

SETS

→ implemented using Binary Search trees (BST)

↓ more precisely

Red-black tree

#1. STL container → stores unique elements.

2. Values are stored in ordered state (increasing / decreasing order)
3. No indexing, elements are identified by their own values.
4. Once value is inserted in a set, it cannot be modified.

↑
pehle pura value remove karo,
then new value daalo.

10	20	30
remove		

Advantages:
1. to store unique values.
2. store data in ordered manner.

3. Dynamic Size.
4. Faster → insertion
deletion
searching
↓
binary search } $O(\log N)$ time complexity

Disadvantages:
1. Cannot access elements using indexing.

2. consumes more memory.

3. not suitable for large data-sets.

↳ insertion/deletion : $O(\log N)$

declaration : #include <set>

set <data-type> set-name;

e.g. → set <int> set2;

initialise → set <int> set2 = {1, 3, 2, 4};

by default, values are stored in increasing order

$\therefore \text{Set } 2 = \{1, 2, 3, 4\}$ // done automatically.

to store values in a set in decreasing order:

```
set <int, greater<int>> set2;
```

insertion in set ⇒

```
Set1.insert (4);  
Set1.insert (1);
```

} [1, 4]
sets

this funcⁿ returns an
iterator to the inserted
pointer, ← value.

Traversal of a Set

using Iterator → `set_name.begin()`; → iterator pointing to 1st element of my set.
→ `set_name.end()`; → iterator pointing to the posⁿ after the last element of my set.

Set A: [1, 2, 3, 4]

`SetA.begin()` `SetA.end()`

```
set<int> set1 = { 4, 1, 5, 3};  
Set1<int>::iterator itr; // or auto itr = set1.begin();
```

```
for (itr = sets.begin(); itr != sets.end(); itr++) {
```

```
cout << *itr << " " ;
```

3

Op → 1 3 4 5

~~deleting~~ → 1. `set1.erase(3);` } $\xrightarrow{O(\log N)}$
 2. `set1.erase(itr);` } `set_name.erase(value / pos^n);`
 iterator.
 3. `set_name.erase(start pos, end pos);`
 deletes elements from $start pos^n$ including it till
 $end pos^n$, excluding it. i.e. $[start, end)$

Eg → set → $[1, \cancel{2}, 3, 4, 5]$ $\xrightarrow[s.erase(s, e)]{O(N)}$ $[1, 4, 5]$
 ↓
 $N = \text{no. of elements in the range.}$

to move iterator by k steps → `advance(iterator_name, k);`

MEMBER FUNCTIONS:

- `size()`, `max_size()` → returns max. no of elements, set can hold.
- `empty()` → returns true if empty else false.
- `clear()` → removes all elements from set.
- `find()` → returns pos. of element if present, else returns end iterator.
- `count()` → returns no. of occurrences of an element.
 agar koi value hogi → return 1
 gf not present in set → return 0
- `lower_bound` → returns element if present, else returns just greater value.

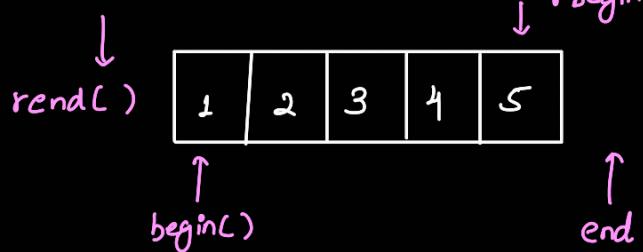
Eg → $s = [10, 20, 30, 40]$
 $s.lower_bound(25); \Rightarrow$ returns 30.

→ `upper_bound` → returns next greater value.

eg → `s.upper_bound(20);` \Rightarrow returns 30.

→ `rbegin()` → returns iterator to first element of set in reverse order.

→ `rend()` → returns iterator to posⁿ after last element in reverse order.



Unordered Set : → values are stored in unordered fashion.
→ rest same as sets.

implemented
using
hashing

advantage over ordered set :

insertion } avg. time complexity:
deletion } $O(1)$ as compared to $O(\log N)$ in ordered set.
search }

1. Insert / → single element $\rightarrow O(1)$ avg.
deletion / find / count $\rightarrow O(N)$ worst → rehashing
↳ when size > capacity.

avg. case:

multiple elements: $O(n)$ → n is no. of elements being inserted.

worst case: $O(n * (N+1))$
↳ size of unordered set. (rehashing)

hashing: load factor

`load_factor()` → $\frac{\text{size of unordered set}}{\text{bucket - count}}$

\Rightarrow `rehash(x)` → sets the no. of buckets to x or more.
($O(N)$ avg.)
 $= O(1^2) \text{ worst}$

($O(N)$ worst)
Multiset: It can store duplicate values. (ordered manner)
↓
rest, same as ordered set.

Member functions:

1. $\text{insert}()$ → $O(\log N)$
2. Deletion → $\text{erase}()$ → $\text{erase}(\text{value}) \rightarrow O(\log N) \rightarrow$ deletes all instances of that value.
 $\text{erase}(\text{position}) \rightarrow$ deletes the value at that iterator/
 $\text{erase}(\text{start}, \text{end}) \rightarrow O(n) \rightarrow$ range.
3. $\text{find}() \rightarrow$ lower bound of element if found else end iterator.
4. $\text{count}() \rightarrow$ returns no. of occurrences. $O(K + \log N)$
↓
no. of occ. of the element
5. $\text{upper-bound}()$ and $\text{lower-bound}() \rightarrow$ same as $O(\log N)$ set.
6. Unordered multiset: allows duplicate values to be stored in unordered manner.
↳ member functions same as unordered-set and similar time complexities.

