

This entrance exam booklet is being re-used for mid-sem exams.
Please ignore the printed material on the sheets.



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Mid Semester Examination, Aug - Nov 2023

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Date:	Subject: <u>Haskell</u>
Course & Year: <u>MSC 1 2023</u>	Total No. of Pages:

<u>Part B</u>		
No.	Marks	Remarks
11		
12		
13		
14		
15		
16		

<u>Part B (ctd.)</u>		
No.	Marks	Remarks
17*		
18*		
19*		
20*		

Further remarks:

Part A	<u>35</u>
Part B	<u>10+2+20+20 = 52</u>
Total	<u>87</u>

Please indicate on the front page the questions in Part B to be marked.

Solution to Question (II)

1) Reverse $[(-10) \dots 10] \rightarrow [(-10) \dots 10]$

$= [10, 9, \dots, (-10)] \rightarrow [(-10), (-9), \dots, 10]$

$= [10, 9, \dots, -10] \downarrow \quad -10 \dots -9, \dots, 10$

$\therefore \text{filter } (\geq 5) (\quad)$

$= [10, 9, 8, 7, 6, 5, 5, 6, 7, 8, 9, 10] \checkmark$

2) (reverse $[0 \dots 10] \rightarrow [(-10) \dots 10]$)

$= [10, 9, 8, \dots, 0, \dots, -10, \dots, -9, \dots, 10]$

\downarrow

$\text{take while } (\geq -6) (\dots)$

$= [10, 9, 8, \dots, 0 \rightarrow]$

length (\downarrow)

= 11 \checkmark

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Solution to Question (12)

3. fold n (11) [] [6..1], [2..5] , [6..13])

=

.....

([6..13] ++ [])

[6..7, 8 .. 13] ++ []

= 6..7 .. :13(0++[]) 1st time will be called

8 times 2nd time recalled

Similarly, [2..5] ++ [6 .. 13]

will call 2nd time length [2..5] times

= 4

1st will call = 2 times as [length [0, 1]] = 2

So, 8 * 4 + 2 = (14) times ✓

Please indicate on the front-page the questions in Part B to be marked.

Solution to Question (13)

(4) $\text{foldl} (\text{++}) [] [0..1], [2..5], [6..13]$

$$= ((([] \downarrow \text{++} [0..1]) \text{++} [2..5]) \text{++} [6..13])$$

0 calls to 2nd lines 0

Now $[0..1] \text{++} [2..5]$

2 calls to 2nd lines 2

$[0..1, 2..5] \text{++} [6..13]$

6 calls to 2nd lines

$$\frac{+ 6}{= 8}$$

So total 8 calls

Solution to Question (14)

Solution to Question (14)

$$(5) \quad \left[\begin{array}{ccccccccc} (1,0), & (2,0), & (3,0), & (4,0), & (5,0), & (6,0), & (7,0), & (8,0), & (9,0), \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ (2,1), & (3,1), & (4,1) \end{array} \right] \dots]$$

So answer should be (11)

- $\text{foldr} \ (\lambda x) \ [v] \left[\begin{array}{c} v \\ [v] \\ v \end{array} \right]$
 \downarrow \downarrow
 $(\text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool}) \quad \text{Bool} \quad [\text{Bool}] \quad \text{Bool}$

∴ type for expression folds (22) will be

Bool → [Bool] → Bool.

~~g.f~~ should have an input type α .

as $(\lambda x : f \text{ type})$ takes two lists

and

$g.f$ should have an input type say a .

such that $f : a \rightarrow [\text{any}]$

throughout

and so

$g.f : [^a] \rightarrow [\text{any}]$

any will be

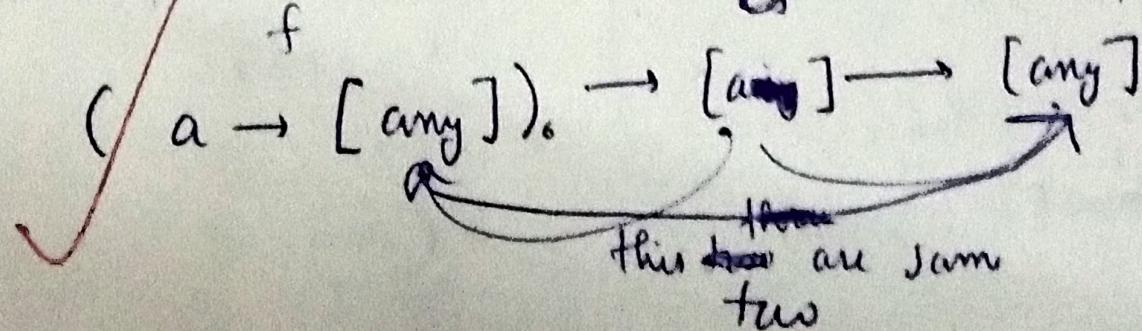
constant

once fixed

: g should have input:

b

output



Part D

Please indicate on the front page the questions in Part D to be marked.

Solution to Question (D)

(1) We will replace $\left(\begin{matrix} \text{fib}(1) & \text{with } 1 \\ \text{fib}(0) & \text{with } 0 \end{matrix} \right)$ whenever we see it
 in the base case.

$$\text{fib}(5) = \text{fib}(3) + \text{fib}(4)$$

$$= (\text{fib}(1) + \text{fib}(2)) + \text{fib}(4)$$

$$= (1 + \text{fib}(1)) + \text{fib}(4)$$

$$= (1 + (1 + (\text{fib}(0) + \text{fib}(1)))) + \text{fib}(4)$$

✓

$$= (1 + 0 + 1) + \text{fib}(4)$$

$$= 2 + \text{fib}(4)$$

$$= 2 + (\text{fib}(2) + \text{fib}(3))$$

$$= 2 + ((\text{fib}(0) + \text{fib}(1)) + \text{fib}(3))$$

✓ (10)

$$= 2 + (1 + (1 + (\text{fib}(0) + \text{fib}(1))))$$

$$= 2 + (1 + (1 + (0 + 1)))$$

$$= 2 + (1 + 1)$$

$$= 2 + 2$$

$$= 2 + 3$$

$$= 5 \quad \text{(ans)}$$

$$= 2 + (1 + (\text{fib}(1) + \text{fib}(2)))$$

$$= 2 + (1 + (\cancel{\text{fib}(1)} + \text{fib}(2)))$$

Head lookup and prepending takes
so, reversing is a good option.

3

~~Implementation~~

Intuition : we will use a helper function

that takes the reverse of the ~~list~~ list

and two empty list as its argument

helper : $[Int] \rightarrow [Int] \rightarrow [[Int]] \rightarrow [[2^{n+1}]]$

\Downarrow
Input reversed
list

\Downarrow
buffer
if we find
a lower number
wht ih Head,

we put this into
the next list of $[2^n]$
and make it null

\Downarrow
stores the
current answer
current
answ

(will proceed like
tail recursion)

Reversing the list will make the logic / algo easier
as everything will go as expected [as head lookup \rightarrow
takes constant time]

Please indicate on the front page the questions in Part B to be marked.

Solution to Question (20)

upRm : [Int] → [[Int]]

upRm $\overset{x}{\bullet}$ = helper (remove x) [] []

helper : [2nt] → [Int] → [[2nt]] → [[2nt]]

helper [] ab = (a:b)

helper (x:x₃) [] b = helper x₃ [x] b

helper (x:x₃) (y:y₃:as) b •

| x > a = helper x₃ • [x] ((a:as):b)

| otherwise = helper x₃ (x:a:as) (b)

Briefly what helper function does on one of the inputs.

upRm [3, 2, 1] = helper [1, 2, 3] [] [] ✓

helper [1, 2, 3] • [] []

- helper [2, 3] [1] []

= helper [3] $\overset{a>1}{[2]} [1]$ []

* helper [] [3] [[2], [1]]

~~= helper [3] [2] [1]~~

Similarly upRm [1, 2, 3] :

helper [3, 2, 1] [] [] = helper [2, 1] [3] []

\Rightarrow helper [1] [2, 3] []

* helper [] [1, 2, 3] []

* [1, 2, 3]

Rough work only. Do not write your final answer here.

getMeSum :: IO()

4. getMeSum = recurse 10 0

recurse :: Int \rightarrow Int \rightarrow IO()

recurse 0 sum = putStrLn ("sum")
(show sum)

recurse n sum = do

cur \leftarrow getLine

recurse
~~return~~ $\frac{(n-1)}{^1}$ $\frac{\text{sum} + (\text{read}(\text{cur})) :: \text{Int}}{20}$

Alternatively we can use instead

Rough work only. Do not write your final answer here.

Q. What say in 'l' we have 3 elements.

So, $(\text{foldl } f \text{ id } l)$



= $\text{foldl } f \text{ id } [a; b; c]$

= ~~$\text{foldl } ((\text{id } f a) f b) f c$~~

= $((((\text{id } f a) f b) f c))$

2. ~~$[\]$~~

reverse l = $((((\text{id } f a) f b) f c)) [\]$

output should be $[c, b, a]$

$f(f(f(f(id, a), b), c), [\])$

'at the last depth it will call recursive'

$f(\text{id}[\]) a = [\] a : [\]$ $(\text{id}[\])$

$f(f(id, a)) b = b : f(id, a)$

$\rightarrow b : [a]$

$\rightarrow [b, a]$

$\therefore f(f(f(id, a), b), c) = [c, b, a]$

$s_0, f: [a] \rightarrow a \rightarrow [a]$

$$f[\]x = [x]$$

~~$f(x:y)$~~

$$f(y:y_s) x = (x:y:y_s)$$

$$f g x l =$$

$$x:(g l)$$

is the correct answer.

~~rough work only. Do not write your final answer here.~~

Part (a) 7

(+) function composition

not is a function that has type

$$\text{Bool} \rightarrow \text{Bool}$$

the resulting function should look like.

$$\text{not}(\cdot) \circ (\cdot) \circ f_2(\cdot) \dots (f_n)$$

Say f_n has arity k_n

$$f_{n-1} \quad " \quad " \quad k_{n-1}$$

f_n will take k_n inputs

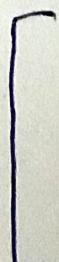
and the return value of f_n will be used in f_{n-1}

So, it will have arity k_{n-1}

and we need to provide it with $k_{n-1}-1$ values
~~the input~~ and other one will be of the
return type f_n

the type expression will be

$$a_k \rightarrow a_2 \rightarrow \dots \rightarrow a_{k-1}$$



$$\left[a_{n_1} \rightarrow a_{n_2} \rightarrow \dots \rightarrow a_{n_{k_{n-1}}} \rightarrow a_{(n-1)} \rightarrow \dots \rightarrow \text{Bool} \right]$$