

29/7/2019

Programming Fundamentals & OOP concepts

Not portable but easier.
⇒ Assembly Level prog.

SKS Sir

Pseudo - code

(Alphanumeric)

Pneumonics .

Low Level Language

- Error prone of Machine Language ⇒ 0s & 1s.
- Debugging is difficult.
- Machine Dependent (Not portable).
- Tedium

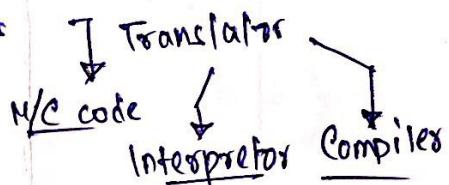
Assembly Level

Assembler ↔ Translator
Machine Lang:

Source code can be copied to any machine
High Level Language object code → Machine-dependent

↓ → Portable
FORTRAN, BASIC, C, ... etc ⇒ English like.

GRAMMAR ⇒ SET OF RULES



Source Program

↓
Translator
Object code

Interpreter → Takes a line at a time from source code and translates it & executes it; stops if no line is left or encounters any error.

Error

→ Rectify & start again from first line.

Compiler → Translates the whole program in a single shot. If no error, creates object code else displays a list of errors.

Programming Approach

②

Paradigm

System is developed

Modity

Changes to be made

1. Data
2. Algorithm

→ Organization of the data

→ Finding efficient algorithm to work on data

Problem

Sub-tasks

→ Write the procedure/fn
for the sub-problem

↓
to focus
on the algorithm

MODULAR PROGRAM

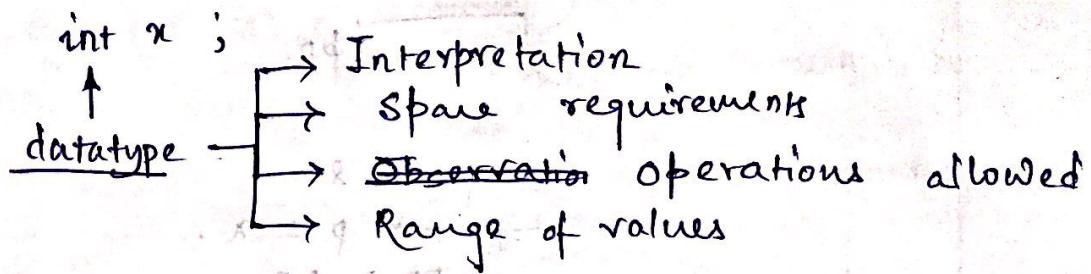
→ RELATED PROCEDURES and data on which they work

→ Organization of data is ignored are put together to form a module.

Data Abstraction

Access control

Object Oriented Approach



int $x = 65$;
 ↓
 binary

char $C = 'A'$; → Stream of 0/1
 ↓ depends on
 ASCII value → binary datatype.

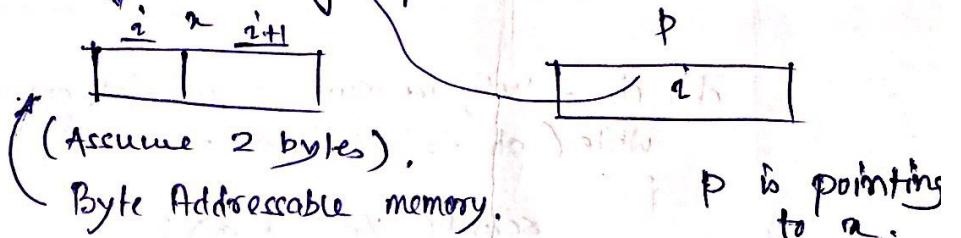
COBOL → Common Business
 Oriented Language.

Space allocation is same.

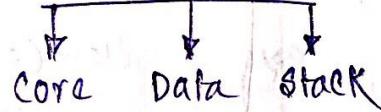
can hold the address of another integer variable.

{ int x ;
 int $\star p$;

When we don't know the value is garbage



Segments



Segmentation Fault

Accessing memory not allocated.

$p = \&x$;

↑
 address of
 x .

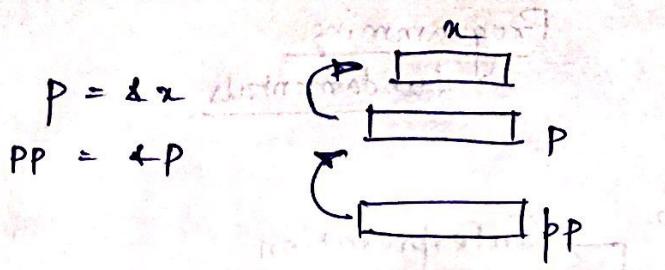
$\star p$ = (content of location to which p points).

$\star p = x$

scanf("%d", & x);

if & is not provided, it will not give any error.

int x ; int $\star p$;
 ↓
 type of variable to which p will point.
 int $\star pp$;
 ↓
 type of variable to which pp will point.
 $pp = \&p$.



②

$$\begin{aligned} \Rightarrow P &\equiv \&x \\ \cancel{\Rightarrow \&P} &\equiv \&x \\ \Rightarrow PP &\equiv \&P \\ \Rightarrow \&PP &\equiv \& \equiv \&x \\ \Rightarrow \&(\&PP) &\equiv \&P \equiv \&x. \end{aligned}$$

~~same~~ Arg of n. numbers
while loop.

char ch = 'y'; int sum = 0; int cnt = 0;
while (ch == 'y' || ch == 'Y').

```
{
    scanf("y.d", &ch);
    sum = sum + a;
    cnt = cnt + 1;
}
```

~~double a~~ ^{float} arg = float(sum)/float(cnt);
(either one).

Mixed type

~~More General type~~

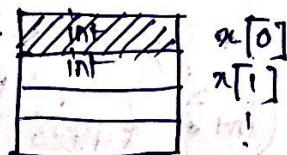
→ Elements in a single variable & be able to
use them → variable name

int x [100];
type size.
of the elements → integer pointer.

#define LENGTH 100
↑
Preprocessor
directive.

Collection of homogeneous elements → A block of memory of contiguous space will be allocated to hold 100 integers.

Array name → m →
points to 1st element



char s1[30], s2[80];

scanf("%s", s1);

strcpy(s2, s1);

$$\boxed{s2 = s1}$$

\Downarrow
assigning one pointer
to another

however they are constant
pointers

(int) const *x;

\downarrow
type to which the constant
pointer is pointing to.

for (i=0; i<n; i++)

{

scanf("%d", &x[i]);

}

for (i=0; i<n; i++)

{

printf("%d", x[i]);

}

$$x \rightarrow x[0]$$

$$x+i \rightarrow x[i]$$

x \rightarrow x[0].
p \Rightarrow pointer.

p+1 \rightarrow 500 \rightarrow 502.
int pointer.

char pointer \rightarrow 500 \rightarrow 501

p+k \rightarrow pointer to kth variable.

\Downarrow
p + (k + size of T).

$\Rightarrow \&x[i]$
can be written as
n+i

$$(n+i) \equiv x[i]$$

$$x[i] \equiv *(n+i),$$

T*p

④

Max 1000 integers

#define LENGTH 1000

Mostly \Rightarrow working ~~with 100 numbers~~ with 100 numbers.

int x[LENGTH];

leads to space wastage.

Dynamic Memory Allocation

n \rightarrow numbers.

int *x;

int *x;

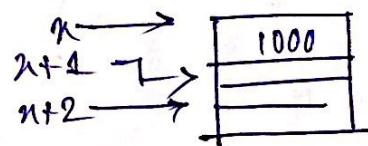


$x = (\text{int} *) \text{malloc}(\text{sizeof}(\text{int}) * n);$

(void *)

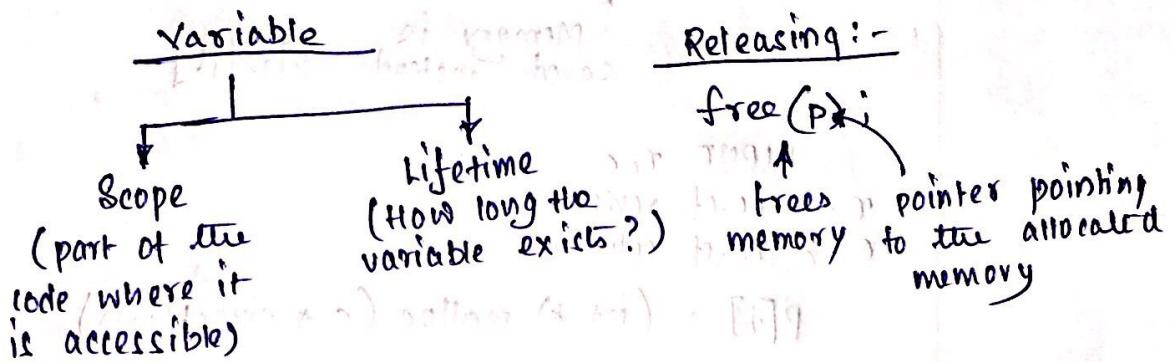
No. of bytes we need.

Try to allocate specified amount of memory, successful or null.



1D Array & Pointer

`int *p; p = malloc(...);`
~~p == m~~ \Downarrow
 dynamically allocated memory has to be released once it's not required



2-D Array

`int n[10][5];`

↑ Rows ↓ Column

In memory

Stored - 1D

$n[0][0], n[0][1], \dots, n[1][0], \dots$

2D Array

A one dimensional array where elements are again 1D array.

$x \rightarrow$ 1st element

$\uparrow \Rightarrow x \rightarrow x[0] \rightarrow x[0][0]$

`int **(a[0])`

$x+1 \rightarrow x[1]$

$x[0] + 1 \rightarrow x[0][1]$

$x[0][0]$

$* (x+i) = x[i]$

General form

$* (x+i) + j = x[i][j]$

$* [*(x+i)+j] = x[i][j]$

$$*(x+i) = x[i][0]$$

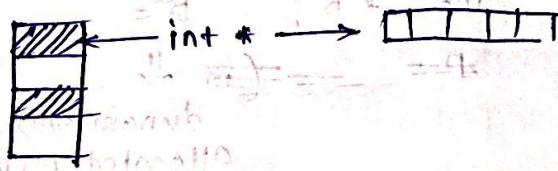
$$*(x[i]+j) = x[i][j]$$

size specified at Execution

No. of Rows \rightarrow (known)

~~No. of columns~~

We need to have an array of integer pointers



int * p[10];

{Memory is saved instead} $p[10][5]$

INPUT r, c

r \rightarrow no. of rows

c \rightarrow no. of columns

$p[i] = (\text{int} *) \text{malloc}(c * \text{sizeof}(\text{int}))$.

int ** p;

INPUT r, c;

$p = (\text{int} **)\text{malloc}(r * \text{sizeof}(\text{int}))$;

for (int i = 0; i < r; i++)

$p[i] = (\text{int} *) \text{malloc}(c * \text{sizeof}(\text{int}))$;

Freesing up memory

for (int i = 0; i < r; i++)

$\text{free}(p[i])$;

$\text{free}(p)$;

$E_1 \leftarrow E_0$

$E_2[0] \leftarrow E_1 \cap E_0$

$[0][0] \leftarrow E_0$

$[0][1] \leftarrow E_1$

$(E_1 \wedge E_2) \wedge$

not maxmin

filled \rightarrow closed

$E_2[1][2] = (E_1 \wedge E_2) \wedge$

$E_2[1][3] = (E_1 \wedge E_2) \wedge$

PTR to ARRAY

`int x[10];`

`int (*p)[10];`

p is ~~an~~ a pointer
to an array of 10 integers.

$$p = \&x$$

$$p \rightarrow \boxed{1000}$$

$$\begin{array}{l} \text{int } * t = x \\ t \rightarrow 1000 \\ \boxed{P++} \rightarrow 1020. \end{array}$$

`int * p[5]`

$t \Rightarrow \text{rows}$

$p = (\text{int}**) \text{malloc} (5 * \text{size of (int)})$

$p + i \rightarrow i^{\text{th}} \text{ row}$.

$\Rightarrow p + 1 \rightarrow 1^{\text{st}} \text{ row}$

function

~~large task~~ \Rightarrow smaller tasks.

calling function

Present
in
stack

main()

{

:

}

called fn.

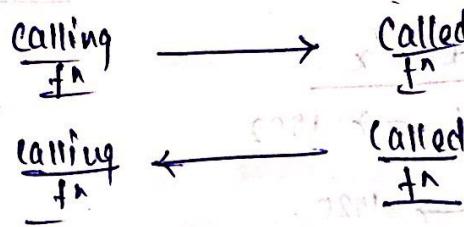
1. When a fn. is called, execution of a calling fn. is temporarily suspended.

2. Control of execution goes to called fn. Once execution is over, the control comes back to the calling fn.

④

Function
creates more stress on the system, however modularity is preferred more.

Call by value & call by Reference



int main()

```
{  
    int a, b;  
    int area, perimeter;  
    calc-parameter(a, b);  
}
```

void calc-parameter
(int a, int b)

→ Use of global variables is not advised

parameter passing

by value by reference/address

calc-parameter (a, b, &area, &perimeter).

$$* \text{area} = a * b$$

$$* \text{perimeter} = (a+b) * 2$$

(5)

char *c = "abcd" ;

f(c)

....

void f(char *p)

variable where change is to be reflected, send its address.

07/08/2019

Programming Practice

SKS
Sir

①

Passing array to function

1D Array

```
#define SIZE -
```

```
int x[SIZE];
```

```
int n;
```

```
INPUT n;
```

```
getnumbers(x, n);
```

prototype void getnumbers(int [], int); → for(...)
 or int *, int
 void printnumbers(*x, n); scanf(...)
&p[i] OR (p+i).

2D Array

```
int mat[row][col];
```

void getnumbers(int **p
 int p[][], int r, int c)
 \uparrow
 $p+i$
 gives ith row.

Recursion

Function calling itself.

```
main()
```

```
{
```

```
f(..)
```

```
}
```

```
int k;
```

```
k = f(0);
```

```
^
```

Not an
error.

void f(int),

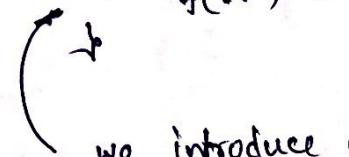
k = f(0);

Error

~~infinitely many calls~~

f(...)
|
f(...)

Never ending.



we introduce a condn

if (...) = True.

f(...)

if (...) = False

break;

else

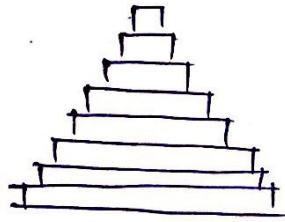
do

...

end

Stack

Recursion may have memory overhead. [Stack overflow]



~~int fact(int n)~~

{ if (n > 1)

return n * fact(n-1);

else
return 1;

Pts to function

(1) int * p; int pointer

int (*p[]); array of ptr.

int (*p)[]; ptr to array

int *f(); fn ret. pointer

int (*f)(); ptr. to fn that returns int.

`int (*pf)(int, char);`

`int f(int n, char c)`

{

:

: return(...);

}

Alt.

`pf = f; pf = &f;`

`*pf = (10, 'A')`

Calling a fn

but here we can skip it.

'&' & '*' can be dropped.

ptr to fn \Rightarrow has the address of

\Downarrow the corre. code of

the fn that it points.

malloc/dealloc

cannot be done.

processing (...)

{

op. . .

f(1) . . .

f(2) . . .

: . . .

f(n) . . .

{ ith op.

f1() . . . | pf1()

f2() . . . | pf2()

`int (*pt[])(int, int)`

`pt[0] = add; add;`

:

`pt[3] = divide.`

Array to function
pointer having 2 integers
as arguments

④

Array

⇒ Collection of homogeneous elements

Structure

not homogeneous.

Collection of diff. element - which may not be of same type.

struct <tag>

{

int roll;

char name[31];

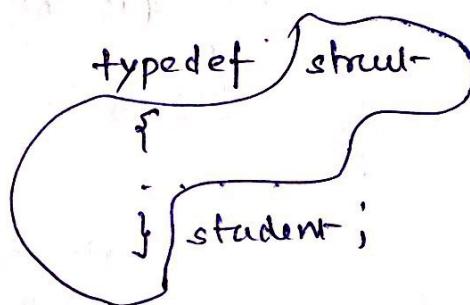
int marks[5];

}

[int(2) char(31) int(10)]

$$31 + 10 = 43$$

struct <tag> x;
struct <tag> y;



typedef oldtype newtype;

scanf("%d",
& a.roll).

$$p = \&a.$$

a.name, &a.marks[i]

$$p \rightarrow \text{roll}$$

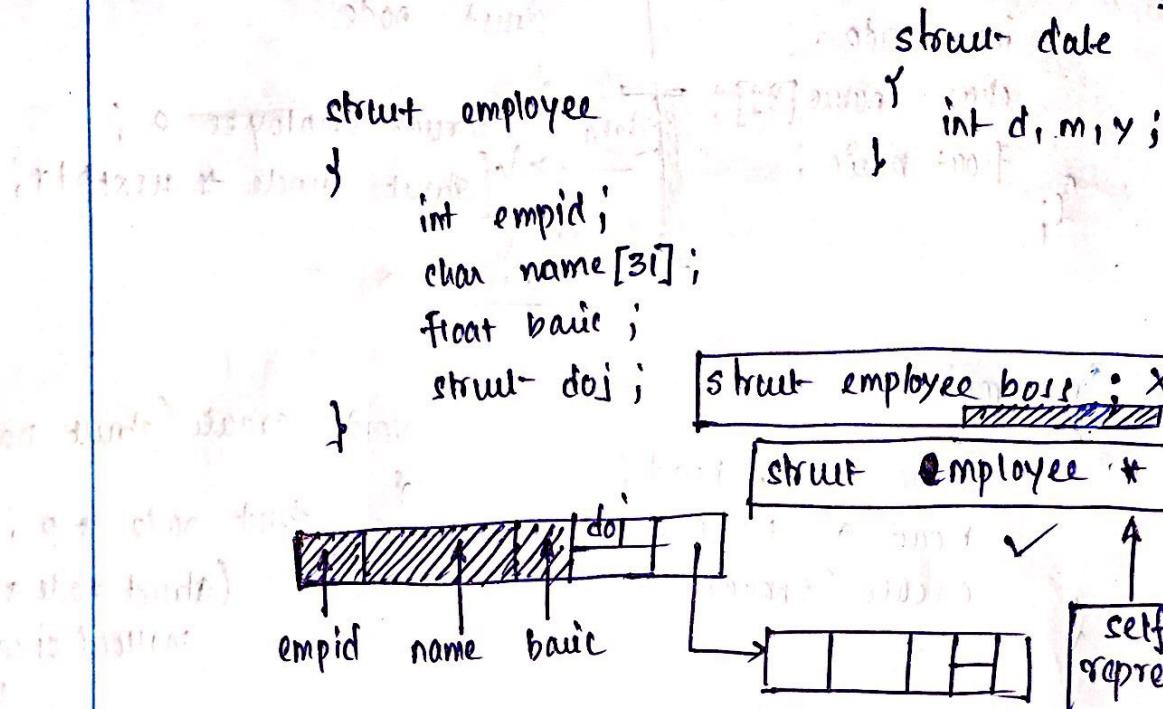
$$p \rightarrow \text{name}$$

$$\&(p \rightarrow \text{marks}[i]).$$

08/8/2019

Programming Practice

Page No. 1
Sanjay Kumar Saha



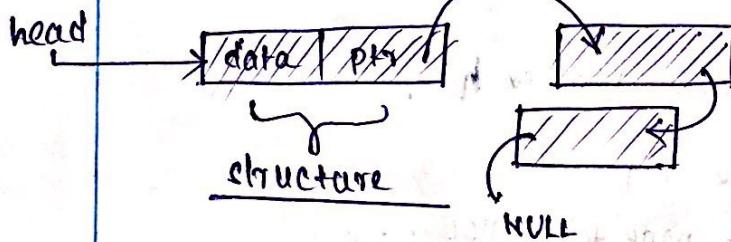
Contiguous Memory Allocation

→ Random Access possible

→ Insertion/Deletion not possible.

Linked List

No contiguous space for all elements required.

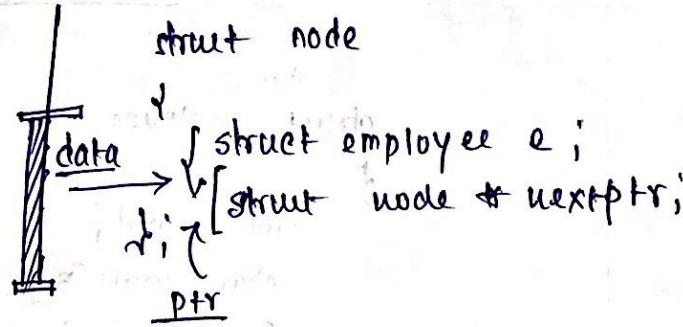


- Insertion & Deletion easier wrt array.
- No predetermined size.
- Sequential access.
- Overhead:- In each node, a pointer to next node is required.

struct employee

{

```
int encode ;  
char name[31] ;  
float basic ;  
};
```



int main()

{

```
struct node* head ;  
head = NULL ;  
create (&head) ;
```

void create (struct node **p)

{

```
struct node *p ;  
p = (struct node *)
```

```
malloc (sizeof(struct  
node)) ;
```

p → nextptr = NULL ;

void get_employee_data (struct employee * ep)
{
 ...
}

get_employee_data (&(p → e))

*hp = p

void append (struct node **h) .

{

```
struct node *p ;  
p = (struct node *) malloc ...  
p → nextptr = NULL ;
```

~~get_employee_data (&(p → e)) ;~~

while (h → nextptr != NULL)

{ h = h → nextptr ;

}

h → nextptr = p .

}

~~void~~ display-employee-data($h \rightarrow e$)

~~void~~ ~~h =~~

~~void~~ ~~search (head)~~

search(head, k);

if ($(h \rightarrow e).ecode == k$)

{
...
}.

if ($h == \text{NULL}$)

{
...
}.

printf("Not found");

~~void~~ insert-before(&head, k)

void insert-before(struc-node **hp, int p)

{

struc-node *p, *prev, *cur;

cur = *hp;

if ($(cur \rightarrow e).ecode == k$)

{
...
}.

p = ... ; get ...;

p \rightarrow nextptr = ~~cur~~; cur = ~~nextptr~~;

~~head~~ \rightarrow ~~nextptr~~ = *hp = p;

else

{

curr = (*hp) \rightarrow nextptr;

prev = (*hp);

while (curr != NULL)

{

if ($(curr \rightarrow e).ecode == k$)

{
break;
}.

else prev = curr;

curr = curr \rightarrow nextptr;

```

if (curr == NULL)
{
    printf("Not found");
}
else
{
    p = curr;
    prev curr → nextptr = p;
    p → nextptr = curr;
}

```

```

void delete(strut_node *hp, int k).
{
    strut_node *p, *prev, *curr;
    curr = hp;
    if ((curr → e).eode == k)
    {
        *hp = (*hp) → nextptr;
        free(curr);
    }
    else
    {
        prev → nextptr = curr → nextptr;
        free(curr);
    }
}
```

File Handling

Array \Rightarrow Large amount of data (IN Primary Memory)

Once program is over, everything is lost.

File \Rightarrow Permanent storage of Data.

Data Structure (only)
with Secondary memory.

To work with a file

In C
 \downarrow
 FILE *fp;
opening mode
 {
 \Rightarrow open a file \rightarrow store data
 \Rightarrow work with file \rightarrow Access data.
 \Rightarrow close the file
 special type of structure (defined in stdio.h)
 store info. regarding the file
 W \Rightarrow write
 R \Rightarrow read
 Pts to the buffer

fopen

open the
file in

desired mode.

If successful,

return file

pointer. If it fails,

return NULL.

if ($fp == \text{NULL}$)
 {
 printf("ERROR");
 exit(-1);

FILE

Binary

* Data elements are
represented by Ascii
value. Can be opened using any editor.

fprintf(fileptr, ...)

②

fscanf()

Buffer reads & stores blocks.

→ inode contains info about ~~file~~ block
out is maintained by the operating system.

fprintf(..)

fscanf(..)

• ASCII aka formatted file.

• Variable length of record ⇒ Maintenance of such file is difficult (Issue arises when jump is happening)
value of an attribute may get changed.

store it in a variable

struct rec

{

int roll;

char name[32];

int scope;

};

fwrite(&s, sizeof(s), 1, fp);

struct rec a[100];

fwrite(a+i, sizeof(s), n, fp)

w → if the file is existing
the entire data will be lost.

To append,
we have 'a' mode.

(append) → added at the end.

Reading data from the file

→ 'r' mode
and the file must exist.

```
struct rec s;
fread(&s, sizeof(s), 1, fp);
```

→ after reading also shifts the file markers

feof(fileptr)

if marker encounters end of file.
without feof
implementation
 while(!feof(...))

fclose(file).

fseek → sets the marker at a posn in file.

fseek(fileptr, offset, ref) [offset (non-negative)]
 ↓
 long integer | → 0, wrt start of the file
 | → 1, wrt current posn
 | → 2, wrt end of file [offset (non-positive)]

fseek(fp, 0, 0) → start of the file

ftell(fp) → specify the current byte posn.

To get kth record

start = (k-1) * sizeof(~~charing point~~)

↓
 struct rec[0]

Modify Records

$\boxed{r+}$ $\boxed{r++}$

\rightarrow old thing
will be lost

modifying ~~roll no.~~ roll no. 10.

$fp = fopen(..., "r+\0")$

```
while( !fread( ... ) )
{
    check. if( ... ) check = 0;
    {
        check = 1;
        break;
    }
    ...
}

if ( check == 0 )
    printf( .... )
```

To modify

```
fseek( fp, -1 * size of( struct rec ), 1 );
...
fwrite( &s, size of( struct rec ), 1, fp );
```

Sorted - order

\rightarrow use dummy records.

Deleting

\rightarrow Take the help of a temporary file.

Logical deletion, Physical deletion

\Downarrow
Assigning dummy
to file to be deleted.

20/8/2019

Object-Oriented

Programming

Prof.

Satyay

Kumar

Lata

①

Low Level

↓

High Level

→ Procedural Approach

→ Modular Approach

→ Object-Oriented Approach

Emphasis is on
~~programming~~ algorithm/
organisation of data
is ignored

To some
extent data org.
is considered
Related Proc. &
data \Rightarrow module

Language which
supports object-oriented
features \Rightarrow object-oriented
prog. lang.

OOP

Object \Rightarrow corresponds
to real world entity
or concept

Description
& Behaviour \leftarrow Similar type of object

Group / Set

Class is the
collection of similar
objects. Class defn
provides the description &
behaviour of the objects,
belonging to the type

Class Student

{

Data Member / Attribute

Roll, Name, Phone-No., dob

Description

admission()

change-details()

Behaviour / Method /

Member fn.

Data Abstraction

Hiding the complexity
of internal representation

student x

object is
the instance
of the class

⑧

As end user, implementation details are not sufficient required to know the methods one can utilize.

→ Polymorphism

Same method can have different implementation in different places.

Implementation to be followed is based on target object.

Logically same operation:

⇒ Actual implementation differs.

→ Inheritance → helps in easy maintenance

Sharing of attributes & Behaviours

Makes the code reusable.

Class (I)

~~f1# f2#~~

f1()

f12()

Class (II)

↓

f2()

f22()

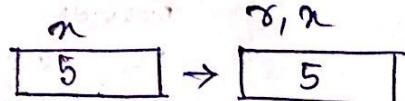
its lifetime
(as per scope).

Its reference variable throughout

In C++, still refers to some variable

Reference Variable

int n; f think as an
int &r = n; alias.



→ r is a reference variable, and it refers to integer variable

C++

Whenever a reference variable is declared, it must refer to an existing variable. No reference space is allocated, or if it is another name for the variable it refers.

2r, 2x ⇒ Same
No separate
memory is allocated

main() ① $n = 5$.
 (After) $n = 5$ ② $n = 5$
 f(n); After $x = 10$

void f(int p) → param by value.

{
 p = p + 5;
 }
 ↓

→ by reference

void f(int &p)
 {
 p = p + 5;

#include <stdio.h> → #include <iostream.h>

To display.

cout << " ";
 cout << "\n";

cout << "Hello World" << "\n";

↑ ↑
 insertion operator concatenation

Reference to ostream class.

cout << "Value is " << a << "\n";

cin >> a;

in put extraction
 operator.

main()

int a, b;

cin >> a;

cin >> b;

int &c = max(a, b);

must have reference

↑ ↓

int &c = max(int a, int b)

do certain operations
 affecting a.

int &c = max(a, b)

↑
 get the reference.

int *c = max(a, b)

↑
 get the value

→ #define max(x,y) ((x)>(y)) ? (x) : (y)

Preprocessor
directive

prior to compilation
macro substitution takes place

No fn call
hence no push/pop
in the stack present

→ There is no type checking in macro
substitution, may create problem
during run time.

~~inline fn~~ inline int max (int a, int b)
request to the compiler to make the fn inline. If the compiler grants, every call to the fn will be replaced by equivalent code during compilation.

{
if (a>b)
return a;
else
return b;

Simple fn \Rightarrow inline.

Compile time \uparrow
Run time \Rightarrow fast -

functions with default parameter(s).

calc-percentage (int marks, int total = 100);

float calc-percentage (int s, int fm)

return ~~(s + fm)~~, ~~(s + fm)/100~~
~~(s + 100.0)/fm;~~

→ Only the rightmost subset of arguments can have default values.

Note

function overloading

A function with same name with different implementation.

Each defn must differ in signature.

No. of ~~types~~ arguments and/or addressing type of arguments

```
void f(int a)
{
}
void f(char c)
```

```
main()
{
    char x;
    int p, float z;
    f(p); f(x); f(z);
}
```

converted
to int

Rules

→ exact match / match after trivial conversion

T	T&
T&	T
T	count T

f(int a, int b)

→ Promotion (char to int), (short int to int),
(float to double)

→ Standard conversion

→ User defined conversion

→ Ellipsis (...)

f(...)

{ }

All rules
are followed
then relaxation
and other is done.

⑥

n arguments in the fn call

for ($i = 0; i < n; i++$)

$s[i] = 'A'$

for ($i = 0; i < n; i++$)

value of $s[i]$ find match with i th

argument

Apply the
Rule - (1) for
all defn.

$s \cap s[i]$

taking the intersection

if no match
{

if ($\#ps$) $\neq 1$.

Apply Rule - (2)

~~repeat~~ Report
ambiguity

it no match
{

else

the defn in s
is executed.

Rule - (3)

$+Cint a, int b)$

Rule (4)

$f(char c, char d)$

Rule (5)

$f(int, float)$

int i ; char n ;

$+ (i, n)$,

[check for
ellipsis].

$s_0 = + (int, float),$
 $+ (int, c)$

$s_1 = + (char, char),$

$s_0 \cap s_1 = \text{NULL}$

structure Student

void Student::get_info() \rightarrow scope operator

{ int roll;

char name[31];

void get_info();

void display();

{

}

cin >> roll;
cin >> name;

Scanned with CamScanner

Student s;
s. get - info();
s. display();

→ However

structures have no
area control

class Student

```
int roll;
char name[31];
public:
void get - info()
{
    :
    inline void display(void)
    {
        :
    }
}
```

inline

↑
Request
to make
it inline
to the
compiler.

for each
object space
allocated

void student:: display(void)

```
{ :
}
```

int main (void)

```
{
    Student s;
    Student t;
}
```

→ Each object
will have its
own copy of data
members.

→ All objects of
the class share
single copy of
member fn code.

s. get - info()

invoking object member

A member
fn is always
invoked by an
object.

(8)

Reference of invoking object -
automatically paired.

public

available
outside the
class also.

private

can be accessed
using member
functions.

→ By default (if nothing is mentioned) \Rightarrow ~~public~~ private

Attributes & methods that
can work with the data
members are bundled
together \rightarrow ENCAPSULATION

An user of a class should know
members only, for public members

Ref into \Rightarrow Rest hidden

public
 \Rightarrow prototype

~~# Another method~~ | int compare (Student x)
| if (scope == x . score)
| return 0 ;
| if (score > x . score)
| return 1 ;
| else
| return -1 ;

student * sp = a . higherScore (t),

student & higherScore (student t);

(Student &) student :: higherScore (Student & t),

if (t . score) \neq (& . score)
if (t . score) > (& . score)
return + ; | else
| * this ;

How many students have taken admission during the session of the program?

:

int[i].admissions;
i++;

Suppose we use static int cnt;

for a class there is only one copy of static data members, all objects share this.

Student s;

Each object will have

their own copy of data members.

→ Declare the static data member outside the class.

~~int student::count = 0;~~

→ int student::cnt = 0;

static void show_count(),

cout << count << "In";

becomes
a part of the
class
(can be invoked
without
the object).

Student :: static-fn(...)

No non static
elements can be referred.
But one can declare local
objects and can work with that.

→ What is class?

→ Encapsulation

→ Abstraction.

→ How to define a class?

→ Data members/Attributes

→ Member fn

↴ ↓
can be can be defined
defined inside outside the class
the class (scope resolution
defn operator).

→ Any member fn has to be invoked by the object of the class

→ Reference of invoking object is automatically passed.

→ Access control → (i) Private → Can be accessed only ~~with~~ in member fn
 (ii) Public → Everywhere ~~with~~ member fn

data members → Private
 & member function → Public.

→ If data members become public, then there would be a compromise in security.

→ Data members → public
 [NOT ADVISABLE]

No. of methods have certain common attributes activity ⇒ repeat in all the methods

private {
 After Make a separate fn with those common parts, and use it in the methods whenever necessary.

INDEPENDENT PUBLIC WILL ALWAYS BE A PART

PRIVATE

To hide sensitive issues
use private else use public.
security purpose.

static data member

attribute of the class
itself not of the individual object

All objects share this attribute

static number fn

Does not depend on any object.

$\text{size of object} \geq \text{sum of size of non-static members}$.

class Student

{

int roll;

char name[31];

int score;

public :

void getdata();

void showdata();

};

void Student::getdata()

{

}

void Student::showdata()

{

}

Student s, t;

~~s = t;~~

~~s.getdata();~~

s.showdata();

t = s; \Rightarrow Byte wise assignment.

Dynamic memory allocation

```
int *p;
p = new int;
delete(p); ← freeing up memory
```

```
int *p;
p = new int[10];
p[0] → first element
of the array.
```

```
Student *s, x;
s → get_data();
(*s).get_data();
x.get_data();
```

In a class,
one can have object of other class.

```
class Date
{
    int y, m, d;
public:
    ...
};
```

class Student

```
{Date dob; ...}
```

If student is above Date

then we put class Date; → Forward Declaration

We cannot use

cin << dob;

Not possible

④ Nesting of class defn.

Defn
is limited
to that
class

```

class A
{
    class myDate
    {
        int dd, mm, yy;
        void setDate(int d, int m, int y)
        {
            dd = d; mm = m; yy = y;
        }
    };
}

```

↑
creating an object → [0]
[0] → [0]

```
void A::myDate::get-data()
```

```
class Item
{
    char icode[5];
    char iName[31];
    float price;
    float qty;
}
```

```

class ItemList
{
    Item list[10];
    void getData();
    void showAll();
    void insertItem();
}
```

~~private:~~ public:

~~Code { ... }~~ void get-data();

~~Code { ... }~~ void show-data();

~~Code { ... }~~ void set-price(); void set-price(float);

~~Code { ... }~~ void change-qty(float);

if trying to access a private member from outside, we need to get the data using a getter ~~method~~ method.

ItemList Interface

```

{
    static void operation()
    {
        ...
    }
}
```

Implement menu in main()

OR ItemListInterface :: operation()

27/8/2014

Programming

Practise

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Sanjay
Kumar
Saha
①

class Income

```
float amt;
```

```
public:
```

```
void set_income(float f)
```

```
{ void show...
```

→ In the class, declare the fn as friend.

→ If required, forward declaration of class.

class Expense

```
float amt;
```

```
public:
```

```
return (ic.amt -
```

```
exp.amt);
```

↳ objects of all the classes of which it is a friend are to be passed as arguments.

↳ In each of the class, declare it as a friend.

If a fn is friend of a class then that

fn can access the ~~public members~~ private members of that class. (e.g. if netBalance becomes friend of that class).

then it can access the private data members.

If a method working with objects of different type.

Global fn

Working as friend of no. of classes.

```
float
```

```
friend void netBalance (Income, Expense)
```

inside the class defn.)

② A fn member of one class acting as friend of others, then we use.

→ friend ^ Income :: netBalance (Expense)

→ I may skip forward declaration.

friend class B;

→ all member fn of class B will act as my friend.

setter fn → mutator / MODifiers
getter fn → accessor

int retr() const

x++;

return x;

→ Mutator

→ Accessor

→ Management

Constructor

Destructor

⇒ special type of member fn which are called automatically at the time of object creation

Can be used to initialize the object

⇒ can have any / arg. with default values no. of arguments

Normally, public / protected.

(in some case, it can be made private)

⇒ fn name is same as className

(constructor) constructor

⇒ No return type (even not void). [implicitly return the object reference]

⇒ Can never be static

→ If user does not provide any constructor, system provides a constructor called default constructor.

→ When user writes a constructor, system withdraws default constructor.

class student
{
 ...
 public:
 Student(int n=0)
 {
 ...
 }
 ...
 }
}

can take a single OR no argument

object declaration must match with exactly one constructor deft.

Constructor with no arguments or all arguments with default value \Rightarrow Default constructor.

Student s(5);
student t = s;
 \Rightarrow Copy constructor System automatically provides unless otherwise copying.

To invoke copy constructor

student t;
t = s;
Not an issue of copy constructor.

student s;
student t = s;
student t(s) \leftarrow copy constructor.

student t(student).

↑
Diff.
are altogether

class Array

{

 int *p;
 int size;

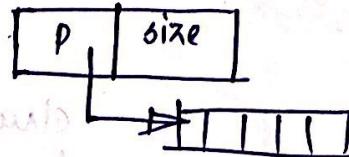
public:

 Array (int s = 0)

 { ... }

 { ... }

 { ... }



 Array (y = n); → Copy constructor

ptr data member
for which memory is dynamically
allocated.

→ Now, system given copy constructor

will make the ptr of new object (in
which copied) to point to the same
location.

→ Two objects will share the same
single copy of ~~referenced~~ referred
location. change made by one
will affect other.

→ If it is undesirable, one must
write own copy constructor to have separate
allocation.

copy constructor → $\text{Array}(\text{Array}, \&+)$

Destructor

~~#~~ ↓
 automatically
 called when the object is destroyed.

→ public → non static
 → no return type → no argument
 → \sim classname()
 → not even void.

\sim student()		stored in stack.
{ ...		1. s
}		2. +

→ We must write our own destructor when memory allocation is made

delete ap;

Array $a(\ast ap);$
 $\text{delete } ap;$

↓
 system will delete
 only ap.
 our destructor will
 delete the entire thing.

02/9/19

Programming

Practical

Part -
Sanjay
Kumar
Saha

Class Test

Syllabus upto (day before 02/9/19),

upto Constructor & Destructor

Date : 16th Sept, 12:00 pm

DATA ABSTRACTION

- To hide the complexity of representation
- Implementation details are hidden from the user.
- Behaviour of similar types should be similar.

e.g. $a * b$, $* p$ | $<<$, $>>$
↓ playing multiple role
Left shift, Right shift.

Operator Overloading

class Complex

```
int r, int i; add = add();  
public:  
Complex(int a=0, int b=0);
```

```
Complex Complex::add  
(Complex t)
```

```
r=a;  
i=b;
```

```
Complex add(Complex
```

$c3 = c1 + c2$

↑ this is overloading example

overload

• (a) review of the `operator` and `friend`

To overload any operator.

fn name

operator symbol()

e.g. operator +()

c3 = c1 + c2 ;

$c1.operator + (c2)$

(operator on the left)

taken on invoking

object

OR operator + (c1, c2)

either write
a member function
or a friend function

cannot have any default
value

~~complex~~

complex Complex::operator + (complex +)

complex c;

return c;

Cannot
discriminate
pre & post

Binary Operator

friend fn

→ operator on the left is 1st arg.

→ operator on the right is 2nd arg.

No implicit

type casting

complex operator * (int k,
complex c)
return (c*k);

Overloading
operator
to handle casts.

$$c_2 = c_1 * k$$

Complex operator * (int t)

friend Complex operator * (int t, complex r).

for e.g. ~~c - c~~

$$c_2 = -c_1$$

$$c_3 = c_1 - c_2$$

operator - (complex)

complext c1;
 istream → cin >> c1;
 ostream ← cout << c1; ostream, complext
object
friend void operator<<()
 + ... return t;
 ostream

friend void operator>>()
friend void

Note :- Destructor
 should be friend for
 a Constructor

byte-wise assignment

↴ c = a
 ↓
 c.operator=(a)

=
 ↑
 overloading
is done by member fn
 and NEVER by
 a friend fn.

operator[](int)

↑
Member fn only.

int operator[](int n)
 { ... }

friend ostream operator>>()
 istream &, complext

class Array

{
 int *p;
 int size;

public:

{
 ↴

}
 + copy
 constructor
 Array(const
 Array &t) ~Array()
 {
 if (p1 == NULL)
 {
 p = new int[t.size];
 for (int i = 0; i < t.size; i++)
 p[i] = t.p[i];

↴
 Array b = a
 ↴ Array(a)
 Array b(a)

To have no. of
reference.

Array

{
 ...
 int *ref_count;

}
 : *ref_count
 else
 ref_count = 1;
 ref_count = NULL;

④

Array (Array *t),

{
int a[10];
}

void operator()

{

int b[10];

~t { ~~Destructor~~

~t { ~~destructor~~

~~~t~~ { ~~deallocate~~

(deallocate), free b[10];

~~~t~~

~t { [deallocate] b[10];

(deallocate)

~t {

void f

{ int a[10];

~a(); }

(11) f();

; a[10]; {

(b) ~~destructor~~ ~~deallocate~~

(b) ~~deallocate~~ b[10];

a = d (point)

(d) points to

(a) d points to

b[10] (point to)

int a[10];

~a(); }

int b[10];

~b(); }

bool f() { t a[10]; }

int b[10];

~b(); }

(a) f();

int c[10];

~c(); }

int d[10];

~d(); }

(b) D ~operator

~operator

~operator

~operator

~operator

~operator

~operator

~operator

~operator

03/09/2019

Assignment Operator Overloading

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Saha

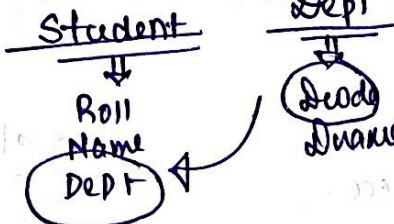
Programming Practice

→ Return type as per requirement

return type void obj1 = obj2
obj1.operator(obj2)

return type
object

obj1 = obj2 = obj3



Don't inherit the entire class, if you need only one part

studentList

Array of student object

DeptList

Array of dept. objects

student

studentList

Student list[100]

Dept

DeptList

display()

to display department,
we need to provide support

student return student
(int rollno)

{ ... }

return student

b student

(int rollno)

student 8987

return student

() { rollno }

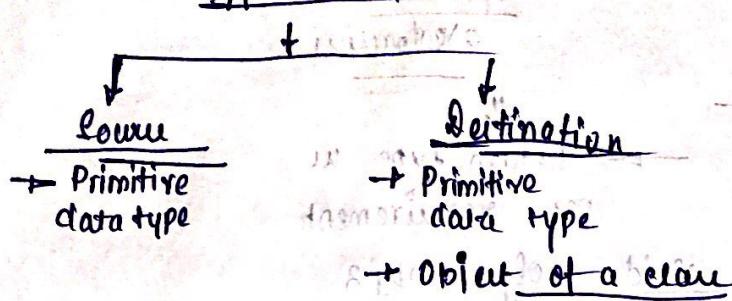
8987

rollno 47

return student

(2)

Type Casting



destination

source

$$x = \underline{y}$$

explicit

(float) n.

or float(n)

time t;

!

t=n;

int
no. of sec.

class Time

{ int min, sec, hr;
public:

Have a constructor
with basic type as argument

Source

Destination

time(int n)

t

n

hr = n/3600;

}

Destination

Source

member fn

n

t

operator int()

||
...
|

Write suitable
type cast operator

operator $\varphi()$

type
in which
casting is
required

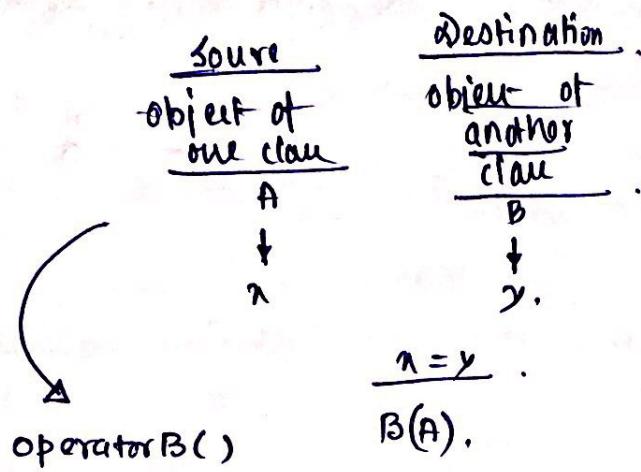
type cast operator

↓
member fn
no return type in
implied

double d

d = float(d)

$C1 = C2 + \textcircled{3}$ → user-defined
function



```
class month-day
{
    int month, day;
public:
```

```
class dayOfYear
{
    int day;
    ...
}
```

```
month-day m;
day-of-year d;
m = 1;
month-day(day-of-year& k);
```

operator month-day()

```
m = —
d = —
return
month-day(m,d)
```

09/09/2019

Programming Practice

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(1)

Inheritance

Base class / Super class / Parent class

To share attributes & behaviour of parent class with the child class.

Derived class / Sub class / Child class

Need to add certain attributes / behaviour to modify in an existing class

↓
direct change in the class may not be possible (source not available) even if accessible some case old behaviour is required.

For some logical order & in some case modified behaviour is required.

OR Redoing old things which are required and designing own class covering all derived features.

Best Soln.

Existing Class

Inherit

child class

Add additional attributes / behaviour or change existing behaviour.

class B : private A

Type / Mode of Inheritance,

class A

{ int x;

public:

getdata(void)

{ cin>x;

show(void)

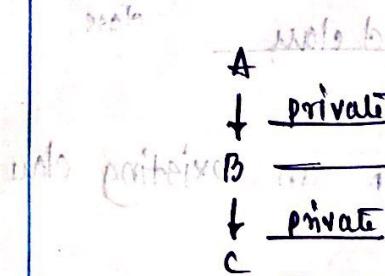
{ cout<x;

}

};

② private data of base class
are not accessible in
derived class
(but those are inherited)

Multi-Level Inheritance



Important

Protected

Protected members
are available in
itself & in derived class

Type of Inheritance

Access Specifier

in base class in derived class

~~Inheritance~~
Inheriting
everything

Private

Private

Not accessible

Public

Accessible

Private

Protected

Protected

Private

Protected

Private

Not accessible

Protected

Protected

Public

Protected

When inherited in private

public base → private in derived class.
members

Selectively the public members
can be inherited as public/private

Public

Private
Protected
Public

Not accessible
Protected
Public.

Use scope resolution operator ::
to access the fn from the base
class having the same name

function overwriting

If a derived class fn has same name (signature may not be same) as a base class fn then it is that base class defn that is overwritten in derived class.
With a derived class object if overwritten fn is invoked, it always goes till derived class.

Lookup for a fn

An object invokes a fn

C++ → Fn overloading is in a single class.

To go to the base class,
use scope resolution
operator.
(check accessibility)

LOOKUP Mechanism

- Find class in the same level
- If not found, go to the class at higher level
- If not found, error.
- If an variable is present in the ~~class~~, fn we that if variable is present in the class , we that if not found go to the higher class
- If not found finally global variable

(b) Inherited
members

static members of base class

Also be shared

by derived class

instance

availability in
derived class to

protected/public

Friendship

(friendship for class B)

friend class B
granted friendship
to B

$A \rightarrow D$
automatically
does not inherit
friendship.

$A \rightarrow Aa$; def.
 $B \rightarrow Bb$; def.
constructor

$B(\text{const } A \& -)$:

$y = +y;$

$A(\text{const } A \& -)$

$A \rightarrow B$
derived class
const. can work with
only additional
attribute

$B(\text{int } a=0, \text{ int } b=0) : A(a)$

call default
const. of base
class.

explicit
call to
base
class

constructor.

base class

base class

base class

base class

base class

base class

20/11/19
17/11/19

Programming Practices

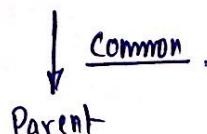
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(1)

→ Inheritance
Sharing of attributes behaviour

Generalization

Child class



→ Mode of inheritance

→ Scope Rule.

Specialization

General class/parent

↓
Subclasses.

Over-riding

fn name same as that of base class in derived class

Signature may/may not be same.

Whenever derived class

obj → calls a fn.

derived class reference is automatically converted to base class reference.

Derived
class to
Base
class

B(const B & +) : A(+).

A a;
B b;

→ | a = b.

Takes base class posn & rest is ignored.

Base class to
Derived class

NOT AUTOMATIC

OR

operator B()

②

A → B.

A * p ;

↓
point to either a base class object
or a derived class object

p = &a ; p = &b ;

p = new A ;

p = new B ;

EMPLOYEE

ename
basic

public:

call-pay()

EMPLOYEE

OFFICER
→ allowance
calc-pay()

runtime / dynamic
linking OR
late binding

EMPLOYEE * ep ;
EMPLOYEE e ; OFFICER o ;
ep → calc-pay()

ep = &O

ep → calc-pay().

Using base
class pointer,
Access is limited
to base class
part

C++

• Base class fn

for which run time polymorphism
is required has to be
overridden in derived class with
same signature.

• Base class is to be declared as
virtual,

hence Base class.

EMPLOYEE

.....

virtual void call-pay() { . . . }

System

Whenever a class has a virtual fn, a table for the class called VTBL is prepared. Maintains an entry for each virtual fn in the class.

fn prototype & starting address of the code.

In the class, system keeps an attribute VPTR.

VPTR \Rightarrow points to the VTBL of the class.

A derived class inherits from the base class.

VTBL of base class is copied.

VPTR of any derived object points to VTBL of its own class.

If any virtual fn is modified then its address in VTBL is updated. If any new virtual function is added, then an entry will be made.

$A \rightarrow B.$
 $ap = new B(\dots);$
 ↓
 constructor.

make * ap ← delete ap;
 call to the
 base class destructor.

→ To handle delete the destructor as virtual.

Why? Why?

derived class destructor
 ↓
 base class destructor.

24/9/2019

Programming Practice

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Class Student, Faculty, Staff

may have common attribute.

which can be used for

generalization. We use
a higher class for this.

Abstract Class

In reality there exists child/sub derived class objects. They have lots of commonalities. To avoid repetition, take out common attribute/behavior and form a base class. If it is so then in reality, an implementation object of the class will not exist then make it abstract class.

A class with atleast one pure virtual fn is an abstract class.

Note Any class derived from an abstract class must over-ride the pure virtual fn.

Same name -
same
signature

class A

public :

virtual void f() = 0;

}

void A::f()

{

A

{}

tail tail + A

b f();
public void b();

void f() { }
Not possible

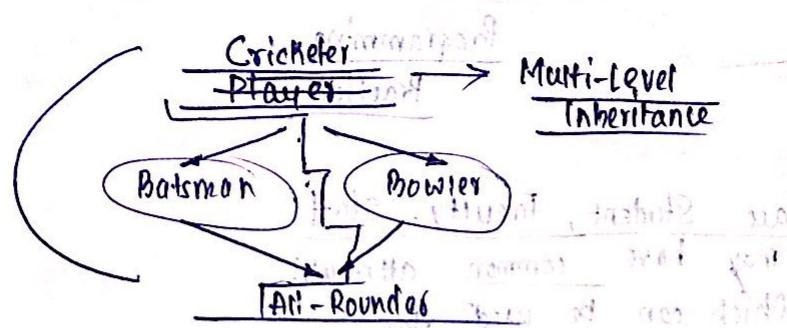
public void a();
void f() { }
Not possible

No direct
instance of
abstract class

A a; → X

A *P → ✓

② ~~for
multiple
inher~~
Multiple
Inheritance



Multi-level Inheritance

Multiple instances of common ancestor via different paths.

Note:- Do not consider a relationship consists of multiple inheritance.

e.g. If ~~person~~ is person & wife & husband & wife & family

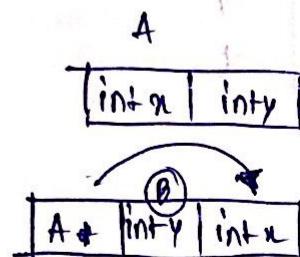
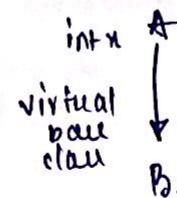
Not inheritance of the teacher who has family who holds A

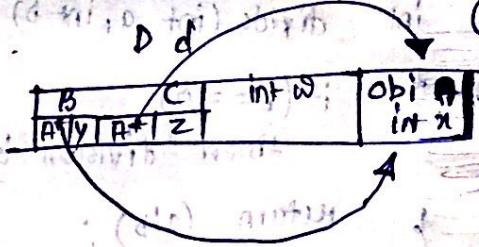
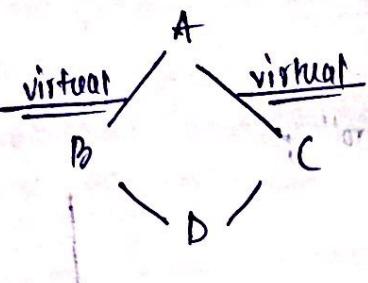
Virtual fn

Common ancestor will be declared as virtual base class for its children

class batsman : virtual public Cricketer

class bowler : virtual public cricketer





If there is any fn. with same name present in both B and C, consider must override this.

D d;

A * p;

p = &d; + limited access to base class.

Exception Handling

int divide(int a, int b)

return a/b;

if (b == 0)
exit(1);

try block

try {
 throw exception
};

To raise
the
exception
abrupt
termination
of the
program

int main(void)

{
char c; int x;
int y;
do
{
cin >> x;
};

cout << divide(x,y) << "\n";

};
while(c == 'x' || c == 'y');

| if (y == 0)
| division
| by zero

exact match/
match after trivial conversion.

No promotion/No standard
conversion

comes out
from subsequent
catch blocks

int divide(int a, int b)

{ if ($b = 0$)
throw "division by zero"; }

} return (a/b);

close my exception object

a < b
throw

my exception object
t(" ", a, b).
} throw b;

char s[];

int x, y;

public:

my exception object(...)

{ strcpy(s, c); x = i;
y = j; throw exception exp r;

}

void show()

Instead of
default
terminator.

fn

#include <except.h>

Set_terminate (function
pointer)
Set_unexpected (...)

return type fn-name(args[]) throw (exp1, exp2, ...)

try {

fn-name()

catch ()

{ ... }

... }

... }

... }

... }

... }

... }

... }

... }

... }

... }

... }

... }

`int max(int, int)` Same code
`float max(float, float)` || repeatedly in all
`char max(char, char)`. defn.

(88,1a) non M.F.U.O.T

Function Template

Used to create an instance of ~~function~~ It is a skeleton/template for defining a function for a specific type. It is a skeleton/template for generating a number of function definitions. Formal arguments may be of generic type. Instead of specific type, one can put placeholder for data type. Actual type will be decided looking into the actual parameters.

max(Ta, Tb) | e = max(a, b)

max(Ta, Tb)

template <fun T>
max(Ta, Tb)

Each generic type must appear at least once in the argument list.

If $a > b$
return a;
else
return b;

int a, b;
float f1, f2;

cout << max(a, b);

An instance of $\max(\text{int}, \text{int})$ exists or not?
If not then create an instance using function template.

Whenever an instance is created for a fn template

for float use: $\max(\text{float}, \text{float})$

if will work only with exact match.

cout << max(f1, f2);

max(float, float)

however

cout << max(a, f2)

→ create an error

① Exact match with normally defined (not template) overloaded version.

If found then

else ② tries to create an instance with template if possible consider that else

looks into overloaded fn with promotion, conversion etc.

②

class Student

• must have operator > (student)

↑
operator >
↑
operator <

Tout <- max (s1, s2)

not. friend ostream & operator << (ostream &, student);
work w.r.t. to reading a file
year, stream & operator <<
to insert the value to <<
in relation to file or file pointer
Note: No. of classes with same attribute &
methods may have to be defined even
if only certain attribute type is different
but role/interpretation is same.

template <class T>

Required for class

(defn. of) elegant

class 'Array - 1D'

{

• T* arr; // note

• int size; // note

public:

 Array - 1D (int u = 0)

 to standard constructor

 no. of rows to fill

 interprets size as u;

 if (size == 0)

 arr = NULL;

else

 arr = new T[size];

 resized

 T add ()

{

 ...

}

};

..

class

array

int

size

arr

int

size

arr

int

size

Array - 1D <int> x;

every member fn

of a class template

in a fn template -

template <class T>

T Array - 1D <T> :: add (vo

of

int i;

int arr;

arr[i] = vo;

return arr;

for template void D

(empty fn) biggest element

template <main> void D

int Array <T> :: eat_size

of

return size;

Template graph after iteration 10

class template

When an object is created for a type and it is the first object for the generic type.

- Instance of the ~~template~~ class defn has to be created.

All static & non-static data members fn & static data members (if at all) are also created.

For the next objects, only the non-static data members are created.

class c1 < class < c2 > >
 +
 space .

if func is a static data member.

template <class T>
T Array-1D<T>::c = 0
char Array-1D <char>::c = ' '
specific initialization

template <class T>
class ~~simpl~~ Spt-Array
{
 ...
 friend class Array-1D<T>;
};

template <class T>
class B : public A<T>
{
 B() : A<T>();
};

do have,
to specify

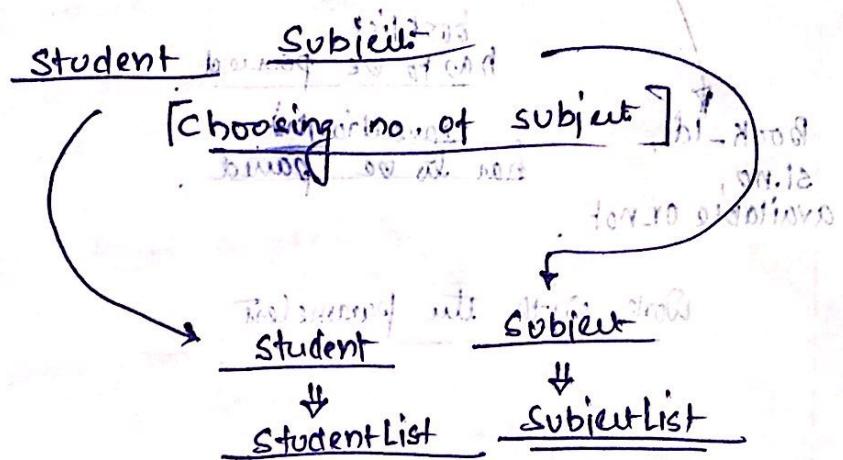
21/10/2019

Programming Practice

(...) week - working. 9/10/2019

Prot.
Sanjay
Kumar
Saha

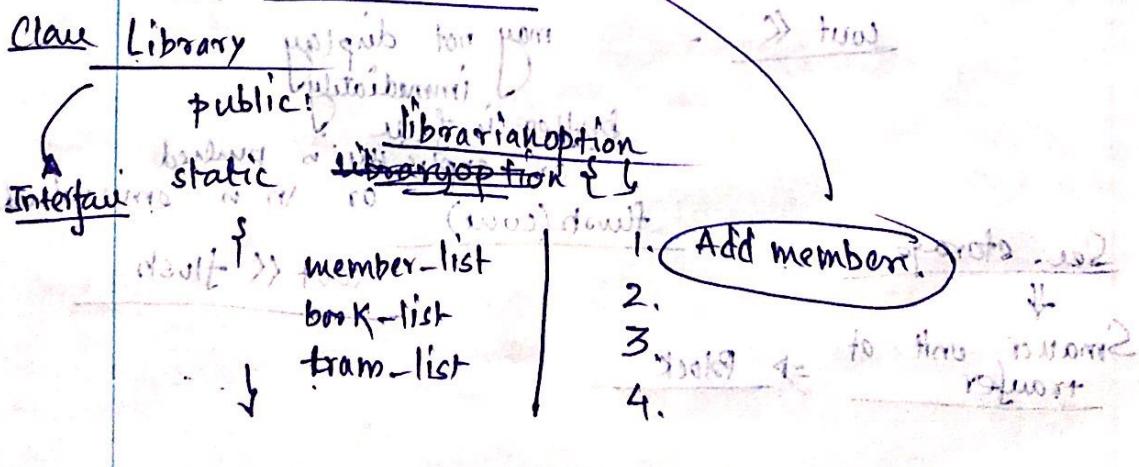
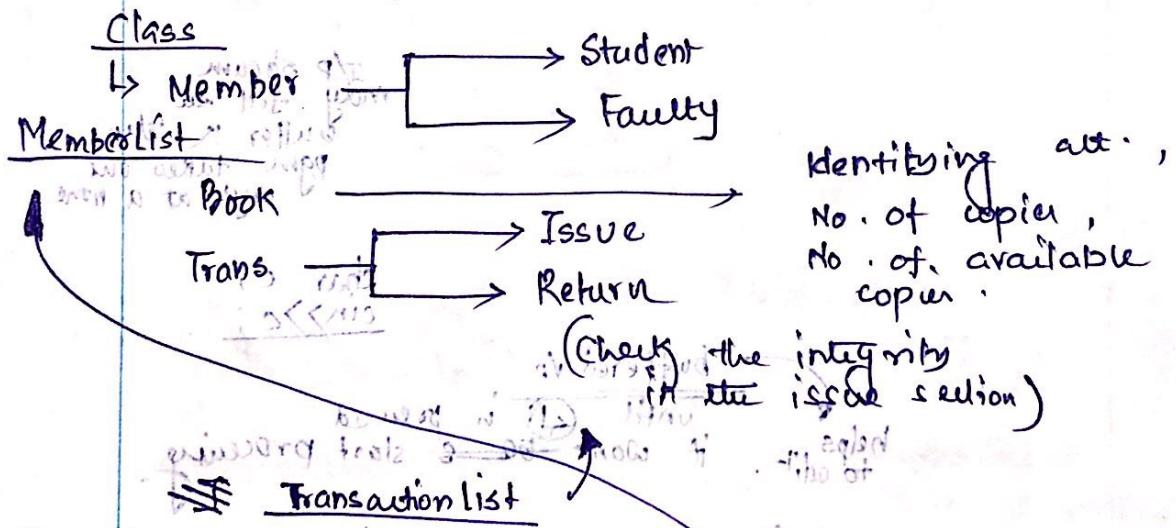
①



3rd & 4th
Collection class
OptionList class

broadens A. respo.
Simplifies A. respo.
Identifying alt.,
No. of copies,
No. of available
copies.

Library



②
2015
KUMAR
SOP

writer's premonition

processes

③

Issue - process - issue (...)

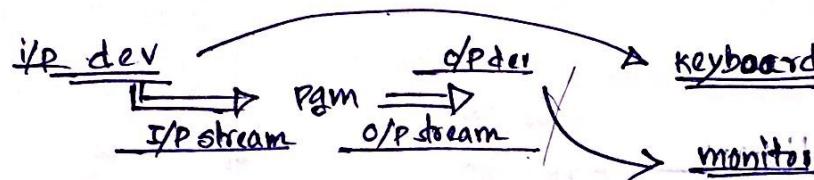
①
booklist has to be passed attribute
Book-Id, transactionlist has to be passed
sl.no,
available or not.

Work with the parameters

file pointer
file pointer

note: I/O & FILE

note: stream of bytes



processable
biggs to on
distances to on
newspaper

attribute
buffer
queue
writer
char c; NOT
cin>c;

I/P stream may fill the buffer rather pgm takes one byte at a time

buffered I/O
helps to edit until it won't be in proceed if won't be start processing

cout <<

may not display immediately

buffer is full or explicitly pushed flush(count) or in environment

See: storage

Smaller unit of transfer ⇒ Block

cout << flush

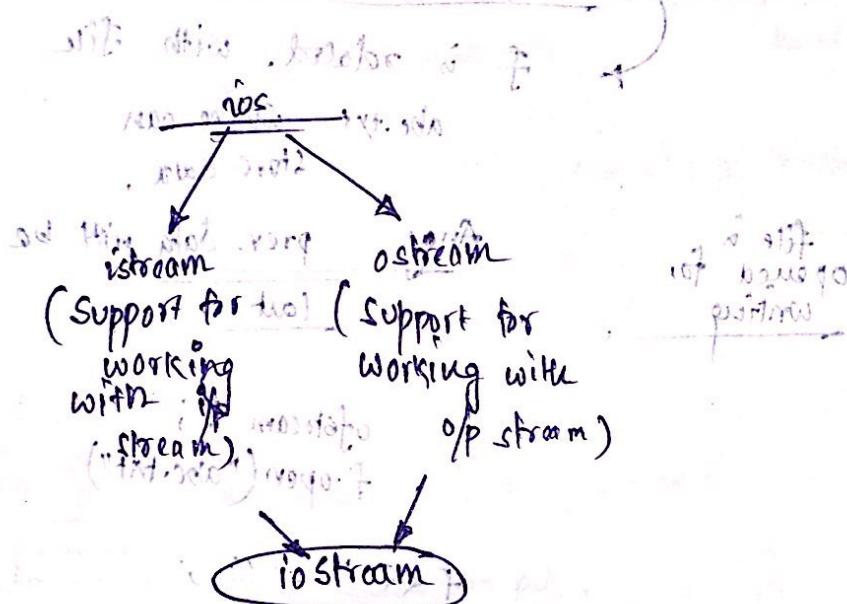
flush

flush

streambuf class

↳ buffer management information about the memory for buffer to support + buffer filling, reading from buffer, etc. (buffer mgmt.)

ios | keeps a ptr to 'streaming object'
keeps basic file info like
binary/text file information open for reading
(reading)



char ch; // ignore tab/space/\n
cin >> ch;
while (ch != '\n')
{
 cout << ch;
 cin >> ch;
}

Go on continue
cin.get(ch).
get(ch).
...

takes at most 50 char
or it encounters a newline which even happens
cin.get(str, 50);
cin.getline(str, 50);
abc. xyz\n abc in will stay in the stream
abc in will be deleted out

noto *loderio*

`cout.put('a');` appends 'a' to the output stream

ifstream class

ifstream ifstream class ofstream

iostream

quidam et usq. secesserunt ut vix videtur.

• $\text{ofstcam} \rightarrow f(\text{abc}, \text{txt})$

f is related with \tilde{f}

~~abc.txt where can
store data,~~

file in
opened for
writing

Anay prev. data will be

at 100% out of time

What you know

ofstream f;
f.open("abc.txt")

和分離的，

One can
work
with f stream of

~~guitar no 32~~

(45) 120-1912

(Wb) def

ifstream f("abc.txt");
ofstream f1; f1 << "Hello";

• (Ab)-ges. 100 → \rightarrow name
imperfektiv modaler auxiliell \rightarrow name
~~if fact fact~~ while(f.get(c))

(1) ~~logistic~~ ~~multiple~~
of outcome

411-47638
100% pure

三

~~closed~~)

f.close();

reports with input 2: line 10

fopen(

two batelets as New York 345 app

Programming Practice

22/10/19

SKS Sir

2nd CT

→ NOV 18

(except FILE)

random fop

random fop

Binary file

fop of small

orderly work

bit by bit

1) ios::in ⇒ A file must exist (to read)

2) ios::out ⇒ to write (old data lost)

3) ios::app ⇒ to add data at the end

class Student
f ... ;

fstream f;

f.open("student.dat", ios::out | ios::binary);

student s;

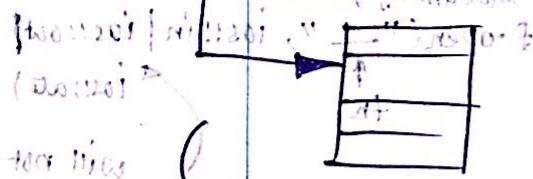
s.getdata();

f.write((char*)&s, sizeof(student));

write(char *, no. of bytes)

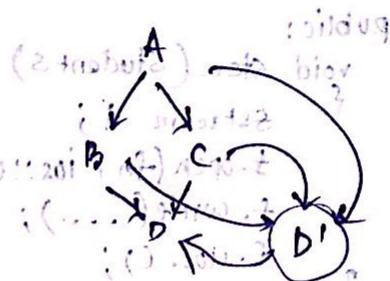
Student

| roll | name | score |
|------|------|-------|
| 5 | 100 | 78 |



If address is stored, then
Who says the address will
be stored,
address will
stay the same.

(q. tri) introduction for



Read handle file

fstream f;
f.open(" ", ios::in | ios::binary);

while(!f.eof()) {
 f.read(s...);
 s.display();
}

white(f.read(..))
s.display(); OR

PI/CS
2
11-2012

wir sind gemeinsam

①

Random Access

std::vector<

get marker

put marker

from where
to read

where to put

std::vector

object. seekg() (fxt object; seekpc);

offset, reference. & p. ios::beg (wrt start),
ios::end (wrt start),
ios::cur (wrt current).

i.e. iostream

f.seekg("file", ios::beg, 0);

f.read(...); gets the first object.

(On encountering eof, it a flag in set.)

Subsequent read operations are ignored.

(Set to on, & so on) f.fail.

Object.fail() => if last operation fails.

f.clear() => clears the flag,

class StudentList

f

char fn[31];

public:

void store(Students)

f

fstream f;

f.open(fn, ios::out | ios::app);

f.write(...);

f.close();

void showall()

{ f.read();

f.read(fn, 31);

now tellg() goes to pos of getmarker

tellpc() goes to pos of putmarker

say we want to

modify (int x)

fstream f;

f.open("c", ios::in | ios::out)

ios::ate)

fn

write f

out, read

Truncation will be suppressed

file unusable

use ate part

void searchstudent(int p)

{

}

((..) boor. f) show

; () kogaiib.

(3)

```
searchstudent(int p).
if (rec!=0)
{
    f.seekg(ios::beg , (rec-1)*sizeof(student)),  

    f.read(...)
    f.seekp( ), );
    f.write( ... )
}
f.close();
```

If we have variable sizes, then we can store the size look up table in another memory structure.

05/10

Programming Practice

Prof.
Sanjay
Kumar Saha

05/11/2019

Namespace is scope for std:: namespace or si
NAMESPACES multilevel namespaces facilitate better organization



- To group named entities within a scope. (Other than Global).

(Identifier, fn, class, etc). safe to have

multiple objects if used in

- To avoid name collision

entities by same name can exist in diff. namespace.

• May span multiple-blocks

↓
maybe in different files also

should not part of main()

namespace ns1 { }

int x; ⇒ available to all

} void display(); ... in this namespace.

→ void... f(void);

Any other global fn.

another function instance

function inside

ns1:: f(void);

↓

use this name space.

functions & void ns1::A::display(..)

functions f(...)

May have diff. files

→ nspace1.cpp

→ nspace2.cpp.

to

with appropriate header

#include "nspace1.cpp" - header

2 ways:

using declaration

using declaration. declarator.

• declaration: signature;

using ns1::abc;

• declare with member, I have
for every member, I have
to provide a declaration

• not triggered since obvously it is
using namespace ns1;

• required till declaration of → too many namespaces
multiple .h2 files

using ns1;

Standard Template Library

is a collection of lots of tools & facilities with efficient implementation that can be reused for developing applications.

Suppose A infinite length binary array of a

List of Nos. (H2, m2, H3, m3, m4)

we have to find the average.

→ Array → Start from 1st element
→ Linked list → add...list n.

→ Calculate average by, definitely
by count.

Taking help of

(STL)

1. container

2. iterator

3. Generic algorithm

4. function object

5. Allocator & Adapter.

provide
limited
interface
in a
container

Vector<int> v;

vector<int> :: iterator p;

Container

→ is an object that holds no. of objects say, vector list.

- Mostly are template & abstract.
- General use of container.

Iterator → Helps to iterate through the container without identifying

the name, type, ..

Container specific

container<type> :: iterator i;

* i → gives the value.

i++ → increments.

iterator is a copy of it.

Generic algorithm → to provide basic support for the container.

Function object → to customize the behaviour of certain std. algorithm.

Stack, Queue

vector

Random access,
 ↑ Array like, expandable.

```
#include <vector>
using namespace std;
```

```
vector<int> v;
```

To add data at the end : (3) push-back (3)

v.push_back(10);

v.size() → Gets the size

v.at(index) → checks the bounds.

v1.swap(v2) → Swap their contents

v1 = v2 (all elements)
 assignment

==, !=, >, <

void pop_back() Remove the last element.

↑
 will not
 return any value

v.front()

& v.back

Both getter & setter.

v.resize()

all.

v.resize()

v.resize(index)

Removes the
 element at the index

(1) begin()
 (2) return
 (3) iterator
 (4) + 1 to 1st
 (5) copy & pack.

end()

last + 1

Iterator

```
vector<int> :: iterator it;
```

```
if = v.begin()
```

```
while(it != v.end())
```

```
{ ... }
```

(1) first &

(2) next &

(3) push - 909

(4) next - 909

(5) loop ...

list()#include<list>List

All features of vector present here.

void push-front()

void pop-front();

void push-back();

void pop-back(); *last add to front than FT*

if container contains user-defined objects

then for comparison on container

they support most of the overloading operators

remove (iterator)

all occurrences will be removed

n. sort() must be overloaded

support for operator <()

l1. merge(l2);

begin(), end(),

#include<stack>stack

Returns the top element

stack<int> stak;
st.push();
st.pop();

stack(adapter),

while(!st.empty())
{ cout<<st.top();
st.pop(); }

#include<queue>queue<int> q

q.front()

q.back()

pop-front()

push-front()

... back();

(Histogram) < stack > class Test-(18/9/18)

Remaining
(Chancery - 1) file except File

11/11/2019

STL

Prof.
Saujan
Kumar
Saha

contains

different

components

sequence container

vector, (list) stack, queue

adapter

stack, queue, etc. class

Iterators

generic algorithm

→ fn objects

→ allocators

(or) stack, set

(or) stack, stack

with respect
with respect to
circular

stack \leftrightarrow st

st.push()

st.pop()

st.top()

size()

empty()

while(!st.empty())

{

wout < st.top()

{

: st.pop()

{

(empty var)

(new var)

can be used for assignment

Queue

empty()

size()

front()

back()

pop()

push()

→ returns 1st element

→ returns last element

push() → appended

Priority Queue

```
#include <priority-queue>
priority-queue<string> pq;
pq.push("Hello");
pq.push("World");
```

- definition

Basic Data Type

higher the value, higher the priority.

pq.push(10);

pq.push(20);

10 < 20 \Rightarrow 20 > 10

pq.top() \rightarrow element with the highest priority.

class myString

{ string w;

public:

myString(string s)

{ w = s;

bool operator < (myString s)

{ if there are issues
return (w > s.w)

friend bool operator <= (myString s1, myString s2)

{ return !(s1 < s2); }

friend bool operator < (myString s1, myString s2)

{ return (s1.w < s2.w); }

`#include <string>`

string s1("abc");
 s1 = "abce";

string s2(s1)

`*s1`

`s1[index]`

`s1.at(index)`

if out of range, return exception is raised.

`s1.find(string s2)`: to find first no of match of s2 in s1.
 if it returns the start index when s2 is found.
 if it cannot find s2, returns string::npos.

`find-first-of(string)`

→ pos of first occurrence of any character in s1 which matches with any of the characters in s2.

`find-last-of(string)`

`find-not-first-of()`

→ find the first occurrence of mismatch.

to get a substring.

`substr()` → returns another copy

→ `(substr(pos))` → return substring from pos to end.

`substr(start, pos, end, pos)`.

`size()`, `empty()`, `swap()`.

smaller in pos to big end.

(pos, return till points, till goes to max)

last = max - r.

~~replace~~ → now many ch. to
↑ be replaced

~~replace~~ → replace (arg₁, arg₂, arg₃, arg₄)

from which
the replacement
starts : ("H")₁₂ Replacing
starting posⁿ
in replacing
string. [ch]₁₂ = D string.

#include <pair>

pair \Rightarrow collection of 2 values.

pair < type₁, type₂ > p

p. first

p. second

p = pair<int, int>(value₁; value₂) \Rightarrow two if

\rightarrow one copy of key (pair₂).₁₂ \Rightarrow allowed in map. all numbers are

\hookrightarrow sorted associative table terms & &

pairs are stored in a sorted manner. on key.

Value is information
associated with the key.

queue, dequeue

#include <queue>

dequeue \hookrightarrow

double ended

m2(m1)

m2(&pa[i], &pa[j])

m2.insert(pair<int, string> :: value-type (pair<int, string>), \downarrow)

Returns a pair
(iterator, bool)

one copy of key is allowed

\hookrightarrow pair <map<int, string> :: iterator, bool>

r.second = true.

`begin()`
`end()`

`map<-, -> :: iterator t;`

~~t = m.begin()~~
`t ++`

`t → / / / |`

`r.first = t`

`value . (r.first) → first`

`m1[key]`

\Downarrow
`m1.at(key)`

OutofRangeException

if datatype not
a basic datatype
then `<, >, ..`
operators must be
overloaded.

hash-map

⇒ associated array.

elements

are stored according
to hashing scheme
on the key.

to access the
elements in a faster way
without the burden
of sorting on key
 \Downarrow
read in
map.

String
`char const * c = c-str(string)`

`hash-map<type1, type2>`

\Downarrow
my \Downarrow value type.
my type

`hash-map<key type, value type,
hash fn, Equality
object>`

class template,
for temp file,
namespace template