

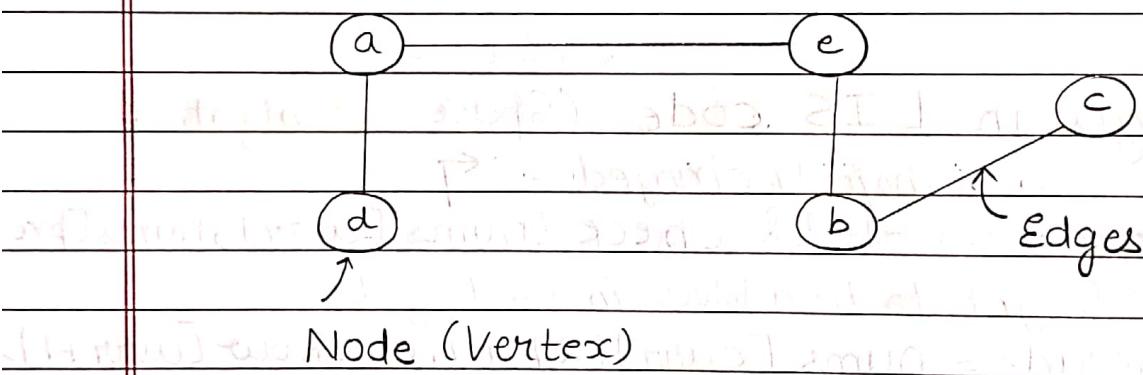
17/06/2023

## Graphs

It is a data structure which is formed by combination of nodes and edges

Node → Stores data

Edge → Connect different nodes



Vertex / nodes → a, b, c, d, e

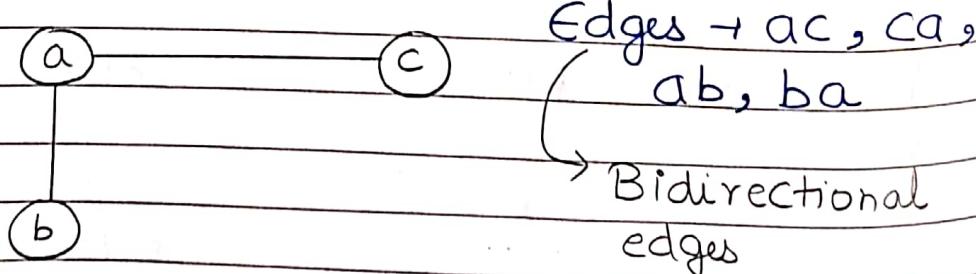
Edges → ae, ad, be, bc

### Applications

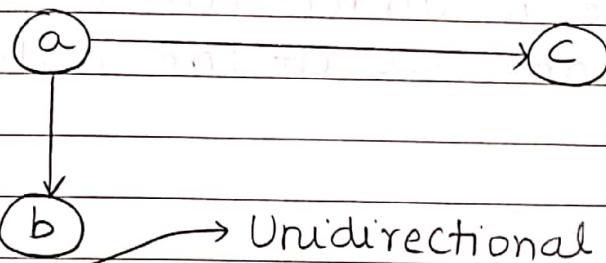
Bus system (route deciding), google maps, facebook etc. → Dijkstra algorithm

### Types of graph

#### 1) Undirected graph



#### 2) Directed graph

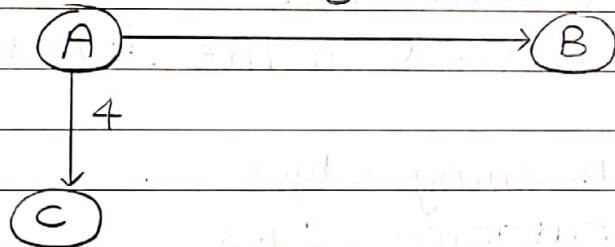


Unidirectional edge

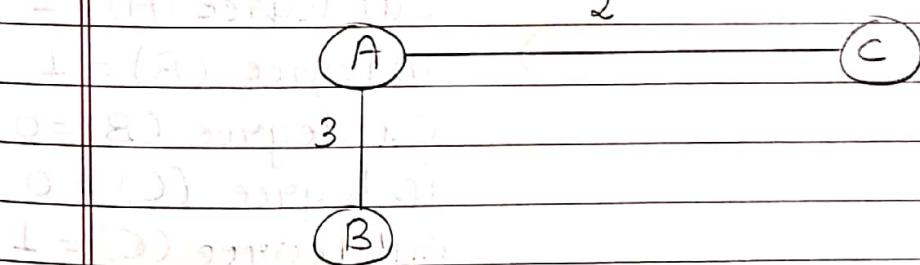
Edges  $\rightarrow ac, ab$

Here ca and ba are not considered as edges.

### 3) Weighted graph



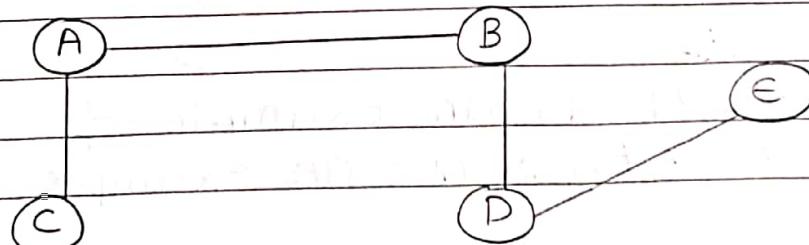
The above graph is an example of weighted directed graph.



The above graph is an example of weighted undirected graph.

Note - If the weights are not given, then we consider 1 as the weight for each edge.

### Degree



Degree of node in undirected graph is the no. of edges connected to the node.

$$\text{degree}(A) = 2$$

$$\text{degree}(B) = 2$$

$$\text{degree}(C) = 1$$

$$\text{degree}(D) = 2$$

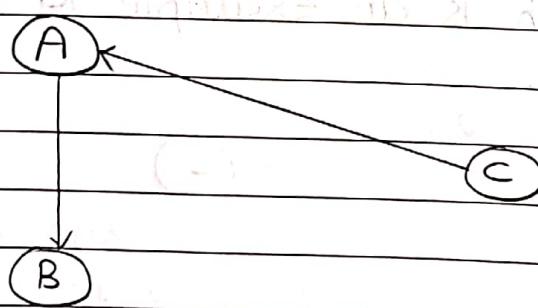
$$\text{degree}(E) = 1$$

Indegree and Outdegree

This concept is used in the directed graph.

Indegree  $\rightarrow$  incoming edges

Outdegree  $\rightarrow$  outgoing edges



$$\text{indegree}(A) = 1$$

$$\text{outdegree}(A) = 2$$

$$\text{indegree}(B) = 1$$

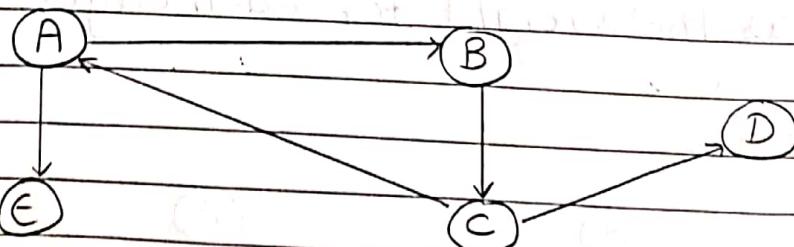
$$\text{outdegree}(B) = 0$$

$$\text{indegree}(C) = 0$$

$$\text{outdegree}(C) = 1$$

Path

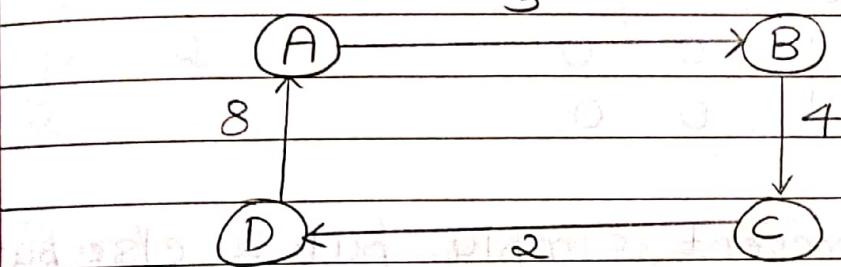
Nodes can not be repeated in a path. Path is defined as sequence of nodes.



$A \rightarrow B \rightarrow C \rightarrow D$  is an example of path

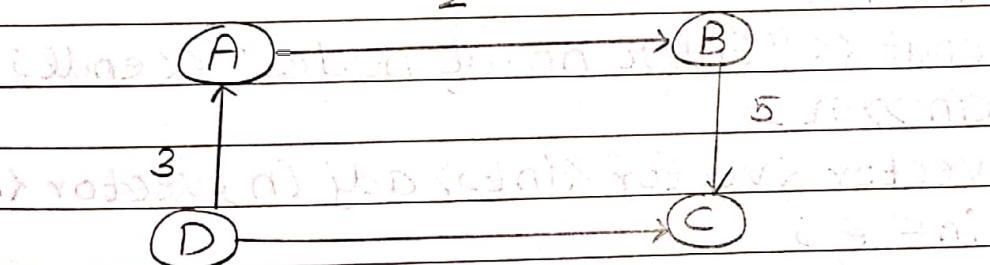
$A \rightarrow B \rightarrow C \rightarrow A$  is not an example of path  
→ repeating

## Cyclic graph



The above graph is an example of weighted cyclic directed graph.

## Acyclic graph

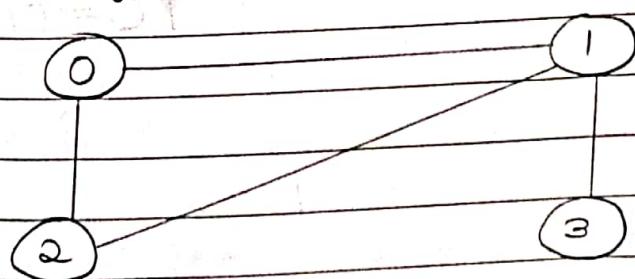


The above graph is an example of weighted acyclic directed graph

## Implementation of graphs

- 1) Adjacency matrix  $\rightarrow$  2D array
- 2) Adjacency list  $\rightarrow$  map/list/vector/Set

### \* Adjacency matrix



No. of nodes = 4

Simply make  $4 \times 4$  matrix

0	1	2	3	
0	0	1	1	0
1	1	0	1	1
2	1	1	0	0
3	0	1	0	0

0	0	1	1	0
1	1	0	1	1
2	1	1	0	0
3	0	1	0	0

If edge is present, simply put 1 else put 0. We simply require the edge list in i/p.

### Code

```

int main () {
    int n ;
    cout << "Enter no. of nodes" << endl ;
    cin >> n ;
    vector <vector <int>> adj (n , vector <int>(n,0));
    int e ;
    cout << "Enter no. of edges " << endl ;
    cin >> e ;
    for (int i = 0 ; i < e ; i++) {
        int u , v ;
        cin >> u >> v ;
        // mark 1
        adj [u][v] = 1 ;
    }
    return 0 ;
}

```

SC =  $O(V^2)$

### Output

3 → nodes (number)

6 → edges (number)

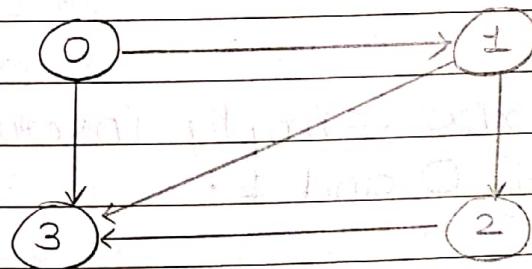
0 1  
1 0

1	2
2	1
0	2
2	0

Print adjacency matrix  $\rightarrow$

$$\begin{matrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{matrix}$$

\* Adjacency list



$$SC = O(V+E) \rightarrow \text{average}$$

$$SC = O(V^2) \rightarrow \text{worst}$$

List of edges  $0 - 1$

$1 - 3$

$1 - 2$

$2 - 3$

$0 - 3$

Adjacency list :-

$0 \rightarrow \{1, 3\}$

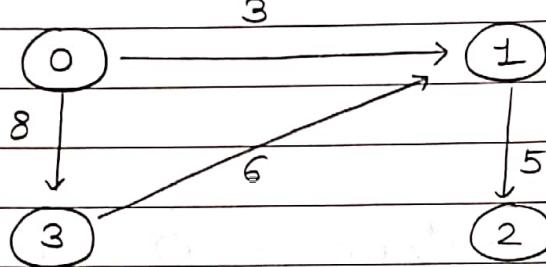
$1 \rightarrow \{2, 3\}$

$2 \rightarrow \{3\}$

$3 \rightarrow \{\}$

Note → Both of the implementation can be used for undirected and directed graph.

## Weighted graph



node ← | → distance

$0 \rightarrow \{(1, 3), (3, 8)\}$

$1 \rightarrow \{(2, 5)\}$

$2 \rightarrow \{3\}$

$3 \rightarrow \{(1, 6)\}$

In adjacency matrix, simply insert the weight instead of 0 and 1.

## Code

```
class Graph {
```

```
public :
```

```
unordered_map <int, list <int>> adjList;
```

```
void addEdge (int u, int v, bool direction) {
```

```
// direction = 0 → undirected
```

```
// direction = 1 → directed
```

```
// Create edge from u to v
```

```
adjList[u].push_back(v);
```

```
// add edge from v to u if undirected  
if (direction == 0)
```

```
adjList[v].push_back(u);
```

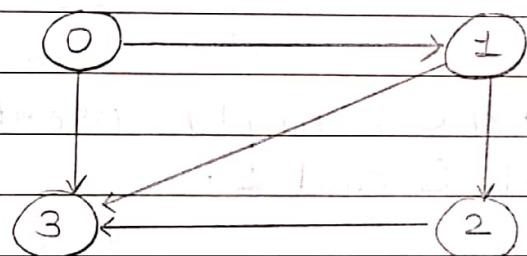
```
void printAdjList () {
```

1	2
2	1
0	2
2	0

Print adjacency matrix  $\rightarrow$

$$\begin{matrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{matrix}$$

\* Adjacency list



$$SC = O(V+E) \rightarrow \text{average}$$

$$SC = O(V^2) \rightarrow \text{worst}$$

List of edges  $0 - 1$

$1 - 3$

$1 - 2$

$2 - 3$

$0 - 3$

Adjacency list :

$$0 \rightarrow \{1, 3\}$$

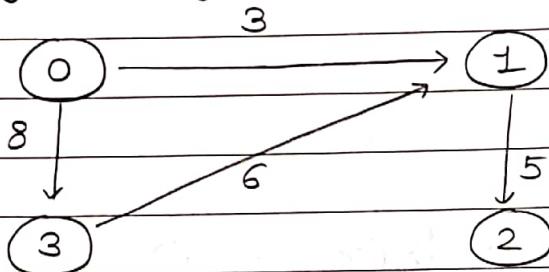
$$1 \rightarrow \{2, 3\}$$

$$2 \rightarrow \{3\}$$

$$3 \rightarrow \{\}$$

Note  $\rightarrow$  Both of the implementation can be used for undirected and directed graph.

## Weighted graph



node ← | → distance

$0 \rightarrow \{(1, 3), (3, 8)\}$

$1 \rightarrow \{(2, 5)\}$

$2 \rightarrow \{\}$

$3 \rightarrow \{(1, 6)\}$

In adjacency matrix, simply insert the weight instead of 0 and 1.

Code

```

class Graph {
public:
    unordered_map<int, list<int>> adjList;
    void addEdge (int u, int v, bool direction) {
        // direction = 0 → undirected
        // direction = 1 → directed
        // Create edge from u to v
        adjList[u].push_back(v);
        // add edge from v to u if undirected
        if (direction == 0)
            adjList[v].push_back(u);
    }
    void printAdjList () {
    }
}
  
```

```

for (auto node : adjList) {
    cout << node.first << " → ";
    for (auto neighbours : node.second) {
        cout << neighbours << ", ";
    }
    cout << endl;
}
}
}

```

Note → In case of weighted graph, changes in code

- 1) unordered\_map <int, list <pair <int, int>>>  
adjList ; Pass weight in addEdge function.
- 2) adjList [u].push-back ({v, weight});  
adjList [v].push-back ({u, weight});
- 3) In print function, inside inner for loop.

```

cout << "(" << neighbours.first << ", " << neighbours.
second << ")" << ", ";

```

Note → In adjacency matrix, instead of marking with 1, simply mark with weight.

Generic graph

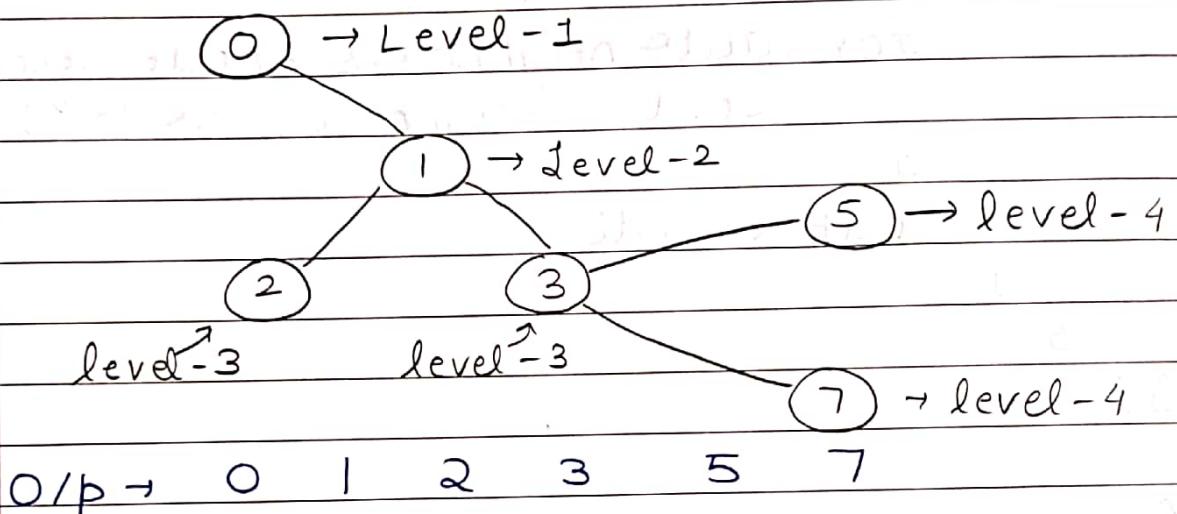
It can be created with the help of template.

template <typename T>

Now replace int with T in the code.

Graph <char> g; → Creation of graph

Breadth first search  $\rightarrow$  Level order traversal



- 1) Simply create edge list and then create the adjacency list.

$0 \rightarrow \{1, 3\}$   
 $1 \rightarrow \{2, 3, 0\}$   
 $2 \rightarrow \{1, 3\}$   
 $3 \rightarrow \{1, 5, 7\}$

$5 \rightarrow \{3\}$   
 $7 \rightarrow \{3\}$

- 2) Make a queue and push the source node in queue.

queue  $\rightarrow \{0\}$

Find front node. Pop it & insert the neighbours.  
Also we need to maintain a visited data

Structure.

$0 \rightarrow F T$   
 $1 \rightarrow F T$   
 $2 \rightarrow F T$   
 $3 \rightarrow F T$

$5 \rightarrow F T$   
 $7 \rightarrow F T$

We have to push the node in queue if it is not visited.

1) O/p → 0

queue → {1 3}

2) O/p → 0 1

queue → {2, 3} (0 not inserted)

3) O/p → 0 1 2

queue → {3} (1 not inserted)

4) O/p → 0 1 2 3

queue → {5, 7} (1 not inserted)

5) O/p → 0 1 2 3 5

queue → {7} (3 not inserted)

6) O/p → 0 1 2 3 5 7

queue → {} → empty & hence stop

## Code

```
void bfs(int src) {
    queue<int> q;
    unordered_map<int, bool> visited;
    q.push(src);
    visited[src] = true;
    while (!q.empty()) {
        // Front node
        int frontNode = q.front();
```

~~q.pop();~~

~~cout << frontNode;~~

~~//insert neighbours~~

~~for (auto neighbours : adjList[frontNode]) {~~

~~//If not visited, then insert~~

~~if (!visited[neighbours]) {~~

~~q.push(neighbours);~~

~~visited[neighbours] = true;~~

3

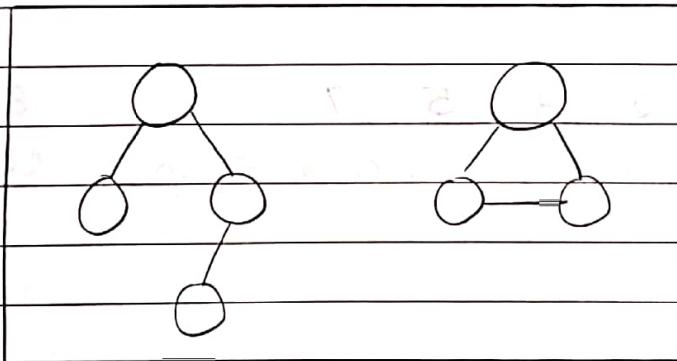
3

3

3

Time complexity =  $O(V + E)$

Components of graph



As no. of components  $> 1$ , this is known as disconnected graph. To handle this run a for loop on number of nodes.

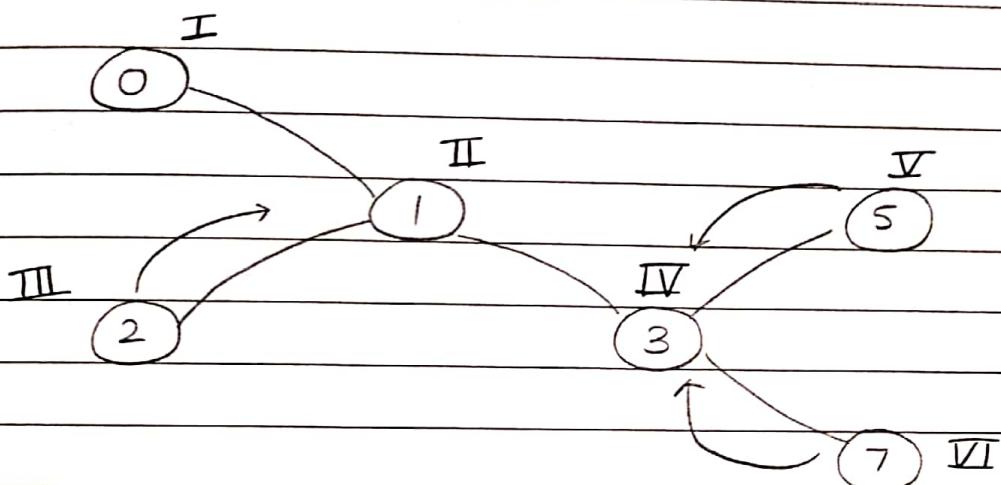
for (int i=0 ; i<=7; i++) {

if (!visited[i])

bfs(i, visited);

3

## Depth first Search



0 1 2 3 5 7

Print source node and then recursive call for neighbours. (if not visited)

### Code

```
void dfs (int src, unordered_map<int, bool>& visited) {
    cout << src << " ";
    visited[src] = true;
    for (auto neighbour : adjList[src]) {
        if (!visited[neighbour])
            dfs(neighbour, visited);
    }
}
```

3  
3  
3

Why no base case required?

Recursive call will go only if not visited but at some point of time all nodes will be visited.

TC = O (V+E)