



BINARY SEARCH TREE



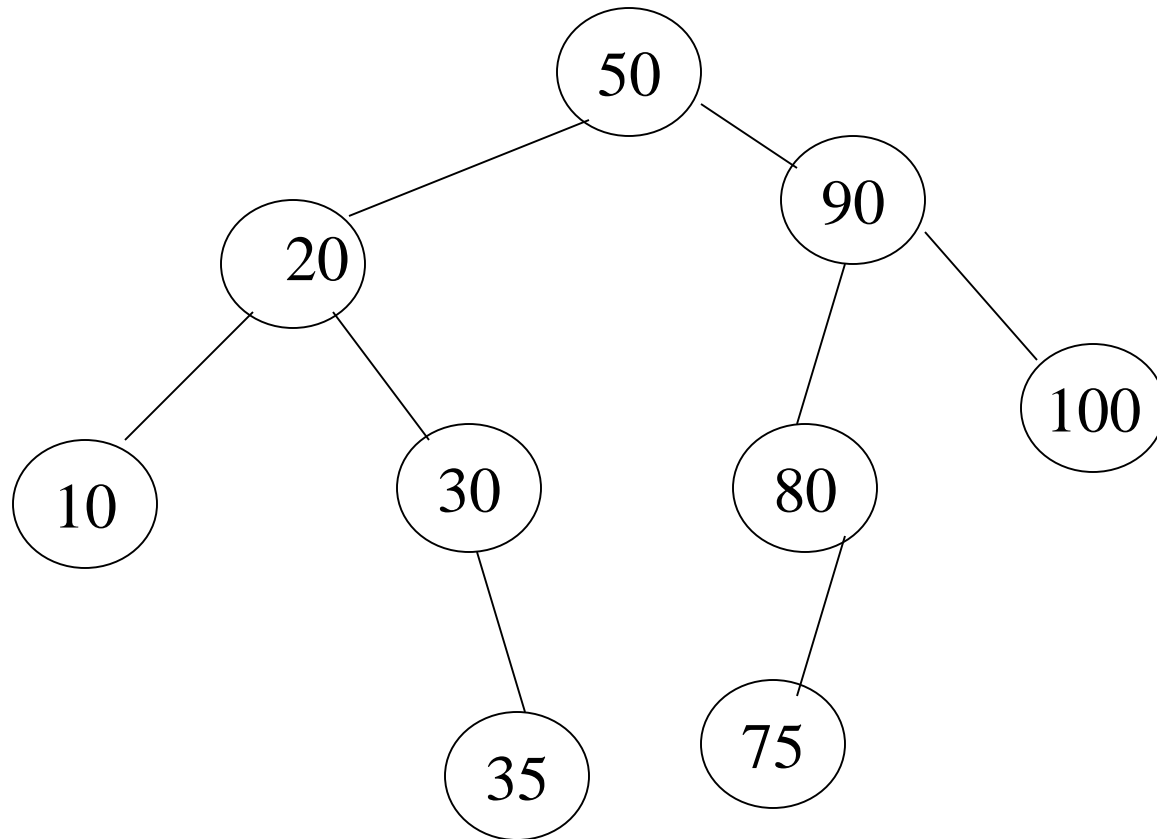
Binary Search Tree (BST)

A BST is a binary tree T with the following conditions:

- a) Key of every node in the right sub-tree of T is greater than the Key at root.
- b) Key of every node in the left sub-tree of T is less than the Key at root .
- c) All Keys are distinct.



An Example





BST Operations (in addition to those of Binary Tree)

1. Search for a key
2. Insert a key
3. Delete a key
4. Findmax & Findmin



Recursive Search

```
BST * search (T key, BST * t){  
    if (empty_t(t))  
        return NULL;  
    else if (key==t→info)  
        return t;  
    else if (key < t→info)  
        return (search (key,t→left));  
    else  
        return (search (key, t→right));  
}
```



Non-recursive Search

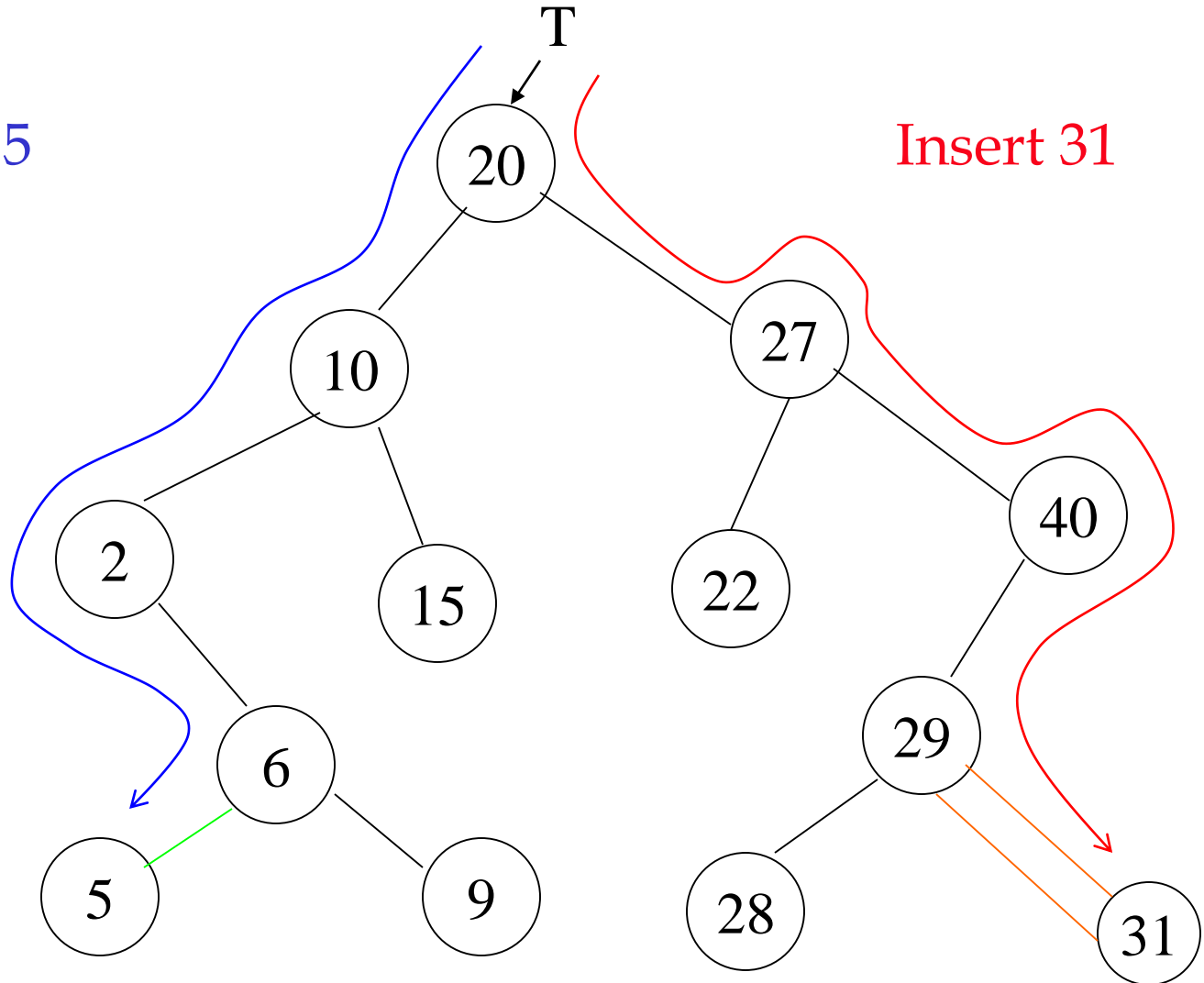
```
BST * search (T key, BST * t) {  
    BST *cur ; int found;  
    if (empty_t(t))  
        return NULL;  
    else{  
        cur=t; found=0;  
        while( (cur!=NULL) & (!(found))){  
            if (key==cur → info) found=1;  
            else if (key < cur→info)  
                cur=cur→left;  
            else cur=cur→right;  
        }  
        return cur;  
    } }
```



Insertion Example

Insert 5

Insert 31





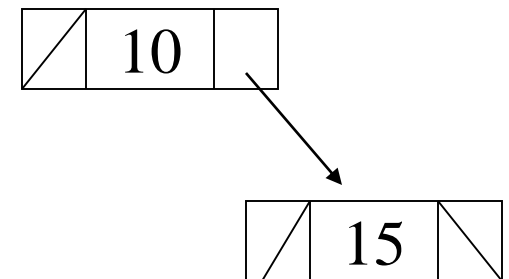
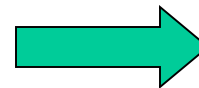
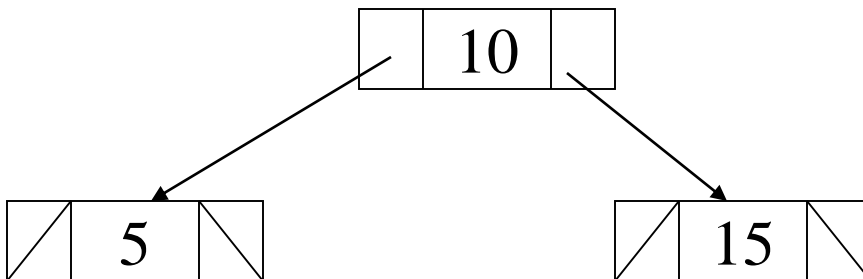
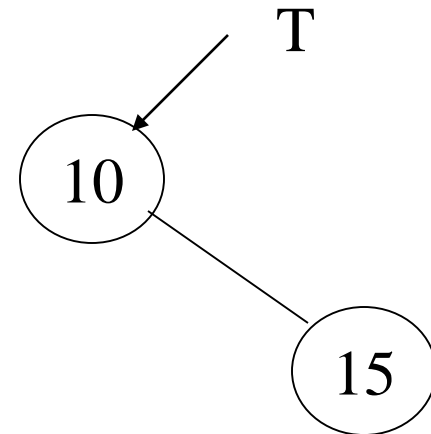
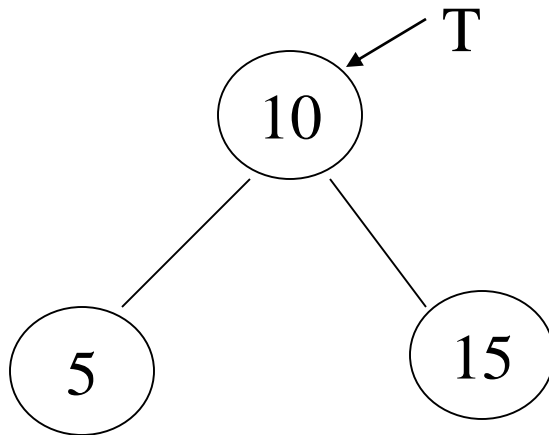
Deletion Example

- Delete 5



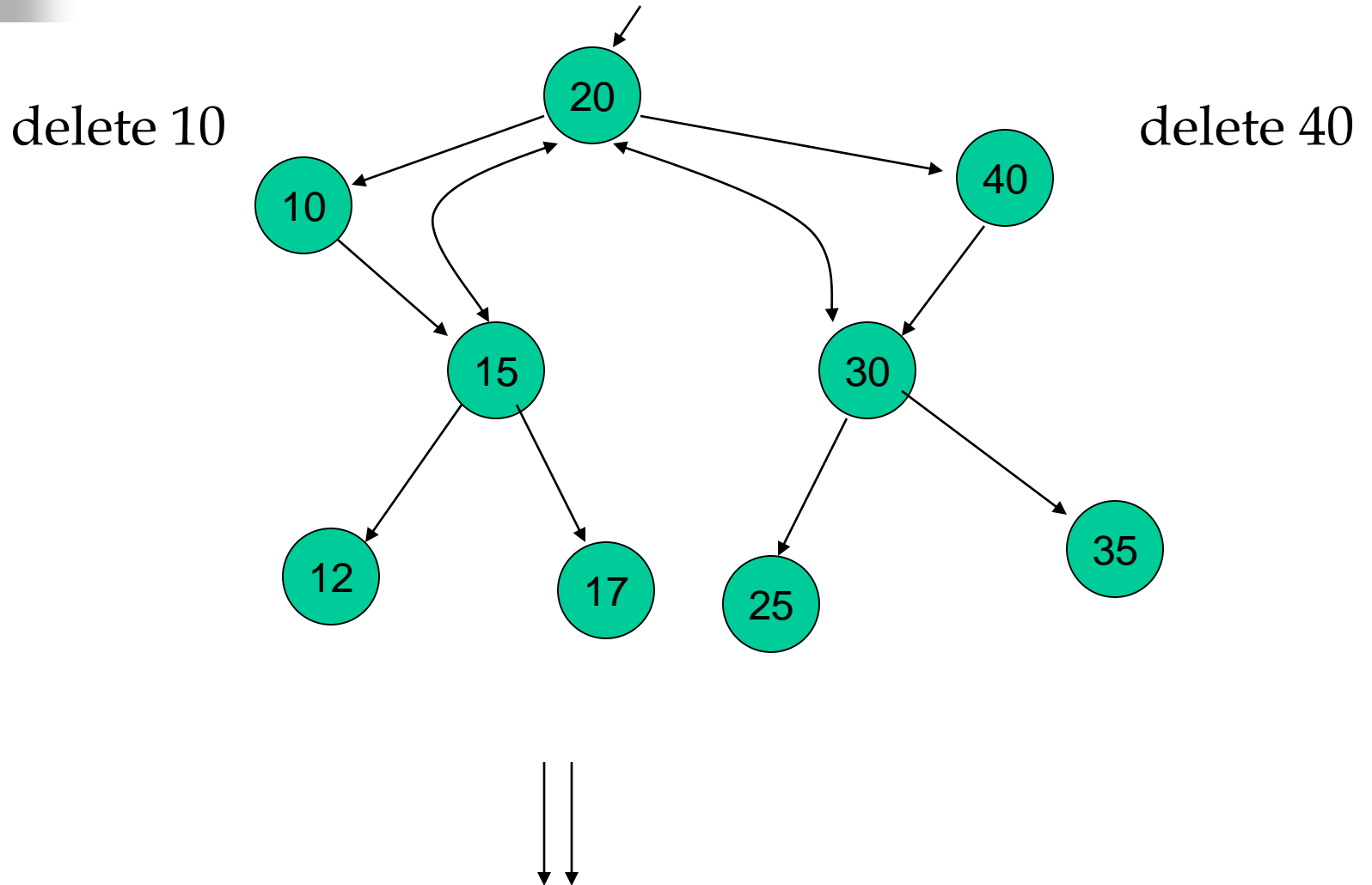
Error

- Delete 5



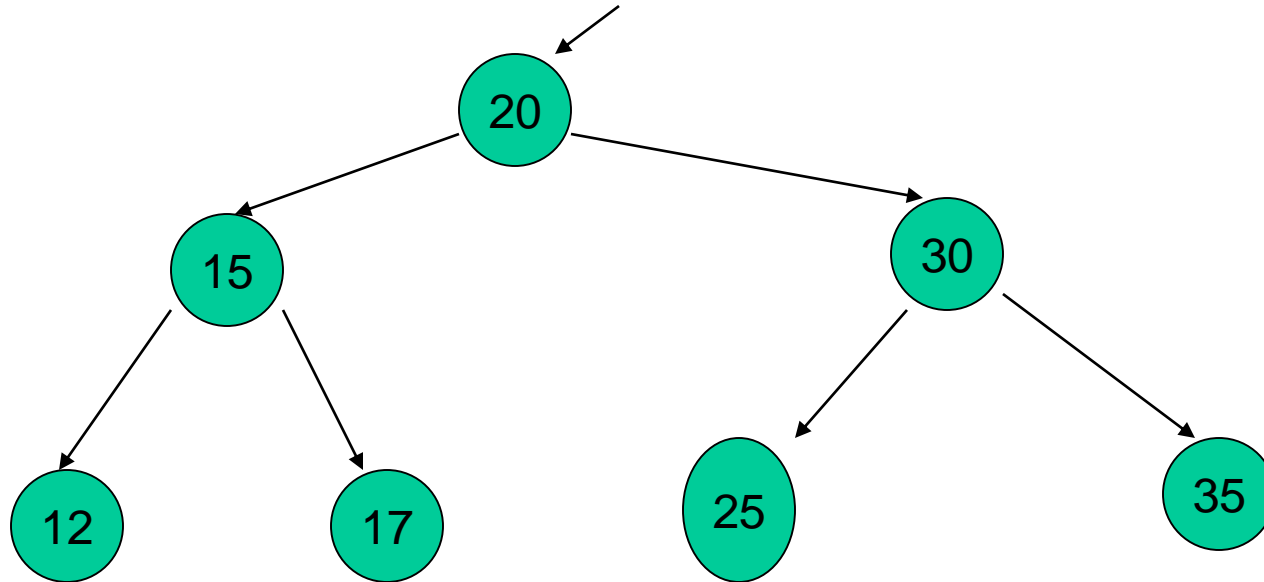


Deletion Example ...



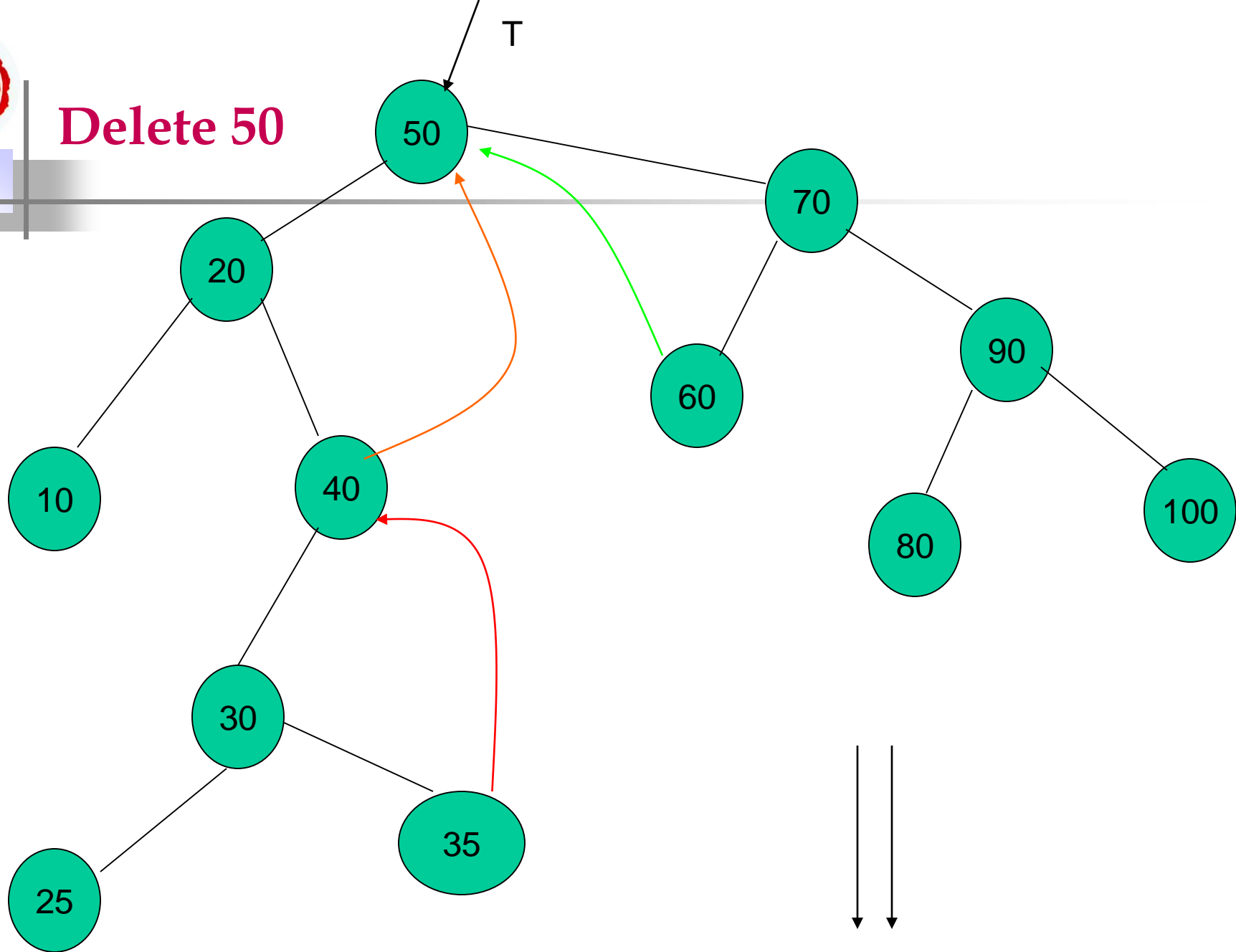


Deletion Example ...



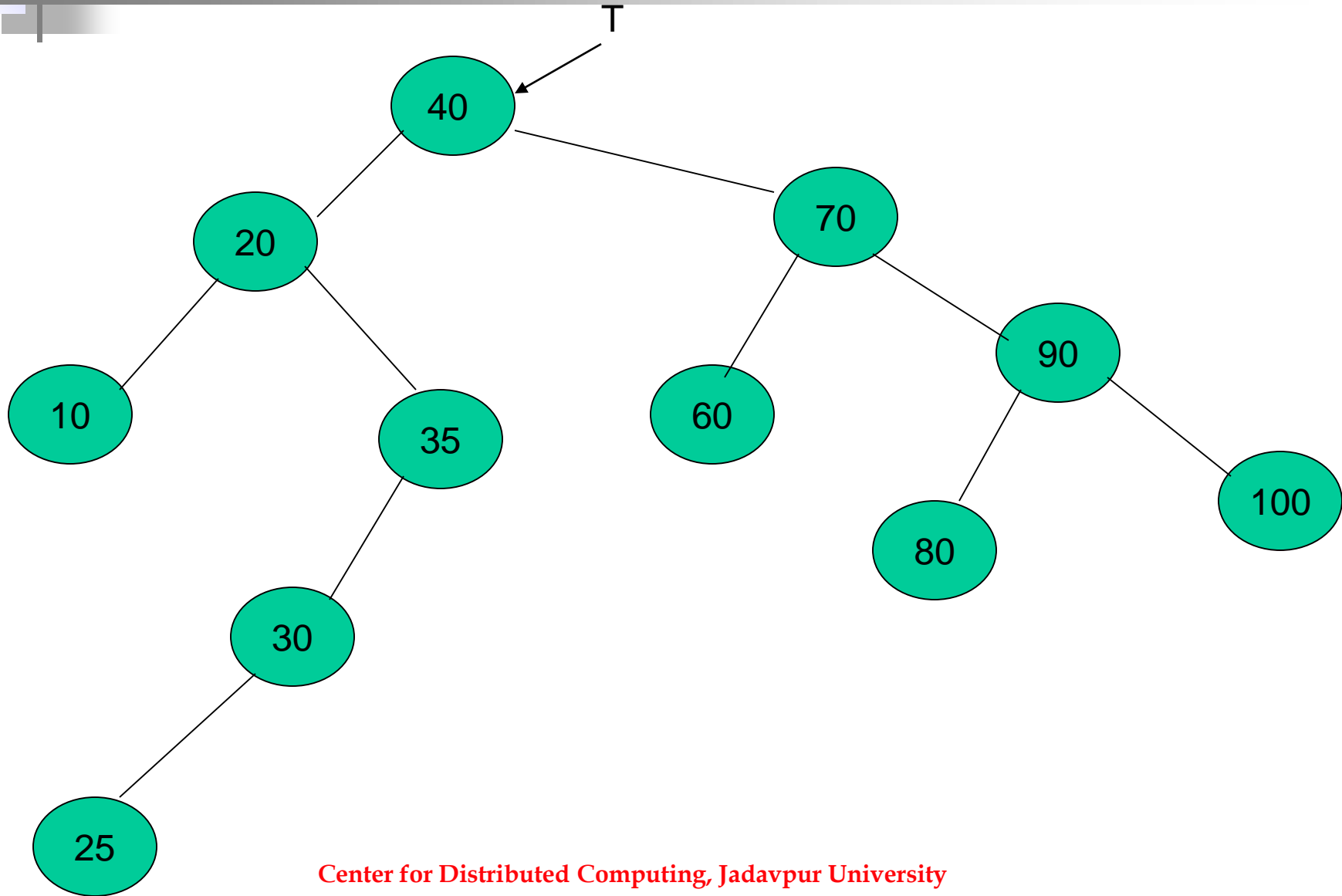


Delete 50



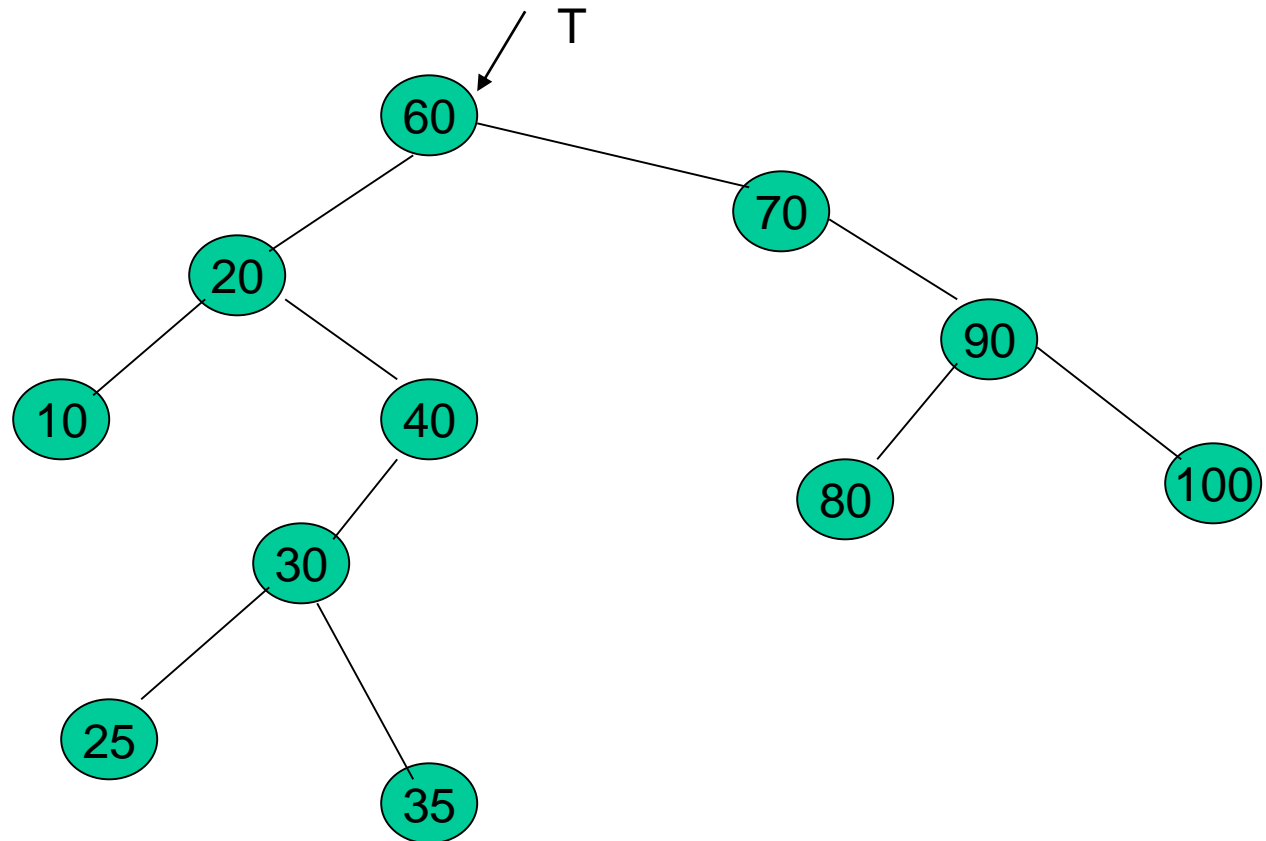


Result 1





Result 2





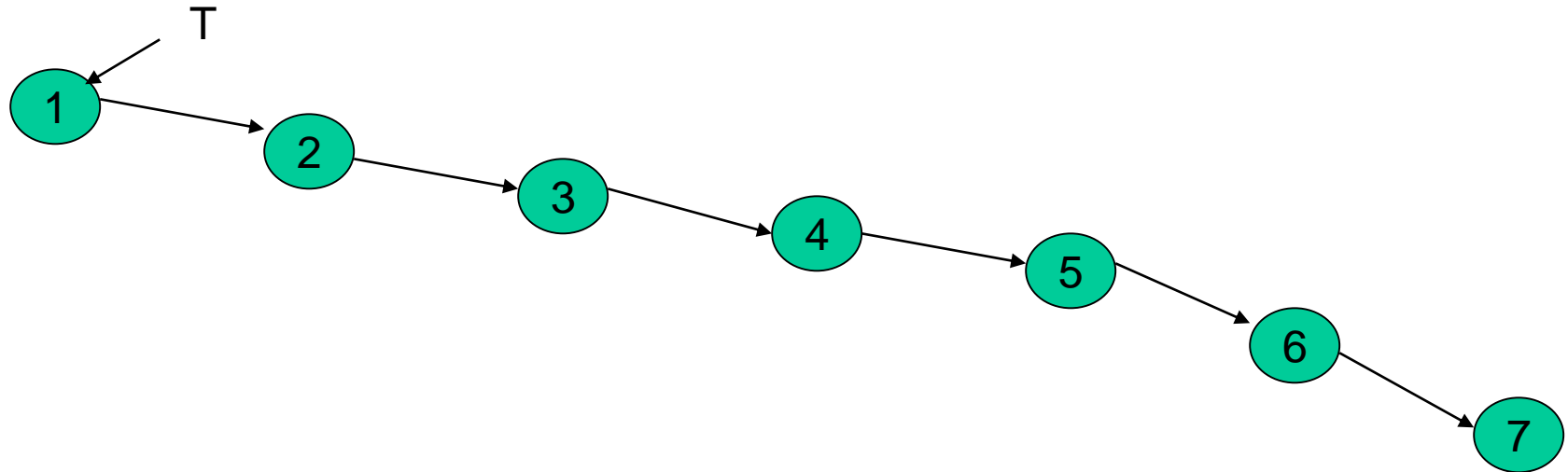
Problem of BST

- Average case complexity of search, insertion and deletion operations is $O(\log_2 n)$, where n is the no of nodes in the tree.
- The height of a BST depends on the sequence of insertion and deletion of keys.
- An extreme case:
Draw a BST for the following sequence of insertions:

1, 2, 3, 4, 5, 6, 7



Problems of BST ...



The tree degenerates into a linked list.

The worst case complexity of search, insertion and deletion are $O(n)$.

Remedy: Balanced tree.



Height Balanced Tree (AVL Tree)

- Invented by Adelson-Velskii, Landis (Russian)
- AVL tree is a BST where at each node (including the root node) the left sub-tree and the right sub-tree do not differ in height by more than one.

$$|h_L - h_R| \leq 1$$



Balance Factor

- Balance Factor (BF) of a node is the difference between the heights of its left and right sub-trees.

$$BF = h_L - h_R$$

$BF = 1$	left high
$BF = -1$	right high
$BF = 0$	equal high



AVL Tree Operations (in addition to those of Binary tree)

1. Search a key

2. Find max & Find min

Same as BST

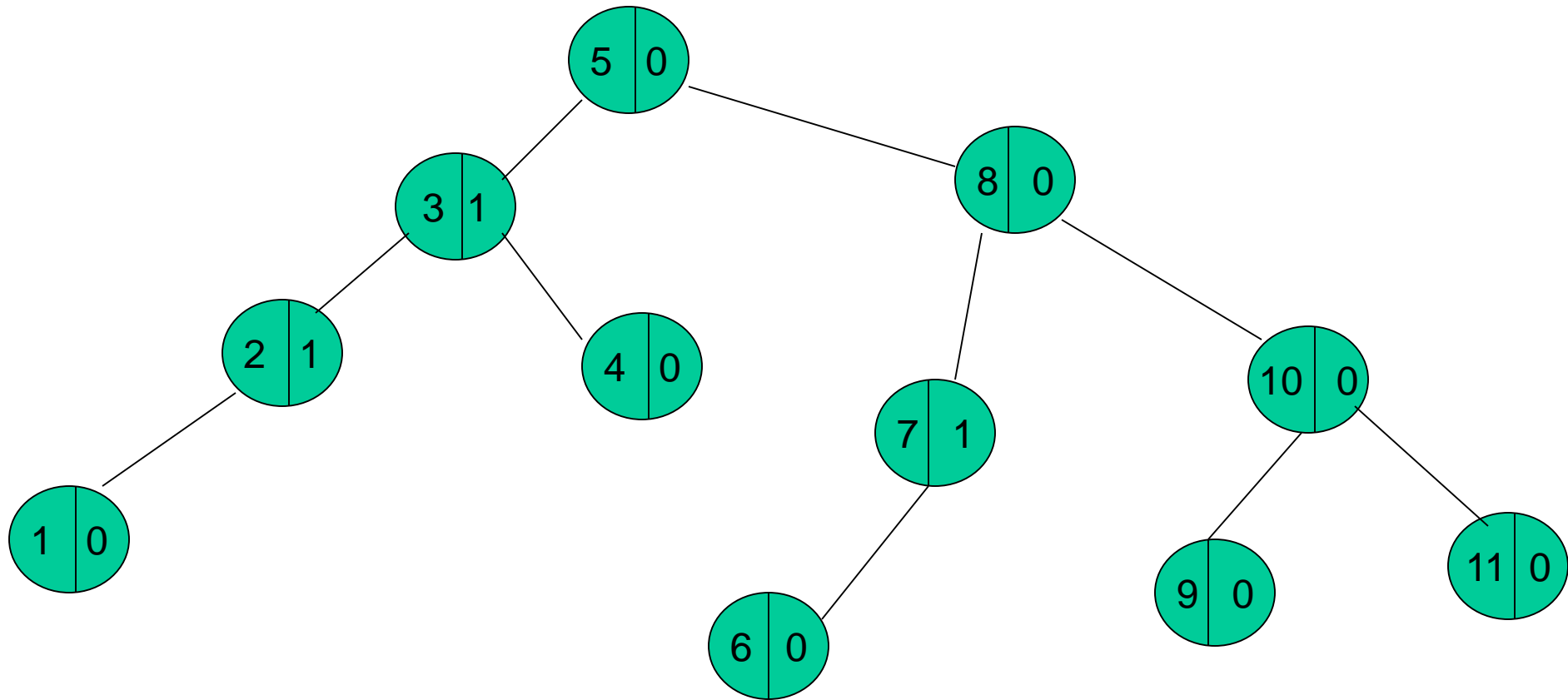
3. Insert a Key

4. Delete a Key

Insert / Delete as in BST;
then rebalance the resultant tree
if necessary

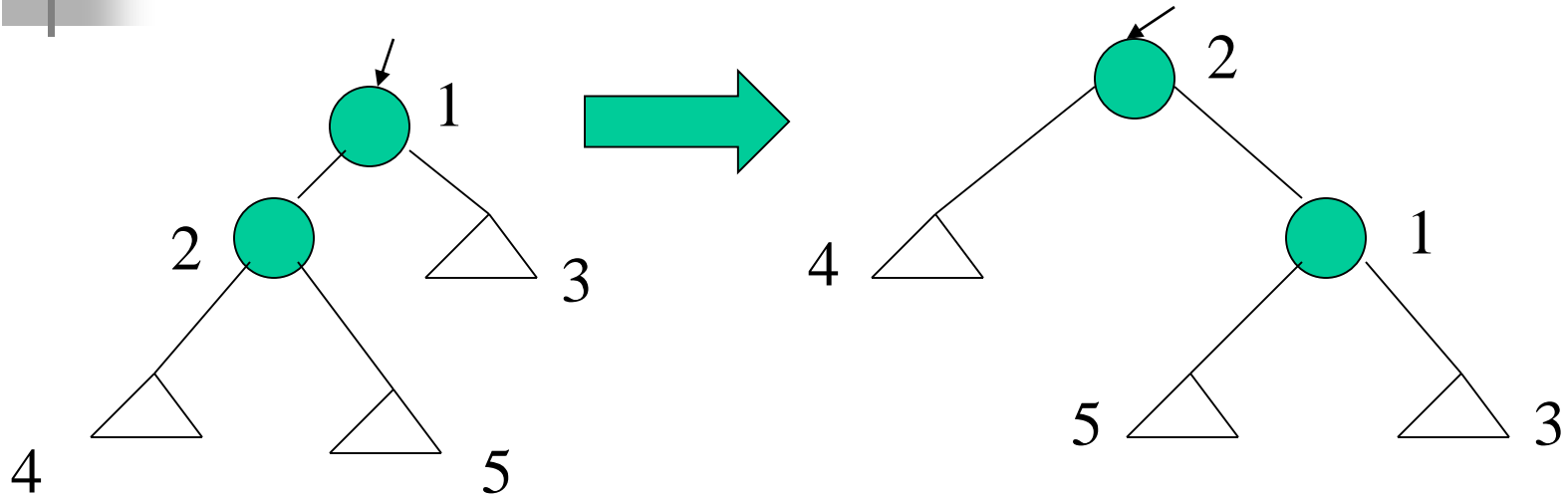


AVL Tree Example





Rebalancing needs Rotation

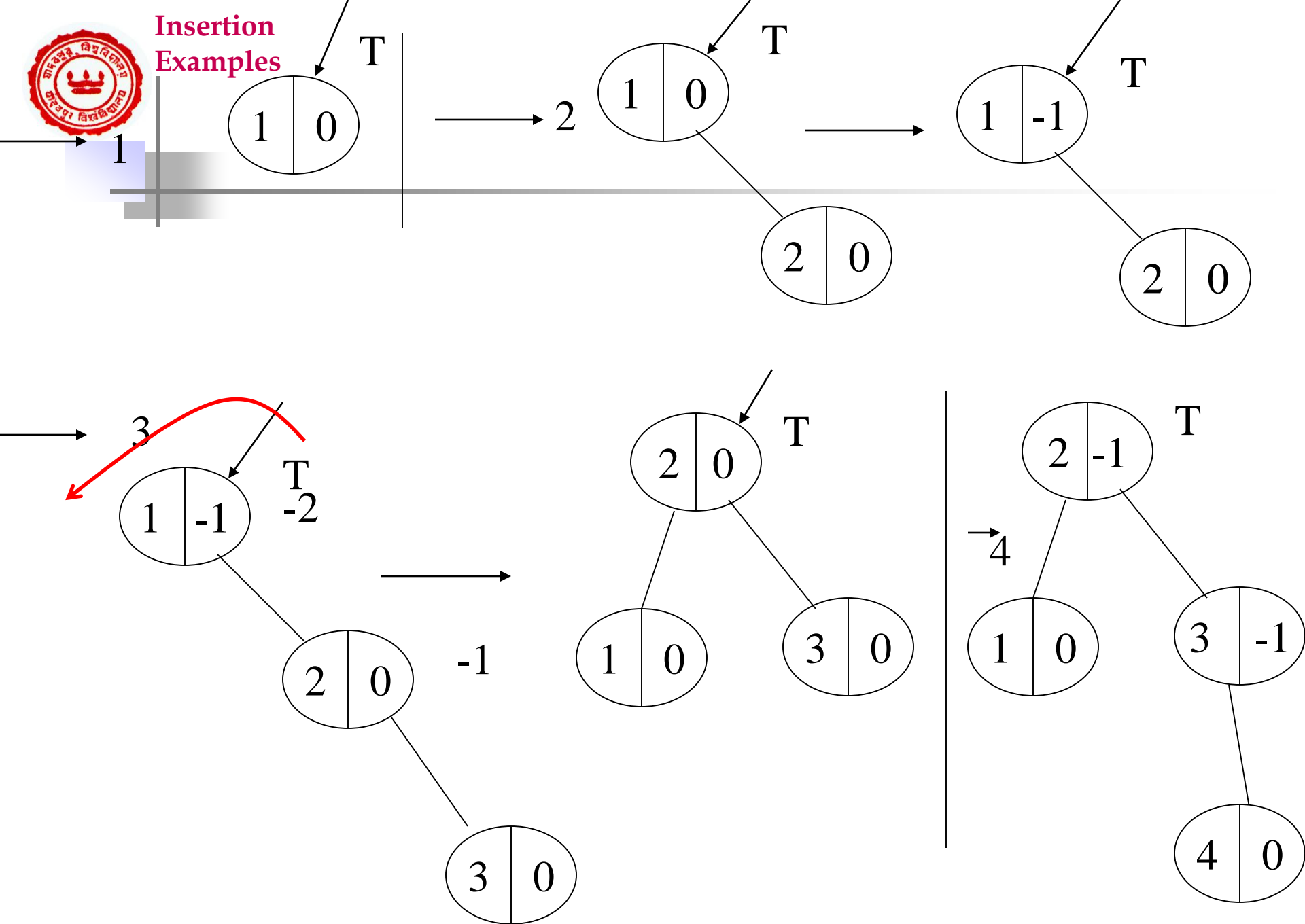




Right Rotation

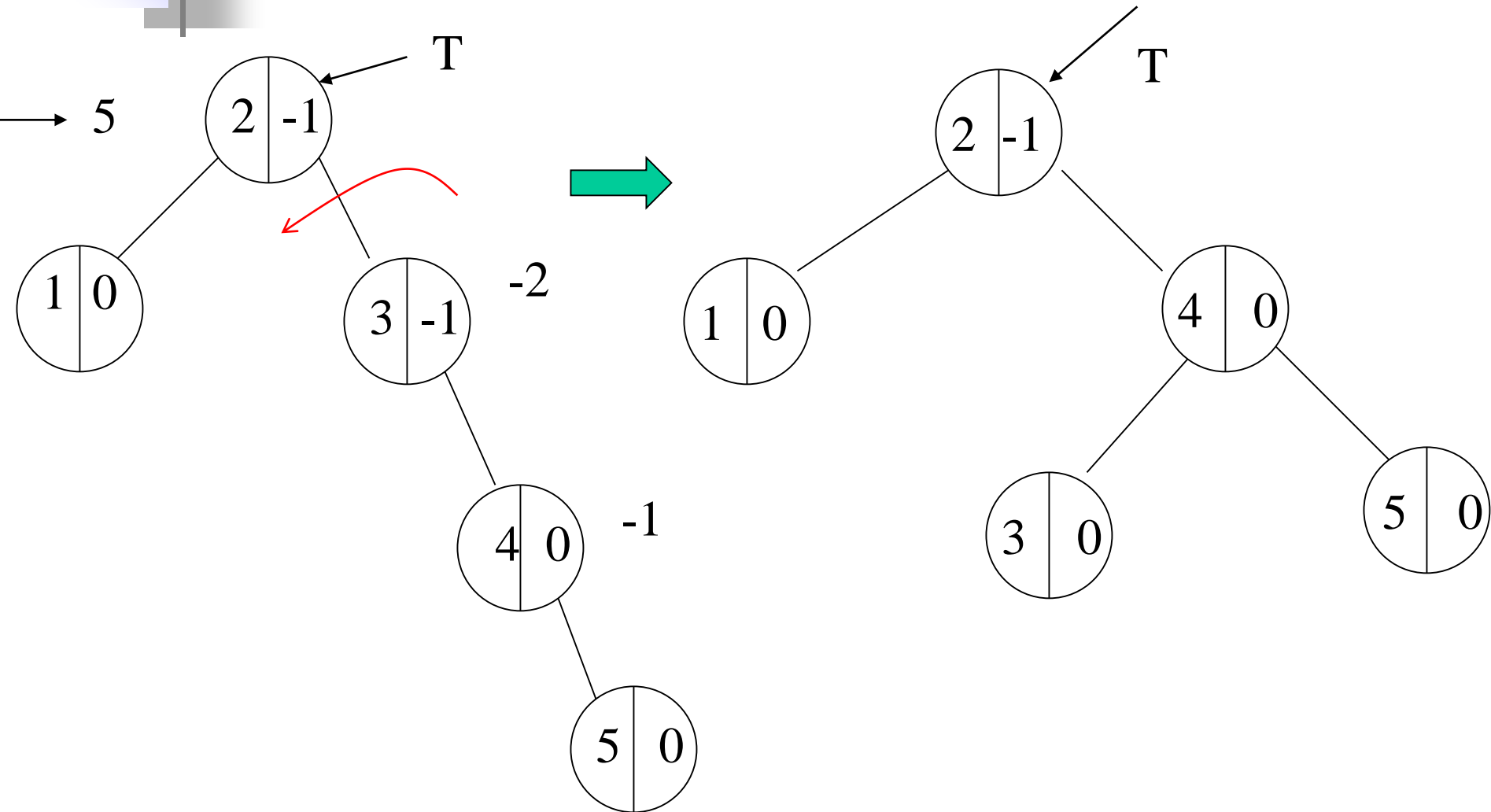
```
avltree * rotate-right (avltree * t) {  
    avltree * temp;  
    temp = t → left;  
    t → left = temp → right;  
    temp → right = t;  
    return temp;  
}
```

Insertion Examples



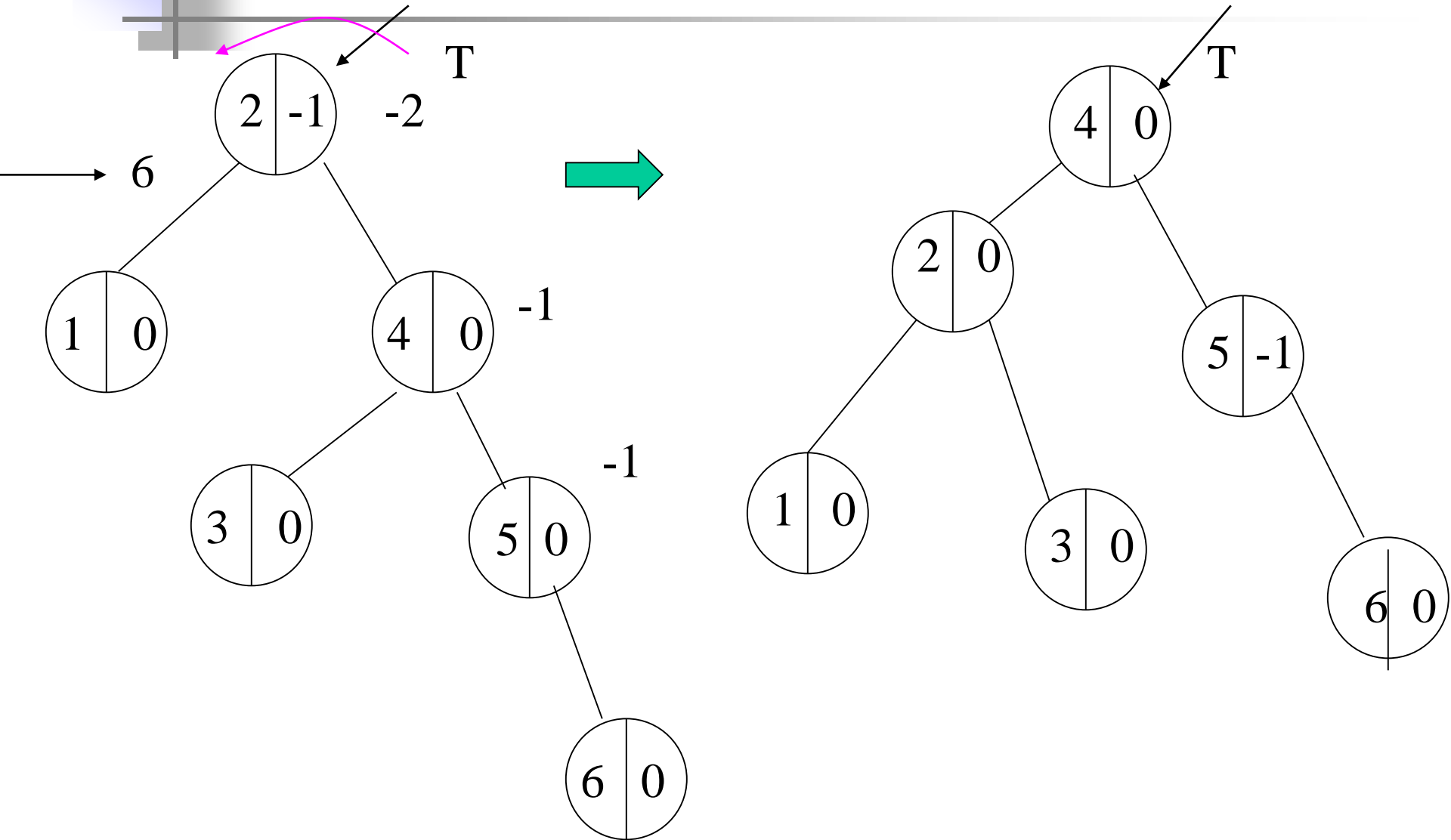


Insertion Examples ...



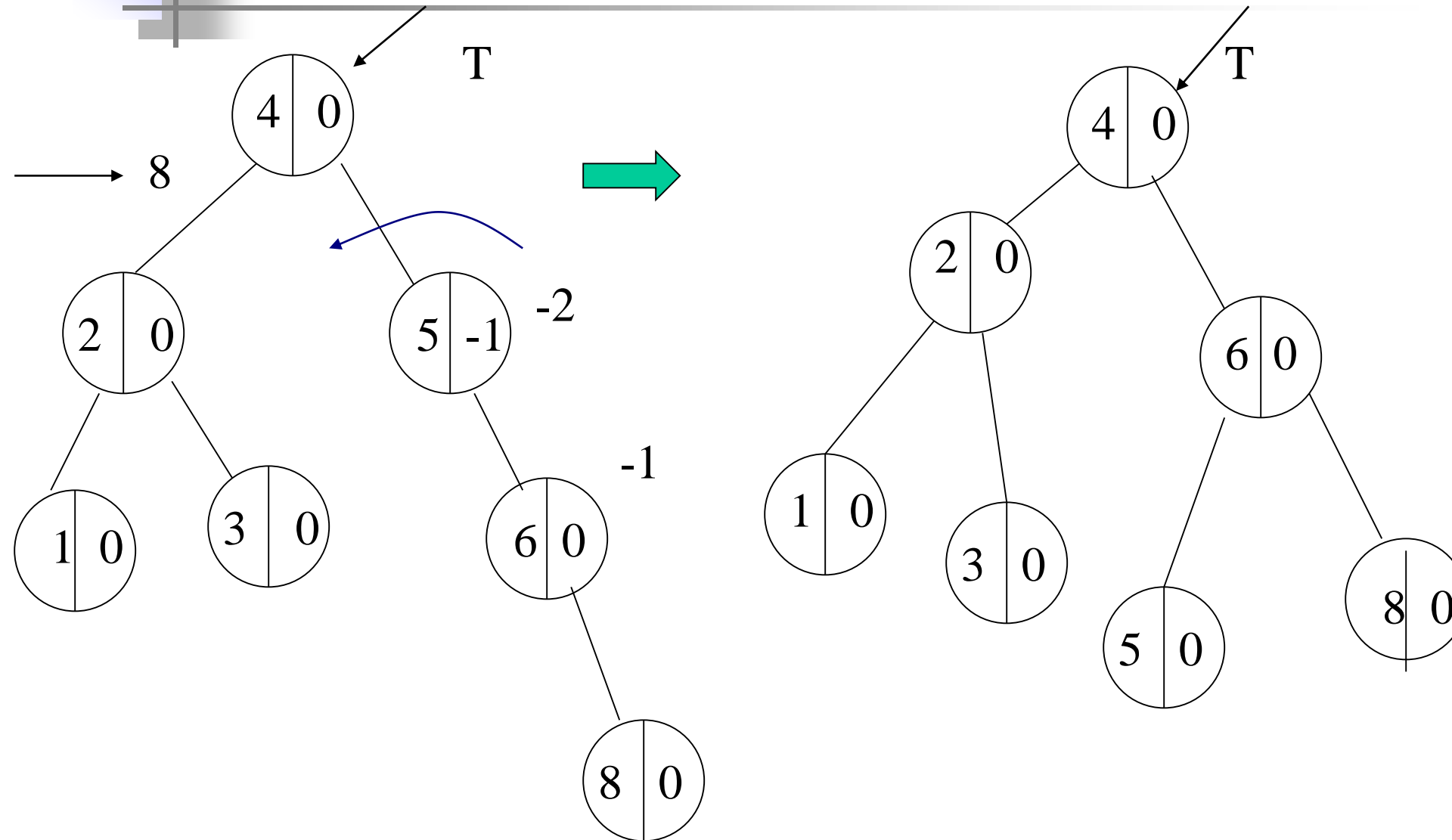


Insertion Examples ...



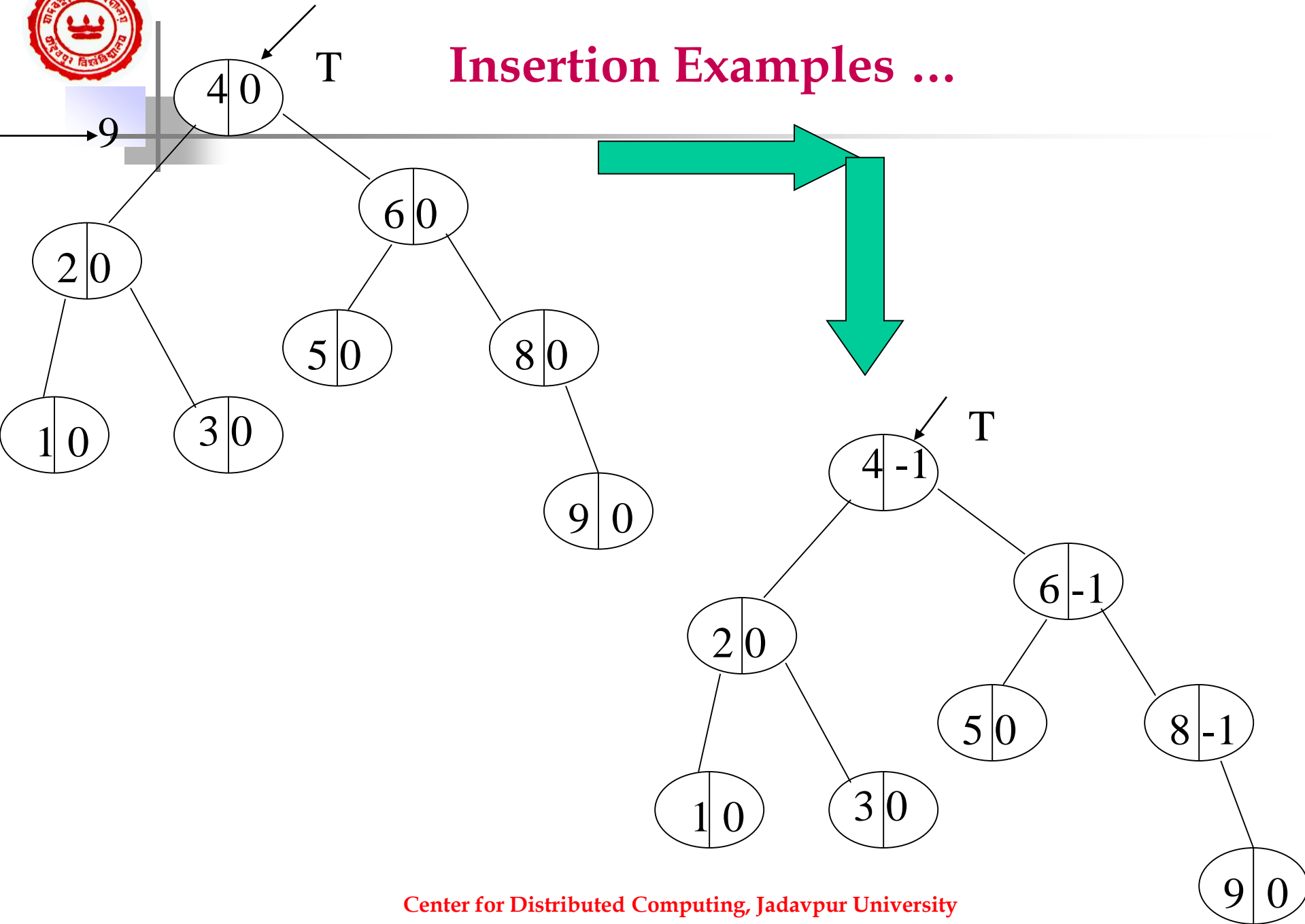


Insertion Examples ...



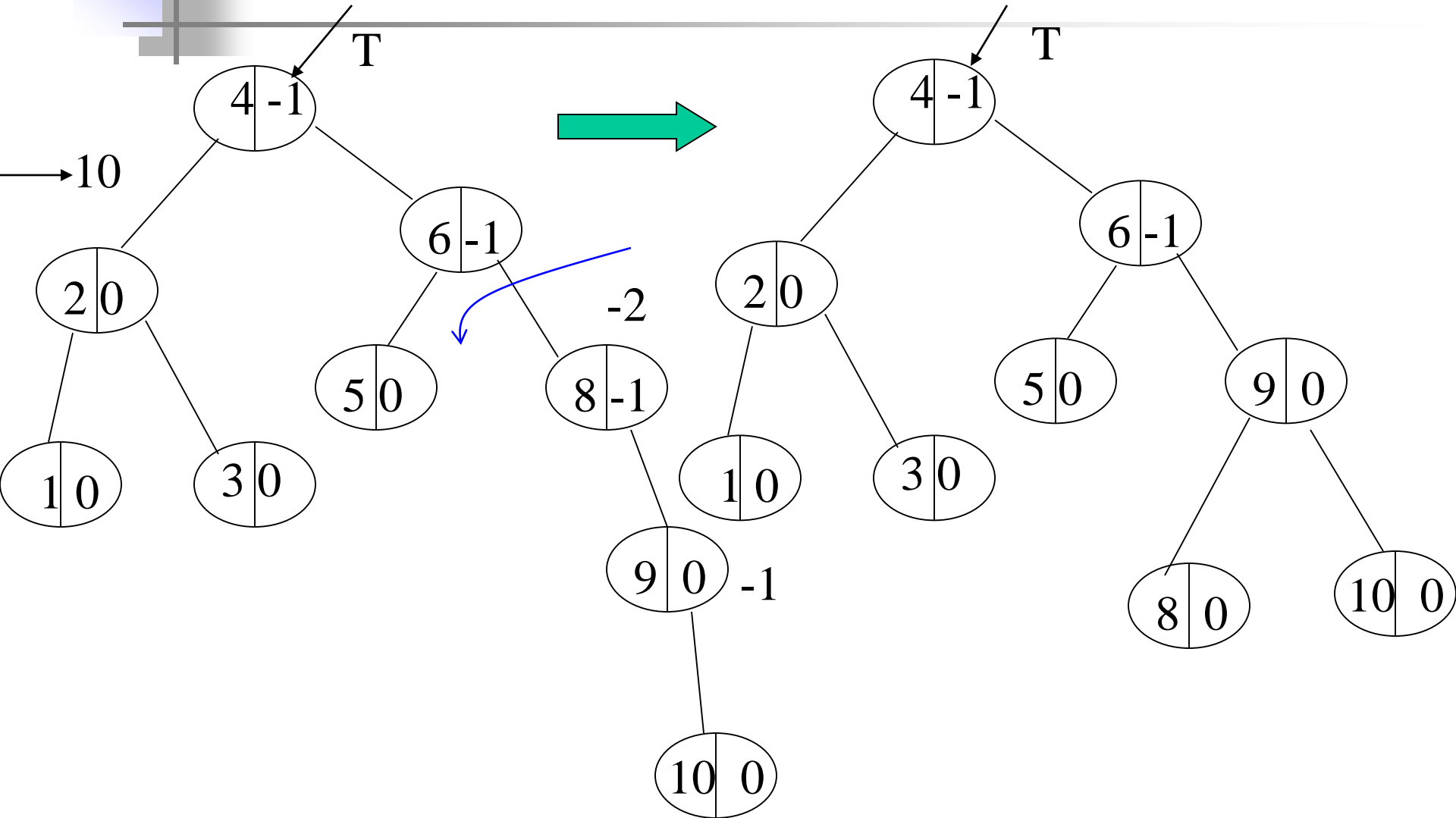


Insertion Examples ...



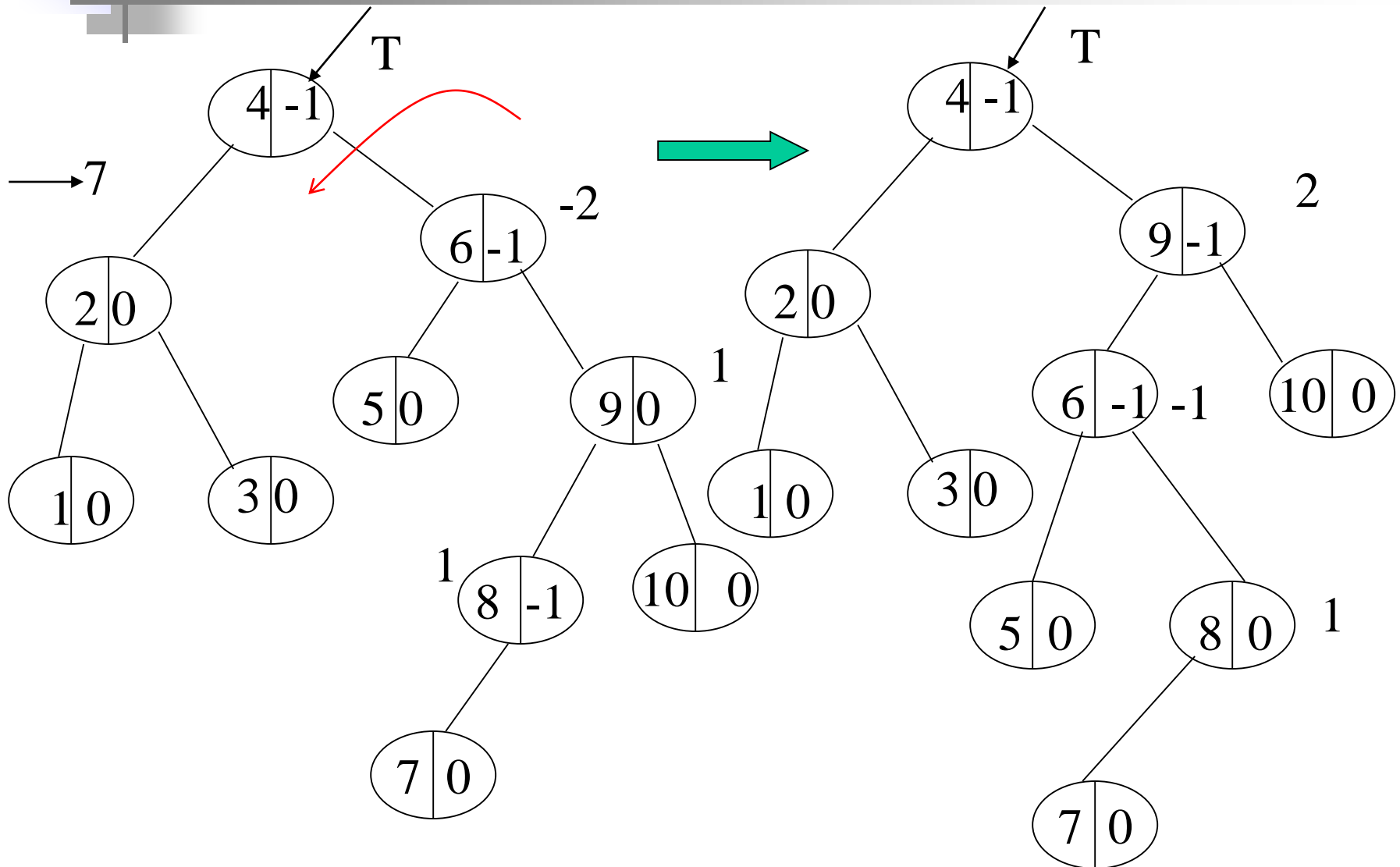


Insertion Examples ...



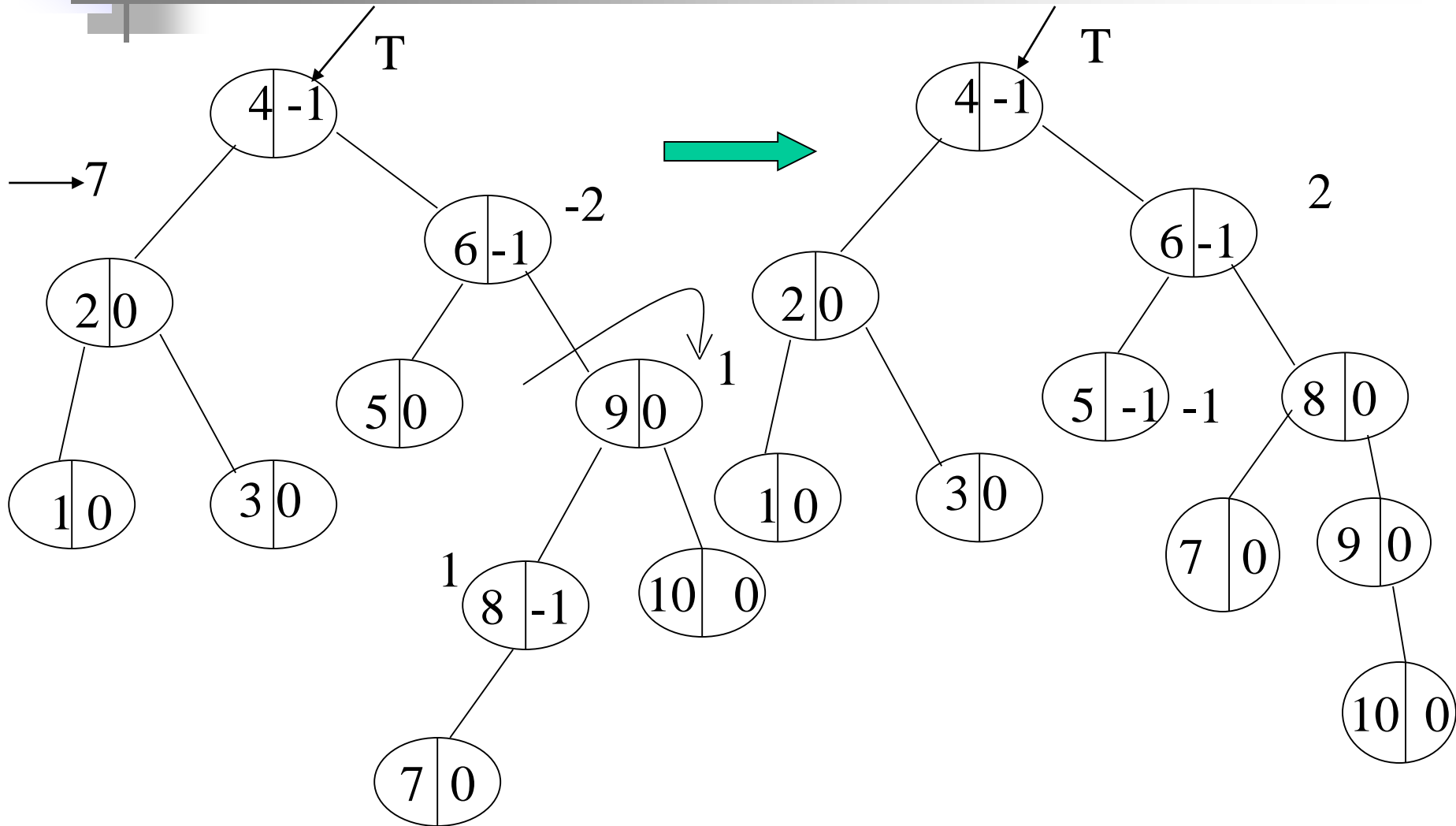


Tree remains unbalanced even after rotation



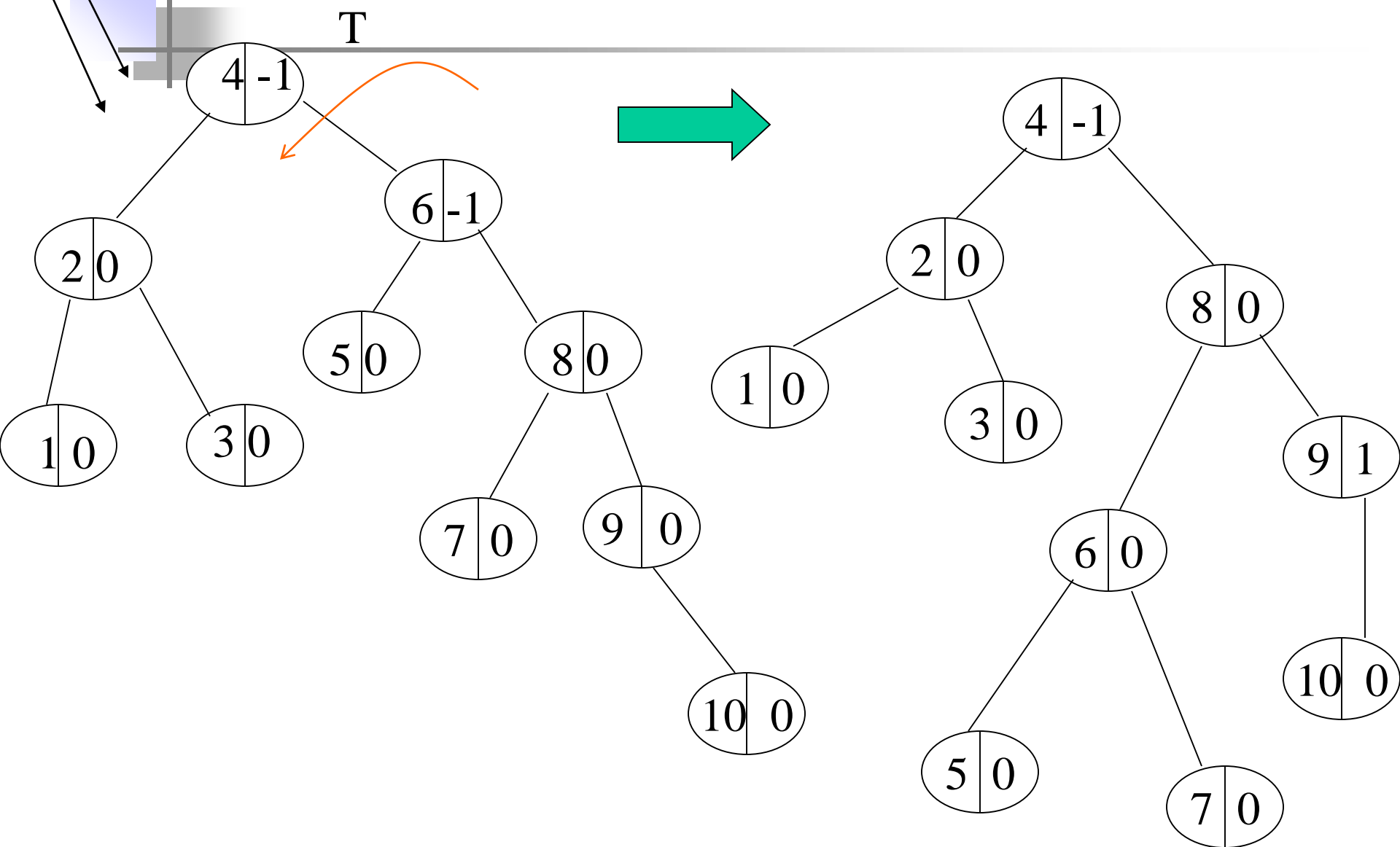


Right rotate the right child



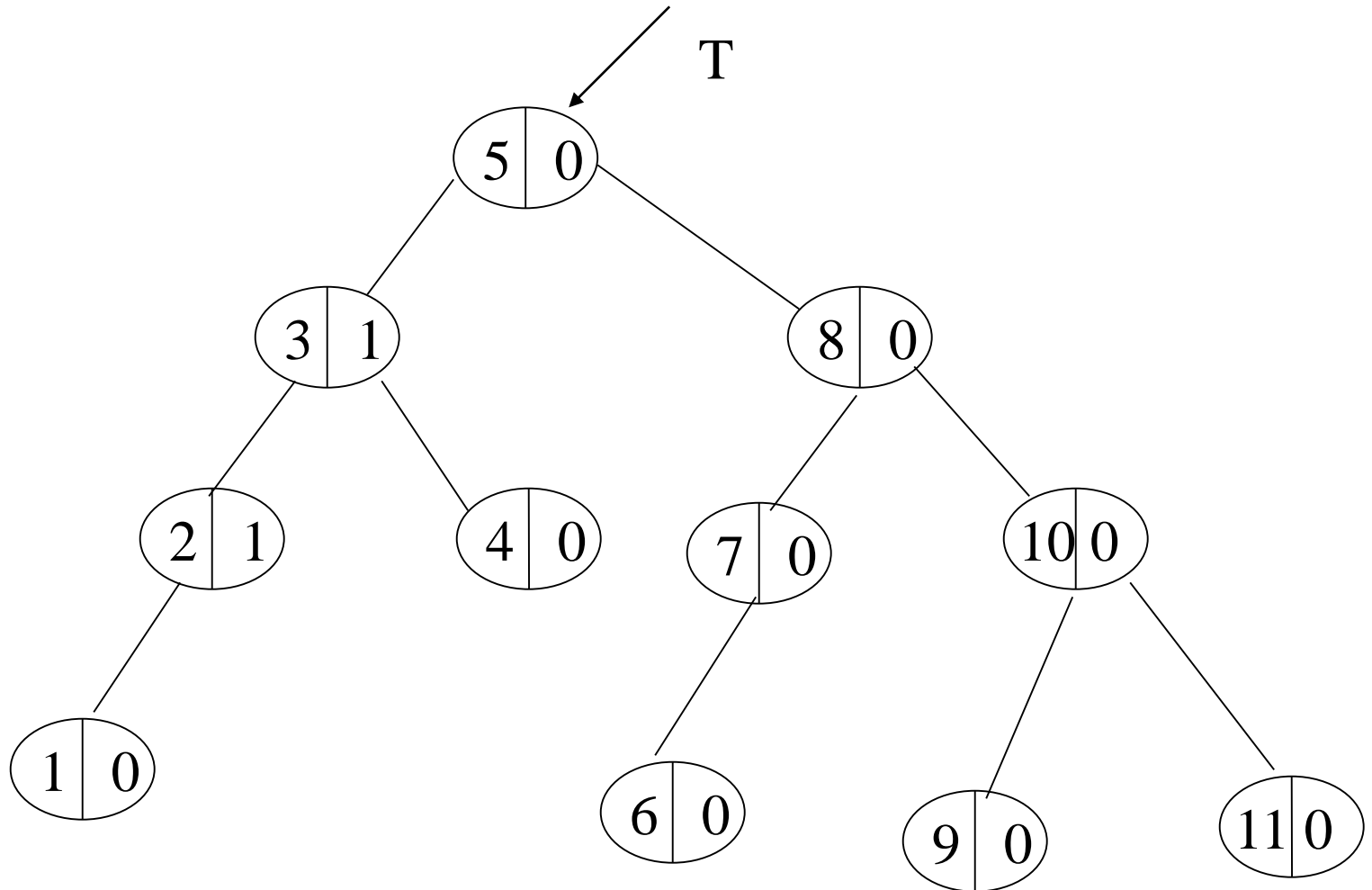


Double rotation: here, right rotation followed by left rotation



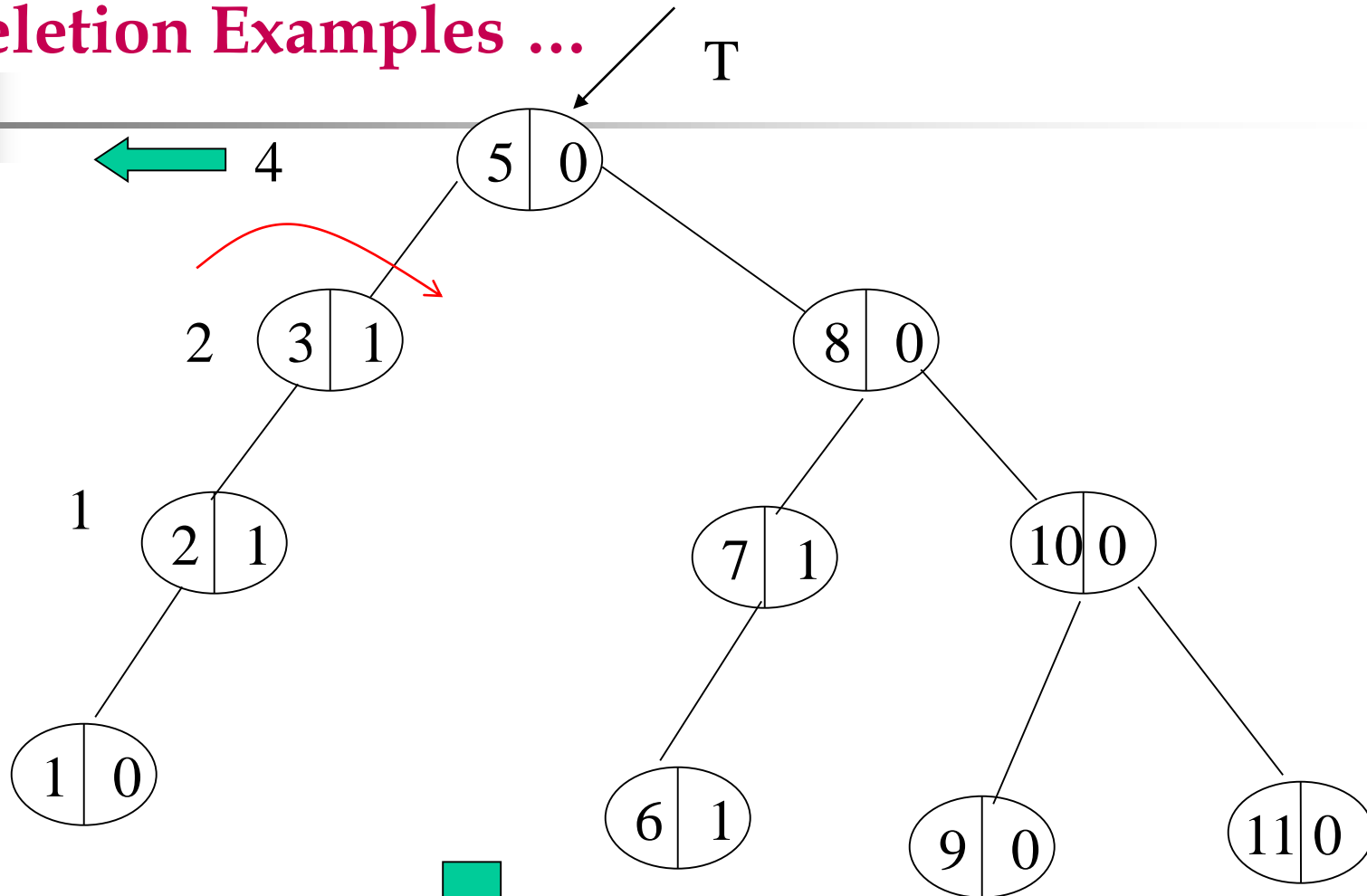


Deletion Examples



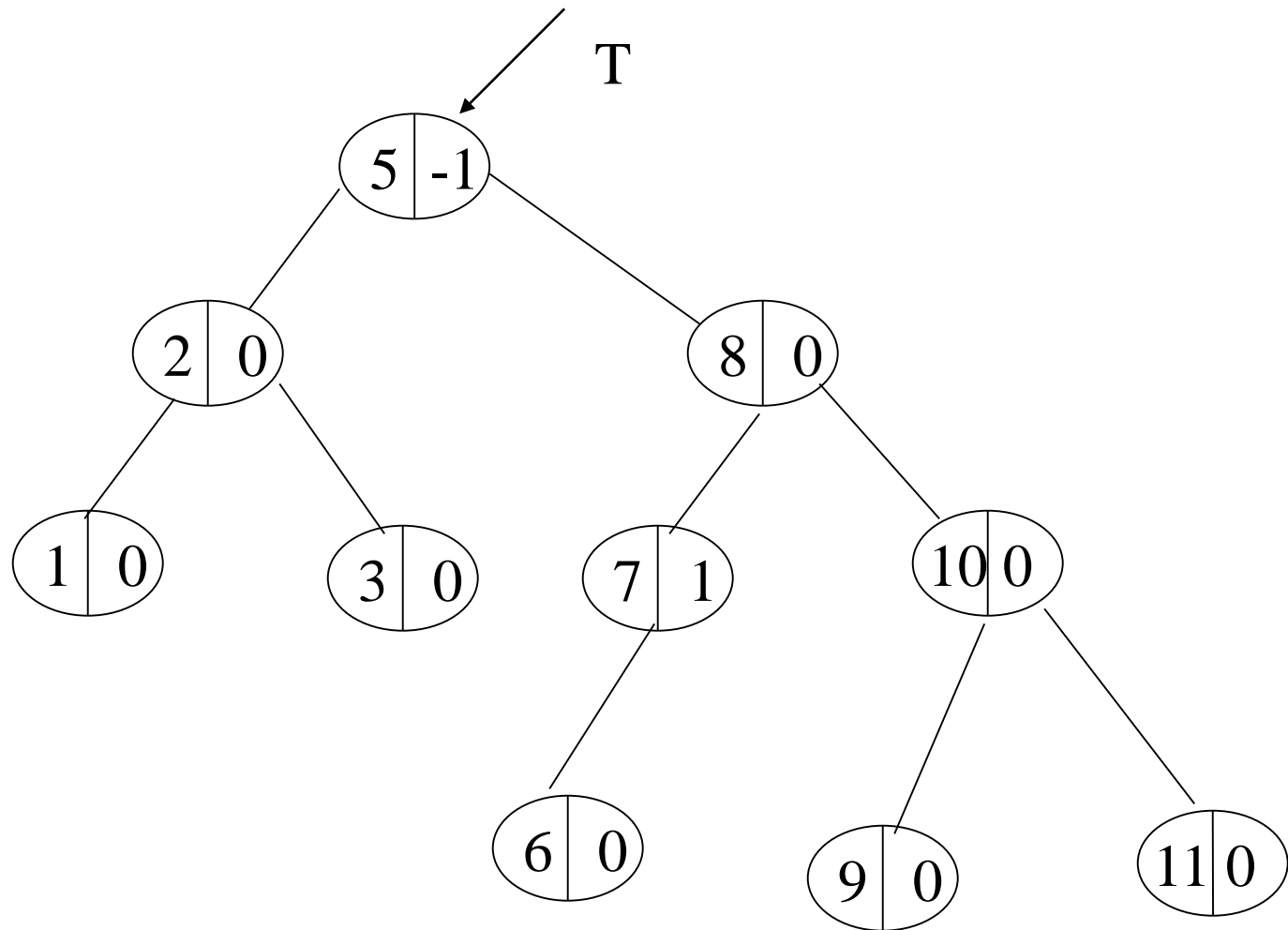


Deletion Examples ...



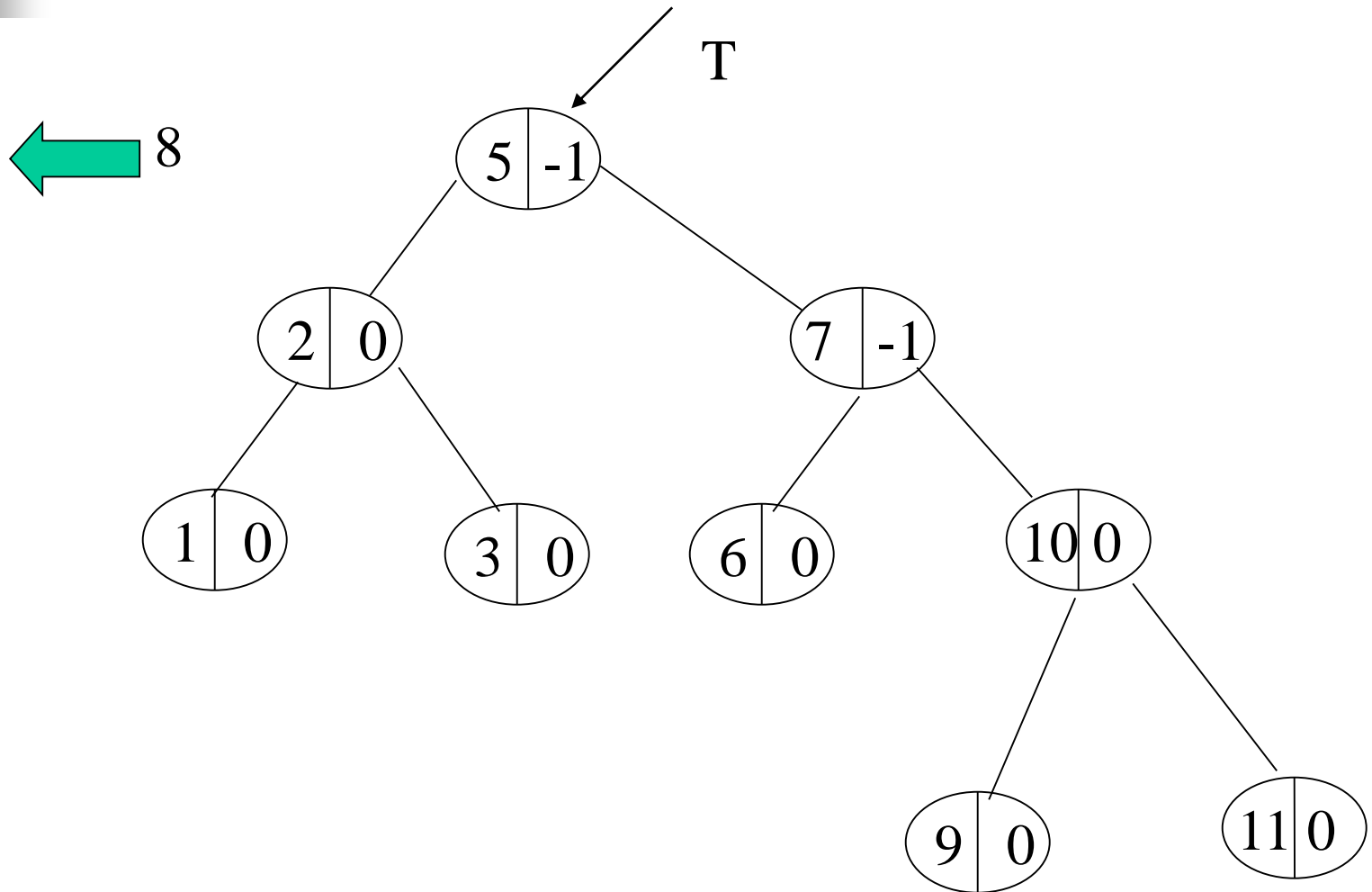


Deletion Examples ...



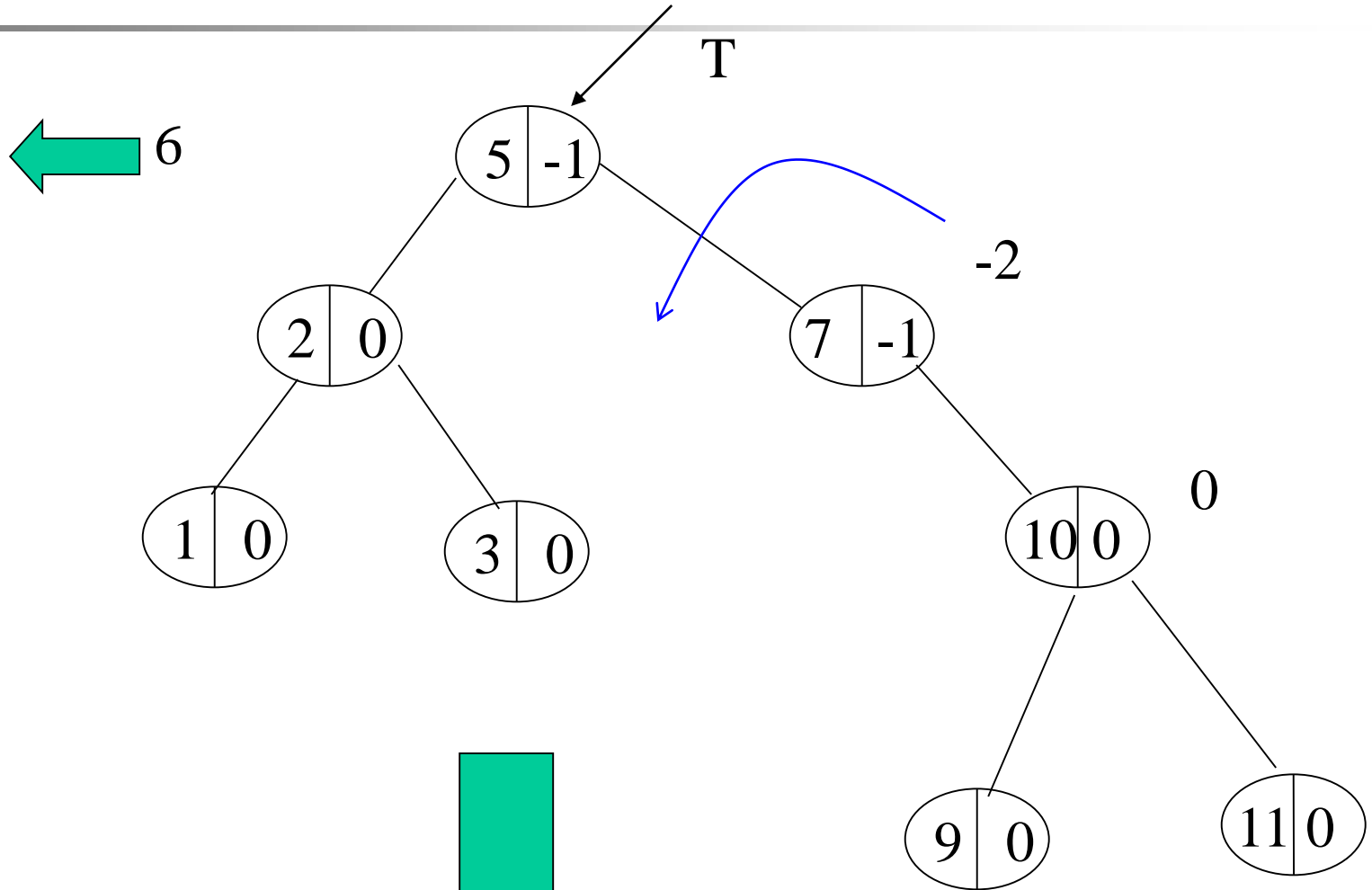


Deletion Examples ...



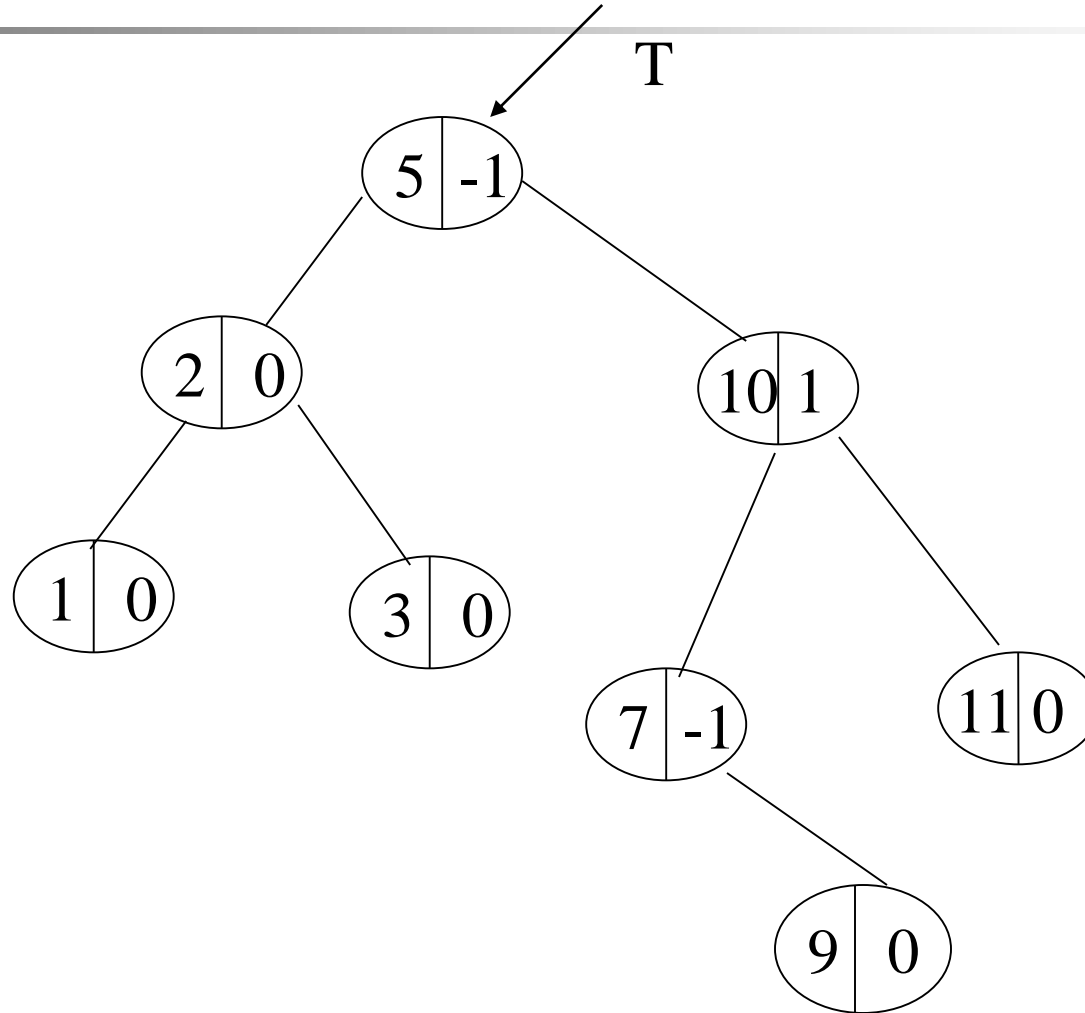


Deletion Examples ...





Deletion Examples ...





Conclusion

- Height of a height-balanced (AVL) Tree is guaranteed to be $O(\log n)$, n being the no. of nodes.
- The insertion/deletion step takes at most $O(\log n)$ time.
- Each rebalancing step, i.e., rotation (possibly double rotation) and updation of BF takes a constant amount of time.
- The rebalancing may go up to the root. Thus, there can be at most $O(\log n)$ rebalancing steps.
- Thus the overall complexity of insertion/deletion is $O(\log n)$.



Various types of trees used in other applications

- Splay Tree
- Red Black Tree
- Trie
- Quad Tree
- Octree
- ...