

Fibonacci Series

This C++ program calculates the nth term of the Fibonacci series using both a simple recursive function and a recursive function with memoization. The Fibonacci series is a sequence of numbers where each number is the sum of the two preceding ones, usually starting with 0 and 1.

Recursive function to calculate the nth Fibonacci number

1. The program starts by including the necessary header files for input/output and an unordered map (used for memoization).
2. The **fibonacciSeries** function is defined, which takes an integer **sequence** as input and returns the nth term of the Fibonacci series.
3. In the **fibonacciSeries** function, there are two cases:
 - Base Case: If **sequence** is 0 or 1, the function returns **sequence**. For **sequence = 0**, the Fibonacci series starts with 0, and for **sequence = 1**, it starts with 0, 1.
 - Recursive Case: If **sequence** is greater than 1, the function calculates the nth term of the Fibonacci series recursively by adding the (**sequence - 1**)th and (**sequence - 2**)th terms.

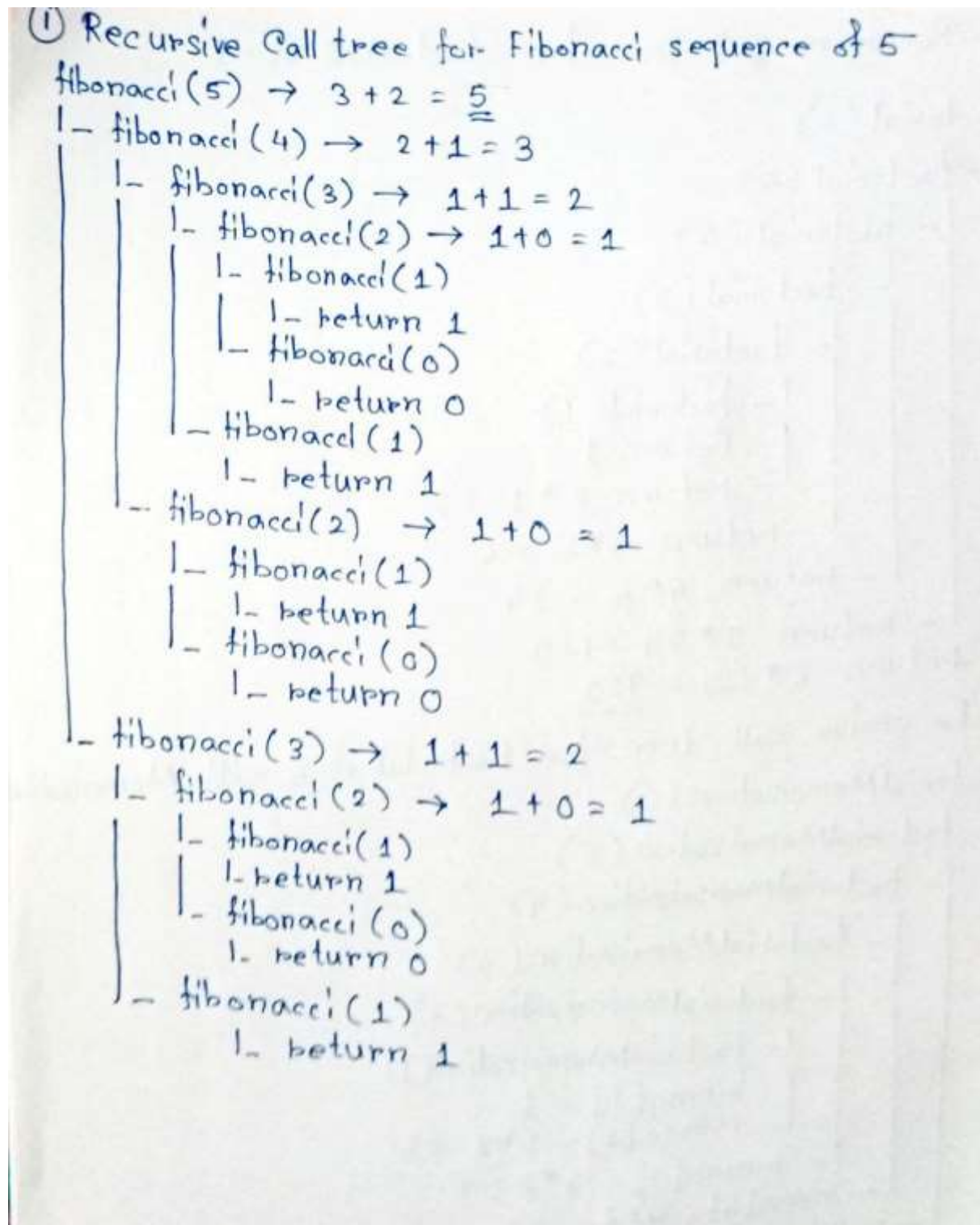
Time Complexity:

- The time complexity of the **fibonacciSeries** function without memoization is exponential, specifically $O(2^n)$. This is because the function makes two recursive calls for each sequence number greater than 1, leading to redundant calculations.

Space Complexity:

- The space complexity of the **fibonacciSeries** function without memoization is $O(n)$. This is because it requires space for n recursive calls on the call stack.

Recursive call stack for the approach:



Recursive function to calculate the nth Fibonacci number with memoization

1. The **fibonacciSeriesMemoization** function is defined, which takes an integer **sequence** as input and returns the nth term of the Fibonacci series using memoization.

2. Inside the **fibonacciSeriesMemoization** function, there is a memoization step using an unordered map named **memo**.
3. In the **fibonacciSeriesMemoization** function, there are three cases:
 - Base Case: If **sequence** is 0 or 1, the function returns **sequence**, same as the **fibonacciSeries** function.
 - Memoization Check: The function checks if the Fibonacci value for the given **sequence** is already calculated and stored in the **memo** map using **memo.count(sequence)**. If it is present, it directly returns the precalculated value.
 - Recursive Case: If the Fibonacci value is not already calculated, the function calculates it recursively by adding the **(sequence - 1)**th and **(sequence - 2)**th terms of the Fibonacci series. The result is then stored in the **memo** map for future use.

Time Complexity:

- The time complexity of the **fibonacciSeriesMemoization** function with memoization is linear, specifically $O(n)$. The memoization approach reduces redundant calculations by storing previously calculated Fibonacci values in the memo map. As a result, each Fibonacci number is computed only once, and subsequent calls retrieve the precalculated values in constant time.

Space Complexity:

- The space complexity of the **fibonacciSeriesMemoization** function with memoization is also $O(n)$. The memoization approach optimizes the recursion and reduces redundant calculations, but it requires additional space to store the memo map, which can hold up to n Fibonacci values.

Recursive call stack for the approach:

② Fibonacci Series of 5 using Memoization

fibonacciMemoization(5) \rightarrow memo[5] = 3 + 2 = 5

├ fibonacciMemoization(4) \rightarrow memo[4] = 2 + 1 = 3

│ └ fibonacciMemoization(3) \rightarrow memo[3] = 1 + 1 = 2

│ │ └ fibonacciMemoization(2) \rightarrow memo[2] = 1 + 0 = 1

│ │ │ └ fibonacciMemoization(1)

│ │ │ │ └ memo[1] = 1

│ │ │ │ └ fibonacciMemoization(0)

│ │ │ │ │ └ memo[0] = 0

│ │ │ └ fibonacciMemoization(1)

│ │ │ │ └ memo[1] = 1

│ │ └ fibonacciMemoization(2)

│ │ │ └ memo[2] = 1

│ └ fibonacciMemoization(3)

│ │ └ memo[3] = 2