

Find All Duplicates in an Array [LeetCode](#)

Given an integer array `nums` of length `n` where all the integers of `nums` are in the range `[1, n]` and each integer appears once or twice, return an array of all the integers that appears twice.

You must write an algorithm that runs in $O(n)$ time and uses only constant extra space.

Input: `nums = [4,3,2,7,8,2,3,1]`

Output: `[2,3]`

Input: `nums = [1,1,2]`

Output: `[1]`

Input: `nums = [1]`

Output: `[]`

Note:

You should not modify the original array.

The order of elements in the output vector does not matter.

Your solution should have a time complexity of $O(n)$ and a space complexity of $O(1)$ (excluding the output vector).

Approach 1: Using Nested Loop to compare array elements with each other.

iterates through the input array `nums` using two nested loops. It compares each element with the remaining elements to count the occurrences of each element. If an element occurs exactly twice, it is considered a duplicate and added to the `ans` vector.

Time Complexity: The outer loop runs `n` times, where `n` is the size of the input array. The inner loop runs `n - i - 1` times for each outer loop iteration. Overall, this approach has a **time complexity of $O(n^2)$, where `n` is the size of the input array.**

Space Complexity: The `ans` vector stores the duplicate elements, **which can have a maximum size of `n/2` (assuming every element appears twice).** Therefore, the space complexity is $O(n)$.

Approach 2: Using Hash map to keep track of the count of the array elements.

uses an unordered map, `count`, to store the frequency of each element in the input array `nums`. It then iterates through the map and adds the elements with a frequency of 2 to the `ans` vector.

Time Complexity: The first loop that counts the frequency of each element in the input array takes $O(n)$ time, where n is the size of the input array. The second loop iterates through the map, which contains at most $n/2$ elements (assuming every element appears twice), **resulting in a time complexity of $O(n)$.**

Space Complexity: **The space complexity is $O(n)$ because the unordered map, count,** can store at most $n/2$ elements (assuming every element appears twice), and the ans vector can also have a maximum size of $n/2$.

Approach 3: Using "Negation Marking" or the "Index as a Hash" approach.

Note: it comes with the constraint that the elements in the array must be within a specific range and allows only non-negative values. If the range or the non-negative constraint is not guaranteed, this approach may not be applicable.

The optimized approach takes advantage of the fact that the elements in the array are in the range of 1 to n , where n is the size of the array. It modifies the array itself to keep track of the duplicates.

Consider the example array `nums = {4, 3, 2, 7, 8, 2, 3, 1}`.

Step 1:

Initialize an empty vector to store the duplicate elements: `duplicates = {}`.

Step 2:

Iterate through each element in the array:

For the first element, `num = 4`, the corresponding index is $\text{abs}(4) - 1 = 3$. Since the value at index 3 is positive (`nums[3] = 7`), we multiply it by -1: `nums[3] = -7`.

`nums = {4, 3, 2, -7, 8, 2, 3, 1}`

For the second element, `num = 3`, the corresponding index is $\text{abs}(3) - 1 = 2$. Since the value at index 2 is positive (`nums[2] = 2`), we multiply it by -1: `nums[2] = -2`.

`nums = {4, 3, -2, -7, 8, 2, 3, 1}`

For the third element, `num = -2`, the corresponding index is $\text{abs}(-2) - 1 = 1$. Since the value at index 1 is positive (`nums[1] = 3`), we multiply it by -1: `nums[1] = -3`.

`nums = {4, -3, -2, -7, 8, 2, 3, 1}`

Continuing this process, we encounter the element 7, which has a corresponding index of 6. The value at index 6 is positive (`nums[6] = 3`), so we multiply it by -1: `nums[6] = -3`.

`nums = {4, -3, -2, -7, 8, 2, -3, 1}`

The element 8 corresponds to index 7, which is positive. We multiply it by -1: $\text{nums}[7] = -1$.

`nums = {4, -3, -2, -7, 8, 2, -3, -1}`

For the element 2, the corresponding index is 1. The value at index 1 is already negative ($\text{nums}[1] = -3$), indicating that 2 is a duplicate. So we add its absolute value to the duplicates vector: $\text{duplicates} = \{2\}$.

Finally, for the element 3, the corresponding index is 2. The value at index 2 is already negative ($\text{nums}[2] = -2$), indicating that 3 is a duplicate. So we add its absolute value to the duplicates vector: $\text{duplicates} = \{2, 3\}$.

`duplicates = {2, 3}`

This problem with a time complexity of $O(n)$ and constant space complexity ($O(1)$).