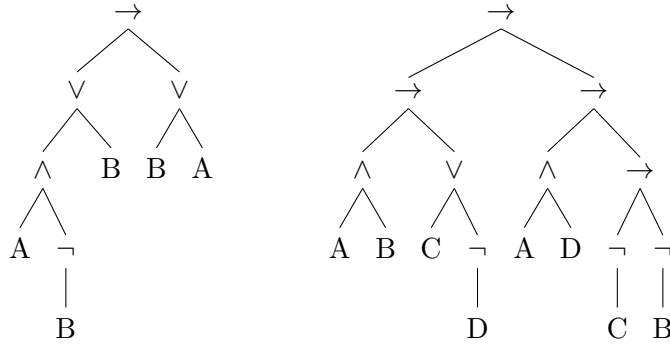


Exercise Sheet 12g - Solutions

Propositional Logic – Natural Deduction

1. The parse trees of F and G are



2. Here is a constructive Natural Deduction proof of $(A \wedge \neg B) \vee B \rightarrow B \vee A$

$$\begin{array}{c}
 \frac{\overline{A \wedge \neg B}^1 \quad \frac{\overline{A \wedge \neg B}^2 \quad \frac{A}{B \vee A} [\vee I_R] \quad \frac{B}{B \vee A} [\vee I_L]}{A \wedge \neg B \rightarrow B \vee A}^2 [\rightarrow I] \quad \frac{B}{B \rightarrow B \vee A}^3 [\rightarrow I]}{(A \wedge \neg B) \vee B \rightarrow B \vee A}^1 [\rightarrow I] \\
 \frac{\overline{A \wedge \neg B}^1 \quad \frac{\overline{A \wedge \neg B}^2 \quad \frac{A}{B \vee A} [\vee I_R] \quad \frac{B}{B \vee A} [\vee I_L]}{A \wedge \neg B \rightarrow B \vee A}^2 [\rightarrow I] \quad \frac{B}{B \rightarrow B \vee A}^3 [\rightarrow I]}{(A \wedge \neg B) \vee B \rightarrow B \vee A}^1 [\rightarrow I]
 \end{array}$$

3. Here is a constructive Natural Deduction proof of $(A \wedge B \rightarrow C \vee \neg D) \rightarrow A \wedge D \rightarrow \neg C \rightarrow \neg B$

$$\begin{array}{c}
 \frac{\overline{A \wedge B \rightarrow C \vee \neg D}^1 \quad \frac{\overline{A \wedge D}^2 \quad \frac{A}{A \wedge B} [\wedge I] \quad \frac{B}{\neg C}^4 [\neg I] \quad \frac{C}{\neg C}^5 [\neg E] \quad \frac{D}{\neg D}^6 [\neg E]}{A \wedge B \rightarrow C \vee \neg D}^1 \quad \frac{\overline{A \wedge D}^2 \quad \frac{A}{A \wedge D} [\wedge I] \quad \frac{D}{\neg D}^6 [\neg E]}{A \wedge D \rightarrow \neg C \rightarrow \neg B}^2 [\rightarrow I] \\
 \frac{\overline{A \wedge B \rightarrow C \vee \neg D}^1 \quad \frac{\overline{A \wedge D}^2 \quad \frac{A}{A \wedge B} [\wedge I] \quad \frac{B}{\neg C}^4 [\neg I] \quad \frac{C}{\neg C}^5 [\neg E] \quad \frac{D}{\neg D}^6 [\neg E]}{A \wedge B \rightarrow C \vee \neg D}^1 \quad \frac{\overline{A \wedge D}^2 \quad \frac{A}{A \wedge D} [\wedge I] \quad \frac{D}{\neg D}^6 [\neg E]}{A \wedge D \rightarrow \neg C \rightarrow \neg B}^2 [\rightarrow I]}{(A \wedge B \rightarrow C \vee \neg D) \rightarrow A \wedge D \rightarrow \neg C \rightarrow \neg B}^1 [\rightarrow I]
 \end{array}$$

4. F is provable and therefore valid by soundness, and so any valuation satisfies the formula such as $A = \mathbf{T}$, $B = \mathbf{T}$.
 G is provable and therefore valid by soundness, and therefore not falsifiable.