

# The Combination Lemma for curves

①

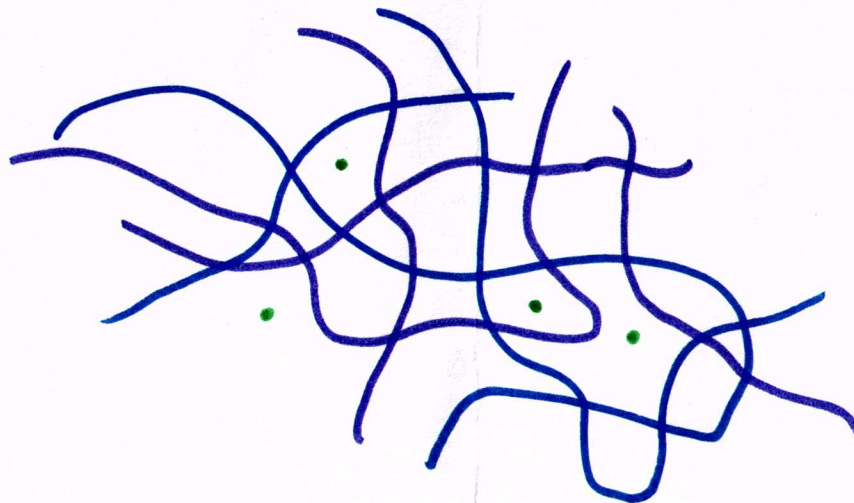
$T_1, T_2$  - sets of Jordan arcs.

$P, |P|=k$  - marking points.

Let  $A(T_1), A(T_2)$  have complexities  $r$  &  $b$  for marked faces in  $A(T_1)$  and  $A(T_2)$ , respectively.

$A(T_1 \cup T_2)$  has complexity  $O(r+b+k)$  for marked faces in  $A(T_1 \cup T_2)$ .

GSS, DCG, 4:491-521,  
Swihart et. al. 1989



## Applications.

Use recursion for estimating "many faces" complexity in an arrangement of Jordan arcs.

③'

$$\text{So, } (s-1) \binom{m}{r} \geq n \binom{\lceil E(G_{m,n})/n \rceil}{r}$$

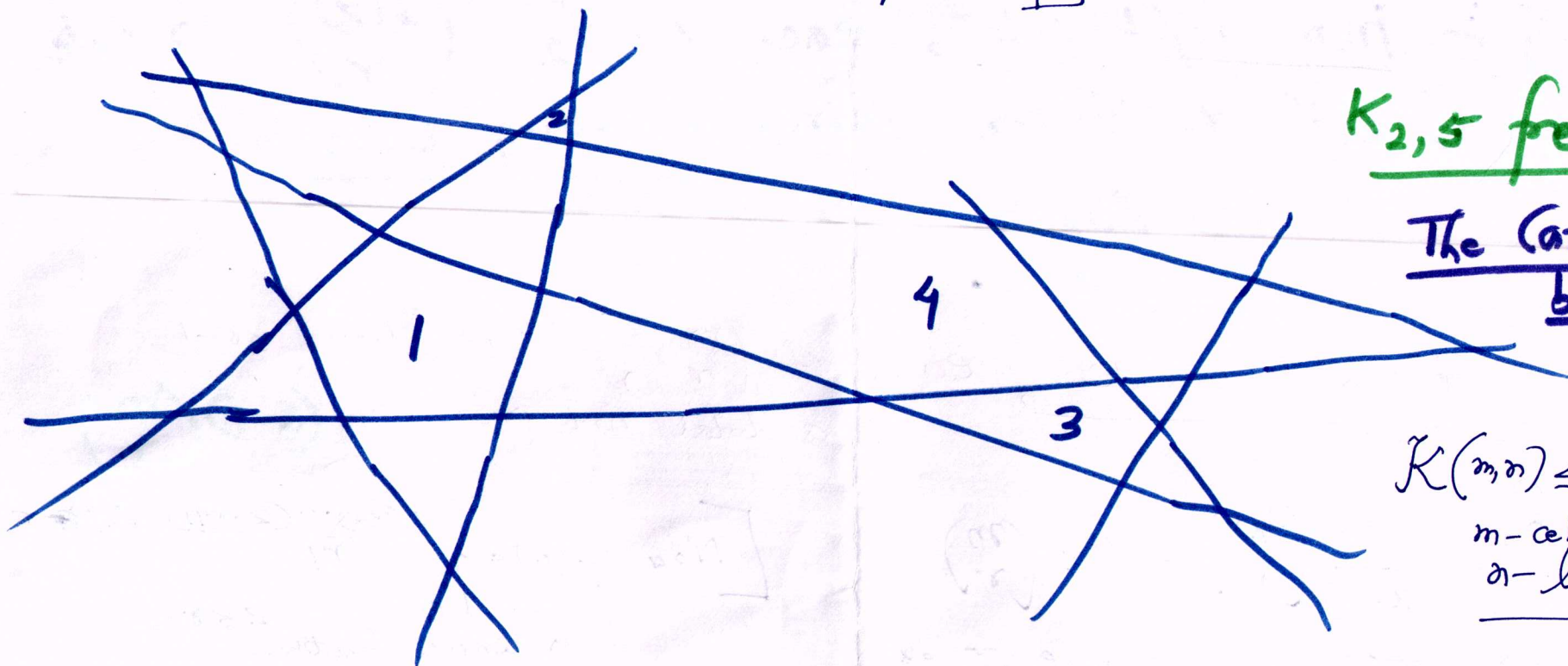
$$\Rightarrow (s-1) m^r \geq n \left( \frac{E(G_{m,n})}{n} - r + 1 \right)^r$$

$$\left[ \begin{array}{c} f(G_{m,n}/n) \\ \leq \end{array} \binom{\lceil G_{m,n} \rceil/n}{r} \right]$$

$$\Rightarrow (s-1)^{1/r} m n^{1-\frac{1}{r}} + (r-1)n \geq |E(G_{m,n})|$$

[denominators  $r!$  get balanced]

□



$K_{2,5}$  free

The Canham  
bound  
(1969)

$$K(m,n) \leq c(m\sqrt{n} + n)$$

$m$  - cells  
 $n$  - lines



# Kővári, Sós, Turán (1954) & Erdős [3]

$G_{m,n}$  bipartite  $m \times n$  vertices, no  $K_{r,s}$  as a subgraph

$$|E(G_{m,n})| \leq c_{r,s} (mn^{1-\frac{1}{r}} + n)$$

$c_{r,s}$  - constant, function of only  $r$  &  $s$ , not  $m$  &  $n$ .

Proof: From right side, each  $x \in V_2$   $\binom{d(x)}{r}$   $r$ -tuple  $W$  with  $x$   $(W, x)$  pairs, totalling exactly

$$\sum_{x \in V_2} \binom{d(x)}{r}.$$

From left side, each of  $\binom{m}{r}$ ,  $(W, x)$  pairs  $\leq s-1$ ,  
so a total not exceeding  $(s-1) \binom{m}{r}$

$$\text{so, } \sum_{x \in V_2} \binom{d(x)}{r} \leq (s-1) \binom{m}{r}$$

$$\left[ \text{Now } f(z) = \begin{cases} \frac{z(z-1)\cdots(z-r+1)}{r!} & z \geq r \\ 0 & z \leq r-1 \end{cases} \right]$$

is a convex function.

$$\Rightarrow \sum_{x \in V_2} f(d(x)) \geq n f\left(\frac{\sum d(x)}{n}\right) \rightarrow \boxed{\text{Jensen}}$$

$$= n f\left(\frac{|E(G_{m,n})|}{n}\right)$$

## Crossing number theorem

⑤

$G(V, E)$ , no multiple edges.

$$cr(G) \geq \frac{1}{6} \cdot \frac{|E|^3}{|V|^2} - |V|.$$

$$cr(K_5) = 1. \quad cr(K_4) = 0.$$

For non-planar graphs  $G$ ,  $cr(G) > 0$ .

$\Rightarrow$   $cr(G) \geq |E| - 3|V|$ . If  $|E| > 3|V|$  and we had  $< |E| - 3|V|$  crossings then for each crossing drop an edge (at most  $|E| - 3|V|$  edges dropped), still having a planar graph with  $|E| > 3|V|$  edges\*, a Contradiction  $\square$

\* A graph with more than  $3|V| - 6$  edges has  $cr(G) > 0$ , non-planar