

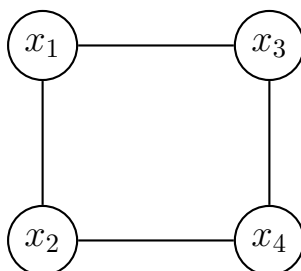
CS 446: Machine Learning

Homework 7

Due on Tuesday, March 6, 2018, 11:59 a.m. Central Time

1. [4 points] Inference Methods for Discrete Markov Random Fields

For this problem, consider the following Markov Random Field, where each node can be assigned one of the values in $\{1, 2, 3, 4, 5\}$:



- (a) To conduct MAP inference on this graph using exhaustive search, how many configurations must be checked?

Your answer:

5^4

- (b) Can MAP inference be run on this graph using a dynamic programming algorithm? Why or why not?

Your answer: No, because the MRF has cycles in it, e.g. it is not a tree.

- (c) To run MAP inference on this graph using loopy belief propagation, how many messages must be computed?

Your answer: Let's start by noticing that

$$\mathcal{R} = \{\{1\}, \{2\}, \{3\}, \{4\}, \{1, 2\}, \{1, 3\}, \{3, 4\}, \{2, 4\}\}$$

Since the number of messages is the number of Lagrangian multipliers and we have one multiplier for each of the constraints $\sum_{\mathbf{y}_p \mathbf{y}_r} b_p(\mathbf{y}_b) = b_r(\mathbf{y}_r)$.

Since $P(r \in \{\{1, 2\}, \{1, 3\}, \{3, 4\}, \{2, 4\}\}) = \emptyset$ only the other values of r will generate a constraint.

$$\mathbf{y}_r \in \{1, 2, 3, 4, 5\} \forall r \in \{\{1\}, \{2\}, \{3\}, \{4\}\}$$

Finally, note that there are two parents for each of r , that is $|P(r)| = 2$.

Then there are $4 \times 5 \times 2 = 40$ constraints and, therefore, 40 messages.

2. [7 points] ILP Inference formulation in Discrete Markov Random Fields

- (a) Suppose we have two variables $x_1 \in \{0, 1\}$ and $x_2 \in \{0, 1\}$ and their local evidence functions $\theta_1(x_1)$ and $\theta_2(x_2)$ as well as a pairwise function $\theta_{1,2}(x_1, x_2)$. Using this setup, inference solves $\arg \max_{x_1, x_2} \theta_1(x_1) + \theta_2(x_2) + \theta_{1,2}(x_1, x_2)$. Using

$$\theta_1(x_1) = \begin{cases} 1 & \text{if } x_1 = 0 \\ 2 & \text{otherwise} \end{cases} \quad \theta_2(x_2) = \begin{cases} 1 & \text{if } x_2 = 0 \\ 2 & \text{otherwise} \end{cases}$$
$$\theta_{1,2}(x_1, x_2) = \begin{cases} -1 & \text{otherwise} \\ 2 & \text{if } x_1 = 0 \text{ \& } x_2 = 1 \end{cases}$$

what is the integer linear programming formulation of the inference task? Make the cost function and constraints explicit for the given problem, i.e., do not use a general formulation.

Your answer: Lets first find the function of b_r that we are trying to maximize.

$$\begin{aligned}
\begin{bmatrix} b_1(0) \\ b_1(1) \\ b_2(0) \\ b_2(1) \\ b_{1,2}(0,0) \\ b_{1,2}(0,1) \\ b_{1,2}(1,0) \\ b_{1,2}(1,1) \end{bmatrix}^\top \begin{bmatrix} \theta_1(0) \\ \theta_1(1) \\ \theta_2(0) \\ \theta_2(1) \\ \theta_{1,2}(0,0) \\ \theta_{1,2}(0,1) \\ \theta_{1,2}(1,0) \\ \theta_{1,2}(1,1) \end{bmatrix} &= \begin{bmatrix} b_1(0) \\ b_1(1) \\ b_2(0) \\ b_2(1) \\ b_{1,2}(0,0) \\ b_{1,2}(0,1) \\ b_{1,2}(1,0) \\ b_{1,2}(1,1) \end{bmatrix}^\top \begin{bmatrix} 1 \\ 2 \\ 1 \\ 2 \\ -1 \\ 2 \\ -1 \\ -1 \end{bmatrix} \\
&= b_1(0) + 2b_1(1) + b_2(0) + 2b_2(1) - \\
&\quad b_{1,2}(0,0) + 2b_{1,2}(0,1) - b_{1,2}(1,0) - b_{1,2}(1,1) \\
&:= A_b
\end{aligned}$$

Then the inference task can be rewritten as

$$\max_b A_b \quad \text{s.t.} \quad \begin{cases} b_r(\mathbf{y}_r) \geq 0 & \forall r, \mathbf{y}_r \\ \sum_{\mathbf{y}_r} b_r(\mathbf{y}_r) = 1 & \forall r \\ \sum_{\mathbf{y}_p/\mathbf{y}_r} b_p(\mathbf{y}_p) = b_r(\mathbf{y}_r) & \forall r, p \in P(r), \mathbf{y}_r \end{cases}$$

Now let's write the constraints explicitly.

$$\max_b A_b \quad \text{s.t.} \quad \begin{cases} b_1(0) \geq 0 \\ b_1(1) \geq 0 \\ b_2(0) \geq 0 \\ b_2(1) \geq 0 \\ b_{1,2}(0,0) \geq 0 \\ b_{1,2}(0,1) \geq 0 \\ b_{1,2}(1,0) \geq 0 \\ b_{1,2}(1,1) \geq 0 \\ b_1(0) + b_1(1) = 1 \\ b_2(0) + b_2(1) = 1 \\ b_{1,2}(0,0) + b_{1,2}(0,1) + b_{1,2}(1,0) + b_{1,2}(1,1) = 1 \\ b_{1,2}(0,0) + b_{1,2}(0,1) = b_1(0) \\ b_{1,2}(1,0) + b_{1,2}(1,1) = b_1(1) \\ b_{1,2}(0,0) + b_{1,2}(1,0) = b_2(0) \\ b_{1,2}(0,1) + b_{1,2}(1,1) = b_2(1) \end{cases}$$

(b) What is the solution (value and argument) to the program in part (a).

Your answer: Note that the new cost function is

$$b_1(0) + 2b_1(1) + b_2(0) + 2b_2(1) - b_{1,2}(0,0) + 2b_{1,2}(0,1) - b_{1,2}(1,0) - b_{1,2}(1,1)$$

By simply looking at the previous item we can see that the maximum value is achieved when letting $b_{1,2}(0,1) = 1$. Then the solution is The solution is:

$$b_1(0) = 1$$

$$b_1(1) = 0$$

$$b_2(0) = 0$$

$$b_2(1) = 1$$

$$b_{1,2}(0,0) = 0$$

$$b_{1,2}(0,1) = 1$$

$$b_{1,2}(1,0) = 0$$

$$b_{1,2}(1,1) = 0$$

And then the value of our cost function is:

$$1 + 2(0) + 0 + 2(1) - 0 + 2(1) - 0 - 0 = 1 + 2 + 2 = 5$$

- (c) Why do we typically not use the integer linear program for reasonably sized MRFs?

Your answer: Because it becomes very slow for large programs. The reason being that it still has to do exhaustive search on the individual restrictions' arguments. e.g. $b_r(\mathbf{y}_r)$. Which, in large programs, may still be too large. Either $|r|$ is too large for some r or there are too many restrictions, $|\mathcal{R}|$, is too large.