

Introduction to Categorical Logic

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Contents

2	First-Order Logic	5
2.1	Theories	7
2.2	Predicates as subobjects	10
2.3	Cartesian logic	14
2.3.1	Subtypes	23
2.4	Quantifiers as adjoints	27
2.4.1	The Beck-Chevalley condition	31
2.5	Regular logic	34
2.5.1	Regular categories	34
2.5.2	Images and existential quantifiers	40
2.5.3	Regular theories	44
2.5.4	Classifying category of a regular theory	48
2.5.5	Coherent logic	54

Chapter 2

First-Order Logic

Having considered equational theories, we now move on to first-order logic. This is the usual predicate logic with propositional connectives like \wedge and \Rightarrow , and quantifiers \forall and \exists . The general approach to studying logic via category theory is to determine categorical structures that model the first-order logical operations, or a suitable fragment of it, and then consider categories with these structures and functors that preserve them. Here adjoint functors play an important role, as the basic logical operations are recognized as adjoints. We again show that the semantics is “functorial”, meaning that models of a theory are functors that preserve suitable categorical structure. We again construct classifying categories representing theories, which are the counterparts of the algebraic theories that we have already met.

Let us demonstrate our approach informally with an example. In section ?? we considered models of algebraic theories in categories with finite products. Recall that e.g. a group is a structure of the form:

$$e : 1 \rightarrow G , \quad m : G \times G \rightarrow G , \quad i : G \rightarrow G .$$

for which, moreover, certain diagrams built from these basic arrows must commute. We can express some properties of groups in terms of further equations, for example commutativity

$$x \cdot y = y \cdot x .$$

As we saw, such equations can be interpreted in any category with finite products. This provides a large scope for categorical semantics of algebraic theories.

However, there are also many significant properties of algebraic structures which cannot be expressed with equations. Consider the statement that a group (G, e, m, i) has no non-trivial square roots of unity,

$$\forall x : G . (x \cdot x = e \Rightarrow x = e) . \quad (2.1)$$

This is a first-order logical statement which cannot be rewritten as a system of equations (proof!). To see what its categorical interpretation ought to be, we look at its usual set-theoretic interpretation. Each of subformulas, $x \cdot x = e$ and $x = e$, determines a subset

$$\begin{array}{ccc} \{x \in G \mid x \cdot x = e\} & \xrightarrow{\quad \quad \quad} & \{x \in G \mid x = e\} \\ & \searrow i \quad \quad \swarrow j & \\ & G & \end{array}$$
$$\begin{array}{c} \{x \in G \mid x \cdot x = e\} \rightrightarrows G \xrightarrow{\langle 1_G, 1_G \rangle} G \times G \xrightarrow{m} G \\ \quad \quad \quad \searrow \text{\scriptsize $e \circ !_G$} \nearrow \\[10pt] \{x \in G \mid x = e\} \rightrightarrows G \xrightleftharpoons[e \circ !_G]{1_G} G \end{array}$$

The aim of this chapter is to analyze how such examples can be treated in general. We want to relate first-order logic and fragments of it to categorical structures that are suitable for the interpretation of the logic. The general outline will be as follows:

1. A language \mathcal{L} for a first-order theory consists, as usual, of some basic relation, function, and constant symbols, say $\mathcal{L} = (R, f, c)$.
2. An \mathcal{L} -structure in a category \mathcal{C} with finite limits is an interpretation of \mathcal{L} in \mathcal{C} as an object M equipped with corresponding relations and operations (of appropriate arities), e.g.

$$\begin{aligned} R^M &\rightharpoonup M \times \cdots \times M \\ f^M &: M \times \cdots \times M \rightarrow M \\ c^M &: 1 \rightarrow M. \end{aligned}$$

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3. Formulas $\varphi(x_1, \dots, x_n)$ in (some fragment of) first-order logic will be interpreted as “generalized subsets”, i.e. subobjects,

$$\llbracket \varphi(x_1, \dots, x_n) \rrbracket \multimap M \times \dots \times M.$$

The interpretation makes use of categorical operations in \mathcal{C} corresponding to the logical ones appearing in the formula $\varphi(x_1, \dots, x_n)$.

4. A theory \mathbb{T} in (a fragment of) first-order logic will consist of a set of (binary) sequents,

$$\varphi(x_1, \dots, x_n) \vdash \psi(x_1, \dots, x_n).$$

5. A model of the theory is then an interpretation M in which the corresponding subobjects satisfy all the sequents of \mathbb{T} , in the sense that

$$\llbracket \varphi(x_1, \dots, x_n) \rrbracket \leq \llbracket \psi(x_1, \dots, x_n) \rrbracket \quad \text{in } \mathbf{Sub}(M^n).$$

6. We shall give a deductive calculus for such sequents, prove that it is sound with respect to categorical models, and then use it to construct a classifying category $\mathcal{C}_{\mathbb{T}}$, with the expected universal property: models of \mathbb{T} correspond to (structure-preserving) functors on $\mathcal{C}_{\mathbb{T}}$.
7. Completeness of the calculus in general follows from classification, and more specialized completeness results from embedding theorems applied to the classifying category.

2.1 Theories

A *first-order theory* \mathbb{T} consists of an underlying *type theory* and a set of formulas in a *fragment of first-order logic*. Anticipating Chapter ??, the type theory is given by a set of basic types, a set of basic constants together with their types, rules for forming types, and rules and axioms for deriving typing judgments

$$x_1 : A_1, \dots, x_n : A_n \mid t : B,$$

expressing that term t has type B in typing context $x_1 : A_1, \dots, x_n : A_n$, and a set of axioms and rules of inference which tell us which equations between terms

$$x_1 : A_1, \dots, x_n : A_n \mid t = u : B,$$

are valid. This part of the theory \mathbb{T} may be regarded as providing the underlying structure, on top of which the logical formulas are defined. For first-order logic, the underlying type theory will be essentially the same as the equational logic that we already met in Chapter ??.

A fragment of first-order logic is then given by a set of *basic relation symbols* together with a specification of which first-order operations are being considered in building formulas. Each basic relation symbol has a *signature* (A_1, \dots, A_n) , which specifies the types of its arguments. The *arity* of a relation symbol is the number of arguments it takes. The judgment²

$$x_1 : A_1, \dots, x_n : A_n \mid \phi \text{ pred}$$

states that ϕ is a well-formed formula in typing context $x_1 : A_1, \dots, x_n : A_n$. For each basic relation symbol R with signature (A_1, \dots, A_n) there is an inference rule

$$\frac{\Gamma \mid t_1 : A_1 \quad \dots \quad \Gamma \mid t_n : A_n}{\Gamma \mid R(t_1, \dots, t_n) \text{ pred}}$$

Depending on what fragment of first-order logic is involved, there may be other rules for forming logical formulas. For example, if equality is present, then for each type A there is a rule

$$\frac{\Gamma \mid t : A \quad \Gamma \mid u : A}{\Gamma \mid t =_A u \text{ pred}}$$

and if conjunction is present, then there is a rule

$$\frac{\Gamma \mid \varphi \text{ pred} \quad \Gamma \mid \psi \text{ pred}}{\Gamma \mid \varphi \wedge \psi \text{ pred}}$$

Other such rules will be given when we come to the study of particular logical operations.

The basic logical judgment of a first-order theory is *logical entailment* between formulas,

$$x_1 : A_1, \dots, x_n : A_n \mid \varphi_1, \dots, \varphi_m \vdash \psi$$

which states that in the typing context $x_1 : A_1, \dots, x_n : A_n$, the hypotheses $\varphi_1, \dots, \varphi_m$ entail ψ . It is understood that the terms appearing in the formulas are well-typed in the typing context, and that formulas $\varphi_1, \dots, \varphi_m, \psi$ are part of the fragment of the logic of \mathbb{T} . When the fragment contains conjunction \wedge it is convenient to restrict attention to *binary* sequents,

$$x_1 : A_1, \dots, x_n : A_n \mid \varphi \vdash \psi,$$

by replacing $\varphi_1, \dots, \varphi_m$ with $\varphi_1 \wedge \dots \wedge \varphi_m$. When the fragment contains equality, we may replace the type-theoretic equality judgments

$$x_1 : A_1, \dots, x_n : A_n \mid t = u : B$$

with the logical statements

$$x_1 : A_1, \dots, x_n : A_n \mid \cdot \vdash t =_B u.$$

²We follow type-theoretic practice here by adding the tag **pred** to turn what would otherwise be an exhibited formula in context into a judgement concerning the formula.

The subscript at the equality sign indicates the type at which the equality is taken. In a theory \mathbb{T} there are basic entailments, or axioms, which together with the inference rules for the operations involved can be used for deriving valid judgments, as usual.

We shall consider several standard fragments of first-order logic, determined by selecting a subset of logical connectives and quantifiers. These are as follows:

1. *Full first-order logic* is built from logical operations

$$= \top \perp \neg \wedge \vee \Rightarrow \forall \exists .$$

2. *Cartesian logic* is the fragment built from

$$= \top \wedge .$$

3. *Regular logic* is the fragment built from

$$= \top \wedge \exists .$$

4. *Coherent logic* is the fragment built from

$$= \top \wedge \exists \perp \vee .$$

5. A *geometric formula* is a formula of the form

$$\forall x : A. (\varphi \Rightarrow \psi) ,$$

where φ and ψ are coherent formulas.

The names for these fragments come from the names of various categorical structures in which they are interpreted.

The well-formed terms and formulas of a first-order theory \mathbb{T} constitute its *language*. It may seem that we are doing things backwards, because we should have spoken of first-order languages before we spoke of first-order theories. While this is possible for simple theories, it becomes difficult to do when we consider more complicated theories in which types and logical formulas are intertwined. In such cases the typing judgments and logical entailments may be given by a mutual recursive definition. In order to find out whether a given term is well-formed, we might have to prove a logical statement. In everyday mathematics this occurs all the time, for example, to show that the term $\int_0^\infty f$ denotes a real number, it may be necessary to prove that $f : \mathbb{R} \rightarrow \mathbb{R}$ is an integrable function and that the integral has a finite value. This is why it does not always make sense to strictly differentiate a language from a theory.³

In order to focus on the logical part of first-order theories, we are going to limit attention to only two very simple kinds of type theory. A *single-sorted* first-order theory has as its underlying type theory a single type A , and for each $k \in \mathbb{N}$ a set of basic k -ary function symbols. The rules for typing judgments are:

³However, it *does* make sense to distinguish syntax from theory. Rules of substitution and the behaviour of free and bound variables are syntactic considerations, for example.

1. Variables in contexts:

$$\overline{x_1 : A, \dots, x_n : A \mid x_i : A}$$

2. For each basic function symbol f of arity k , there is an inference rule

$$\frac{\Gamma \mid t_1 : A \quad \dots \quad \Gamma \mid t_n : A}{\Gamma \mid f(t_1, \dots, t_n) : A}$$

This much is essentially an algebraic theory. In addition, however, a single-sorted first-order theory may contain relation symbols, formulas, axioms, and rules of inference which an algebraic theory does not.

A slight generalization of a single-sorted theory is a *multi-sorted* one. Its underlying type theory is given by a set of types, and a set of basic function symbols. Each function symbol f has a *signature* $(A_1, \dots, A_n; B)$, where n is the arity of f and A_1, \dots, A_n, B are types. The rules for typing judgments are:

1. Variables in contexts:

$$\overline{x_1 : A_1, \dots, x_n : A_n \mid x_i : A_i}$$

2. For each basic function symbol f with signature $(A_1, \dots, A_n; B)$, there is an inference rule

$$\frac{\Gamma \mid t_1 : A_1 \quad \dots \quad \Gamma \mid t_n : A_n}{\Gamma \mid f(t_1, \dots, t_n) : B}$$

We often write suggestively $f : A_1 \times \dots \times A_n \rightarrow B$ to indicate that $(A_1, \dots, A_n; B)$ is the signature of f . However, this does not mean that $A_1 \times \dots \times A_n \rightarrow B$ is a type! A multi-sorted first-order theory does *not* have any type forming operations, such as \times and \rightarrow .

2.2 Predicates as subobjects

Formulas of first-order logic will be interpreted as “generalized subsets”, i.e. subobjects. We therefore need to review some of the basic theory of these.

Let A be an object in a category \mathcal{C} . If $i : I \rightarrow A$ and $j : J \rightarrow A$ are monos into A , we say that i is smaller than j , and write $i \leq j$, when there exists a morphism $k : I \rightarrow J$ such that the following diagram commutes:

$$\begin{array}{ccc} I & \xrightarrow{\quad k \quad} & J \\ & \searrow i \quad \swarrow j & \\ & A & \end{array}$$

If such a k exists then it, too, is monic, since i is, and it is unique, since j is monic. The class $\mathbf{Mono}(A)$ of all monos into A is this preordered by this relation \leq , it is the same as

the slice category $\mathbf{Mono}(\mathcal{C})/A$ of all monos in \mathcal{C} , sliced over the object A . Let $\mathbf{Sub}(A)$ be the poset reflection of this preorder. Thus the elements of $\mathbf{Sub}(A)$ are equivalence classes of monos into A , where monos $i : I \rightarrowtail A$ and $j : J \rightarrowtail A$ are equivalent when $i \leq j$ and $j \leq i$ (note that then $I \cong J$). The induced relation \leq on $\mathbf{Sub}(A)$ is then a partial order.

We have to be a bit careful with the formation of $\mathbf{Sub}(A)$, since it is defined as a quotient of a *class* $\mathbf{Mono}(A)$. In many particular cases the general construction by quotients can be avoided. If we can demonstrate that the preorder $\mathbf{Mono}(A)$ is equivalent, as a category, to a poset P then we can simply take $\mathbf{Sub}(A) = P$. At any rate, we usually require that $\mathbf{Sub}(A)$ is small.

Definition 2.2.1. A category \mathcal{C} is *well-powered* when, for all $A \in \mathcal{C}$, there is at most a *set* of subobjects of A , so that the category $\mathbf{Mono}(A)$ is equivalent to a small poset. In other words, for every $A \in \mathcal{C}$, $\mathbf{Sub}(A)$ is a small category.

We shall often speak of subobjects as if they were monos rather than equivalence classes of monos. It is understood that we mean the subobjects represented by monos and not the monos themselves. Sometimes we refer to a mono $i : I \rightarrowtail A$ by its domain I only, even though the object I itself does not determine the morphism i . Hopefully this will not cause confusion, as it is always going to be clear which mono is meant to go along with the object I .

In a category \mathcal{C} with finite limits the assignment $A \mapsto \mathbf{Sub}(A)$ is the object part of the *subobject functor*

$$\mathbf{Sub} : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Poset} .$$

The morphism part of \mathbf{Sub} is pullback. More precisely, given a morphism $f : A \rightarrow B$, let $\mathbf{Sub}(f) = f^* : \mathbf{Sub}(B) \rightarrow \mathbf{Sub}(A)$ be the monotone map which maps the subobject $[i : I \rightarrowtail B]$ to the subobject $[f^*i : f^*I \rightarrowtail A]$, where $f^*i : f^*I \rightarrowtail A$ is a pullback of i along f :

$$\begin{array}{ccc} f^*I & \xrightarrow{\quad} & I \\ \downarrow f^*i & \lrcorner & \downarrow i \\ A & \xrightarrow{\quad f \quad} & B \end{array}$$

Recall that a pullback of a mono is again mono, so this definition makes sense. We need to verify (1) that if two monos $i : I \rightarrowtail A$ and $j : J \rightarrowtail A$ are equivalent, then their pullbacks are so as well; and (2) that $\mathbf{Sub}(1_A) = 1_{\mathbf{Sub}(A)}$ and $\mathbf{Sub}(g \circ f) = \mathbf{Sub}(f) \circ \mathbf{Sub}(g)$. These all follow easily from the fact that pullback is a functor $\mathcal{C}/B \rightarrow \mathcal{C}/A$, which reduces to the familiar “2-pullbacks” lemma:

Lemma 2.2.2. *Suppose both squares in the following diagram are pullbacks:*

$$\begin{array}{ccccc} \cdot & \xrightarrow{\quad} & \cdot & \xrightarrow{\quad} & \cdot \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ \cdot & \xrightarrow{\quad} & \cdot & \xrightarrow{\quad} & \cdot \end{array}$$

Then the outer rectangle is a pullback diagram as well. Moreover, if the outer rectangle and the right square are pullbacks, then so is the left square.

Proof. This is left as an exercise in diagram chasing. \square

Of course, pullbacks are really only determined up to isomorphism, but this does not cause any problems because isomorphic monos represent the same subobject.

In the semantics to be given below, a formula

$$x : A \mid \varphi \text{ pred}$$

will be interpreted as a subobject

$$\llbracket x : A \mid \varphi \rrbracket \rightrightarrows \llbracket A \rrbracket.$$

Thus $\mathbf{Sub}(A)$ can be regarded as the poset of “predicates” on A , generalizing the powerset of a set A . Logical operations on formulas then correspond to operations on $\mathbf{Sub}(A)$. The structure of $\mathbf{Sub}(A)$ therefore determines which logical connectives can be interpreted. If $\mathbf{Sub}(A)$ is a Heyting algebra, then we can interpret the full intuitionistic propositional calculus (cf. Subsection ??), but if $\mathbf{Sub}(A)$ only has binary meets then all that can be interpreted are \top and \wedge . We will work out details of different operations in the following sections, but one common aspect that we require is the “stability” of the interpretation of the logical operations, in a sense that we now make clear.

Substitution and stability

Let us consider the interpretation of substitution of terms for variables. There are two kinds of substitution, into a term, and into a formula. We may substitute a term $x : A \mid t : B$ for a variable y in a term $y : B \mid u : C$ to obtain a new term $x : A \mid u[t/y] : C$. If t and u are interpreted as morphisms

$$\llbracket A \rrbracket \xrightarrow{\llbracket t \rrbracket} \llbracket B \rrbracket \xrightarrow{\llbracket u \rrbracket} \llbracket C \rrbracket$$

then $u[t/y]$ is interpreted as their composition:

$$\llbracket x : A \mid u[t/y] : C \rrbracket = \llbracket y : B \mid u : C \rrbracket \circ \llbracket x : A \mid t : B \rrbracket.$$

Thus, *substitution into a term is composition*.

The second kind of substitution occurs when we substitute a term $x : A \mid t : B$ for a variable y in a formula $y : B \mid \varphi$ to obtain a new formula $x : A \mid \varphi[t/y]$. If t is interpreted as a morphism $\llbracket t \rrbracket : \llbracket A \rrbracket \rightarrow \llbracket B \rrbracket$ and φ is interpreted as a subobject $\llbracket \varphi \rrbracket \rightrightarrows \llbracket B \rrbracket$ then the interpretation of $\varphi[t/y]$ is the pullback of $\llbracket \varphi \rrbracket$ along $\llbracket t \rrbracket$:

$$\begin{array}{ccc} \llbracket \varphi[t/y] \rrbracket = \llbracket t \rrbracket^* \llbracket \varphi \rrbracket & \xrightarrow{\quad} & \llbracket \varphi \rrbracket \\ \downarrow & \lrcorner & \downarrow \\ \llbracket A \rrbracket & \xrightarrow{\llbracket t \rrbracket} & \llbracket B \rrbracket \end{array}$$

Thus, *substitution into a formula is pullback*,

$$\llbracket x : A \mid \varphi[t/y] \rrbracket = \llbracket x : A \mid t : B \rrbracket^* \llbracket y : B \mid \varphi \rrbracket.$$

Now, because substitution respects the syntactical, logical operations, e.g.

$$(\varphi \wedge \psi)[t/x] = \varphi[t/x] \wedge \psi[t/x],$$

the categorical structures used to interpret the various logical operations such as \wedge must also behave well with respect to the interpretation of substitution, i.e. pullback. We say that a categorical property or structure is *stable (under pullbacks)* if it is preserved by pullbacks.

For example, a category \mathcal{C} has *stable meets* if each poset $\mathbf{Sub}(A)$ has binary meets, and the pullback of a meet $I \wedge J \rightarrowtail A$ along any map $f : B \rightarrow A$ is the meet $f^*I \wedge f^*J \rightarrowtail A$ of the respective pullbacks,

$$f^*(I \wedge J) = f^*I \wedge f^*J.$$

This means that the meet operation,

$$\wedge : \mathbf{Sub}(A) \times \mathbf{Sub}(A) \longrightarrow \mathbf{Sub}(A)$$

is natural in A , in the sense that for any map $f : B \rightarrow A$ the following diagram commutes.

$$\begin{array}{ccc} \mathbf{Sub}(A) \times \mathbf{Sub}(A) & \xrightarrow{\wedge_A} & \mathbf{Sub}(A) \\ f^* \times f^* \downarrow & & \downarrow f^* \\ \mathbf{Sub}(B) \times \mathbf{Sub}(B) & \xrightarrow{\wedge_B} & \mathbf{Sub}(B) \end{array}$$

Exercise 2.2.3. Show that any category \mathcal{C} with finite limits has stable meets in the foregoing sense: each poset $\mathbf{Sub}(A)$ has all finite meets (i.e. including the “empty meet” 1), and these are stable under pullbacks. Conclude that $\mathbf{Sub} : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Posets}$ factors through the subcategory of \wedge -semi-lattices.

Generalized elements

In any category, we sometimes consider arbitrary arrows $x : X \rightarrow C$ as *generalized elements* of C , thinking thereby of variable elements or parameters. With respect to a subobject $S \rightarrowtail C$, such an element is said to be *in the subobject*, written

$$x \in_C S,$$

if it factors through (any mono representing) the subobject,

$$\begin{array}{ccc} & & S \\ & \nearrow & \downarrow \\ X & \xrightarrow{x} & C \end{array}$$

which, observe, it then does uniquely. The following “generalized element semantics” can be a useful technique for “externalizing” the operations on subobjects into statements about generalized elements.

Proposition 2.2.4. *Let C be any object in a category \mathcal{C} with finite limits.*

1. *for the top element $1 \in \mathbf{Sub}(C)$, and for all $x : X \rightarrow C$,*

$$x \in_C 1.$$

2. *for any $S, T \in \mathbf{Sub}(C)$,*

$$S \leq T \iff x \in_C S \text{ implies } x \in_C T, \text{ for all } x : X \rightarrow C.$$

3. *for any $S, T \in \mathbf{Sub}(C)$, and for all $x : X \rightarrow C$,*

$$x \in_C S \wedge T \iff x \in_C S \text{ and } x \in_C T.$$

4. *for the subobject $\Delta = [\langle 1_C, 1_C \rangle] \in \mathbf{Sub}(C \times C)$, and for all $x, y : X \rightarrow C$,*

$$\langle x, y \rangle \in \Delta \iff x = y.$$

5. *for the equalizer $E_{(f,g)} \rightharpoonup A$ of a pair of arrows $f, g : A \rightrightarrows B$, and for all $x : X \rightarrow A$,*

$$x \in_A E_{(f,g)} \iff fx = gx.$$

6. *for the pullback $f^*S \rightharpoonup A$ of a subobject $S \rightharpoonup B$ along any arrow $f : A \rightarrow B$, and for all $x : X \rightarrow A$,*

$$x \in_A f^*S \iff fx \in_B S.$$

Exercise 2.2.5. Prove the proposition.

2.3 Cartesian logic

As a first example we look at the logic of *cartesian categories*, which are categories with finite limits, to be called *cartesian logic*. This is a logic of formulas over a multi-sorted type theory with unit type 1. (See section ?? for multi-sorted type theories and the axioms for the unit type). The logical operations are $=$, \top , and \wedge .

Formation rules for cartesian logic

Given a basic language consisting of a stock of relation and function symbols (with arities), the terms are built up as usual from the basic function symbols and variables (we take “constants” to be 0-ary function symbols). The rules for constructing formulas are as follows:

1. The 0-ary relation symbol \top is a formula:

$$\overline{\Gamma \mid \top \text{ pred}}$$

2. For each basic relation symbol R with signature (A_1, \dots, A_n) there is a rule

$$\frac{\Gamma \mid t_1 : A_1 \quad \dots \quad \Gamma \mid t_n : A_n}{\Gamma \mid R(t_1, \dots, t_n) \text{ pred}}$$

3. For each type A , there is a rule

$$\frac{\Gamma \mid s : A \quad \Gamma \mid t : A}{\Gamma \mid s =_A t \text{ pred}}$$

4. Conjunction:

$$\frac{\Gamma \mid \varphi \text{ pred} \quad \Gamma \mid \psi \text{ pred}}{\Gamma \mid \varphi \wedge \psi \text{ pred}}$$

5. Weakening:

$$\frac{\Gamma \mid \varphi \text{ pred}}{\Gamma, x : A \mid \varphi \text{ pred}}$$

Observe that, as usual, there is then a derived operation of substitution of terms for variables into formulas, defined by structural recursion on the above specification of formulas:

Substitution:

$$\frac{\Gamma \mid t : A \quad \Gamma, x : A \mid \varphi \text{ pred}}{\Gamma \mid \varphi[t/x] \text{ pred}}$$

Inference rules for cartesian logic

Although we do not yet need them, we state the rules of inference here, too, for the convenience of having the entire specification of cartesian logic in one place. As already mentioned, we can conveniently state this deductive calculus entirely in terms of *binary* sequents,

$$\Gamma \mid \psi \vdash \varphi.$$

We omit mention of the context Γ when it is the same in the premisses and conclusion of a rule.

1. Weakening:

$$\frac{\Gamma \mid \psi \vdash \varphi}{\Gamma, x : A \mid \psi \vdash \varphi}$$

2. Substitution:

$$\frac{\Gamma \mid t : A \quad \Gamma, x : A \mid \psi \vdash \varphi}{\Gamma \mid \psi[t/x] \vdash \varphi[t/x]}$$

3. Identity:

$$\overline{\varphi \vdash \varphi}$$

4. Cut:

$$\frac{\psi \vdash \theta \quad \theta \vdash \varphi}{\psi \vdash \varphi}$$

5. Equality:

$$\frac{}{\psi \vdash t =_A t} \quad \frac{\psi \vdash t =_A u \quad \psi \vdash \varphi[t/z]}{\psi \vdash \varphi[u/z]}$$

6. Truth:

$$\overline{\psi \vdash \top}$$

7. Conjunction:

$$\frac{\vartheta \vdash \varphi \quad \vartheta \vdash \psi}{\vartheta \vdash \varphi \wedge \psi} \quad \frac{\vartheta \vdash \varphi \wedge \psi}{\vartheta \vdash \psi} \quad \frac{\vartheta \vdash \varphi \wedge \psi}{\vartheta \vdash \varphi}$$

Exercise 2.3.1. Derive symmetry and transitivity of equality:

$$\frac{\Gamma \mid \psi \vdash t =_A u}{\Gamma \mid \psi \vdash u =_A t} \quad \frac{\Gamma \mid \psi \vdash t =_A u \quad \Gamma \mid \psi \vdash u =_A v}{\Gamma \mid \psi \vdash t =_A v}$$

Example 2.3.2. The theory of a poset is a cartesian theory. There is one basic sort \mathbf{P} and one binary relation symbol \leq with signature (\mathbf{P}, \mathbf{P}) . The axioms are the familiar axioms for reflexivity, transitivity, and antisymmetry:

$$\begin{aligned} & x : \mathbf{P} \mid \cdot \vdash x \leq x \\ & x : \mathbf{P}, y : \mathbf{P}, z : \mathbf{P} \mid x \leq y, y \leq z \vdash x \leq z \\ & x : \mathbf{P}, y : \mathbf{P} \mid x \leq y, y \leq x \vdash x =_{\mathbf{P}} y \end{aligned}$$

There are also many examples, such as ordered groups, ordered fields, etc., that are posets with further algebraic structure.

Example 2.3.3. An *equivalence relation* in a cartesian category is a model of the corresponding theory with one basic sort \mathbf{A} and one binary relation symbol \sim with signature (\mathbf{A}, \mathbf{A}) . The axioms are the familiar axioms for reflexivity, symmetry, and transitivity:

$$\begin{aligned} & x : \mathbf{A} \mid \cdot \vdash x \sim x \\ & x : \mathbf{A}, y : \mathbf{A} \mid x \sim y \vdash y \sim x \\ & x : \mathbf{A}, y : \mathbf{A}, z : \mathbf{A} \mid x \sim y \wedge y \sim z \vdash x \sim z \end{aligned}$$

Semantics for cartesian logic

In order to give the semantics of cartesian logic, we require a couple of useful propositions regarding cartesian categories.

Proposition 2.3.4. *If a category \mathcal{C} has pullbacks then, for every $A \in \mathcal{C}$, $\mathbf{Sub}(A)$ has finite limits. Moreover, these are stable under pullback.*

Proof. The poset $\mathbf{Sub}(A)$ has finite limits if it has a top object and binary meets. The top object of $\mathbf{Sub}(A)$ is the subobject $[1_A : A \rightarrow A]$. The meet of subobjects $i : I \rightarrow A$ and $j : J \rightarrow A$ is the subobject $i \wedge j = i \circ (i^*j) = j \circ (j^*i) : I \wedge J \rightarrow A$ obtained by pullback, as in the following diagram:

$$\begin{array}{ccc} I \wedge J & \xrightarrow{j^*i} & J \\ \downarrow i^*j & \lrcorner & \downarrow j \\ I & \xrightarrow{i} & A \end{array}$$

It is easy to verify that $I \wedge J$ is the infimum of I and J . Finally, stability follows from a familiar diagram chase on a cube of pullbacks. \square

Proposition 2.3.5. *If a category has finite products and pullbacks of monos along monos then it has all finite limits.*

Proof. It is sufficient to show that the category has equalizers. To construct the equalizer of parallel arrows $f : A \rightarrow B$ and $g : A \rightarrow B$, first observe that the arrows

$$A \xrightarrow{\langle 1_A, f \rangle} A \times B \qquad A \xrightarrow{\langle 1_A, g \rangle} A \times B$$

are monos because the projection $\pi_0 : A \times B \rightarrow A$ is their left inverse. Therefore, we may construct the pullback

$$\begin{array}{ccc} P & \xrightarrow{p} & A \\ \downarrow q & \lrcorner & \downarrow \langle 1_A, f \rangle \\ A & \xrightarrow{\langle 1_A, g \rangle} & A \times B \end{array}$$

The morphisms p and q coincide because $\langle 1_A, f \rangle$ and $\langle 1_A, g \rangle$ have a common left inverse π_0 :

$$p = 1_A \circ p = \pi_0 \circ \langle 1_A, f \rangle \circ p = \pi_0 \circ \langle 1_A, f \rangle \circ q = 1_A \circ q = q .$$

Let us show that $p : P \rightarrow A$ is the equalizer of f and g . First, p equalizes f and g ,

$$f \circ p = \pi_1 \circ \langle 1_A, f \rangle \circ p = \pi_1 \circ \langle 1_A, g \rangle \circ q = g \circ q = g \circ p .$$

If $k : K \rightarrow A$ also equalizes f and g then

$$\langle 1_A, f \rangle \circ k = \langle k, f \circ k \rangle = \langle k, g \circ k \rangle = \langle 1_A, g \rangle \circ k,$$

therefore by the universal property of the constructed pullback there exists a unique factorization $\bar{k} : K \rightarrow P$ such that $k = p \circ \bar{k}$, as required. \square

We now explain how cartesian logic is interpreted in a cartesian category \mathcal{C} (i.e. \mathcal{C} is finitely complete). Let \mathbb{T} be a multi-sorted cartesian theory. Recall that the type theory of \mathbb{T} is specified by a set of sorts (types) A, \dots and a set of basic function symbols f, \dots together with their signatures, while the logic is given by a set of basic relation symbols R, \dots with their signatures, and a set of axioms in the form of logical entailments between formulas in context,

$$\Gamma \mid \psi \vdash \varphi.$$

Definition 2.3.6. An *interpretation* of \mathbb{T} in \mathcal{C} is given by the following data, where Γ stands for a typing context $x_1 : A_1, \dots, x_n : A_n$, and ψ and φ are formulas:

1. Each sort A is interpreted as an object $\llbracket A \rrbracket$, with the unit sort 1 being interpreted as the terminal object 1 .
2. A typing context $x_1 : A_1, \dots, x_n : A_n$ is interpreted as the product $\llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket$. The empty context is interpreted as the terminal object 1 .
3. A basic function symbol f with signature $(A_1, \dots, A_m; B)$ is interpreted as a morphism $\llbracket f \rrbracket : \llbracket A_1 \rrbracket \times \dots \times \llbracket A_m \rrbracket \rightarrow \llbracket B \rrbracket$.
4. A basic relation symbol R with signature (A_1, \dots, A_n) is interpreted as a subobject $\llbracket R \rrbracket \in \mathbf{Sub}(\llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket)$.

We then extend the interpretation to all terms and formulas as follows:

1. A term in context $\Gamma \mid t : B$ is interpreted as a morphism

$$\llbracket \Gamma \mid t : B \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket B \rrbracket$$

according to the following specification.

- A variable $x_0 : A_1, \dots, x_n : A_n \mid x_i : A_i$ is interpreted as the i -th projection $\pi_i : \llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket \rightarrow \llbracket A_i \rrbracket$.
- The interpretation of $\Gamma \mid * : 1$ is the unique morphism $!_{\llbracket \Gamma \rrbracket} : \llbracket \Gamma \rrbracket \rightarrow 1$.
- A composite term $\Gamma \mid f(t_1, \dots, t_m) : B$, where f is a basic function symbol with signature $(A_1, \dots, A_m; B)$, is interpreted as the composition

$$\llbracket \Gamma \rrbracket \xrightarrow{\langle \llbracket t_1 \rrbracket, \dots, \llbracket t_m \rrbracket \rangle} \llbracket A_1 \rrbracket \times \dots \times \llbracket A_m \rrbracket \xrightarrow{\llbracket f \rrbracket} \llbracket B \rrbracket$$

Here $\llbracket t_i \rrbracket$ is shorthand for $\llbracket \Gamma \mid t_i : A_i \rrbracket$.

2. A formula in a context $\Gamma \mid \varphi$ is interpreted as a subobject $\llbracket \Gamma \mid \varphi \rrbracket \in \mathbf{Sub}(\llbracket \Gamma \rrbracket)$ according to the following specification.

- The logical constant \top is interpreted as the maximal subobject, represented by the identity arrow:

$$\llbracket \Gamma \mid \top \rrbracket = [1_{\llbracket \Gamma \rrbracket} : \llbracket \Gamma \rrbracket \rightarrow \llbracket \Gamma \rrbracket]$$

- An atomic formula $\Gamma \mid R(t_1, \dots, t_m)$, where R is a basic relation symbol with signature (A_1, \dots, A_m) is interpreted as the left-hand side of the pullback:

$$\begin{array}{ccc} \llbracket \Gamma \mid R(t_1, \dots, t_m) \rrbracket & \xrightarrow{\quad} & \llbracket R \rrbracket \\ \downarrow & & \downarrow \\ \llbracket \Gamma \rrbracket & \xrightarrow{\langle \llbracket t_1 \rrbracket, \dots, \llbracket t_m \rrbracket \rangle} & \llbracket A_1 \rrbracket \times \dots \times \llbracket A_m \rrbracket \end{array}$$

- An equation $\Gamma \mid t =_A u$ **pred** is interpreted as the subobject represented by the equalizer of $\llbracket \Gamma \mid t : A \rrbracket$ and $\llbracket \Gamma \mid u : A \rrbracket$:

$$\llbracket \Gamma \mid t =_A u \rrbracket \rightrightarrows \llbracket \Gamma \rrbracket \begin{array}{c} \xrightarrow{\llbracket t \rrbracket} \\ \xrightarrow{\llbracket u \rrbracket} \end{array} \llbracket A \rrbracket$$

- By Proposition 2.3.4, each $\mathbf{Sub}(A)$ is a poset with binary meets. Thus we interpret a conjunction $\Gamma \mid \varphi \wedge \psi$ **pred** as the meet of subobjects

$$\llbracket \Gamma \mid \varphi \wedge \psi \rrbracket = \llbracket \Gamma \mid \varphi \rrbracket \wedge \llbracket \Gamma \mid \psi \rrbracket .$$

- A formula formed by weakening is interpreted as pullback along a projection:

$$\begin{array}{ccc} \llbracket \Gamma, x : A \mid \varphi \rrbracket & \xrightarrow{\quad} & \llbracket \Gamma \mid \varphi \rrbracket \\ \downarrow & & \downarrow i \\ \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket & \xrightarrow{\pi} & \llbracket \Gamma \rrbracket \end{array}$$

Computing this pullback one sees that the interpretation of $\llbracket \Gamma, x : A \mid \varphi \rrbracket$ turns out to be the subobject

$$\llbracket \Gamma \mid \varphi \rrbracket \times \llbracket A \rrbracket \xrightarrow{i \times 1_A} \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket$$

This concludes the definition of an interpretation of a cartesian theory \mathbb{T} in a cartesian category \mathcal{C} .

As was explained in the previous section, the operation of substitution of terms into formulas is interpreted as pullback:

Lemma 2.3.7. *Let the formula $\Gamma, x : A \mid \varphi$ and the term $\Gamma \mid t : A$ be given. Then the substituted formula $\Gamma \mid \varphi[t/x]$ is interpreted as the pullback indicated in the following diagram:*

$$\begin{array}{ccc} \llbracket \Gamma \mid \varphi[t/x] \rrbracket & \longrightarrow & \llbracket \Gamma, x : A \mid \varphi \rrbracket \\ \downarrow & & \downarrow \\ \llbracket \Gamma \rrbracket & \xrightarrow{\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle} & \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket \end{array}$$

Proof. A simple induction on the structure of φ . We do the case where φ is an atomic formula $R(t_1, \dots, t_m)$. Let $\Gamma = x_1 : A_1, \dots, x_n : A_n$ and $\Gamma, x : A \mid t_i : B_i$ for $i = 1, \dots, m$, where (B_1, \dots, B_m) is the signature of R . For the interpretation of $\Gamma, x : A \mid R(t_1, \dots, t_m)$, by Definition 2.3.6 we have a pullback diagram:

$$\begin{array}{ccc} \llbracket \Gamma \mid R(t_1, \dots, t_m) \rrbracket & \longrightarrow & \llbracket R \rrbracket \\ \downarrow & & \downarrow \\ \llbracket \Gamma, x : A \rrbracket & \xrightarrow{\langle \llbracket t_1 \rrbracket, \dots, \llbracket t_m \rrbracket \rangle} & \llbracket B_1 \rrbracket \times \dots \times \llbracket B_m \rrbracket \end{array}$$

Now suppose $\Gamma \mid t : A$, and consider the substitution

$$\Gamma \mid R(t_1, \dots, t_m)[t/x] = \Gamma \mid R(t_1[t/x], \dots, t_m[t/x])$$

For the interpretations of the substituted terms $t_i[t/x]$ we have the composites

$$\llbracket t_i[t/x] \rrbracket = \llbracket t_i \rrbracket \circ \langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle : \llbracket \Gamma \rrbracket \longrightarrow \llbracket \Gamma, x : A \rrbracket \longrightarrow \llbracket B_i \rrbracket$$

by (associativity of composition and) the definition of the interpretation of terms. Thus for the interpretation of $\Gamma \mid R(t_1, \dots, t_m)[t/x]$ we have the outer pullback rectangle below.

$$\begin{array}{ccccc} & & \text{---} & & \\ & & \text{---} & & \\ \llbracket \Gamma \mid R(t_1, \dots, t_m)[t/x] \rrbracket & \dashrightarrow & \llbracket \Gamma, x : A \mid R(t_1, \dots, t_m) \rrbracket & \longrightarrow & \llbracket R \rrbracket \\ \downarrow & & \downarrow & & \downarrow \\ \llbracket \Gamma \rrbracket & \xrightarrow{\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle} & \llbracket \Gamma, x : A \rrbracket & \xrightarrow{\langle \llbracket t_1 \rrbracket, \dots, \llbracket t_m \rrbracket \rangle} & \llbracket B_1 \rrbracket \times \dots \times \llbracket B_m \rrbracket \\ & \searrow & & \nearrow & \\ & & \langle \llbracket t_1[t/x] \rrbracket, \dots, \llbracket t_m[t/x] \rrbracket \rangle & & \end{array}$$

But since the righthand square is a pullback, there is a unique dotted arrow as indicated. By the 2-pullbacks lemma, the lefthand square is then also a pullback, as required. \square

Exercise 2.3.8. Complete the proof.

When we deal with several different interpretations at once we may name them M, N, \dots , and subscript the semantic brackets accordingly, $\llbracket \Gamma \rrbracket_M, \llbracket \Gamma \rrbracket_N, \dots$

Definition 2.3.9. If $\Gamma \mid \psi \vdash \varphi$ is one of the logical entailment axioms of \mathbb{T} and

$$\llbracket \Gamma \mid \psi \rrbracket_M \leq \llbracket \Gamma \mid \varphi \rrbracket_M$$

holds in an interpretation M , then we say that M *satisfies* or *models* $\Gamma \mid \psi \vdash \varphi$ and write

$$M \models \Gamma \mid \psi \vdash \varphi.$$

An interpretation M is a *model* of \mathbb{T} if it satisfies all the axioms of \mathbb{T} .

Theorem 2.3.10 (Soundness of cartesian logic). *If a cartesian theory \mathbb{T} proves an entailment*

$$\Gamma \mid \psi \vdash \varphi$$

then every model M of \mathbb{T} satisfies the entailment:

$$M \models \Gamma \mid \psi \vdash \varphi.$$

Proof. The proof proceeds by induction on the proof of the entailment. In the following we often omit the typing context Γ to simplify notation, and all inequalities are interpreted in $\mathbf{Sub}(\llbracket \Gamma \rrbracket)$. We consider all possible last steps in the proof of the entailment:

1. Weakening: if $\llbracket \Gamma \mid \psi \rrbracket \leq \llbracket \Gamma \mid \varphi \rrbracket$ in $\mathbf{Sub}(\llbracket \Gamma \rrbracket)$ then

$$\llbracket \Gamma, x : A \mid \psi \rrbracket = \llbracket \Gamma \mid \psi \rrbracket \times A \leq \llbracket \Gamma \mid \varphi \rrbracket \times A = \llbracket \Gamma, x : A \mid \varphi \rrbracket \quad \text{in } \mathbf{Sub}(\llbracket \Gamma, x : A \rrbracket).$$

2. Substitution: by lemma 2.3.7, substitution is interpreted by pullback so that $\llbracket \varphi[t/x] \rrbracket = \langle \mathbf{1}_{\llbracket \psi \rrbracket}, \llbracket t \rrbracket \rangle^* \llbracket \varphi \rrbracket$ and $\llbracket \psi[t/x] \rrbracket = \langle \mathbf{1}_{\llbracket \psi \rrbracket}, \llbracket t \rrbracket \rangle^* \llbracket \psi \rrbracket$. Because

$$\langle \mathbf{1}_{\llbracket \psi \rrbracket}, \llbracket t \rrbracket \rangle^* : \mathbf{Sub}(\llbracket \psi \rrbracket) \rightarrow \mathbf{Sub}(\llbracket \psi \rrbracket \times \llbracket A \rrbracket)$$

is a functor it is a monotone map, therefore $\llbracket \psi \rrbracket \leq \llbracket \varphi \rrbracket$ implies

$$\langle \mathbf{1}_{\llbracket \psi \rrbracket}, \llbracket t \rrbracket \rangle^* \llbracket \psi \rrbracket \leq \langle \mathbf{1}_{\llbracket \psi \rrbracket}, \llbracket t \rrbracket \rangle^* \llbracket \varphi \rrbracket.$$

3. Identity: trivially

$$\llbracket \varphi \rrbracket \leq \llbracket \varphi \rrbracket.$$

4. Cut: if $\llbracket \psi \rrbracket \leq \llbracket \theta \rrbracket$ and $\llbracket \theta \rrbracket \leq \llbracket \varphi \rrbracket$ then clearly $\llbracket \psi \rrbracket \leq \llbracket \varphi \rrbracket$, since $\mathbf{Sub}(\llbracket \Gamma, x : A \rrbracket)$ is a poset.

5. Truth: trivially $\llbracket \psi \rrbracket \leq \llbracket \top \rrbracket$.
6. The rules for conjunction clearly hold because by the definition of infimum $\llbracket \vartheta \rrbracket \leq \llbracket \varphi \wedge \psi \rrbracket$ if, and only if, $\llbracket \vartheta \rrbracket \leq \llbracket \varphi \rrbracket$ and $\llbracket \vartheta \rrbracket \leq \llbracket \psi \rrbracket$.
7. Equality: the axiom $t =_A t$ is satisfied because an equalizer of $\llbracket t \rrbracket$ with itself is the maximal subobject:

$$\llbracket \psi \rrbracket \leq [1_{\llbracket \Gamma \rrbracket} : \llbracket \Gamma \rrbracket \rightarrow \llbracket \Gamma \rrbracket] = \llbracket t =_A t \rrbracket .$$

For the other axiom, suppose $\llbracket \psi \rrbracket \leq \llbracket t =_A u \rrbracket$ and $\llbracket \psi \rrbracket \leq \llbracket \varphi[t/z] \rrbracket$. It suffices to show $\llbracket t =_A u \rrbracket \wedge \llbracket \varphi[t/z] \rrbracket \leq \llbracket \varphi[u/z] \rrbracket$ for then

$$\llbracket \psi \rrbracket \leq \llbracket t =_A u \rrbracket \wedge \llbracket \varphi[t/z] \rrbracket \leq \llbracket \varphi[u/z] \rrbracket .$$

The interpretation of $P = \llbracket t =_A u \rrbracket \wedge \llbracket \varphi[t/z] \rrbracket$ is obtained by two successive pullbacks, as in the following diagram:

$$\begin{array}{ccccc} P & \xrightarrow{\quad} & \llbracket \varphi[t/z] \rrbracket & \xrightarrow{\quad} & \llbracket \varphi \rrbracket \\ \downarrow \lrcorner & & \downarrow \lrcorner & & \downarrow \lrcorner \\ \llbracket t =_A u \rrbracket & \xrightarrow{\quad e \quad} & \llbracket \Gamma \rrbracket & \xrightarrow{\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle} & \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket \end{array}$$

Here e is the equalizer of $\llbracket u \rrbracket$ and $\llbracket t \rrbracket$. Observe that e equalizes $\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle$ and $\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket u \rrbracket \rangle$ as well:

$$\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle \circ e = \langle e, \llbracket t \rrbracket \circ e \rangle = \langle e, \llbracket u \rrbracket \circ e \rangle = \langle 1_{\llbracket \Gamma \rrbracket}, \llbracket u \rrbracket \rangle \circ e .$$

Therefore, if we replace $\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket t \rrbracket \rangle$ with $\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket u \rrbracket \rangle$ in the above diagram, the outer rectangle still commutes. By the universal property of the pullback

$$\begin{array}{ccc} \llbracket \varphi[u/z] \rrbracket & \xrightarrow{\quad} & \llbracket \varphi \rrbracket \\ \downarrow \lrcorner & & \downarrow \lrcorner \\ \llbracket \Gamma \rrbracket & \xrightarrow{\langle 1_{\llbracket \Gamma \rrbracket}, \llbracket u \rrbracket \rangle} & \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket \end{array}$$

it follows that P also factors through $\llbracket \varphi[u/z] \rrbracket$, as required.

□

Example 2.3.11. Recall the cartesian theory of posets (example 2.3.2). There is one basic sort \mathbf{P} and one binary relation symbol \leq with signature (\mathbf{P}, \mathbf{P}) and the axioms of reflexivity, transitivity, and antisymmetry. A poset in a cartesian category \mathcal{C} is thus given by an object P , which is the interpretation of the sort \mathbf{P} , and a subobject $r : R \rightarrowtail P \times P$, which the interpretation of \leq , such that the axioms are satisfied. As an example we spell

out when the reflexivity axiom is satisfied. The interpretation of $x : P \mid x \leq x$ is obtained by the following pullback:

$$\begin{array}{ccc} \llbracket x \leq x \rrbracket & \xrightarrow{\quad} & R \\ \downarrow \lrcorner & \nearrow \rho & \downarrow r \\ P & \xrightarrow{\Delta} & P \times P \end{array}$$

where $\Delta = \langle 1_P, 1_P \rangle$ is the diagonal. The first axiom is satisfied when $\llbracket x \leq x \rrbracket = 1_P$, which happens if, and only if, Δ factors through r , as indicated. Therefore, reflexivity can be expressed as follows: there exists a “reflexivity” morphism $\rho : P \rightarrow R$ such that $r \circ \rho = \Delta$. Equivalently, the morphisms $\pi_0 \circ r$ and $\pi_1 \circ r$ have a common right inverse ρ .

As an example, of a poset in a cartesian category other than **Set**, observe that since the definition is stated entirely in terms of finite limits, and these are computed pointwise in functor categories $\mathbf{Set}^{\mathbb{C}}$, it follows that a poset P in $\mathbf{Set}^{\mathbb{C}}$ is the same thing as a functor $P : \mathbb{C} \rightarrow \mathbf{Poset}$. Indeed, as was the case for algebraic theories, we have an equivalence (an isomorphism, actually) of categories,

$$\mathbf{Poset}(\mathbf{Set}^{\mathbb{C}}) \cong \mathbf{Poset}(\mathbf{Set})^{\mathbb{C}} \cong \mathbf{Poset}^{\mathbb{C}}.$$

2.3.1 Subtypes

Let us consider whether the theory of a category is a cartesian theory. We begin by expressing the definition of a category so that it can be interpreted in any cartesian category \mathcal{C} . An *internal category* in \mathcal{C} consists of an *object of morphisms* C_1 , an *object of objects* C_0 , and *domain*, *codomain*, and *identity* morphisms,

$$\text{dom} : C_1 \rightarrow C_0, \quad \text{cod} : C_1 \rightarrow C_0, \quad \text{id} : C_0 \rightarrow C_1.$$

There is also a *composition* morphism $c : C_2 \rightarrow C_1$, where C_2 is obtained by the pullback

$$\begin{array}{ccc} C_2 & \xrightarrow{p_1} & C_1 \\ \downarrow p_0 \lrcorner & & \downarrow \text{dom} \\ C_1 & \xrightarrow{\text{cod}} & C_0 \end{array}$$

The following equations must hold:

$$\begin{aligned} \text{dom} \circ i &= 1_{C_0} = \text{cod} \circ i, \\ \text{cod} \circ p_1 &= \text{cod} \circ c, \quad \text{dom} \circ p_0 = \text{dom} \circ c, \\ c \circ \langle 1_{C_1}, i \circ \text{dom} \rangle &= 1_{C_1} = c \circ \langle i \circ \text{cod}, 1_{C_1} \rangle, \end{aligned}$$

The first two equations state that the domain and codomain of an identity morphism 1_A are both A . The second equation states that $\text{cod}(f \circ g) = \text{cod} f$ and the third one that

$\text{dom}(f \circ g) = \text{dom } g$. The fourth equation states that $f \circ 1_{\text{dom } f} = f = 1_{\text{cod } f} \circ f$. It remains to express associativity of composition. For this purpose we construct the pullback

$$\begin{array}{ccc} C_3 & \xrightarrow{q_2} & C_1 \\ q_{01} \downarrow & \lrcorner & \downarrow \text{dom} \\ C_2 & \xrightarrow{\text{cod} \circ p_1} & C_0 \end{array}$$

The object C_3 can be thought of as the set of triples of morphisms (f, g, h) such that $\text{cod } f = \text{dom } g$ and $\text{cod } g = \text{dom } h$. We denote $q_0 = p_0 \circ q_{01}$ and $q_1 = p_1 \circ q_{01}$. The morphisms $q_0, q_1, q_2 : C_3 \rightarrow C_1$ are like three projections which select the first, second, and third element of a triple, respectively. With this notation we can write $q_{01} = \langle q_0, q_1 \rangle_{C_2}$ because q_{01} is the unique morphism such that $p_0 \circ q_{01} = q_0$ and $p_1 \circ q_{01} = q_1$. The subscript C_2 reminds us that the “pair” $\langle q_0, q_1 \rangle_{C_2}$ is obtained by the universal property of the pullback C_2 .

Morphisms $c \circ q_{01} : C_3 \rightarrow C_1$ and $q_2 : C_3 \rightarrow C_1$ factor through the pullback C_2 because

$$\text{cod} \circ c \circ q_{01} = \text{cod} \circ p_1 \circ q_0 = \text{dom} \circ q_2 .$$

Thus let $r : C_3 \rightarrow C_2$ be the unique factorization for which $p_0 \circ r = c \circ q_{01}$ and $p_1 \circ r = q_2$. Because p_0 and p_1 are like projections from C_2 to C_1 , morphism r can be thought of as a pair of morphisms, so we write $r = \langle c \circ q_{01}, q_2 \rangle_{C_2}$. Morphism $c \circ \langle c \circ q_{01}, q_2 \rangle_{C_2} : C_3 \rightarrow C_1$ corresponds to the operations $\langle f, g, h \rangle \mapsto (f, g) \circ h$, whereas the morphism corresponding to $\langle f, g, h \rangle \mapsto f \circ (g \circ h)$ is obtained in a similar way and is equal to

$$c \circ \langle q_0, c \circ \langle q_1, q_2 \rangle_{C_2} \rangle_{C_2} : C_3 \rightarrow C_1 .$$

Thus associativity is expressed by the equation

$$c \circ \langle c \circ \langle q_0, q_1 \rangle_{C_2}, q_2 \rangle_{C_2} = c \circ \langle q_0, c \circ \langle q_1, q_2 \rangle_{C_2} \rangle_{C_2} .$$

Example 2.3.12. An internal category in \mathbf{Set} is a small category.

Example 2.3.13. An internal category in $\mathbf{Set}^{\mathbb{C}}$ is a functor $\mathbb{C} \rightarrow \mathbf{Cat}$. Indeed, as in previous examples of cartesian theories we have an equivalence of categories,

$$\mathbf{Cat}(\mathbf{Set}^{\mathbb{C}}) \cong \mathbf{Cat}(\mathbf{Set})^{\mathbb{C}} \cong \mathbf{Cat}^{\mathbb{C}} .$$

We have successfully formulated the theory of a category so that it makes sense in any cartesian category. In fact, the definition of an internal category refers only to certain pullbacks, hence the notion of an internal category makes sense in any category with pullbacks. However, if we try to formulate it as a multi-sorted cartesian theory, there is

a problem. Obviously, there ought to be a basic sort of objects \mathbf{C}_0 and a basic sort of morphisms \mathbf{C}_1 . There are also basic function symbols with signatures

$$\text{dom} : (\mathbf{C}_1; \mathbf{C}_0) \qquad \text{cod} : (\mathbf{C}_1; \mathbf{C}_0) \qquad \text{id} : (\mathbf{C}_0, \mathbf{C}_1) .$$

However, it is not clear what the signature for composition should be. It is not $(\mathbf{C}_1, \mathbf{C}_1; \mathbf{C}_1)$ because composition is undefined for non-composable pairs of morphisms. We might be tempted to postulate another basic sort \mathbf{C}_2 but then we would have no way of stating that \mathbf{C}_2 is the pullback of dom and cod . And even if we somehow axiomatized the fact that \mathbf{C}_2 is a pullback, we would then still have to formalize the object \mathbf{C}_3 of composable triples, \mathbf{C}_4 of composable quadruples, and so on. What we lack is the ability to define the type \mathbf{C}_2 as a *subtype* of $\mathbf{C}_1 \times \mathbf{C}_1$.

One way to remedy the situation is to use a richer underlying type theory; in Chapter ?? we will consider the system of *dependent type theory*, which provides the means to capture such notions as the theory of categories (and much more). Here we consider a small step in that direction, namely *simple subtypes*. The formation rule for simple subtypes is

$$\frac{x : A \mid \varphi \text{ pred}}{\{x : A \mid \varphi\} \text{ type}}$$

We can think of $\{x : A \mid \varphi\}$ as the subobject of all those $x : A$ that satisfy φ . Note that we did not allow an arbitrary context Γ to be present. This means that we cannot define subtypes that depend on parameters, which is why they are called “simple”.

Inference rules for subtypes are as follows:

$$\frac{\Gamma \mid t : \{x : A \mid \varphi\}}{\Gamma \mid \text{in}_\varphi t : A} \qquad \frac{\Gamma \mid t : \{x : A \mid \varphi\}}{\Gamma \mid \cdot \vdash \varphi[\text{in}_\varphi t/x]} \qquad \frac{\Gamma \mid t : A \quad \Gamma \mid \cdot \vdash \varphi[t/x]}{\Gamma \mid \text{rs}_\varphi t : \{x : A \mid \varphi\}}$$

$$\frac{\Gamma, x : A \mid \varphi, \psi \vdash \theta}{\Gamma, y : \{x : A \mid \varphi\} \mid \psi[\text{in}_\varphi y/x] \vdash \theta[\text{in}_\varphi y/x]}$$

The first rule states that a term t of subtype $\{x : A \mid \varphi\}$ can be converted to a term $\text{in}_\varphi t$ of type A . We can think of the constant in_φ as the *inclusion* $\text{in}_\varphi : \{x : A \mid \varphi\} \rightarrow A$. The second rule states that every term of a subtype $\{x : A \mid \varphi\}$ satisfies the defining predicate φ . The third rule states that a term t of type A which satisfies φ can be converted to a term $\text{rs}_\varphi t$ of type $\{x : A \mid \varphi\}$. A good way to think of the constant rs_φ is as a partially defined *restriction*, or a type-casting operations, $\text{rs}_\varphi : A \rightarrow \{x : A \mid \varphi\}$.⁴ The last rule tells us how to replace a variable x of type A and an assumption φ about it with a variable y of type $\{x : A \mid \varphi\}$ and remove the assumption. Note that this is a two-way rule.

There are two more axioms that relate inclusions and restrictions:

$$\frac{\Gamma \mid t : \{x : A \mid \varphi\}}{\Gamma \mid \cdot \vdash \text{rs}_\varphi(\text{in}_\varphi t) = t} \qquad \frac{\Gamma \mid t : A \quad \Gamma \mid \cdot \vdash \varphi[t/x]}{\Gamma \mid \cdot \vdash \text{in}_\varphi(\text{rs}_\varphi t) = t} .$$

⁴Inclusions and restrictions are like type-casting operations in some programming languages. For example in Java, an inclusion corresponds to an (implicit) type cast from a class to its superclass, whereas a restriction corresponds to a type cast from a class to a subclass. Must I write that Java is a registered trademark of Sun Microsystems?

In an informal discussion it is customary for the inclusions and restrictions to be omitted, or at least for the subscript φ to be missing.⁵

Exercise 2.3.14. Suppose $x : A \mid \psi$ and $x : A \mid \varphi$ are formulas. Show that

$$x : A \mid \psi \vdash \varphi$$

is provable if, and only if, $\{x : A \mid \psi\}$ factors through $\{x : A \mid \varphi\}$, which means that there exists a term k ,

$$y : \{x : A \mid \psi\} \mid k : \{x : A \mid \varphi\} ,$$

such that

$$y : \{x : A \mid \psi\} \mid \cdot \vdash \text{in}_\psi y =_A \text{in}_\varphi k$$

is provable. Show also that k is determined uniquely up to provable equality.

Example 2.3.15. We are now able to formulate the theory of a category as a cartesian theory whose underlying type theory has product types and subset types. The basic types are the type of objects \mathbf{C}_0 and the type of morphisms \mathbf{C}_1 . We define the type \mathbf{C}_2 to be

$$\mathbf{C}_2 \equiv \{p : \mathbf{C}_1 \times \mathbf{C}_1 \mid \text{cod}(\text{fst } p) = \text{dom}(\text{snd } p)\} .$$

The basic function symbols and their signatures are:

$$\text{dom} : \mathbf{C}_1 \rightarrow \mathbf{C}_0 , \quad \text{cod} : \mathbf{C}_1 \rightarrow \mathbf{C}_0 , \quad \text{id} : \mathbf{C}_0 \rightarrow \mathbf{C}_1 , \quad \text{c} : \mathbf{C}_2 \rightarrow \mathbf{C}_1 .$$

The axioms are:

$$\begin{aligned} a : \mathbf{C}_0 \mid \cdot &\vdash \text{dom}(\text{id}(a)) = a \\ a : \mathbf{C}_0 \mid \cdot &\vdash \text{cod}(\text{id}(a)) = a \\ f : \mathbf{C}_1, g : \mathbf{C}_1 \mid \text{cod}(f) = \text{dom}(g) &\vdash \text{dom}(\text{c}(\text{rs } \langle f, g \rangle)) = f \\ f : \mathbf{C}_1, g : \mathbf{C}_1 \mid \text{cod}(f) = \text{dom}(g) &\vdash \text{cod}(\text{c}(\text{rs } \langle f, g \rangle)) = g \\ f : \mathbf{C}_1 \mid \cdot &\vdash \text{c}(\text{rs } \langle \text{id}(\text{dom}(f)), f \rangle) = f \\ f : \mathbf{C}_1 \mid \cdot &\vdash \text{c}(\text{rs } \langle f, \text{id}(\text{cod}(f)) \rangle) = f \end{aligned}$$

Lastly, the associativity axiom is

$$\begin{aligned} f : \mathbf{C}_1, g : \mathbf{C}_1, h : \mathbf{C}_1 \mid \text{cod}(f) = \text{dom}(g), \text{cod}(g) = \text{dom}(h) &\vdash \\ \text{c}(\text{rs } \langle \text{c}(\text{rs } \langle f, g \rangle), h \rangle) &= \text{c}(\text{rs } \langle f, \text{c}(\text{rs } \langle g, h \rangle) \rangle) . \end{aligned}$$

This notation is quite unreadable. If we write $g \circ f$ instead of $\text{c}(\text{rs } \langle f, g \rangle)$ then the axioms take on a more familiar form. For example, associativity is just $h \circ (g \circ f) = (h \circ g) \circ f$. However, we need to remember that we may form the term $g \circ f$ only if we first prove $\text{dom}(g) = \text{cod}(f)$.

⁵Strictly speaking, even the notation $\text{in}_\varphi t$ is imprecise because it does not indicate that ϕ stands in the context $x : A$. The correct notation would be $\text{in}_{(x:A \mid \varphi)} t$, where x is bound in the subscript. A similar remark holds for $\text{rs}_\varphi t$.

A subtype $\{x : A \mid \varphi\}$ is interpreted as the domain of a monomorphism representing $x : A \mid \varphi$:

$$\llbracket \{x : A \mid \varphi\} \rrbracket \xrightarrow{\llbracket x : A \mid \varphi \rrbracket} \llbracket A \rrbracket$$

Some care must be taken here because monos representing a given subobject are only determined up to isomorphism. We assume that a suitable canonical choice of monos can be made.

An inclusion $\Gamma \mid \text{in}_\varphi t : A$ is interpreted as the composition

$$\llbracket \Gamma \rrbracket \xrightarrow{\llbracket t \rrbracket} \llbracket \{x : A \mid \varphi\} \rrbracket \xrightarrow{\llbracket x : A \mid \varphi \rrbracket} \llbracket A \rrbracket$$

A restriction $\Gamma \mid \text{rs}_\varphi t : \{x : A \mid \varphi\}$ is interpreted as the unique $\llbracket \overline{t} \rrbracket$ which makes the following diagram commute:

$$\begin{array}{ccc} \llbracket \Gamma \rrbracket & \xrightarrow{\llbracket \overline{t} \rrbracket} & \llbracket \{x : A \mid \varphi\} \rrbracket \\ & \searrow \llbracket t \rrbracket & \downarrow \\ & & \llbracket A \rrbracket \end{array}$$

Exercise 2.3.16. Formulate and prove a soundness theorem for subtypes. Pay attention to the interpretation of restrictions, where you need to show unique existence of $\llbracket \overline{t} \rrbracket$.

Remark 2.3.17. Another approach to the logic of cartesian categories that captures the theory of categories and related notions involving partial operations is that of *essentially algebraic theories*, due to P. Freyd; see [?, ?]. A third approach is that of *dependent type theory* to be developed in ?? below.

2.4 Quantifiers as adjoints

The categorical semantics of quantification is one of the central features of the subject, and quite possibly one of the nicest contributions of categorical logic to the field of logic. You might expect that the quantifiers \forall and \exists are “just a big conjunction and disjunction”, respectively. In fact the Polish school of algebraic logicians worked to realize this point of view—but categorical logic shows how quantifiers are treated algebraically as adjoint functors to give a much more satisfactory theory. The original treatment can be found in the classic paper [?].

Let us first recall the rules of inference for quantifiers. The formation rules are:

$$\frac{\Gamma, x : A \mid \varphi \text{ pred}}{\Gamma \mid (\forall x : A. \varphi) \text{ pred}} \qquad \frac{\Gamma, x : A \mid \varphi \text{ pred}}{\Gamma \mid (\exists x : A. \varphi) \text{ pred}}$$

The variable x is bound in $\forall x : A . \varphi$ and $\exists x : A . \varphi$. If x and y are distinct variables and x does not occur freely in the term t then substitution of t for y commutes with quantification over x :

$$\begin{aligned} (\exists x : A . \varphi)[t/y] &= \exists x : A . (\varphi[t/y]) , \\ (\forall x : A . \varphi)[t/y] &= \forall x : A . (\varphi[t/y]) . \end{aligned}$$

For each quantifier we have a two-way rule of inference:

$$\frac{\Gamma, x : A \mid \psi \vdash \varphi}{\Gamma \mid \psi \vdash \forall x : A . \varphi} \qquad \frac{\Gamma, x : A \mid \varphi \vdash \vartheta}{\Gamma \mid (\exists x : A . \varphi) \vdash \vartheta}$$

Note that these rules implicitly impose the usual condition that x must not occur freely in ψ and ϑ , because ψ and ϑ are supposed to be well formed in context Γ , which does not contain x .

Exercise 2.4.1. A common way of stating the inference rules for quantifiers is as follows. For the universal quantifier, the introduction and elimination rules are

$$\frac{\Gamma, x : A \mid \psi \vdash \varphi}{\Gamma \mid \psi \vdash \forall x : A . \varphi} \qquad \frac{\Gamma \mid t : A \quad \Gamma \mid \psi \vdash \forall x : A . \varphi}{\Gamma \mid \psi \vdash \varphi[t/x]}$$

The introduction rule for existential quantifier is

$$\frac{\Gamma \mid t : A \quad \Gamma \mid \psi \vdash \varphi[t/x]}{\Gamma \mid \psi \vdash \exists x : A . \varphi}$$

and the elimination rule is

$$\frac{\Gamma \mid \psi \vdash \exists x : A . \varphi \quad \Gamma, x : A \mid \varphi \vdash \vartheta}{\Gamma \mid \psi \vdash \vartheta}$$

Note that these rules implicitly impose a requirement that x does not occur in Γ and that it does not occur freely in ψ because the context $\Gamma, x : A$ must be well formed and the hypotheses ψ must be well formed in context Γ . Show that these rules can be derived from the ones above, and vice versa. Of course, you may also use the inference rules for cartesian logic, cf. page 15.

In order to discover what the semantics of existential quantifier ought to be, we look at the following instance of the two-way rule for quantifiers:

$$\frac{y : B, x : A \mid \varphi \vdash \vartheta}{y : B \mid \exists x : A . \varphi \vdash \vartheta} \tag{2.2}$$

First observe that this rule implicitly requires

$$y : B, x : A \mid \varphi \text{ pred} \qquad y : B \mid \vartheta \text{ pred} \qquad y : B \mid (\exists x : A . \varphi) \text{ pred}$$

This is required for the entailments to be well-formed. The fourth judgement

$$y : B, x : A \mid \vartheta \text{ pred}$$

follows from the second one above by weakening,

$$\frac{y : B \mid \vartheta \text{ pred}}{y : B, x : A \mid \vartheta \text{ pred}}$$

The interpretations of φ , ϑ , and $\exists x : A. \varphi$ are therefore subobjects

$$\begin{aligned} \llbracket y : B, x : A \mid \varphi \rrbracket &\in \mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket) , \\ \llbracket y : B \mid \vartheta \rrbracket &\in \mathbf{Sub}(\llbracket B \rrbracket) , \\ \llbracket y : B \mid \exists x : A. \varphi \rrbracket &\in \mathbf{Sub}(\llbracket B \rrbracket) . \end{aligned}$$

And the weakened instance of ϑ in the context $y : B, x : A$ is interpreted by pullback along a projection, cf. page 19, as in the following pullback diagram:

$$\begin{array}{ccc} \llbracket y : B, x : A \mid \vartheta \rrbracket & \xrightarrow{\quad} & \llbracket B \rrbracket \times \llbracket A \rrbracket \\ \downarrow \lrcorner & & \downarrow \pi \\ \llbracket y : B \mid \vartheta \rrbracket & \xrightarrow{\quad} & \llbracket B \rrbracket \end{array}$$

Thus we have

$$\llbracket y : B, x : A \mid \vartheta \rrbracket = \pi^* \llbracket y : B \mid \vartheta \rrbracket ,$$

with weakening interpreted as the pullback functor

$$\pi^* : \mathbf{Sub}(\llbracket B \rrbracket) \rightarrow \mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket) .$$

We will interpret existential quantification $\exists x : A$ as a suitable functor

$$\exists_A : \mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket) \rightarrow \mathbf{Sub}(\llbracket B \rrbracket)$$

so that

$$\llbracket y : B \mid \exists x : A. \varphi \rrbracket = \exists_A \llbracket y : B, x : A \mid \varphi \rrbracket .$$

The interpretation of the two-way rule (2.2) then becomes a two-way inequality rule

$$\frac{\llbracket y : B, x : A \mid \varphi \rrbracket \leq \pi^* \llbracket y : B \mid \vartheta \rrbracket}{\exists_A \llbracket y : B, x : A \mid \varphi \rrbracket \leq \llbracket y : B \mid \vartheta \rrbracket}$$

Replacing the interpretations of φ and ϑ by general subobjects $S \in \mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket)$ and $T \in \mathbf{Sub}(\llbracket B \rrbracket)$, we obtain the more suggestive formulation

$$\frac{S \leq \pi^* T}{\exists_A S \leq T} \tag{2.3}$$

This is nothing but an adjunction between \exists_A and π^* ! Indeed, the operations \exists_A and π^* are functors on the subjects posets $\mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket)$ and $\mathbf{Sub}(\llbracket B \rrbracket)$, and the bijection of hom-sets (2.3) is exactly the statement of an adjunction between them. Thus *existential quantification is left-adjoint to weakening*:

$$\exists_A \dashv \pi^*$$

An exactly dual argument shows that *universal quantification is right-adjoint to weakening*:

$$\pi^* \dashv \forall_A$$

Thus, in sum, we have shown that the rules of inference require the quantifiers to be interpreted as operations that are adjoints to the interpretation of weakening, i.e. pullback π^* along the projection $\pi : \llbracket B \rrbracket \times \llbracket A \rrbracket \rightarrow \llbracket B \rrbracket$.

$$\begin{array}{ccc} & \mathbf{Sub}(\llbracket B \rrbracket \times \llbracket A \rrbracket) & \\ \exists \swarrow & \uparrow \pi^* & \searrow \forall \\ & \dashv \quad \dashv & \\ & \mathbf{Sub}(\llbracket B \rrbracket) & \end{array}$$

Note that the familiar side-conditions on the conventional rules for the quantifiers, to the effect that “ x cannot occur freely in ψ ”, etc., which may seem like tiresome bookkeeping, are actually of the essence, since they express the weakening operation to which the quantifiers themselves are adjoints.

Let us see how this works for the usual interpretation in **Set**. A predicate $y : B, x : A \mid \varphi$ corresponds to a subset $\Phi \subseteq B \times A$, and $y : B \mid \vartheta$ corresponds to a subset $\Theta \subseteq B$. Weakening of Θ is the subset $\pi^*\Theta = \Theta \times A \subseteq B \times A$. Then we have

$$\begin{aligned} \exists_A \Phi &= \{y \in B \mid \exists x : A. \langle x, y \rangle \in \Phi\} \subseteq B, \\ \forall_A \Phi &= \{y \in B \mid \forall x : A. \langle x, y \rangle \in \Phi\} \subseteq B. \end{aligned}$$

A moment’s thought convinces us that with this interpretation we do indeed have

$$\frac{\Phi \subseteq \Theta \times A}{\exists_A \Phi \subseteq \Theta} \qquad \frac{\Theta \times A \subseteq \Phi}{\Theta \subseteq \forall_A \Phi}$$

The unit of the adjunction $\exists_A \dashv \pi^*$ amounts to the inequality

$$\Phi \subseteq (\exists_A \Phi) \times A, \tag{2.4}$$

and the universal property of the unit says that $\exists_A \Phi$ is the smallest set satisfying (2.4). Similarly, the counit of the adjunction $\pi^* \dashv \forall_A$ is just the inequality

$$(\forall_A \Phi) \times A \subseteq \Phi, \quad (2.5)$$

and the universal property of the counit says that $\forall_A \Phi$ is the largest set satisfying (2.5). Figure 2.1 shows the geometric meaning of existential and universal quantification.

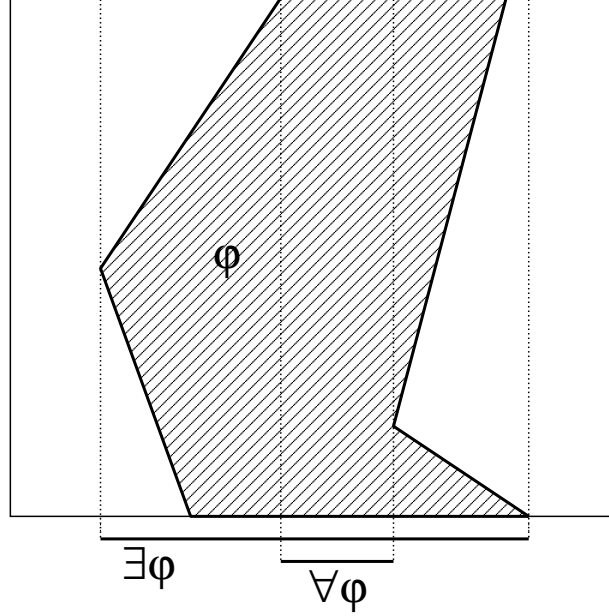


Figure 2.1: $\exists \varphi$ and $\forall \varphi$

Exercise 2.4.2. What do the universal properties of the counit of $\exists_A \dashv \pi^*$ and the unit of $\pi^* \dashv \forall_A$ say?

The weakening functor π^* is a special case of a pullback functor $f^* : \mathbf{Sub}(B) \rightarrow \mathbf{Sub}(A)$ for a morphism $f : B \rightarrow A$. This gives us the idea that we may regard the left and the right adjoint to f^* as a kind of generalized existential and universal quantifier.

We may also be tempted to *define* quantifiers as left and right adjoints to pullback functors. However there is a bit more to quantifiers than that—we are still missing the important *Beck-Chevalley condition*.

2.4.1 The Beck-Chevalley condition

Recall that quantification commutes with substitution, as long as no variables are captured by the quantifier. Thus if $\Gamma \mid t : B$ and $\Gamma, y : B, x : A \mid \varphi$ pred then

$$\begin{aligned} (\exists x : A. \varphi)[t/y] &= \exists x : A. (\varphi[t/y]) \quad . \\ (\forall x : A. \varphi)[t/y] &= \forall x : A. (\varphi[t/y]) \quad . \end{aligned}$$

If semantics of quantifiers is to be sound, the interpretation of these equations must be valid. Because substitution of a term in a formula is interpreted as pullback this means that quantifiers must be *stable* under pullbacks. This is known as the *Beck-Chevalley condition*.

Definition 2.4.3. A family of functors $F_f : \mathbf{Sub}(A) \rightarrow \mathbf{Sub}(B)$ parametrized by morphisms $f : A \rightarrow B$ is said to satisfy the *Beck-Chevalley condition* when for every pullback on the left-hand side, the right-hand square commutes:

$$\begin{array}{ccc} C & \xrightarrow{h} & A \\ k \downarrow \lrcorner & & \downarrow f \\ D & \xrightarrow{g} & B \end{array} \qquad \begin{array}{ccc} \mathbf{Sub}(C) & \xleftarrow{h^*} & \mathbf{Sub}(A) \\ F_k \downarrow & & \downarrow F_f \\ \mathbf{Sub}(D) & \xleftarrow{g^*} & \mathbf{Sub}(B) \end{array}$$

To convince ourselves that Beck-Chevalley condition is what we want, we spell it out explicitly in the case of a substitution into an existentially quantified formula. In order to keep the notation simple we omit the semantic brackets $\llbracket - \rrbracket$. Suppose we have a term $\Gamma \mid t : B$ and a formula $\Gamma, y : B, x : A \mid \varphi$ pred. The diagram

$$\begin{array}{ccc} \Gamma \times A & \xrightarrow{\langle \pi_0, t \circ \pi_0, \pi_1 \rangle} & \Gamma \times B \times A \\ \pi_0^{\Gamma, A} \downarrow \lrcorner & & \downarrow \pi_0^{\Gamma, B, A} \\ \Gamma & \xrightarrow{\langle 1_\Gamma, t \rangle} & \Gamma \times B \end{array}$$

is a pullback. By Beck-Chevalley condition for \exists , the following square commutes:

$$\begin{array}{ccc} \mathbf{Sub}(\Gamma \times A) & \xleftarrow{\langle \pi_0, t \circ \pi_0, \pi_1 \rangle^*} & \mathbf{Sub}(\Gamma \times B \times A) \\ \exists_A^{\Gamma, A} \downarrow & & \downarrow \exists_A^{\Gamma, B, A} \\ \mathbf{Sub}(\Gamma) & \xleftarrow{\langle 1_\Gamma, t \rangle^*} & \mathbf{Sub}(\Gamma \times B) \end{array}$$

Therefore, for $\Gamma, y : B, x : A \mid \varphi$ pred,

$$\begin{aligned} \llbracket (\exists x : A. \varphi)[t/y] \rrbracket &= \langle 1_\Gamma, t \rangle^* (\exists_A^{\Gamma, B, A} \llbracket \varphi \rrbracket) = \\ &= \exists_A^{\Gamma, A} (\langle \pi_0, t \circ \pi_0, \pi_1 \rangle^* \llbracket \varphi \rrbracket) = \llbracket \exists x : A. (\varphi[t/y]) \rrbracket. \end{aligned}$$

This is precisely the equation we wanted. The Beck-Chevalley condition says that (interpretations of) the quantifiers commute with pullbacks, in just the way that the syntactic operations of applying quantifiers to formulas commute with substitutions of terms, which are interpreted as pullbacks.

Definition 2.4.4. A cartesian category \mathcal{C} has *existential quantifiers* if, for every $f : A \rightarrow B$, the left adjoint $\exists_f \dashv f^*$ exists and it satisfies the Beck-Chevalley condition. Similarly, \mathcal{C} has *universal quantifiers* if the right adjoints $f^* \dashv \forall_f$ exist and they satisfy the Beck-Chevalley condition.

Given both adjoints $\exists_f \dashv f^* \dashv \forall_f$, it actually suffices to have the Beck-Chevalley condition for either one in order to infer it for both:

Proposition 2.4.5. *If for every $f : A \rightarrow B$, both the left and right adjoints exist*

$$\exists_f \dashv f^* \dashv \forall_f$$

then the left adjoint satisfies the Beck-Chevalley condition iff the right adjoint does.

Proof. Suppose we have the Beck-Chevalley condition for the left adjoints \exists , and that we are given a pullback square as on the left below. We want to check the Beck-Chevalley square for the right adjoints \forall , as indicated on the right below.

$$\begin{array}{ccc} C & \xrightarrow{h} & A \\ \downarrow k & \lrcorner & \downarrow f \\ D & \xrightarrow{g} & B \end{array} \qquad \begin{array}{ccc} \text{Sub}(C) & \xleftarrow{h^*} & \text{Sub}(A) \\ \downarrow \forall_k & & \downarrow \forall_f \\ \text{Sub}(D) & \xleftarrow{g^*} & \text{Sub}(B) \end{array}$$

Swapping all the functors in the righthand diagram for their left adjoints we obtain the following.

$$\begin{array}{ccc} \text{Sub}(C) & \xrightarrow{\exists_h} & \text{Sub}(A) \\ \uparrow k^* & & \uparrow f^* \\ \text{Sub}(D) & \xrightarrow{\exists_g} & \text{Sub}(B) \end{array}$$

But this is a Beck-Chevalley square for (the “transpose” of) the original pullback diagram, and therefore commutes by the Beck-Chevalley condition for the left adjoints \exists . The original diagram of right adjoints therefore also commutes, by uniqueness of adjoints.

The argument for the dual case is, well, dual. \square

Exercise 2.4.6. In **Set** we can identify $\text{Sub}(-)$ with powersets because $\text{Sub}(X) \cong \mathcal{P}X$. Then quantifiers along a function $f : A \rightarrow B$ are functions

$$\exists_f : \mathcal{P}A \rightarrow \mathcal{P}B, \qquad \forall_f : \mathcal{P}A \rightarrow \mathcal{P}B.$$

Verify that

$$\begin{aligned} \exists_f U &= \{b \in B \mid \exists a : A. (fa = b \wedge a \in U)\}, \\ \forall_f U &= \{b \in B \mid \forall a : A. (fa = b \Rightarrow a \in U)\}. \end{aligned}$$

Thus $\exists_f U$ is just the usual direct image of U by f , sometimes written $f_!(U)$, or simply $f(U)$. But have you seen $\forall_f U$ before? It can also be written as $\forall_f U = \{b \in B \mid f^* \{b\} \subseteq U\}$. What is the meaning of \exists_q and \forall_q when $q : A \rightarrow A/\sim$ is a canonical quotient map that maps an element $x \in A$ to its equivalence class $qx = [x]$ under an equivalence relation \sim on A ?

2.5 Regular logic

We next consider the question of when a cartesian category has existential quantifiers. It turns out that this is closely related to the notion of a *regular category*, a concept which first arose in the context of abelian categories and axiomatic homology theory, quite independently of categorical logic. We will see for instance that all algebraic categories, in the sense of Chapter ??, are regular.

2.5.1 Regular categories

Throughout this section we work in a cartesian category \mathcal{C} . We begin with some general definitions. The *kernel pair* of a morphism $f : A \rightarrow B$ is the pair of morphisms $k_1, k_2 : K \rightrightarrows A$ obtained as in the following pullback

$$\begin{array}{ccc} K & \xrightarrow{k_2} & A \\ k_1 \downarrow & \lrcorner & \downarrow f \\ A & \xrightarrow{f} & B \end{array}$$

Note that a kernel pair determines an equivalence relation $\langle k_1, k_2 \rangle : K \rightrightarrows A \times A$, in the sense that the map $\langle k_1, k_2 \rangle$ is a mono that satisfies the reflexivity, symmetry and transitivity conditions. In **Set** the mono $\langle k_1, k_2 \rangle : K \rightrightarrows A \times A$ is the equivalence relation \sim on A defined by

$$x \sim y \iff fx = fy.$$

Indeed, an equivalence relation in a general cartesian category is precisely a model of the cartesian theory of an equivalence relation, in the sense of example 2.3.3.

Exercise 2.5.1. Prove this.

In general, the *quotient* by the equivalence relation determined by the kernel pair k_1, k_2 is their coequalizer $q : A \rightarrow Q$, if it exists,

$$K \begin{array}{c} \xrightarrow{k_1} \\ \xrightarrow{k_2} \end{array} A \xrightarrow{q} Q$$

Such a coequalizer is called a *kernel quotient*.

Because $f \circ k_1 = f \circ k_2$, we see that f factors through q by a unique morphism $m : Q \rightarrow B$,

$$\begin{array}{ccccc} K & \xrightarrow[k_2]{k_1} & A & \xrightarrow{f} & B \\ & & \searrow q & & \nearrow m \\ & & Q & & \end{array} \quad (2.6)$$

As a coequalizer, $q : A \rightarrow Q$ is always epic; indeed, epis that are coequalizers will be called *regular epimorphisms* and will be denoted by arrows with triangular heads:

$$e : A \longrightarrow\triangleright B$$

It is of some interest to know when the second factor $m : Q \rightarrow B$ in (2.6) is guaranteed to be a mono. For example, in **Set** the function $m : Q \rightarrow B$ is defined by $m[x] = fx$, where $Q = A/\sim$ as above. In this case m is indeed injective, because $m[x] = m[y]$ implies $fx = fy$, hence $x \sim y$ and $[x] = [y]$.

Definition 2.5.2. A category with finite limits is *regular* when it has kernel quotients, and regular epis are stable under pullback. Thus, in detail:

1. the kernel pair of any map has a coequalizer, and
2. any pullback of a regular epi is a regular epi.

Exercise 2.5.3. Suppose $e : A \longrightarrow\triangleright B$ is a regular epi. Prove that it is the coequalizer of its own kernel pair.

Let us return to (2.6) and show that m is monic in any regular category. Consider the following diagram, in which h_1, h_2 are constructed as the kernel pair of m , and the other three squares are constructed as pullbacks:

$$\begin{array}{ccccc} K & \xrightarrow{p_2} & \cdot & \longrightarrow & A \\ \downarrow p_1 & \searrow r & \downarrow s_2 & \lrcorner & \downarrow q \\ \cdot & \xrightarrow{s_1} & H & \xrightarrow{h_2} & Q \\ \downarrow \lrcorner & & \downarrow h_1 & \lrcorner & \downarrow m \\ A & \xrightarrow{q} & Q & \xrightarrow{m} & B \end{array}$$

Because all the smaller squares are pullbacks the large square is a pullback as well, therefore the left-hand vertical morphism is $k_1 : K \rightarrow A$, and the morphism across the top is $k_2 : K \rightarrow A$, and we have the kernel pair $k_1, k_2 : K \rightrightarrows A$ of $f = m \circ q$. The morphisms $s_1,$

s_2 , p_1 , and p_2 are all regular epis because they are pullbacks of the regular epi q . The morphism $r = s_2 \circ p_2 = s_1 \circ p_1$ is epic because it is a composition of regular epis. Observe that

$$h_1 \circ r = q \circ k_1 = q \circ k_2 = h_2 \circ r ,$$

and so, because r is epic, $h_1 = h_2$. But this means that m is monic, since the maps in its kernel pair are equal; indeed, given any $u, v : U \rightarrow Q$ with $m \circ u = m \circ v$, there exists a $w : U \rightarrow H$ such that $u = w \circ h_1 = w \circ h_2 = v$.

Proposition 2.5.4. *In a regular category every morphism $f : A \rightarrow B$ factors as a composition of a regular epi q followed by a mono m ,*

$$\begin{array}{ccccc} & & f & & \\ & \curvearrowright & & \curvearrowright & \\ A & \xrightarrow{q} & Q & \xrightarrow{m} & B \end{array}$$

The factorization is unique up to isomorphism.

Proof. By uniqueness of the factorization we mean that if

$$\begin{array}{ccccc} & & f & & \\ & \curvearrowright & & \curvearrowright & \\ A & \xrightarrow{q'} & Q' & \xrightarrow{m'} & B \end{array}$$

is another such factorization, then there exists an isomorphism $i : Q \rightarrow Q'$ such that $q' = i \circ q$ and $m = m' \circ i$.

$$\begin{array}{ccc} A & \xrightarrow{q'} & Q' \\ \downarrow q & \nearrow i & \downarrow m' \\ Q & \xrightarrow{m} & B \end{array}$$

As the factorization of f we take the one constructed in (2.6). Then q is a regular epi by construction, and we have just shown that m is monic. So it only remains to show that the factorization is unique. Suppose f also factors as $f = m' \circ q'$ where q' is a regular epi and m' is monic. Consider the following diagram, in which k_1, k_2 is the kernel pair of f , q is the coequalizer of k_1 and k_2 , and h_1, h_2 is the kernel pair of q' so that q' is the coequalizer

of h_1 and h_2 :

$$\begin{array}{ccccc}
 & & H & & \\
 & & \downarrow h_1 & & \downarrow h_2 \\
 K & \xrightarrow[k_2]{k_1} & A & \xrightarrow{q} & Q \\
 & & \downarrow q' & \nearrow i & \downarrow m \\
 & & Q' & \xrightarrow{m'} & B \\
 & & & \nwarrow j & \\
 & & & &
 \end{array}$$

Because $m' \circ q' \circ k_1 = m \circ q \circ k_1 = m \circ q \circ k_2 = m' \circ q' \circ k_2$ and m' is monic, $q' \circ k_1 = q' \circ k_2$. So there exists a unique $i : Q \rightarrow Q'$ such that $q' = i \circ q$. But then $m' \circ i \circ q = m' \circ q' = f = m \circ q$ and because q is epi, $m' \circ i = m$.

We prove that i is iso by constructing its inverse j . Because $m \circ q \circ h_1 = m' \circ q \circ h_1 = m' \circ q \circ h_2 = m \circ q \circ h_2$ and m is monic, $q \circ h_1 = q \circ h_2$. So there exists a unique $j : Q' \rightarrow Q$ such that $q = j \circ q'$. Now we have $i \circ j \circ q' = i \circ q = 1_{Q'} \circ q'$, from which we conclude that $i \circ j = 1_{Q'}$ because q' is epi. Similarly, $j \circ i \circ q = j \circ q' = 1_Q \circ q$, therefore $j \circ i = 1_Q$. \square

Corollary 2.5.5. *A map $f : A \rightarrow B$ that is both regular epi and a mono is an iso.*

Proof. Consider the following outer square, regarded as two different reg-epi/mono factorizations.

$$\begin{array}{ccc}
 A & \xrightarrow{1_A} & A \\
 f \downarrow & \nearrow d & \downarrow f \\
 B & \xrightarrow{1_B} & B
 \end{array}$$

A diagonal d is then an inverse of f . \square

A factorization $f = m \circ q$ as in Proposition 2.5.4 determines a subobject

$$\text{im}(f) = [m : Q \rightarrowtail B] \in \text{Sub}(B),$$

called the *image of f* . It is characterized as the least subobject of B through which f factors.

Proposition 2.5.6. *For a morphism $f : A \rightarrow B$ in a regular category \mathcal{C} , the image $\text{im}(f) \rightarrowtail B$ is the least subobject $U \rightarrowtail B$ of B through which f factors.*

Proof. Suppose f factors through $v : V \rightarrowtail B$ as

$$\begin{array}{ccccc}
 & & f & & \\
 & \curvearrowright & & \curvearrowright & \\
 A & \xrightarrow{g} & V & \xrightarrow{v} & B
 \end{array}$$

and consider the factorization of f , as in (2.6). Since $v \circ g \circ k_1 = f \circ k_1 = f \circ k_2 = v \circ g \circ k_2$ and v is mono, $g \circ k_1 = g \circ k_2$, therefore there exists a unique $\bar{g} : Q \rightarrow V$ such that $g = \bar{g} \circ q$. Now $v \circ \bar{g} \circ q = v \circ g = f = m \circ q$ and because q is epic, $v \circ \bar{g} = m$ as required. (The reader should draw the corresponding diagram.) \square

Definition 2.5.7. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is *regular* if it preserves finite limits and regular epis. It follows that F preserves image factorizations. The category of regular functors $\mathcal{C} \rightarrow \mathcal{D}$ and natural transformations is denoted by $\text{Reg}(\mathcal{C}, \mathcal{D})$.

Examples of regular categories

Let us consider some examples of regular categories.

1. The category **Set** is regular. It is complete and cocomplete, so it has in particular all finite limits and coequalizers. To show that the pullback of a regular epi is again a regular epi, note that in **Set** the epis are exactly the surjections, and a surjection is a quotient of its kernel pair, and thus a regular epi. It therefore suffices to show that the pullback of a surjection is a surjection, which is easy.
2. More generally, any presheaf category $\widehat{\mathcal{C}}$ is also regular, because it is complete and cocomplete, with (co)limits computed pointwise. Thus, again, every epi is regular, and epis are stable under pullbacks.
3. (“Fuzzy logic”) Let H be a complete Heyting algebra; thus H is a cartesian closed poset with all small joins $\bigvee_i p_i$. The category of H -presets has as objects all pairs $(X, e_X : X \rightarrow H)$ where X is a set and e_X is a function, called the *existence predicate* of X . For $x \in X$, $e_X(x)$ can be thought of as “the amount by which x exists”. A *morphism of presets* is a function $f : X \rightarrow Y$ satisfying, for all $x \in X$,

$$e_X(x) \leq e_Y(fx).$$

This is a regular category, with the following structure.

- the terminal object is $\top : 1 \rightarrow H$,
- the product of $e_A : A \rightarrow H$ and $e_B : B \rightarrow H$ is

$$e_A \wedge e_B : A \times B \rightarrow H,$$

where $(e_A \wedge e_B)(a, b) = e_A(a) \wedge e_B(b)$,

- the equalizer of two maps $f, g : A \rightarrow B$ is their equalizer as functions, $A' = \{a \mid f(a) = g(a)\} \hookrightarrow A$, with the restriction of $e_A : A \rightarrow H$ to $A' \subseteq A$.
- a map $f : A \rightarrow B$ is a regular epi if and only if it is a surjective function and for all $b \in B$:

$$e_B(b) = \bigvee_{f(a)=b} e_A(a)$$

Exercise 2.5.8. Verify that H -presets form a regular category, and compute the regular epi-mono factorization of a map.

The next example deserves to be a proposition.

Proposition 2.5.9. *The category $\mathbf{Mod}(\mathbb{A}, \mathbf{Set})$ of set-theoretic models of an algebraic theory \mathbb{A} is regular.*

Proof. We sketch a proof, for details see [?, Theorem 3.5.4]. Recall that the objects of $\mathbf{Mod}(\mathbb{A}) = \mathbf{Mod}(\mathbb{A}, \mathbf{Set})$ are \mathbb{A} -algebras, which are structures $A = (|A|, f_1, f_2, \dots)$ where $|A|$ is the underlying set and f_1, f_2, \dots are the basic operations on $|A|$. Every such \mathbb{A} -algebra is also required to satisfy the equational axioms of \mathbb{A} . A morphism $h : A \rightarrow B$ is a function $h : |A| \rightarrow |B|$ that preserves the basic operations.

The category $\mathbf{Mod}(\mathbb{A})$ of \mathbb{A} -algebras has small limits, which are created by the forgetful functor $U : \mathbf{Mod}(\mathbb{A}) \rightarrow \mathbf{Set}$. Thus the product of \mathbb{A} -algebras A and B has as its underlying set $|A \times B| = |A| \times |B|$, and the basic operations of $A \times B$ are computed separately on each factor, and similarly for products of arbitrary (small) families $\prod_i A_i$. An equalizer of morphisms $g, h : A \rightarrow B$ has as its underlying set the equalizer of $g, h : |A| \rightarrow |B|$, and the basic operations inherited from A .

To see that coequalizers of kernel pairs exist, consider a morphism $h : A \rightarrow B$. We can form the quotient \mathbb{A} -algebra Q whose underlying set is $|Q| = |A|/\sim$, where \sim is the relation defined by

$$x \sim y \iff hx = hy ,$$

which is just the kernel quotient of the underlying function h . A basic operation $f_Q : |Q|^k \rightarrow |Q|$ is induced by the basic operation $f_A : |A|^k \rightarrow |A|$ by

$$f_Q\langle [x_1], \dots, [x_k] \rangle = [f_A\langle x_1, \dots, x_k \rangle] .$$

It is easily verified that this is well-defined, that Q is an \mathbb{A} -algebra, and that the canonical quotient map $q : A \rightarrow Q$ is the coequalizer of the kernel pair of h .

Lastly regular epis in $\mathbf{Mod}(\mathbb{A})$ are stable because pullbacks and kernel pairs are computed as in \mathbf{Set} , and a morphism $h : A \rightarrow B$ is a regular epi in $\mathbf{Mod}(\mathbb{A})$ if, and only if, the underlying function $h : |A| \rightarrow |B|$ is a regular epi in \mathbf{Set} , which is therefore stable under pullback. \square

We now know that categories of groups, rings, modules, \mathcal{C}^∞ -rings and other algebraic categories are regular. The preceding proposition is useful also for showing that certain structures *cannot* be axiomatized by algebraic theories. The category of posets is an example of a category that is not regular; therefore the theory of partial orders cannot be axiomatized solely by equations.

Exercise 2.5.10. Show that \mathbf{Poset} is not regular. (Hint: find a regular epi that is not stable under pullback.) Conclude that there is no purely equational reformulation of the cartesian theory of posets.

Exercise* 2.5.11. Is \mathbf{Top} regular? Hint: is there is a topological quotient map $q : X \twoheadrightarrow X'$ and a space Y such that $q \times 1_Y : X \times Y \twoheadrightarrow X' \times Y$ is not a quotient map?

Remark 2.5.12 (Exactness). A regular category \mathcal{C} is said to be *exact* [?] if *every* equivalence relation (not just those arising as kernel pairs) has a quotient. It can be shown fairly easily that categories of algebras are not just regular but also exact: an equivalence relation in such a category is a congruence relation with respect to the algebraic operations, and its (underlying set) quotient is then necessarily also a homomorphism, and thus a coequalizer of algebras.

Exercise 2.5.13. Prove that the regular epis and monos in a regular category \mathcal{C} form the two classes $(\mathcal{L}, \mathcal{R})$, respectively, of an *orthogonal factorization system* in the following sense:

1. every arrow $f : A \rightarrow B$ factors as $f = r \circ l$ with $l \in \mathcal{L}$ and $r \in \mathcal{R}$,
2. \mathcal{L} is the class of all arrows left-orthogonal to all maps in \mathcal{R} , and \mathcal{R} is the class of all arrows right-orthogonal to all maps in \mathcal{L} , where $l : A \rightarrow B$ is said to be *left-orthogonal* to $r : X \rightarrow Y$, and r is said to be *right-orthogonal* to l , if for every commutative square as on the outside below,

$$\begin{array}{ccc} A & \xrightarrow{\quad} & X \\ l \downarrow & \nearrow d & \downarrow r \\ B & \xrightarrow{\quad} & Y, \end{array}$$

there is a unique diagonal arrow d as indicated making both triangles commute.

2.5.2 Images and existential quantifiers

Recall that the poset $\mathbf{Sub}(A)$ is equivalent to the preordered category $\mathbf{Mono}(A)$ of monos into A . If we compose an equivalence functor $\mathbf{Sub}(A) \rightarrow \mathbf{Mono}(A)$ with the inclusion $\mathbf{Mono}(A) \rightarrow \mathcal{C}/A$ we obtain a (full and faithful) inclusion functor

$$I : \mathbf{Sub}(A) \hookrightarrow \mathcal{C}/A. \quad (2.7)$$

In the other direction we have the “image functor” $\mathbf{im} : \mathcal{C}/A \rightarrow \mathbf{Sub}(A)$, which maps an object $f : B \rightarrow A$ in \mathcal{C}/A to the subobject $\mathbf{im}(f) \rightarrowtail A$.

Exercise 2.5.14. In order to show that \mathbf{im} is in fact a functor, prove that $f = g \circ h$ implies $\mathbf{im}(f) \leq \mathbf{im}(g)$.

Proposition 2.5.6 says that the image functor is left adjoint to the inclusion functor (2.7),

$$\mathbf{im} \dashv I.$$

Furthermore, images are stable in the sense that the following diagram commutes for all $f : A \rightarrow B$ (as does the corresponding one with the inclusion I in place of im).

$$\begin{array}{ccc}
C/A & \xleftarrow{f^*} & C/B \\
\text{im}_A \downarrow & & \downarrow \text{im}_B \\
\text{Sub}(A) & \xleftarrow{f^*} & \text{Sub}(B)
\end{array} \tag{2.8}$$

The functor f^* on the top is the “change of base” functor given by pullback of an arbitrary map, and the functor f^* on the bottom is the pullback functor acting on subjects. To see that (2.8) commutes, consider $g : C \rightarrow B$ and the following diagram:

$$\begin{array}{ccc}
f^*C & \xrightarrow{\quad} & C \\
\downarrow & & \downarrow \\
\textstyle\bigcap & \xrightarrow{\quad} & \text{im}(g) \\
\downarrow & & \downarrow \\
A & \xrightarrow{\quad f \quad} & B
\end{array}$$

On the right-hand side we have the factorization of g , which is then pulled back along f . Because monos and regular epis are both stable, this gives a factorization of the pullback f^*g , hence (by the uniqueness of factorizations, Proposition 2.5.4) the claimed equality

$$\text{im}(f^*g) = f^*(\text{im}(g)) \text{ .}$$

Proposition 2.5.15. *A regular category has existential quantifiers. The existential quantifier along $f : A \rightarrow B$,*

$$\exists_f : \text{Sub}(A) \longrightarrow \text{Sub}(B),$$

is given by

$$\exists_f[m : M \multimap A] = \text{im}(f \circ m) ,$$

as indicated below.

$$\begin{array}{ccc} M & \xrightarrow{\quad} & \text{im}(f \circ m) \\ m \downarrow & & \downarrow \\ A & \xrightarrow{\quad f \quad} & B \end{array}$$

Proof. Recall that composition

$$\Sigma_f : \mathcal{C}/A \longrightarrow \mathcal{C}/B$$

by a map $f : A \rightarrow B$ is left adjoint to pullback f^* along f . Thus we are defining $\exists_f = \text{im} \circ \Sigma_f \circ I$ as shown below.

$$\begin{array}{ccc} \text{Sub}(A) & \xrightarrow{\exists_f} & \text{Sub}(B) \\ I \downarrow & & \uparrow \text{im} \\ \mathcal{C}/A & \xrightarrow{\Sigma_f} & \mathcal{C}/B \end{array}$$

First we verify that $\exists_f \dashv f^*$ on subobjects. For $U \rightarrowtail A$ and $V \rightarrowtail B$:

$$\begin{array}{c} \frac{\exists_f U \leq V \quad \text{in } \text{Sub}(B)}{\text{im} \circ \Sigma_f \circ I(U) \leq V \quad \text{in } \text{Sub}(B)} \\ \frac{\Sigma_f \circ I(U) \leq I(V) \quad \text{in } \mathcal{C}/B}{I(U) \rightarrow f^* I(V) \quad \text{in } \mathcal{C}/A} \\ \frac{I(U) \rightarrow I(f^* V) \quad \text{in } \mathcal{C}/A}{U \leq f^* V \quad \text{in } \text{Sub}(A)} \end{array}$$

In the second step in the above derivation we used the adjunction between $\text{im} : \mathcal{C}/B \rightarrow \text{Sub}(B)$ and the inclusion $\text{Sub}(B) \rightarrow \mathcal{C}/B$.

The Beck-Chevalley condition follows from stability of image factorizations. Indeed, given a pullback

$$\begin{array}{ccc} D & \xrightarrow{h} & C \\ k \downarrow \lrcorner & & \downarrow g \\ A & \xrightarrow{f} & B \end{array}$$

and a subobject $U \rightarrowtail C$, (2.8) gives

$$\begin{aligned} f^*(\exists_g U) &= f^* \circ \text{im} \circ \Sigma_g \circ I(U) = \text{im} \circ f^* \circ \Sigma_g \circ I(U) = \text{im} \circ \Sigma_k \circ h^* \circ I(U) \\ &= \text{im} \circ \Sigma_k \circ I \circ h^*(U) = \exists_k(h^* U) \end{aligned}$$

$$\begin{array}{ccccc}
& f^*U & \longrightarrow & U & \\
& \downarrow \lrcorner & & \downarrow & \\
\exists_k f^*U & \xrightarrow{\quad} & D & \xrightarrow{h} & C & \xrightarrow{\quad} & \exists_g U \\
& \downarrow k & \downarrow \lrcorner & & \downarrow g & \\
& A & \xrightarrow{f} & B &
\end{array}$$

as required. \square

Summarizing the results of this section, we have the following.

Proposition 2.5.16. *In any regular category, for every map $f : A \rightarrow B$ we have the following situation, where f^* is pullback:*

$$\begin{array}{ccc}
\text{Sub}(A) & \xleftarrow{f^*} & \text{Sub}(B) \\
\uparrow \text{im} \downarrow I & \xrightarrow{\exists_f} & \uparrow \text{im} \downarrow I \\
C/A & \xleftarrow{f^*} & C/B \\
& \xrightarrow{\Sigma_f} &
\end{array}$$

with adjunctions

$$\exists_f \dashv f^*, \quad \text{im} \dashv I, \quad \Sigma_f \dashv f^*$$

and natural isos

$$f^* \circ \text{im} \cong \text{im} \circ f^*, \quad f^* \circ I \cong I \circ f^*.$$

Note, moreover, that

$$\exists_f \circ \text{im} \cong \text{im} \circ \Sigma_f$$

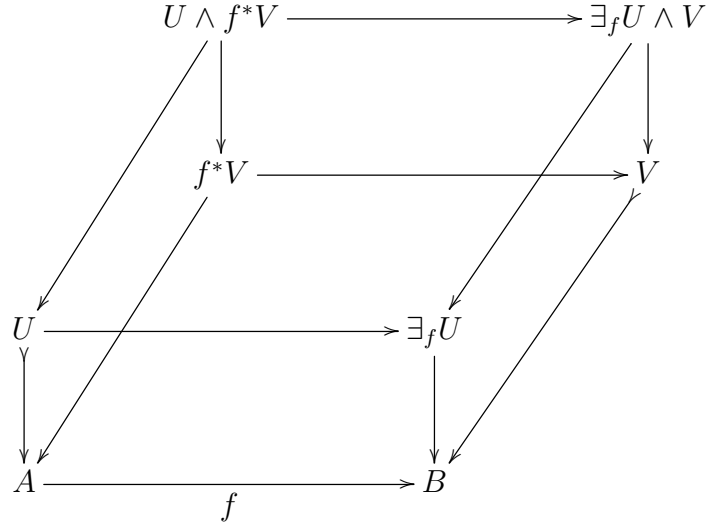
then follows.

Finally, we call attention to the following special fact.

Proposition 2.5.17 (Frobenius Reciprocity). *Given a map $f : A \rightarrow B$ and subobjects $U \leq A$ and $V \leq B$, the following equation holds in $\text{Sub}(B)$.*

$$\exists_f(U \wedge f^*V) = \exists_f U \wedge V$$

Exercise 2.5.18. Prove Frobenius reciprocity, using the following diagram.



2.5.3 Regular theories

A regular category has finite limits and image factorizations, therefore it allows us to interpret a type theory with the terminal type and binary products, and a logic with equality, conjunction, and existential quantifiers. This system is called *regular logic*.

Definition 2.5.19. A (many-sorted) *regular theory* \mathbb{T} is a (many-sorted) type theory together with a set of axioms expressed in the fragment of logic built from $=$, \top , \wedge , and \exists .

In more detail, a regular theory consists of the following data, extending the notion of cartesian theory from section 2.3.

- basic type symbols A_1, \dots, A_k ,
- basic function symbols f, \dots (with signature) $(A_1, \dots, A_m; B)$,
- basic relation symbols R, \dots (with signature) (A_1, \dots, A_n) .

We then define by induction the set of terms in context,

$$\Gamma \mid t : A,$$

as well as the formulas in context,

$$\Gamma \mid \varphi \text{ pred.}$$

Here is the first place where things differ from cartesian logic; we extend the formation rules for cartesian formulas (section 2.3) by the further clause:

6. Existential Quantifier:

$$\frac{\Gamma, x : A \mid \varphi \text{ pred}}{\Gamma \mid \exists x : A. \varphi \text{ pred}}$$

(We also add the evident additional clause for substitution of terms into existentially quantified formulas, namely $(\exists x : A. \varphi)[t/y] = \exists x : A. (\varphi[t/y])$.) This defines the notion of a *regular formula*, i.e. ones built from the atomic formulas $s = t$ and $R(t_1, \dots, t_n)$ using the logical operations \top , \wedge , and \exists .

A regular theory then includes, finally, a set of axioms of the form

$$\Gamma \mid \varphi \vdash \psi$$

where φ, ψ are *regular formulas*.

Example 2.5.20. A ring A (with unit 1) is called *von Neumann regular* if for every element a there is at least one element x for which $a = a \cdot x \cdot a$. Such an x may be thought of as a “weak inverse” of a . The theory of *von Neumann regular rings* is thus an extension of the usual theory of rings with unit by adding the single axiom

$$a : A \mid \top \vdash \exists x : A. a = a \cdot x \cdot a$$

The *rules of inference* of regular logic are those of cartesian logic (section 2.3), with an additional rule for the existential quantifier:

8. Existential Quantifier:

$$\frac{y : B, x : A \mid \varphi \vdash \vartheta}{y : B \mid \exists x : A. \varphi \vdash \vartheta}$$

Note that the lower judgement is well-formed only if $x : A$ does not occur freely in ϑ .

We also add a rule corresponding to Frobenius reciprocity, Proposition 2.5.17, in the form

9. Frobenius:

$$x : A \mid (\exists y : B. \varphi) \wedge \psi \vdash \exists y : B. (\varphi \wedge \psi)$$

provided the variable $y : B$ does not occur freely in ψ .

Note that the converse of Frobenius is easily derivable, so we have the interderivability of $(\exists y : B. \varphi) \wedge \psi$ and $\exists y : B. (\varphi \wedge \psi)$ when $y : B$ is not free in ψ . The Frobenius rule will be derivable in the extended system of Heyting logic (see Proposition ??), and could be made derivable in a suitably formulated system of regular logic using multi-sequents $\Gamma \mid \varphi_1, \dots, \varphi_n \vdash \psi$.

Semantics of regular theories

Turning to semantics, an *interpretation* of a regular theory \mathbb{T} in a regular category \mathcal{C} extends the notion for cartesian logic (section 2.3), and is given by the following data:

1. Each basic sort A is interpreted as an object $\llbracket A \rrbracket$.
2. Each basic constant f with signature $(A_1, \dots, A_n; B)$ is interpreted as a morphism $\llbracket f \rrbracket : \llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket \rightarrow \llbracket B \rrbracket$.

3. Each basic relation symbol R with signature (A_1, \dots, A_n) is interpreted as a subobject $\llbracket R \rrbracket \in \mathbf{Sub}(\llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket)$.

This is the same as for cartesian logic, as is the extension of the interpretation to all terms,

$$\llbracket \Gamma \mid t : A \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \llbracket A \rrbracket$$

For the formulas, we extended the interpretation to cartesian formulas as before (section ??),

$$\llbracket \Gamma \mid \varphi \rrbracket \rightharpoonup \llbracket \Gamma \rrbracket.$$

Finally, existential formulas $\exists x : A. \varphi$ are interpreted by the existential quantifiers in the regular category,

$$\llbracket \Gamma \mid \exists x : A. \varphi \rrbracket = \exists_A \llbracket \Gamma, x : A \mid \varphi \rrbracket,$$

where

$$\exists_A = \exists_\pi : \mathbf{Sub}(\llbracket \Gamma \rrbracket \times \llbracket A \rrbracket) \rightarrow \mathbf{Sub}(\llbracket \Gamma \rrbracket)$$

is the existential quantifier along the projection $\pi : \llbracket \Gamma \rrbracket \times \llbracket A \rrbracket \rightarrow \llbracket \Gamma \rrbracket$.

The following is immediate from these definitions, and the considerations in section 2.4.

Proposition 2.5.21. *The rules of regular logic are sound with respect to the interpretation in regular categories.*

Exercise 2.5.22. Prove this.

If all the axioms of \mathbb{T} are valid in a given interpretation, then we again say that the interpretation is a *model* of the theory \mathbb{T} . Morphisms of models are just morphisms of the underlying cartesian structures. Thus for any regular theory \mathbb{T} and regular category \mathcal{C} , there is a *category of models*,

$$\mathbf{Mod}(\mathbb{T}, \mathcal{C}).$$

Moreover, this semantic category is functorial in \mathcal{C} with respect to regular functors $\mathcal{C} \rightarrow \mathcal{D}$, which, recall, preserve finite limits and regular epis. Indeed, if $F : \mathcal{C} \rightarrow \mathcal{D}$ is regular then given a model M in \mathcal{C} with underlying cartesian structure $\llbracket A \rrbracket_M, \llbracket f \rrbracket_M, \llbracket R \rrbracket_M$, etc., we can determine an interpretation FM in \mathcal{D} by setting:

$$\llbracket A \rrbracket_{FM} = F(\llbracket A \rrbracket_M), \llbracket f \rrbracket_{FM} = F(\llbracket f \rrbracket_M), \llbracket R \rrbracket_{FM} = F(\llbracket R \rrbracket_M)$$

etc., and these will have the correct types (up to isomorphism). To show that FM is a \mathbb{T} -model, if M is one and F is regular, consider an axiom of \mathbb{T} of the form $\Gamma \mid \varphi \vdash \psi$. Satisfaction by M means that $\llbracket \Gamma \mid \varphi \rrbracket_M \leq \llbracket \Gamma \mid \psi \rrbracket_M$ in $\mathbf{Sub}(\llbracket \Gamma \rrbracket_M)$, which in turn means that there is a (necessarily unique) factorization,

$$\begin{array}{ccc} \llbracket \Gamma \mid \varphi \rrbracket_M & \xrightarrow{\quad \quad \quad} & \llbracket \Gamma \mid \psi \rrbracket_M \\ & \searrow \quad \quad \swarrow & \\ & \llbracket \Gamma \rrbracket_M & \end{array},$$

Applying the cartesian functor F will result in an inclusion of subobjects $F\llbracket\Gamma \mid \varphi\rrbracket_M \leq F\llbracket\Gamma \mid \psi\rrbracket_M$ in $\mathbf{Sub}(F\llbracket\Gamma\rrbracket_M) = \mathbf{Sub}(\llbracket\Gamma\rrbracket_{FM})$. Thus it clearly suffices to show that for any regular formula φ ,

$$F\llbracket\Gamma \mid \varphi\rrbracket_M = \llbracket\Gamma \mid \varphi\rrbracket_{FM}.$$

This is an easy induction on φ , using the regularity of F .

Proposition 2.5.23. *Given a regular functor $F : \mathcal{C} \rightarrow \mathcal{D}$, taking images determines a functor*

$$F_* : \mathbf{Mod}(\mathbb{T}, \mathcal{C}) \longrightarrow \mathbf{Mod}(\mathbb{T}, \mathcal{D}).$$

Proof. It only remains show the effect of F_* on morphisms of models. But these are just homomorphisms of the underlying cartesian structure, so they are clearly preserved by the cartesian functor F . \square

An associated result, which we will need, is the following.

Proposition 2.5.24. *Given regular categories \mathcal{C} and \mathcal{D} and a model M in \mathcal{C} , evaluation at M determines a functor*

$$\mathrm{eval}_M : \mathbf{Reg}(\mathcal{C}, \mathcal{D}) \longrightarrow \mathbf{Mod}(\mathbb{T}, \mathcal{D}),$$

which is natural in \mathcal{D} .

The proof is straightforward and can be left as an exercise. The naturality means that for any a regular functor $G : \mathcal{D} \rightarrow \mathcal{D}'$, the following commutes (up to natural isomorphism, as usual):

$$\begin{array}{ccc} \mathbf{Reg}(\mathcal{C}, \mathcal{D}) & \xrightarrow{\mathrm{eval}_M} & \mathbf{Mod}(\mathbb{T}, \mathcal{D}) \\ \mathbf{Reg}(\mathcal{C}, G) \downarrow & & \downarrow G_* \\ \mathbf{Reg}(\mathcal{C}, \mathcal{D}') & \xrightarrow{\mathrm{eval}_M} & \mathbf{Mod}(\mathbb{T}, \mathcal{D}') \end{array}$$

Exercise 2.5.25. Prove this.

Exercise 2.5.26. Show that for any small category \mathbb{C} and regular theory \mathbb{T} , there is an equivalence between models in the functor category and functors into the category of models,

$$\mathbf{Mod}(\mathbb{T}, \mathbf{Set}^{\mathbb{C}}) \simeq \mathbf{Mod}(\mathbb{T})^{\mathbb{C}}.$$

Hint: this is just as for the algebraic and cartesian cases.

2.5.4 Classifying category of a regular theory

We will next show that the framework of *functorial semantics* applies to regular logic and regular categories: there is a *classifying category* $\mathcal{C}_{\mathbb{T}}$ for \mathbb{T} -models, for which there is an equivalence, natural in \mathcal{C} ,

$$\mathbf{Mod}(\mathbb{T}, \mathcal{C}) \simeq \mathbf{Reg}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) ,$$

where $\mathbf{Reg}(-, -)$ is the category of regular functors and natural transformations.

We first sketch the construction of the classifying category $\mathcal{C}_{\mathbb{T}}$ of an arbitrary regular theory \mathbb{T} (a more detailed account can be found in [?, ?]). An object of $\mathcal{C}_{\mathbb{T}}$ is represented by a formula in context,

$$[\Gamma \mid \varphi],$$

where $\Gamma \mid \varphi$ **pred.** Two such objects $[\Gamma \mid \varphi]$ and $[\Gamma \mid \psi]$ are equal if \mathbb{T} proves both

$$\Gamma \mid \varphi \vdash \psi , \quad \Gamma \mid \psi \vdash \varphi .$$

Objects which differ only in the names of free variables are also considered equal. A morphism

$$[x : A \mid \varphi] \xrightarrow{\rho} [y : B \mid \psi]$$

is represented by a formula $x : A, y : B \mid \rho$ such that \mathbb{T} proves that ρ is a *functional relation* from φ to ψ :

$$\begin{array}{ll} x : A \mid \varphi \vdash \exists y : B . \rho & \text{(total)} \\ x : A, y : B, z : B \mid \rho \wedge \rho[z/y] \vdash y = z & \text{(single-valued)} \\ x : A, y : B \mid \rho \vdash \varphi \wedge \psi & \text{(well-typed)} \end{array}$$

Two functional relations ρ and σ represent the same morphism if \mathbb{T} proves both

$$x : A, y : B \mid \rho \vdash \sigma , \quad x : A, y : B \mid \sigma \vdash \rho .$$

Relations which only differ in the names of free variables are considered equal.

The identity morphism on $[x : A \mid \varphi]$ is represented by the relation

$$x : A, y : A \mid (x = y) \wedge \varphi .$$

Composition of morphisms

$$[x : A \mid \varphi] \xrightarrow{\rho} [y : B \mid \psi] \xrightarrow{\tau} [z : C \mid \theta]$$

is given by the relational product,

$$\tau \circ \rho = (\exists y : B . (\rho \wedge \tau)) .$$

Of course, one needs to check that this *is* a morphism from φ to ϑ , i.e. that it is total, single-valued, and well-typed. We leave the detailed proof that $\mathcal{C}_{\mathbb{T}}$ is a category as an exercise; let us just show that composition of morphisms is associative. Given morphisms

$$[x : A \mid \varphi] \xrightarrow{\rho} [y : B \mid \psi] \xrightarrow{\tau} [z : C \mid \theta] \xrightarrow{\sigma} [u : D \mid \zeta]$$

we need to derive in context $x : A, u : D$

$$\exists z : C . ((\exists y : B . (\rho \wedge \tau)) \wedge \sigma) \vdash \exists y : B . (\rho \wedge (\exists z : C . (\tau \wedge \sigma)))$$

But this follows easily by repeated application of the Frobenius rule (Section 2.5.3).

Exercise 2.5.27. Extend the definition of $\mathcal{C}_{\mathbb{T}}$ to morphisms between objects with arbitrary contexts,

$$[\Gamma \mid \varphi] \xrightarrow{\rho} [\Delta \mid \psi]$$

(use relations $\Gamma, \Delta \mid \rho$), and provide a proof that $\mathcal{C}_{\mathbb{T}}$ is a category.

Proposition 2.5.28. *The category $\mathcal{C}_{\mathbb{T}}$ is regular.*

Proof. We sketch the constructions required for regularity.

- The terminal object is $[\cdot \mid \top]$.
- The product of $[x : A \mid \varphi]$ and $[y : B \mid \psi]$, where x and y are distinct variables, is the object

$$[x : A, y : B \mid \varphi \wedge \psi] .$$

The first projection from the product is

$$x : A, y : B, x' : A \mid x = x' \wedge \varphi \wedge \psi ,$$

and the second projection is

$$x : A, y : B, y' : B \mid y = y' \wedge \varphi \wedge \psi ,$$

where we rename the codomains of the projections $[x : A \mid \varphi] = [x' : A \mid \varphi[x'/x]]$, etc., to make the context variables distinct.

- An equalizer of morphisms

$$[x : A \mid \varphi] \begin{smallmatrix} \xrightarrow{\rho} \\ \xrightarrow{\tau} \end{smallmatrix} [y : B \mid \psi]$$

is

$$[x : A \mid \varphi \wedge \exists y : B . (\rho \wedge \tau)] \xrightarrow{\varepsilon} [x' : A \mid \varphi[x'/x]]$$

where ε is the morphism

$$x : A, x' : A \mid (x = x') \wedge \varphi \wedge \exists y : B . (\rho \wedge \tau) .$$

- Finally, let us consider coequalizers of kernel pairs. The kernel pair of a map

$$\rho : [x : A \mid \varphi] \longrightarrow [y : B \mid \psi]$$

is

$$K \begin{array}{c} \xrightarrow{\kappa_1} \\ \xrightarrow{\kappa_2} \end{array} [x : A \mid \varphi]$$

where K is the object

$$[u : A, v : A \mid \varphi[u/x] \wedge \varphi[v/x] \wedge \exists y : B . (\rho[u/x] \wedge \rho[v/x])] ,$$

the morphism κ_1 is

$$u : A, v : A, x : A \mid (u = x) \wedge \varphi$$

and κ_2 is

$$u : A, v : A, x : A \mid (v = x) \wedge \varphi .$$

Now the coequalizer of κ_1 and κ_2 can be shown to be the morphism

$$[x : A \mid \varphi] \xrightarrow{\rho} [y : B \mid \exists x : A . \rho] ,$$

where $[y : B \mid \exists x : A . \rho]$ is the image of ρ , as a subobject of $[y : B \mid \psi]$.

The following lemma shows that regular epis are stable under pullback. □

Lemma 2.5.29. 1. A map $\rho : [x : A \mid \varphi] \longrightarrow [y : B \mid \psi]$ is a regular epi if and only if

$$y : B \mid \psi \vdash \exists x : A . (\rho)$$

2. Regular epis are stable under pullback in $\mathcal{C}_{\mathbb{T}}$.

Proof. ... □

Exercise 2.5.30. Show that in $\mathcal{C}_{\mathbb{T}}$ the regular-epi mono factorization of a morphism $\rho : [x : A \mid \varphi] \rightarrow [y : B \mid \psi]$ is given by

$$[x : A \mid \varphi] \xrightarrow{\rho} [y : B \mid \exists x : A . \rho] \xrightarrow{\iota} [z : B \mid \psi[z/y]]$$

where ι is the morphism

$$y : B, z : B \mid (y = z) \wedge (\exists x : A . \rho) \wedge \psi[z/y] .$$

Theorem 2.5.31 (Functorial semantics for regular logic). *For any regular theory \mathbb{T} , the syntactic category $\mathcal{C}_{\mathbb{T}}$ classifies \mathbb{T} -models in regular categories. Thus for any regular category \mathcal{C} , there is an equivalence of categories*

$$\text{Reg}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) \simeq \text{Mod}(\mathbb{T}, \mathcal{C}). \quad (2.9)$$

which is natural in \mathcal{C} . In particular, there is a universal model U in $\mathcal{C}_{\mathbb{T}}$.

Proof. We have just constructed $\mathcal{C}_{\mathbb{T}}$ and shown that it is regular.

The universal model U , corresponding to the identity functor $\mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}_{\mathbb{T}}$ under (2.9), is determined as follows:

- Each sort A is interpreted by the object $[x : A \mid \top]$
- A basic constant f with signature $(A_1, \dots, A_n; B)$ is interpreted by the formula

$$x_1 : A_1, \dots, x_n : A_n, y : B \mid f(x_1, \dots, x_n) = y .$$

which is plainly a functional relation and thus a morphism $\llbracket A_1 \rrbracket \times \dots \times \llbracket A_n \rrbracket \longrightarrow \llbracket B \rrbracket$.

- A relation symbol R with signature (A_1, \dots, A_n) is interpreted by the subobject represented by the morphism

$$\rho : [x_1 : A_1, \dots, x_n : A_n \mid R(x_1, \dots, x_n)] \longrightarrow [y_1 : A_1, \dots, y_n : A_n \mid \top]$$

where ρ is the formula

$$R(x_1, \dots, x_n) \wedge x_1 = y_1 \wedge \dots \wedge x_n = y_n .$$

It is straightforward to show that this structure has the property

$$U \models \Gamma \mid \varphi \vdash \psi \quad \Longleftrightarrow \quad \mathbb{T} \text{ proves } \Gamma \mid \varphi \vdash \psi .$$

Thus in particular U is indeed a \mathbb{T} -model.

If \mathcal{C} is regular and $F : \mathcal{C}_{\mathbb{T}} \longrightarrow \mathcal{C}$ a regular functor, then applying F to U determines a model FU in \mathcal{C} with

$$\llbracket A \rrbracket_{FU} = F(\llbracket A \rrbracket_U) ,$$

and similarly for the other parts of the structure f , R , etc. Satisfaction of an entailment $\Gamma \mid \varphi \vdash \psi$ is preserved, because the interpretation of the logical operations is determined by the regular structure: pullbacks, images, etc., so that $\llbracket \varphi \rrbracket_U \leq \llbracket \psi \rrbracket_U$ in $\mathbf{Sub}(\llbracket \Gamma \rrbracket)$ implies

$$\llbracket \varphi \rrbracket_{FU} = F(\llbracket \varphi \rrbracket_U) \leq F(\llbracket \psi \rrbracket_U) = \llbracket \psi \rrbracket_{FU}$$

in $\mathbf{Sub}(\llbracket \Gamma \rrbracket_{FU})$.

Moreover, just as for algebraic structures, every natural transformation between regular functors $\vartheta : F \Rightarrow G$ determines a homomorphism of the evaluated models by taking components $\vartheta_U : FU \rightarrow GU$. In this way, as in Proposition 2.5.24, evaluation at U is a functor

$$\text{eval}_U : \mathbf{Reg}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) \longrightarrow \mathbf{Mod}(\mathbb{T}, \mathcal{C}) .$$

We claim that this functor, which is the one mentioned in (2.9), is full and faithful and essentially surjective. The naturality in \mathcal{C} of the equivalence then follows directly from its determination by evaluation at U and Proposition 2.5.24.

To see that eval_U is essentially surjective, let M be a model in \mathcal{C} . There are then objects $\llbracket A \rrbracket_M$ for each type A , as well as interpretations

$$\llbracket \Gamma \mid \varphi \rrbracket \mapsto \llbracket \Gamma \rrbracket$$

for all formulas and

$$\llbracket \Gamma \mid t : B \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \llbracket B \rrbracket$$

for all terms. Using these, we determine a functor

$$M^\sharp : \mathcal{C}_\mathbb{T} \longrightarrow \mathcal{C},$$

taking an object $\llbracket \Gamma \mid \varphi \rrbracket$ to $\llbracket \Gamma \mid \varphi \rrbracket_M$, i.e. the domain of a mono representing the subobject $\llbracket \Gamma \mid \varphi \rrbracket_M \mapsto \llbracket \Gamma \rrbracket_M$. Thus, for the record,

$$M^\sharp[\Gamma \mid \varphi] = \llbracket \Gamma \mid \varphi \rrbracket_M.$$

In the verification that those formulas in context $\llbracket \Gamma \mid \varphi \rrbracket$ that are identified in $\mathcal{C}_\mathbb{T}$ are also identified in \mathcal{C} , we make use of the fact that the rules of inference for regular logic are sound in the regular category \mathcal{C} . Note in particular that for each basic type A , we then have

$$M^\sharp(\llbracket A \rrbracket_U) = M^\sharp([x : A \mid \top]) \cong [x : A \mid \top]_M \cong \llbracket A \rrbracket_M,$$

so that $M^\sharp(U) \cong M$ as required.

Functional relations in $\mathcal{C}_\mathbb{T}$ determine functional relations in \mathcal{C} , also by soundness, which determines the action of M^\sharp on arrows, as well as the functoriality of these assignments.

Finally, to show that eval_U is full and faithful, let $F, G : \mathcal{C}_\mathbb{T} \longrightarrow \mathcal{C}$ be regular functors classifying models FU and GU , and let $h : FU \rightarrow GU$ be a model homomorphism. We then have maps

$$h_{[x:A|\top]} : F([x : A \mid \top]) \longrightarrow G([x : A \mid \top])$$

for all basic types A , and these commute with the interpretations of the function symbols f , and preserve the basic relations R , in the obvious sense, because h is a homomorphism. It only remains to determine the components

$$h_{[\Gamma|\varphi]} : F([\Gamma \mid \varphi]) \rightarrow G([\Gamma \mid \varphi]), \tag{2.10}$$

and to show that they commute with all maps $\rho : [\Gamma \mid \varphi] \rightarrow [\Delta \mid \psi]$. Define

$$h_{[\Gamma|\varphi]} : F[\Gamma \mid \varphi] = \llbracket \Gamma \mid \varphi \rrbracket_{FU} \longrightarrow \llbracket \Gamma \mid \varphi \rrbracket_{GU} = G[\Gamma \mid \varphi]$$

by induction on the structure of φ . The base cases involving the primitive relations R, \dots and equality of terms are given by the assumption that $h : FU \rightarrow GU$ is a model homomorphism, so we just need to check that for every definable subobject

$$\llbracket \Gamma \mid \varphi \rrbracket_{FU} \mapsto \llbracket \Gamma \mid \top \rrbracket_{FU}$$

the following diagram can be filled in as indicated below.

$$\begin{array}{ccc}
 \llbracket \Gamma \mid \varphi \rrbracket_{FU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \top \rrbracket_{FU} \\
 \downarrow h_{[\Gamma \mid \varphi]} & & \downarrow h_{[\Gamma \mid \top]} \\
 \llbracket \Gamma \mid \varphi \rrbracket_{GU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \top \rrbracket_{GU}
 \end{array}$$

Suppose we have e.g. $\varphi = \exists x : A. \psi$, and we have already determined

$$h_{[\Gamma, x:A \mid \psi]} : \llbracket \Gamma, x : A \mid \psi \rrbracket_{FU} \longrightarrow \llbracket \Gamma, x : A \mid \psi \rrbracket_{GU}.$$

An easy diagram chase shows that there is a unique $h_{[\Gamma \mid \exists x:A. \psi]}$ determined by the image factorizations indicated below.

$$\begin{array}{ccccc}
 \llbracket \Gamma, x : A \mid \psi \rrbracket_{FU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \varphi \rrbracket_{FU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \top \rrbracket_{FU} \\
 \downarrow h_{[\Gamma, x:A \mid \psi]} & & \downarrow h_{[\Gamma \mid \exists x:A. \psi]} & & \downarrow h_{[\Gamma \mid \top]} \\
 \llbracket \Gamma, x : A \mid \psi \rrbracket_{GU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \varphi \rrbracket_{GU} & \xrightarrow{\quad} & \llbracket \Gamma \mid \top \rrbracket_{GU}
 \end{array}$$

The other cases are even more direct. Thus we have defined the components (2.10); we leave the required naturality with respect to all maps $\rho : [\Gamma \mid \varphi] \rightarrow [\Delta \mid \psi]$ as an exercise. \square

Exercise 2.5.32. Prove the naturality of the maps (2.10), using the following trick. In any category with finite products, suppose we have objects and arrows

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \alpha \downarrow & & \downarrow \beta \\
 C & \xrightarrow{g} & D
 \end{array} \tag{2.11}$$

Let $\hat{f} = \langle 1_A, f \rangle : A \rightarrowtail A \times B$ be the graph of f , and similarly for $\hat{g} : C \rightarrowtail C \times D$. Then the diagram (2.11) commutes iff the following one does.

$$\begin{array}{ccc}
 A & \xrightarrow{\hat{f}} & A \times B \\
 \alpha \downarrow & & \downarrow \alpha \times \beta \\
 C & \xrightarrow{\hat{g}} & C \times D
 \end{array}$$

Corollary 2.5.33. *The rules of regular logic are sound and complete with respect to semantics in regular categories: a regular theory \mathbb{T} proves an entailment*

$$\Gamma \mid \varphi \vdash \psi \quad (2.12)$$

if, and only if, every model of \mathbb{T} satisfies it.

Proof. As for algebraic logic, soundness follows from classification (although we have of course already proved it separately in Proposition 2.5.21, and made use of it in the proof of the theorem!): if (2.12) is provable from \mathbb{T} , then it holds in the universal model U in $\mathcal{C}_{\mathbb{T}}$ by the construction of U ,

$$U \models \Gamma \mid \varphi \vdash \psi.$$

But since regular functors preserve the interpretations of regular formulas $\llbracket \varphi \rrbracket, \llbracket \psi \rrbracket$ (as well as entailments between them), the entailment (2.12) then holds also in any model M in any regular \mathcal{C} , since there is a classifying functor $M^\sharp : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$ taking U to M , for which

$$M^\sharp(\llbracket \varphi \rrbracket_U) \cong \llbracket \varphi \rrbracket_M.$$

Completeness follows from the syntactic construction of the universal model U in $\mathcal{C}_{\mathbb{T}}$. The model U is logically generic, in the sense that

$$U \models \Gamma \mid \varphi \vdash \psi \quad \Longleftrightarrow \quad \mathbb{T} \text{ proves } \Gamma \mid \varphi \vdash \psi.$$

Thus if $\Gamma \mid \varphi \vdash \psi$ holds in all models, then it holds in particular in U , and is therefore provable. \square

2.5.5 Coherent logic

A regular category is coherent if all the subobject posets are distributive lattices, and that structure is stable under pullback. We add rules to regular logic to describe this further structure, show that the rules are sound in coherent categories, and extend the results on functorial semantics of the previous section to the coherent case, including the completeness theorem.

Definition 2.5.34. A cartesian category \mathcal{C} is *coherent* if:

1. \mathcal{C} is regular, i.e. it has coequalizers of kernel pairs, and regular epimorphisms are stable under pullback,
2. each subobject poset $\mathbf{Sub}(A)$ has all finite joins, in particular 0 and $U \vee V$,
3. for each map $f : A \rightarrow B$, the pullback functor $f^* : \mathbf{Sub}(B) \rightarrow \mathbf{Sub}(A)$ preserves the joins:

$$f^*0_B = 0_A, \quad f^*(U \vee V) = f^*U \vee f^*V.$$

Note that since joins are stable under pullback in a coherent category, the meets distribute over the joins,

$$U \wedge (V \vee W) = (U \wedge V) \vee (U \wedge W), \quad (2.13)$$

so that the posets $\mathbf{Sub}(A)$ are distributive lattices. Indeed, this follows from the fact that $U \wedge V$ may be written as

$$U \wedge V = \Sigma_U \circ U^*(V)$$

where $\Sigma_U : \mathbf{Sub}(U) \rightarrow \mathbf{Sub}(A)$ is the left adjoint (composition) of the pullback functor $U^* : \mathbf{Sub}(A) \rightarrow \mathbf{Sub}(U)$ along the inclusion $U \hookrightarrow A$. Since left adjoints preserve colimits, and thus joins, we therefore have

$$U \wedge (V \vee W) = \Sigma_U \circ U^*(V \vee W) = \Sigma_U \circ U^*(V) \vee \Sigma_U \circ U^*(W) = (U \wedge V) \vee (U \wedge W).$$

A category is said to have *stable sums* if it has all finite coproducts, in particular an initial object 0 and binary coproducts $A+B$, and these are stable under pullback, in the expected sense. The following simple observation provides plenty of examples of coherent categories.

Proposition 2.5.35. *Regular categories with stable sums are coherent.*

Proof. Given subobjects $U, V \hookrightarrow A$, let $U \vee V$ be the image of the canonical map $U+V \rightarrow A$ as indicated below.

$$\begin{array}{ccccc} & & U+V & & \\ & \nearrow & \downarrow & \nwarrow & \\ U & & U \vee V & & V \\ & \searrow & \downarrow & \swarrow & \\ & & A & & \end{array}$$

This easily seems to be the supremum of U and V in $\mathbf{Sub}(A)$. Since the unique map $0 \rightarrow A$ is always monic, it determines the subobject $0 \hookrightarrow A$. Thus $\mathbf{Sub}(A)$ has all finite joins, and they are stable by stability of the coproducts and image factorizations. \square

As examples of coherent categories we thus have \mathbf{Set} and $\mathbf{Set}_{\text{fin}}$, as well as all functor categories $\mathbf{Set}^{\mathbf{C}}$ since limits and colimits (and thus image factorizations) there are computed pointwise.

Coherent logic is the extension of regular logic by adding rules corresponding to joins.

Definition 2.5.36. A *coherent theory* \mathbb{T} is (a type theory together with) a set of axioms expressed in the fragment of logic built from $=$, \top , \perp , \wedge , \vee , and \exists .

We thus extend the formation rules for formulas in context by two additional clauses:

7. The 0-ary relation symbol \perp (pronounced “false”) is a formula :

$$\frac{\cdot}{\Gamma \mid \perp \text{ pred}}$$

6. Disjunction:

$$\frac{\Gamma \mid \varphi \text{ pred} \quad \Gamma \mid \psi \text{ pred}}{\Gamma \mid \varphi \vee \psi \text{ pred}}$$

(We also again add the evident additional clauses for substitution of terms into formulas.)
A *coherent theory* then consists of axioms of the form

$$\Gamma \mid \varphi \vdash \psi$$

where φ, ψ are *coherent formulas*. Coherent logic not only allows for disjunctions $\varphi \vee \psi$ on both side of the \vdash , but the presence of the symbol \perp allows for a certain amount of negation, in the form $\varphi \vdash \perp$, as the following classical example illustrates.

Example 2.5.37. A ring A (with unit 1) is called *local* if it has a unique maximal ideal. This can be captured with two coherent axioms of the form $0 = 1 \vdash \perp$ (to ensure that $0 \neq 1$), and

$$x : A, y : A \mid \exists z : A. z(x + y) = 1 \vdash (\exists z : A. zx = 1) \vee (\exists z : A. zy = 1)$$

Another example is the theory of *fields*, which can be axiomatized by again adding to the theory of rings the law $0 = 1 \vdash \perp$, together with the following:

$$x : A \mid \top \vdash x = 0 \vee (\exists y : A. xy = 1)$$

which is a clever way of saying that every non-zero element has a multiplicative inverse.

The *rules of inference of coherent logic* are those of regular logic (Section 2.5.3), with additional rules for falsehood the disjunctions:

10. Falsehood:

$$\overline{\perp \vdash \psi}$$

11. Disjunction:

$$\frac{\varphi \vdash \vartheta \quad \psi \vdash \vartheta}{\varphi \vee \psi \vdash \vartheta} \quad \frac{\varphi \vee \psi \vdash \vartheta}{\varphi \vdash \vartheta} \quad \frac{\varphi \vee \psi \vdash \vartheta}{\psi \vdash \vartheta}$$

12. Distributivity:

$$\varphi \wedge (\psi \vee \vartheta) \vdash (\varphi \wedge \psi) \vee (\varphi \wedge \vartheta)$$

The latter of course corresponds to the distributive law (2.13); note that the converse can be derived. Like the Frobenius rule, this will be derivable in the extended system of Heyting logic (see Proposition ??), and could also be made derivable in a suitably formulated system of coherent logic using multi-sequents $\Gamma \mid \varphi_1, \dots, \varphi_n \vdash \psi$.

The *semantics for coherent logic* extends that for regular logic in the expected way: the disjunctive formulas are interpreted as the corresponding joins in the subobject lattices,

$$\llbracket \Gamma \mid \perp \rrbracket = 0, \quad \llbracket \Gamma \mid \varphi \vee \psi \rrbracket = \llbracket \Gamma \mid \varphi \rrbracket \vee \llbracket \Gamma \mid \psi \rrbracket.$$

The additional clauses in the proof of soundness are routine, as is the extension of the functorial semantics, and the resulting completeness proof. We leave these routine details to the reader.

Exercise 2.5.38. Extend the functorial semantics theorem 2.5.31 from regular to coherent logic. Specifically, one must determine the components (2.10) of a natural transformation for the extended language of coherent logic.