

15/03

File  $\rightarrow$  fixed length record

unspanned, bfr

$\rightarrow$  Block 0, 1, 2, ..., bfr - 1

$i^{th}$  record:

Block & offset

$\lfloor i/bfr \rfloor$

$i \bmod bfr$

$$71 / 50 \Rightarrow 1$$
$$71 \% 50 = 21$$

File

unordered / head

ordered

Storing records doesn't follow any ordering  
(no ordering on data)  
(Stored in the order of insertion) (APPEND)

ordered based on certain field (ordering field)

student - result

item - stock

Static

$\rightarrow$  unlikely to be changed

dynamic

$\rightarrow$  can

ordered

$\rightarrow$  ordering on key

$\rightarrow$  ordering on nonkey

$\rightarrow$  Binary search on ordered key

$\rightarrow$  Roll  $\geq 50$  & Roll  $\leq 100$

$\rightarrow$  Search 50 by bin search  
 $\rightarrow$  then sequential

update ordered val

$\rightarrow$  delete followed by insert followed by insert

Unordered file

ADV: Insertion

Retrieval  $\rightarrow$  sequential

(worst case  $\rightarrow$  read all)

$(O(n)) \rightarrow n = \text{no of blocks}$

no of accesses = n

average =  $n/2$

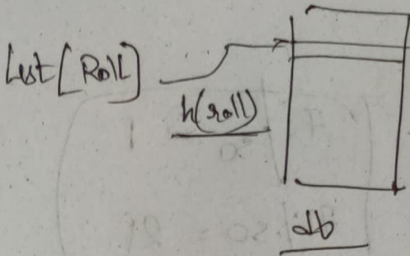
Deletion

$\rightarrow$  seq search  
 $\rightarrow$  reuse is problem

$\rightarrow$  reorganisation required.  
(difficult)

$\therefore$  logical deletion

Internal hashing → in primary memory  
table/array



Hash key

on which hash function is applied  
→ index to table

Disadv: only go for equality of field  
not  $>$ ,  $<$  etc.

Collision

→ linear probing: problem is deletion (what if the record, which collided, is deleted)

→ chaining is better: easy deletion/insertion/search.

External hashing

Disk

→ no of blocks

hash (hash key field)

↓

Bucket no.

(collection of no of blocks)

→ collision prob is less

↓  
As the bucket can hold number of record

↓  
No of hash collision  $\leq$  Bucket capacity

↓  
No prob.

Indexing: To make the search with the file faster

→ overhead: index file

→ Dense index: For every data rec, there is an entry record in index file

→ Sparse index

→ else sparse (index rec  $<$  data rec)

Index

→ ordered file (ordered on index field)

→ mostly fixed length record.

→ 2 fields: < index field value, Record/block pointer >

Insertion in index file is problem



# Data file

ordered

~~unordered~~

Indexed  
ordering

key (primary index)  
nonkey (cluster index)

sparse keys

unordered

(Datafile not ordered on  
index field)

Secondary  
index

→ key  
→ nonkey

Dense sometimes (mostly)  
Sparse rarely