#### Part 1:

Before either banded/unrestricted are called, there is a cost of slicing each string to make it the size of the align length, but since that is a constant, it can be dropped from the overall time and space complexity.

For the unrestricted algorithm, the space complexity is simply O(nm) where n is the length of one string and m is the length of the other string. Since each cell of the table stores a tuple with an integer value and a short string dictating which direction that score came from, it does not affect the asymptotic space complexity of the algorithm.

There are two for loops in the unrestricted algorithm, the outer one being O(n) where n is the length of the string or the align length, and the inner loop being O(m) where m is the length of the other string or align length, whichever is shorter.

The backtrace time complexity for unrestricted is O(n), since it simply constructs the alignments from the directions stored in the table. This would make the overall time complexity O(nm + n), but asymptotically just O(nm)

The above adds another O(n + m) to the space complexity, but again does not contribute asymptotically.

For the banded algorithm, the space complexity is O(kn), where k is the bandwidth (in this case 7), and n is the length of the shorter string. The backtrace adds another O(n + m) to the space complexity, but again does not contribute asymptotically.

The first for loop in the banded algorithm is O(n) time complexity, where n is the minimum of the shorter string's length or the align length. The inner for loop contributes an effectively constant time, in this case O(7) which is asymptotically O(1), since the bandwidth is 7 and the inner loop only goes as long as the bandwidth.

The backtrace for unbanded then is also O(n), as it traces the path to the start from the end and constructs the alignments while doing so. It is O(n) because it is the length of the string (or the align length).

# Part 2:

My algorithm uses the Needleman-Wunsch cost function, and goes through the shorter string character by character, comparing each one to every character of the longer (or horizontally-oriented) string. It formed a DAG because no higher row was dependent on a lower row, and no value was dependent on the value to its right. I stored the value and direction the value came from in tuples in each cell of the table, so that the backtrace could follow the directions in reverse to construct the alignment and arrive at the beginning of the string.

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## Part 3:



Part 4:

#### Unrestricted:

attgcgagcgatttgcgtgcgtgcatcccgcttc-actg--atctcttgttagatcttttcataatctaaactttataaaaacatccactccctgta-g

ataa-gagtgattggcgtccgtacctttctactctcaaactcttgttagtttaaatc-taatctaaactttataaa--cggc-acttcctgtgtg

sequences 3 and 10
attgcgagcgatttgcgtgcgtgcatcccgcttc-actg--at-ctcttgttagatcttttcataatctaaactttataaaaacatccactccctgta-g
ataa-gagtgattggcgtccgtacgtaccctttctactctcaaactcttgttagtttaaatc-taatctaaactttataaa--cggc-acttcctgtgtg

#### Banded:

gattgcgagcgatttgcgtgcgtgcat-ccc--gcttcactgatctcttgttagatcttttcataatctaaactttataaaaacatccactccctgt-a

-a-taagagtgattggcgtccgtacgtaccctttctactctcaaactcttgttagtttaaatc-taatctaaactttat--aaac-ggcacttcctgtgt

sequences 3 and 10 for banded gattgcgagcgatttgcgtgcgtgcat-ccc-gcttcact-gatctcttgttagatcttttcataatctaaactttataaaacatccactcctgt-a -a-taagagtgattggcgtccgtacgtaccctttctactctcaaactcttgttagtttaaatc-taatctaaactttat--aaac-ggcacttcctgtgt

### Part 5:

```
def unrestricted(self, a, b, align_length): # a is vertical, along the side and b is horizontal along the top
    b_end = min(len(b), align_length + 1)
            alignment2 = "-" + alignment2
            alignment1 = "-" + alignment1
            alignment1 = a[i] + alignment1
            alignment2 = b[j] + alignment2
    return score, alignment1, alignment2
def score direction(self, top, left, diagonal):
```

```
def banded_algorithm(self, a, b, align_length):
    for i in range(min(len(a), align_length + 1)):
        row_start = MAXINDELS - i
        for j in range(max(0, row_start), z):
            elif j == row_start and i < MAXINDELS + 1:</pre>
                 left = row[j-1][0] + INDEL if j > 0 else math.inf
                 \label{eq:diagonal} \mbox{diagonal = table[i-1][j+1][0] + INDEL \ \mbox{if} \ \ j < (2 * MAXINDELS) \ \mbox{else math.inf}
                value = self.score_direction(table[i-1][j][0] + top, left, diagonal)
            row[j] = value
        table.append(row)
    score_index = (-MAXINDELS-1) + (len(b) - len(a))
    score = table[-1][score_index][0]
    alignment1 = ""
    alignment2 = ""
    i = min(len(a), align_length - 1) - 1
    j = z + score_index
        if table[i][j][1] == "top":
            alignment1 = a[i] + alignment1
            alignment2 = b[i + j - MAXINDELS] + alignment2
        elif table[i][j][1] == "left":
            alignment1 = "-" + alignment1
            alignment2 = b[i + j - MAXINDELS] + alignment2
            alignment1 = a[i] + alignment1
            alignment2 = "-" + alignment2
    return score, alignment1, alignment2
```