

# Theory of Computation

## Quiz 2

November 2, 2023  
Time: 11:50am to 1.00pm  
Total marks: 50

Write clear and precise answers.

(1) Recall that a context-free language  $L \subseteq \Sigma^*$ , for some finite alphabet  $\Sigma$ , is *linear* if it has a context-free grammar  $G = (V, \Sigma, P, S)$  with all productions of the form  $A \rightarrow \alpha B \beta$  or  $A \rightarrow \alpha$  for variables  $A, B \in V$  and terminal strings  $\alpha, \beta \in \Sigma^*$ .

For a pushdown automaton (PDA) a transition  $(p, \gamma) \in \delta(q, a, X)$  is a *pop operation* if  $\gamma = \epsilon$ . The other transitions are *push operations*.

A *one-turn* pushdown automaton (PDA) is a PDA with the property that, for any input, on any computation path once it does a pop operation it will never use push operations in the rest of the computation.

Show that every linear context-free language can be accepted by a one-turn PDA. 15 marks

(2) Suppose  $L$  is a linear context-free language and  $R$  is a regular language. Show that  $L \cap R$  is a linear context-free language. 10 marks

(3) Write the definition of a deterministic PDA. Suppose  $L = L(M)$  for a deterministic PDA. Let  $L' = \{w \in L \mid \text{no proper prefix of } w \text{ is in } L\}$ . Is  $L'$  accepted by a deterministic PDA? Justify answer. 10 marks

(4) Construct a Turing machine  $M$  that enumerates  $\{0, 1\}^*$  in canonical order. Specifically,  $M$  will output  $x_1 \# x_2 \# \dots \# x_i \# x_{i+1} \# \dots$ , where the  $x_i$  appear in canonical order on the output tape. It will use as subroutine a Turing machine  $M'$  that computes  $x_{i+1}$  from  $x_i$ . Give the transition function of  $M'$  in detail, and a high level description for  $M$ . 15 marks

$\alpha \beta \beta$

$\epsilon, 0, 1, 10, 11, 1^*$   
 $00, 01, 10, 11, 000$

$S \rightarrow a_1 B_1$   
 $B_1 \rightarrow a_2 B_2$

$A \rightarrow \alpha$

Reverse  $\rightarrow$

$\alpha \beta \beta$   
 $\alpha \alpha_1 \beta \beta$   
 $\gamma_1$

$S \rightarrow$

$\alpha$   
 $\beta$   
 $\beta$

$\alpha$   
 $\beta$

$\gamma \omega$

$S \rightarrow a \gamma$

$(a_1)$

$(a_0)$