## BigData – Spark – Processing

discover Spark Core & Sql

#### Outline

Example RAW to LAKE transformations

Explanation step-by-step

**Dataset** 

**Parallel Distribution** 

Reminder: Spark RAW to LAKE samples

# fead transform write

# Typical RAW to LAKE as Spark Java code

```
spark.read
  .format(« csv »)
  .option(«schema », « col1 type1, ... colN typeN »)
  .load(« hdfs://raw/team/domain/table/date=2022-10-12 »)
  .as(Encoder.bean(Bean.Class)
 .map(bean -> transformBean(bean) )
  .toDF()
  .repartition(2, « col1 »)
 .sortWithinPartition(« col1, col2, col3 »)
  .write
 .format(« parquet »)
  .save(« hdfs://lake/team/domain/table/date=2022-10-22 »);
```

#### Typical RAW to LAKE processing with Spark as SQL code **INSERT OVERWRITE** lake\_team\_domain.table SELECT /\* +REPARTITION(col1, 2) \*/ col1, col2, udf\_func1(col3, col4) as col3, transform udf\_func2(col4, col5) as col4, **FROM** read raw\_team\_domain.table **JOIN** transform lake\_anotherTeam\_domain.anotherTable x ON x.ID=id read WHERE date='2022-10-22' AND ... SORT BY col1, col2, col3 -- idem sortWithinPartition write

#### Example of LAKE Aggregation

```
INSERT OVERWRITE
   lake_team_domain.table
SELECT * FROM (
  SELECT * FROM table 1 WHERE ...
 UNION
  SELECT * FROM table 2 WHERE ...
 UNION
  SELECT * FROM table3 WHERE ..
 UNION
  SELECT * FROM table4 WHERE ..
SORT BY col1, col2, col3 -- idem sortWithinPartition
```

# 

# Example of « latest value » cristalisation analytical query « over(partition by) »

```
INSERT OVERWRITE
   lake team domain.table
SELECT
 col1,col2,.... colN -- idem * EXCEPT rank (cf issue SPARK-33164)
FROM (
 SELECT *,
   RANK() OVER (PARTITION BY id ORDER BY update_time DESC) as rank
 FROM lake_team_domain.event_table
WHERE rank=1
SORT BY col1, col2, col3
                        -- idem sortWithinPartition
```

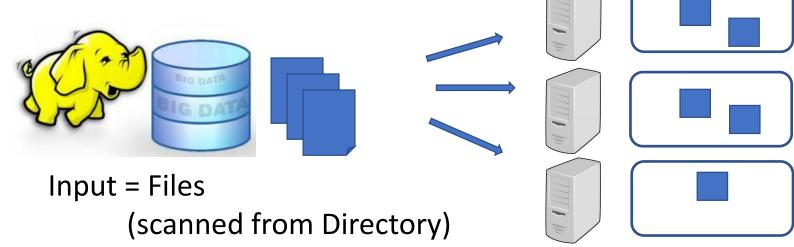
Step-by-Step explained

#### Typical RAW to LAKE as Spark Java code spark.read .format(« csv ») read Step 1/4 .option(«schema », « col1 type1, ... colN typeN ») .load(« hdfs://raw/team/domain/table/date=2022-10-12 ») .as(Encoder.bean(Bean.Class) transform .map(bean -> transformBean(bean) ) Step **2/4** .toDF() .repartition(3, « col1 ») Step **3/4** .sortWithinPartition(« col1, col2, col3 ») write .write Step **4/4** .format(« parquet ») .save(« hdfs://lake/team/domain/table/date=2022-10-22 »);

### RAW to LAKE – Step 1/4: read to Dataset

```
read

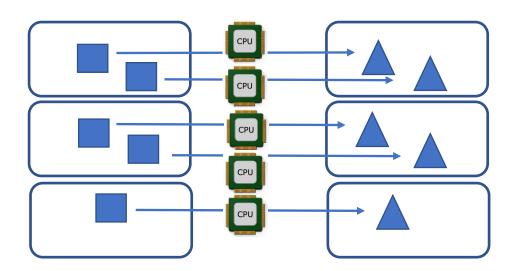
Dataset<Row> ds =
    spark.read
    .format(« csv »)
    .option(«schema », « col1 type1, ... colN typeN »)
    .load(« hdfs://raw/team/domain/table/date=2022-10-12 »)
```



Distributed Read from Storage

Result = Distributed Parts in-memory

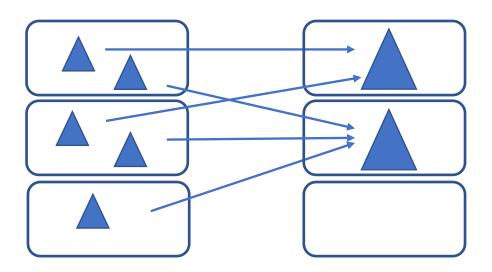
#### RAW to LAKE – Step 2/4 : Transform Dataset



Distributed Processing to compute each new part

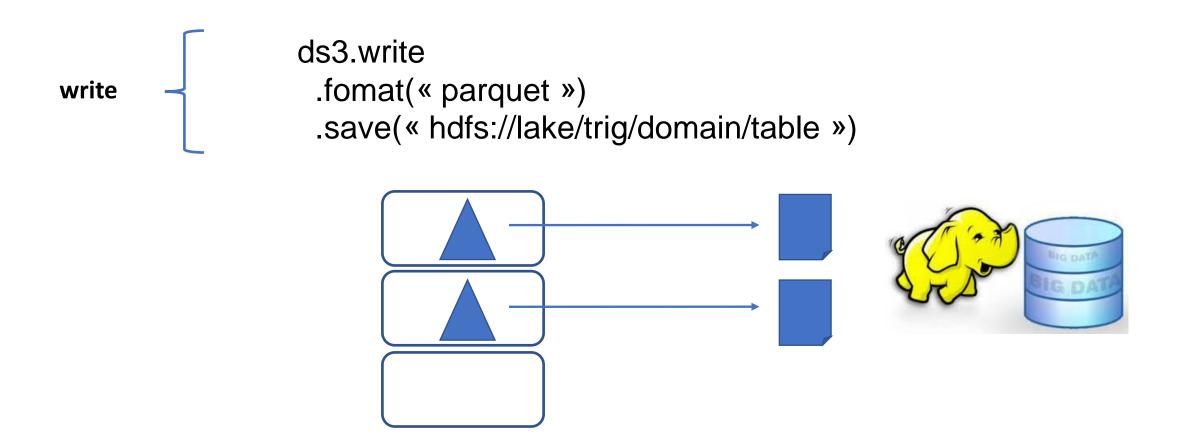
### RAW to LAKE – Step 3/4 : Repartition Dataset

```
transform Dataset<Row> ds3 = ds2
.repartition(2, « col1 »)
.sortWithinPartition(« col1, col2, col3 »)
```



Network Shuffle to distribute / group / sort data

#### RAW to LAKE – Step 4/4 : Write Dataset



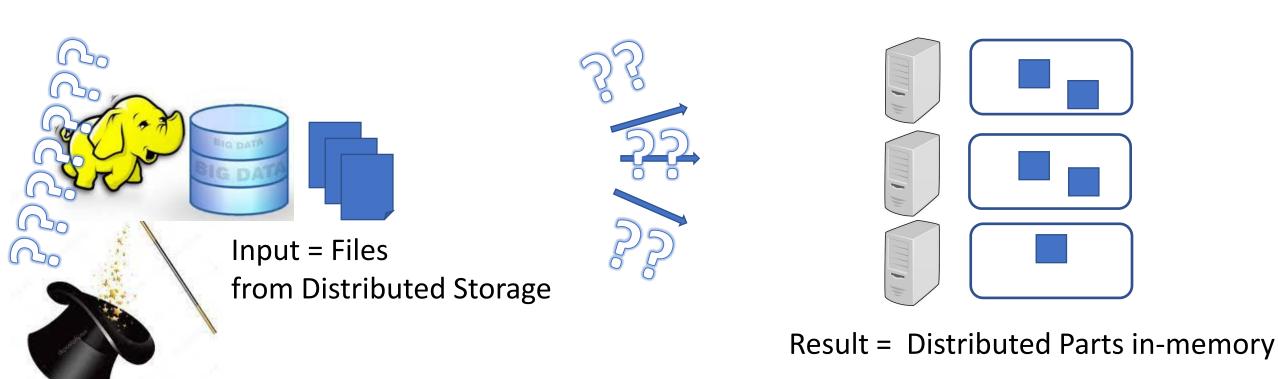
Distributed Write Dataset to Storage

## How it works?



Zooming more ..

#### RAW to LAKE – Step 1/4: read to Dataset



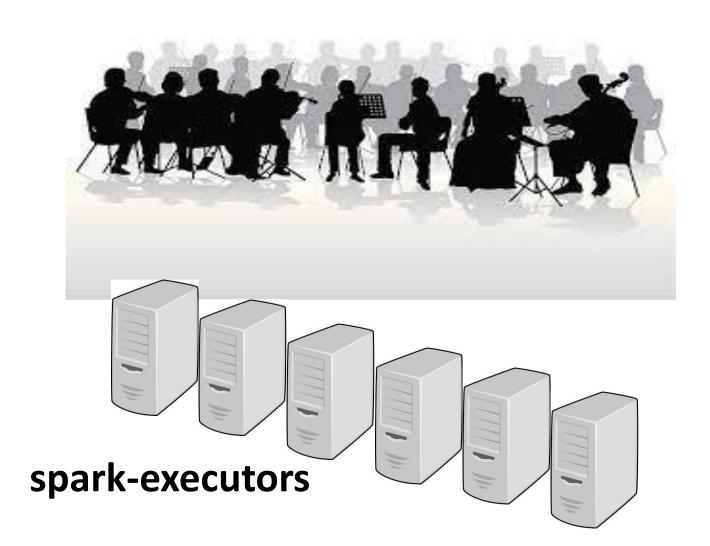
How to Assign N x Files – P x blocks ⇔ to Q x Executors ??
+ Retry on Error ?? + Communicate more ??

## Analogy: How to play music? ( N musicians without 1 Conductor != 1 Orchestra )



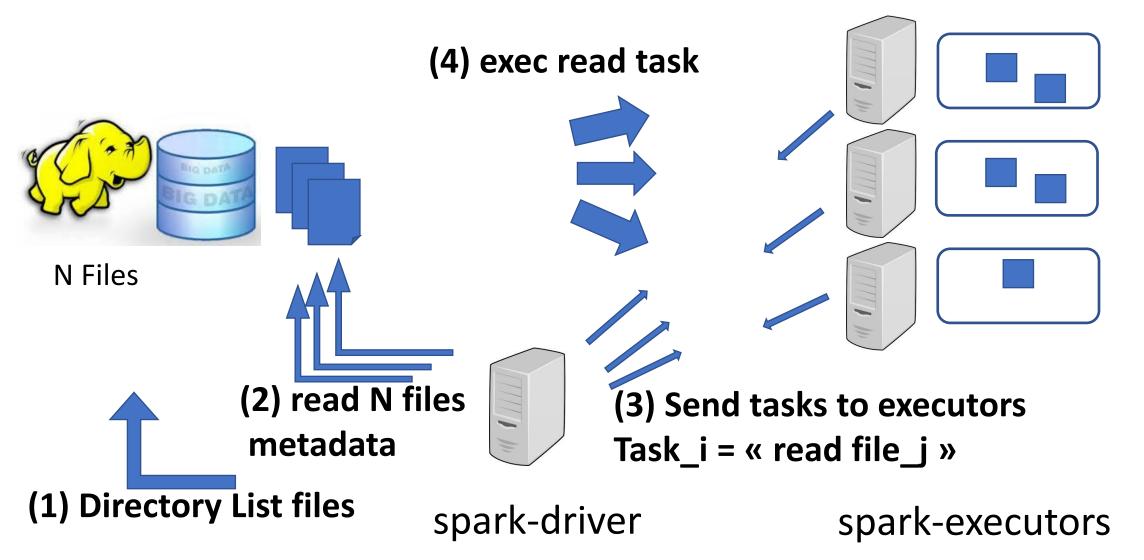
spark-driver



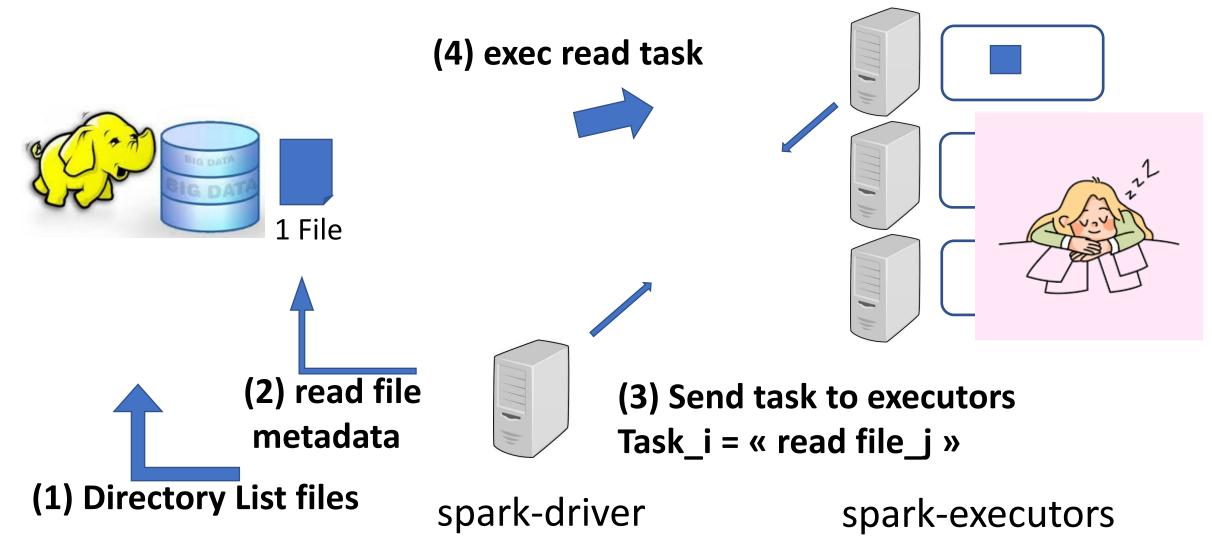


#### Read N Files – assign Tasks to Executors

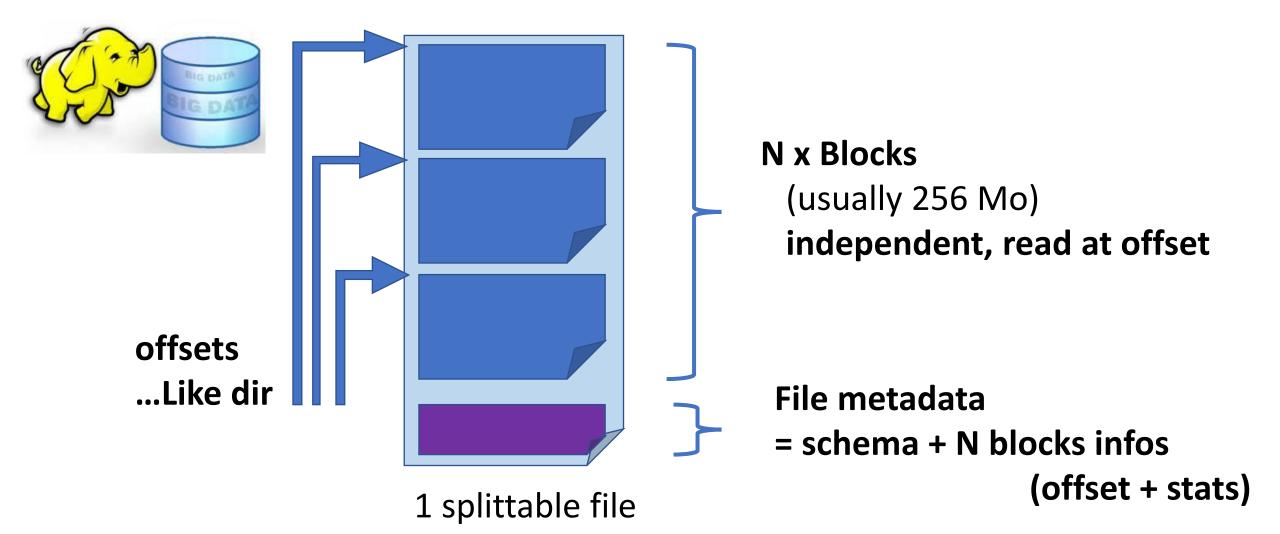
#### (5) Send tasks progress/status



# Remark [1/2] on Parallelism only 1 File -> only 1 Task



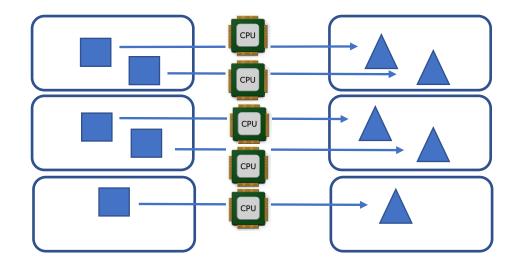
## Remark [2/2] on Parallelism Splittable File format (parquet).. Like dir



## Zooming RAW to LAKE – Step 2/4: Transform Dataset



Dataset<Row> ds2 = ds.map(row -> transformData(row) )

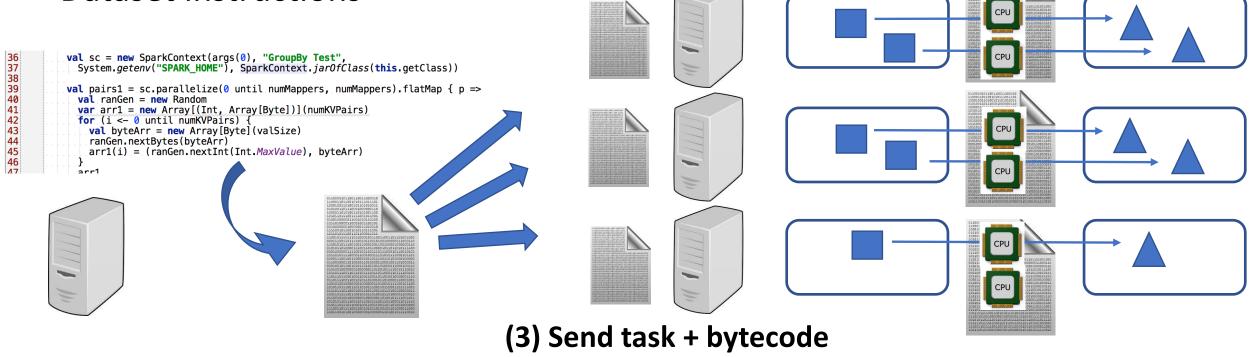


Distributed Processing to compute each new part

#### WholeStageCodeGen

#### Program

= Dataset instructions



to spark-executors

(4) Execute tasks

(1) Generate java code

(RDD Spark sub-class « WholeStageCodeGen\$i »)

(2) Compile Bytecode

# Advanced Transform ... using Row -> Java -> map()-> Java -> Row

```
ds.as( Encoders.bean(InputBean.class) )
                                                               class OutputBean {
class InputBean {
                      OutputBean transformBean(InputBean b) {
                        // complex transform in java
                        return new OutputBean(...);
```

#### Explained as().map().toDF()

```
Dataset<Row> ds = ...
                                 // convert Row->Bean
                                  Dataset<InputBean> dsInputBean =
                                    ds.as(Encoder.bean(InputBean.class))
ds.as( Encoders.bean(InputBean.Class) )
 .map(bean -> transformBean(bean) )
 .toDF()
                                 // map
                                  Dataset<OutputBean> dsOut =
                                    dsInputBean.map(bean -> transformBean(bean))
                                 // convert OutputBean -> Row
                                  Dataset<Row> df = dsOut.toDF();
```

#### Converting Tabular SQL Row to Java Beans

```
CREATE TABLE MyTable (
field1 Int,
field2 String

public class MyBean {
public int field1;
public String field2;
}
```

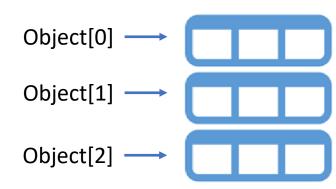
encoder = Encoders.bean(MyBean.class)

Dataset<Row> df = ...



df.as(encoder)
ds.toDF()

Dataset<MyBean> ds = ...



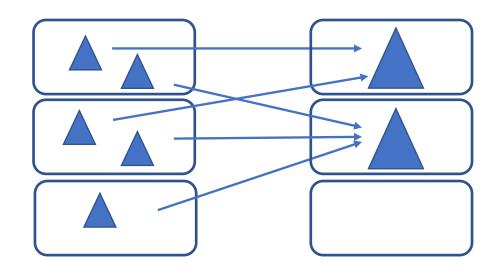
#### RAW to LAKE – Step 3/4 : Repartition Dataset



Dataset<Row> ds3 = ds2

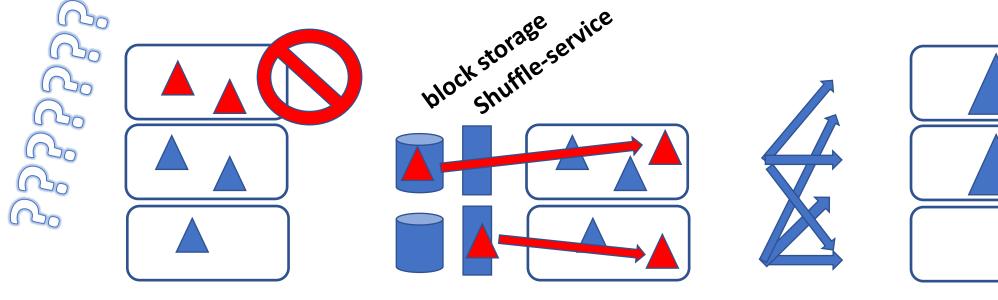
.repartition(2, « col1 »)

.sortWithinPartition(« col1, col2, col3 »)



Network Shuffle to distribute / group / sort data

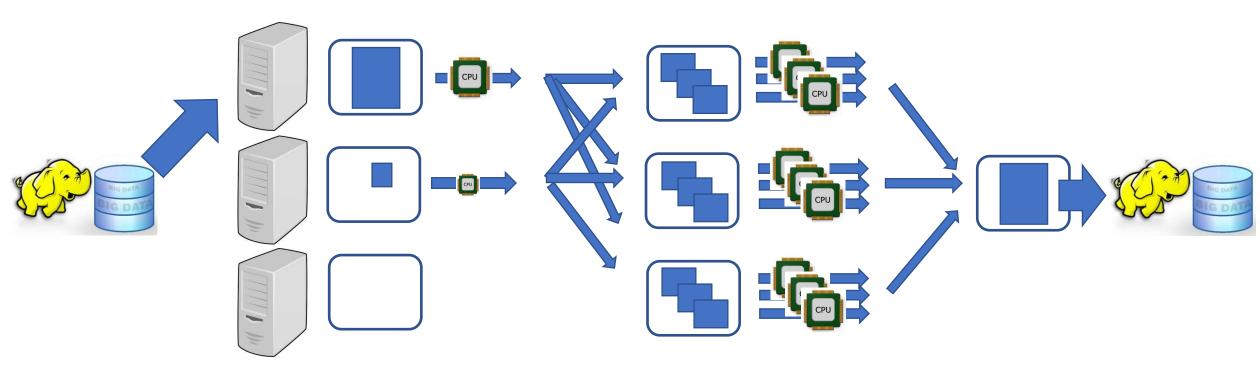
### What if a Spark-Executor crash / Node stop?



Recompute/recover lost blocks
From mem or disk
...using ShuffleService + storage

after shuffle operation,
=> Replicate (peer-to-peer)
 / Persist Blocks (mem and/or disk)
...using ShuffleService + storage

## Example usage: repartition(N).map(..).repartition(1)



INNEFFICIENT !!!Undistributed/ Badly Skewed Data

repartition(N)

**EFFICIENT** 

repartition(1)

**Avoid too many small files** 

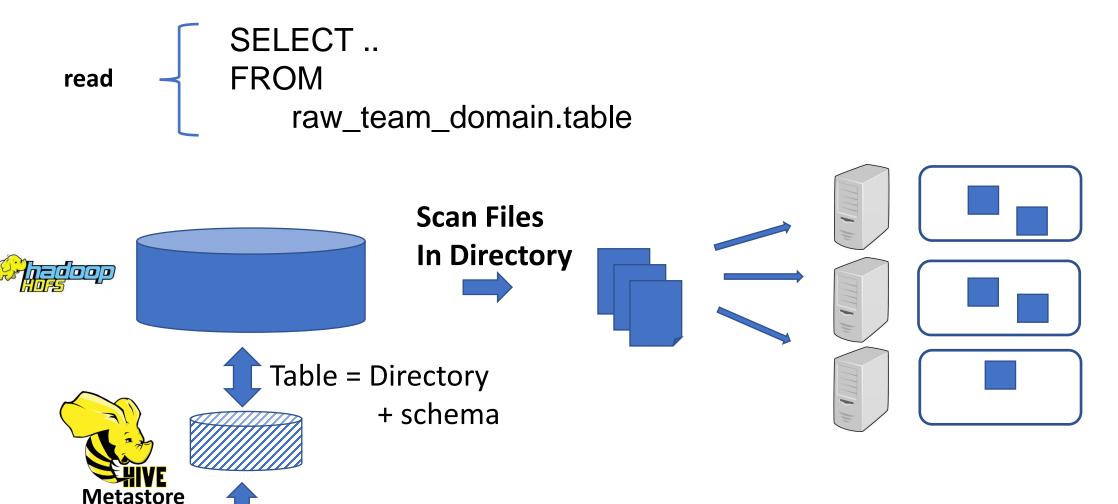
Example transformation ... in SQL

# Neminder

# Typical RAW to LAKE as Spark SQL

```
INSERT OVERWRITE
 write
                    lake team_domain.table
                SELECT /* +REPARTITION(col1, 2) */
                  col1, col2,
                  udf_func1(col3, col4) as col3,
transform
                  udf_func2(col4, col5) as col4,
                FROM
 read
                    raw_team_domain.table
                JOIN
transform
                    lake_anotherTeam_domain.anotherTable x ON x.ID=id
 read
                WHERE date='2022-10-22' AND ...
                SORT BY col1, col2, col3 -- idem sortWithinPartition
 write
```

#### Explained ... SQL (-> Files) -> Dataset



Input = SELECT .. from SQL Table

#### (Hive) MetaStore

**Store ONLY metadatas (DDL + partitions)** 

#### Mapping SQL – Dirs+Files



#### **DDL**:

CREATE EXTERNAL TABLE students (socialSeculd: Int, firstName string, lastName string, birth: Date, ...

) PARTITIONED BY (promo: Int)
STORED AS parquet

LOCATION 'hdfs://lake/students'





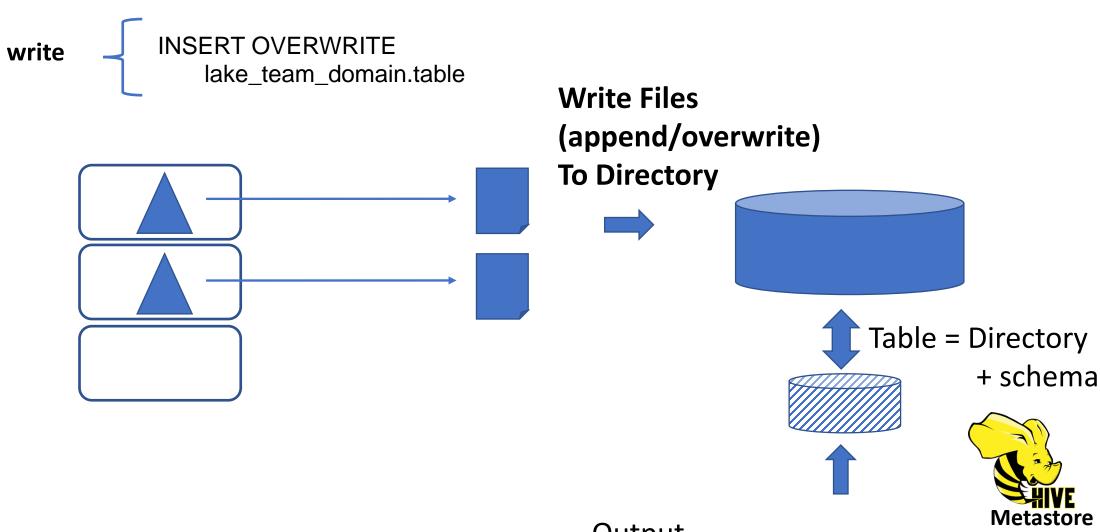




#### Location Dir + Partitions

```
« /students » (table storage dir)
```

#### Dataset -> INSERT SQL Table (-> Files)



Output

**SQL: INSERT INTO/OVERWRITE Table** 

More Java <-> Sql Interactions

#### Executing SQL from Java

=> OK in java code : if, for(), ...

```
for( int i = 0; i < 10; i++) {
 String sql = « SELECT * from db.table » + i;
  Dataset[Row] ds = spark.sql(sql);
                   NO imperative in SQL (cf PL/Sql extensions)
```

#### Java DataSet as SQL View



Dataset[Row] ds = ..

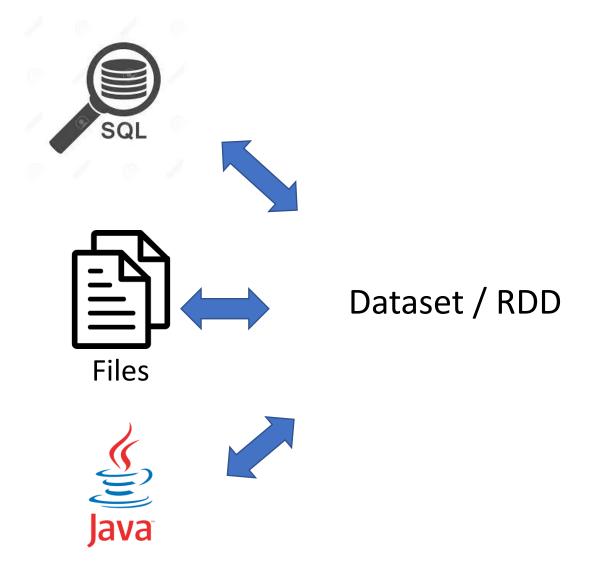
ds.createTemporaryView(« myview1 »)

spark.sql(« SELECT \* FROM myview1 »)

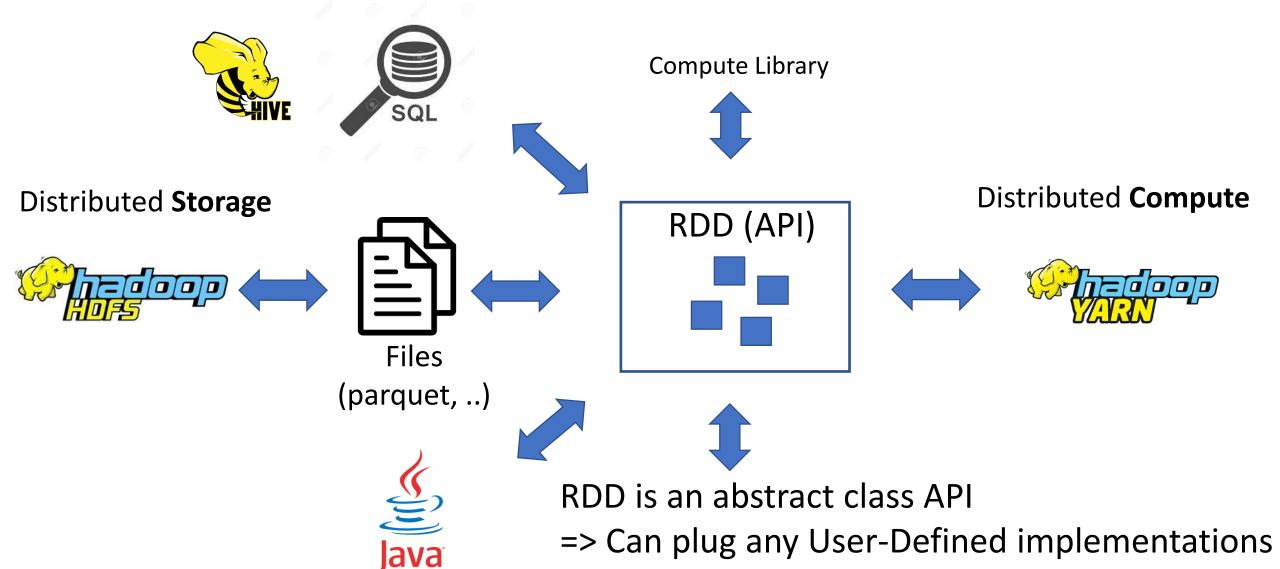
## Calling Java from SQL: User-Defined Function

```
SELECT ..
transform
                      udf_func1(col3, col4) as col3,
                      udf_func2(col4, col5) as col4,
                           int func1(int x, int y) { return x+y; }
                           spark.udf().register(« udf func1",
                               (UDF2<Integer,Integer, Integer>)::func1,
                               DataTypes.IntegerType);
```

# Spark = Unified Sql-Files-Java



# Spark: Unified Engine (Distributed Storage, Distributed Compute)

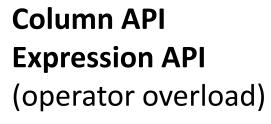


#### Dataset API .... SQL Extensions

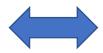


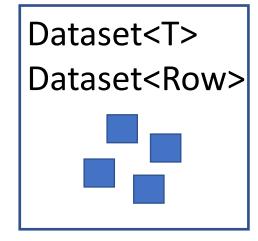




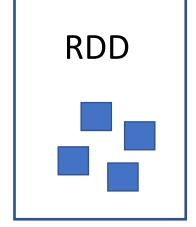












Lamba /Function + runtime compile to bytecode ( WholeStageCodeGen )





#### More Extensions: Hadoop FileSystem API



**HDFS** implements FileSystem



java.io.File adapter







More adapters



••••



#### **Distributed Storage API**







```
abstract class FileSystem {
    ..read, write, list,
}
```

Spark rely on API
=> Can plug any implementations

#### More Extensions: Cluster Scheduler API

Cluster Manager Scheduler API



#### TaskScheduler (SPI)

abstract class TaskScheduler {
 ...start,stop,
 submitTasks,cancelTasks,
 notify Host-Executor-Task changes
}







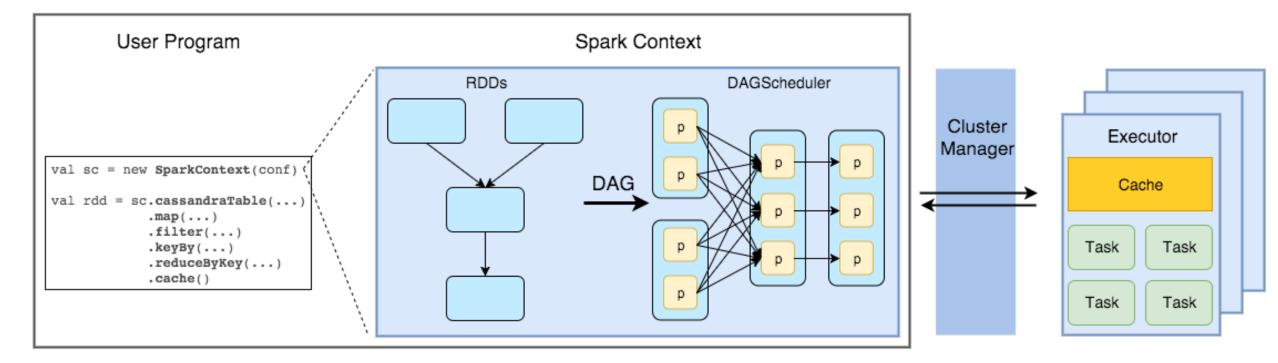




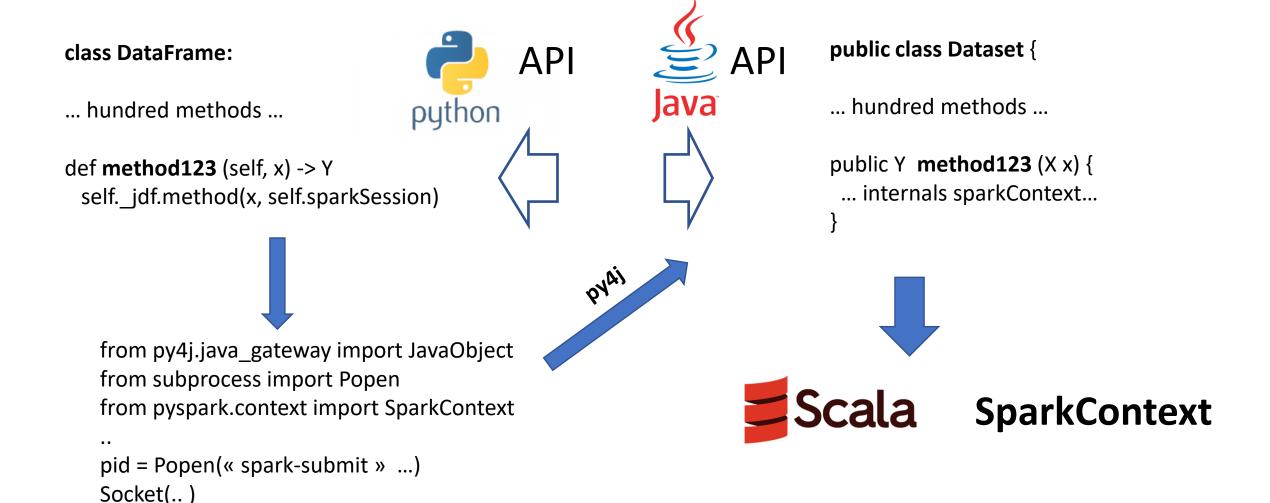




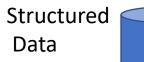
Spark Application Workers

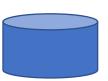


# More « Extensions » backport API to other Langages (Python, R)



#### Spark-Core + ...







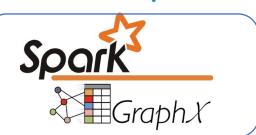














Modules











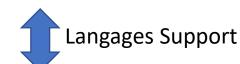
















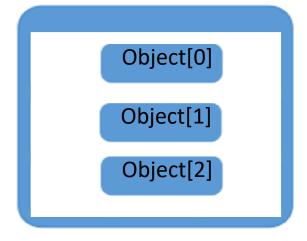




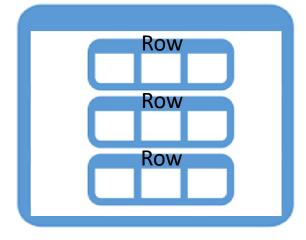
What are Dataset/RDD?
How they work?

#### DataSet ~ Set of <Data>

#### Dataset<UserDefinedClass>



Dataset<Row> = « DataFrame »



Internal RDD



#### DataSet = sql view for RDD

module spark-sql class org.apache.spark.sql.Dataset module spark core class org.apache.spark.rdd.RDD RDD = ?

Resilient

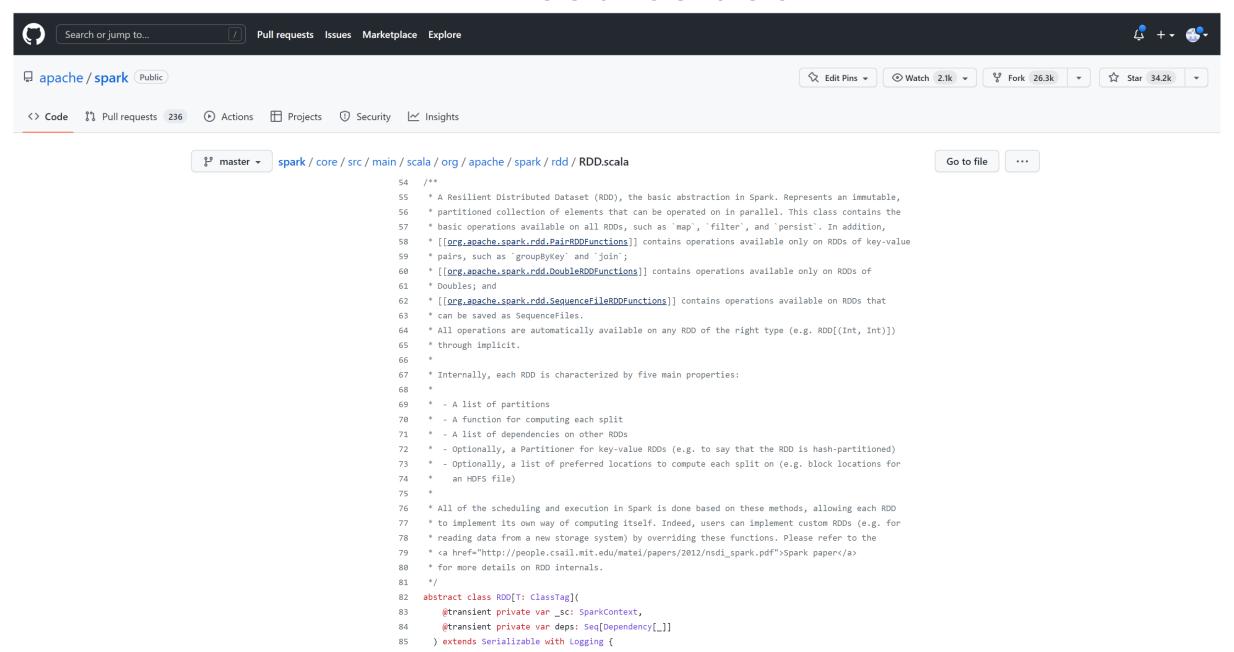
Resist to failure

**D**istributed

Horyzontal scale

**D**ataset

#### RDD source doc



## RDD Doc (1/3)

A Resilient Distributed Dataset (RDD), the basic abstraction in Spark.

Represents
an immutable,
partitioned collection of elements
that can be operated on in parallel.

This class contains the basic operations available on all RDDs, such as 'map', 'filter', and 'persist'.
In addition, PairRDDFunctions (..) of key-value pairs, (..contains) 'groupByKey' and 'join' (..)

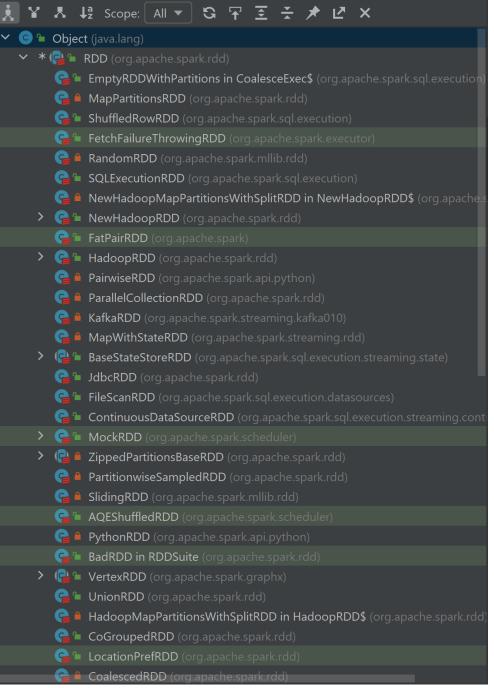
## RDD Doc (2/3)

Internally, each RDD is characterized by :

- A list of partitions
- A function for computing each split
- A list of dependencies on other RDDs
- Optionally, a Partitioner
- Optionally, a list of preferred locations

#### RDD Abstract methods

```
105
        // Methods that should be implemented by subclasses of RDD
106
107
        108
109
110
        * :: DeveloperApi ::
        * Implemented by subclasses to compute a given partition.
111
112
113
       @DeveloperApi
114
       def compute(split: Partition, context: TaskContext): Iterator[T]
115
       /**
116
        * Implemented by subclasses to return the set of partitions in this RDD. This method will only
117
118
        * be called once, so it is safe to implement a time-consuming computation in it.
119
        * The partitions in this array must satisfy the following property:
120
121
        * `rdd.partitions.zipWithIndex.forall { case (partition, index) => partition.index == index }`
122
       protected def getPartitions: Array[Partition]
123
124
125
126
        * Implemented by subclasses to return how this RDD depends on parent RDDs. This method will only
127
        * be called once, so it is safe to implement a time-consuming computation in it.
128
       protected def getDependencies: Seq[Dependency[_]] = deps
129
130
       /**
131
        * Optionally overridden by subclasses to specify placement preferences.
132
133
        */
       protected def getPreferredLocations(split: Partition): Seq[String] = Nil
134
135
136
       /** Optionally overridden by subclasses to specify how they are partitioned. */
137
       @transient val partitioner: Option[Partitioner] = None
```



Class RDD

# Abstract RDD class => (many) concrete sub-classes

```
CoalescedKDD (org.apache.spark.rdd)
   PipedRDD (org.apache.spark.rdd)
> ( EdgeRDD (org.apache.spark.graphx)
   Q ■ DataSourceRDD (org.apache.spark.sql.execution.datasources.v2)
        BlockRDD (org.apache.spark.rdd)
        BaseRRDD (org.apache.spark.api.r)
      ZippedWithIndexRDD (org.apache.spark.rdd)
   Q ■ JDBCRDD (org.apache.spark.sql.execution.datasources.jdbc)
  G ■ CartesianRDD (org.apache.spark.rdd)
        RandomVectorRDD (org.apache.spark.mllib.rdd)
   PartitionerAwareUnionRDD (org.apache.spark.rdd)
   G ShuffledRDD (org.apache.spark.rdd)
   😭 🕒 ContinuousWriteRDD (org.apache.spark.sql.execution.streaming.continuou
   😘 🕒 PartitionPruningRDD (org.apache.spark.rdd)
        CheckpointRDD (org.apache.spark.rdd)
        CyclicalDependencyRDD in RDDSuite (org.apache.spark.rdd)
       FetchFailureHidingRDD (org.apache.spark.executor)
   🚰 🖺 SubtractedRDD (org.apache.spark.rdd)
```

# 1 algorithm / transformation => 1 RDD Sub-class

Example: rdd.map(func) or rdd.flatMap(func)

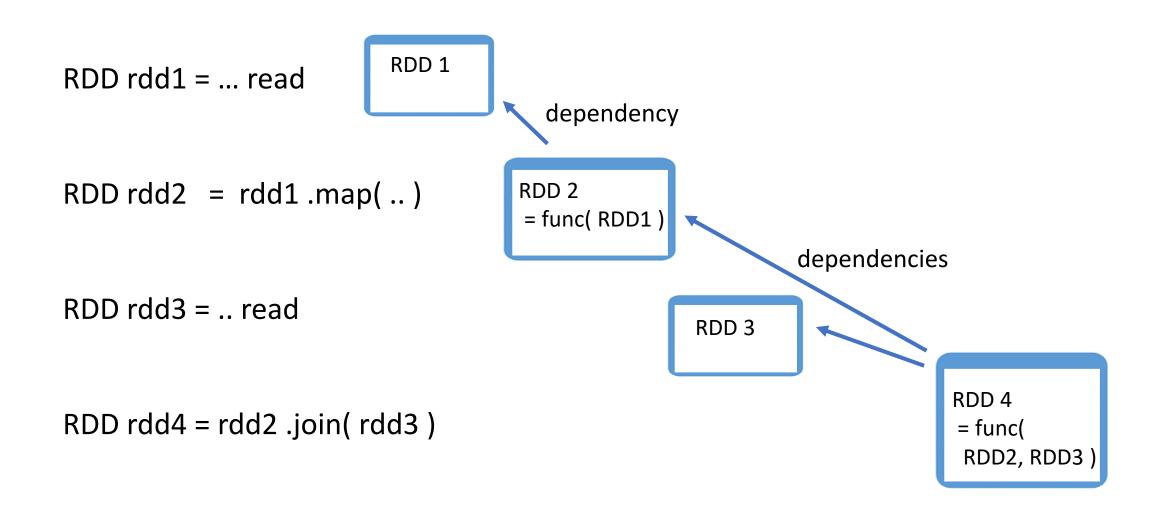
```
/**
  * Return a new RDD by applying a function to all elements of this RDD.

*/
def map[U: ClassTag](f: T => U): RDD[U] = withScope {
  val cleanF = sc.clean(f)
  new MapPartitionsRDD[U, T]( prev = this, (_, _, iter) => iter.map(cleanF))
}

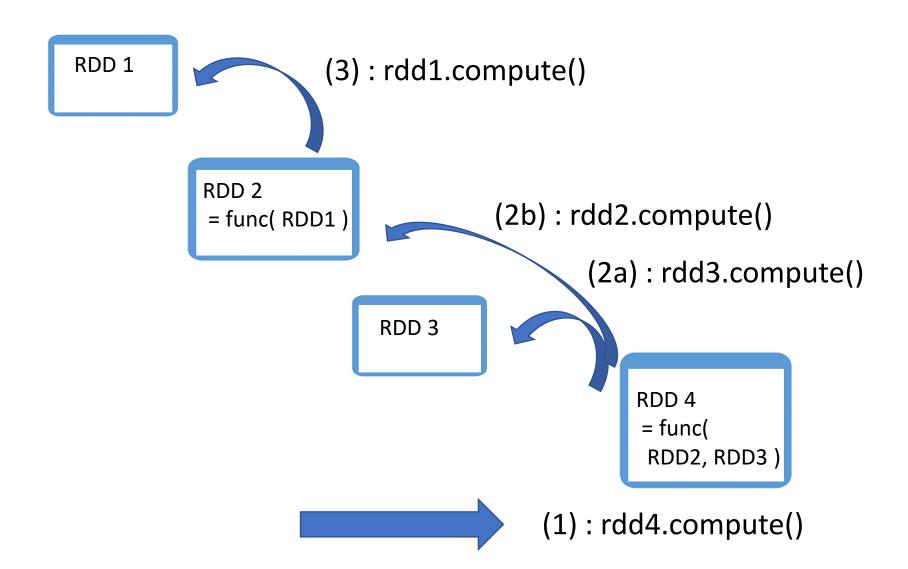
/**
  * Return a new RDD by first applying a function to all elements of this
  * RDD, and then flattening the results.
  */
def flatMap[U: ClassTag](f: T => TraversableOnce[U]): RDD[U] = withScope {
  val cleanF = sc.clean(f)
  new MapPartitionsRDD[U, T]( prev = this, (_, _, iter) => iter.flatMap(cleanF))
}
```

```
st An RDD that applies the provided function to every partition of the parent RDD.
 @param prev the parent RDD.
  Oparam f The function used to map a tuple of (TaskContext, partition index, input iterator) to
  Oparam preservesPartitioning Whether the input function preserves the partitioner, which should
 Oparam isFromBarrier Indicates whether this RDD is transformed from an RDDBarrier, a stage
  Oparam isOrderSensitive whether or not the function is order-sensitive. If it's order
var prev: RDD[T],
  f: (TaskContext, Int, Iterator[T]) => Iterator[U], // (TaskContext, partition index, iterator)
   preservesPartitioning: Boolean = false,
   isFromBarrier: Boolean = false,
   isOrderSensitive: Boolean = false)
 extends RDD[U](prev) {
```

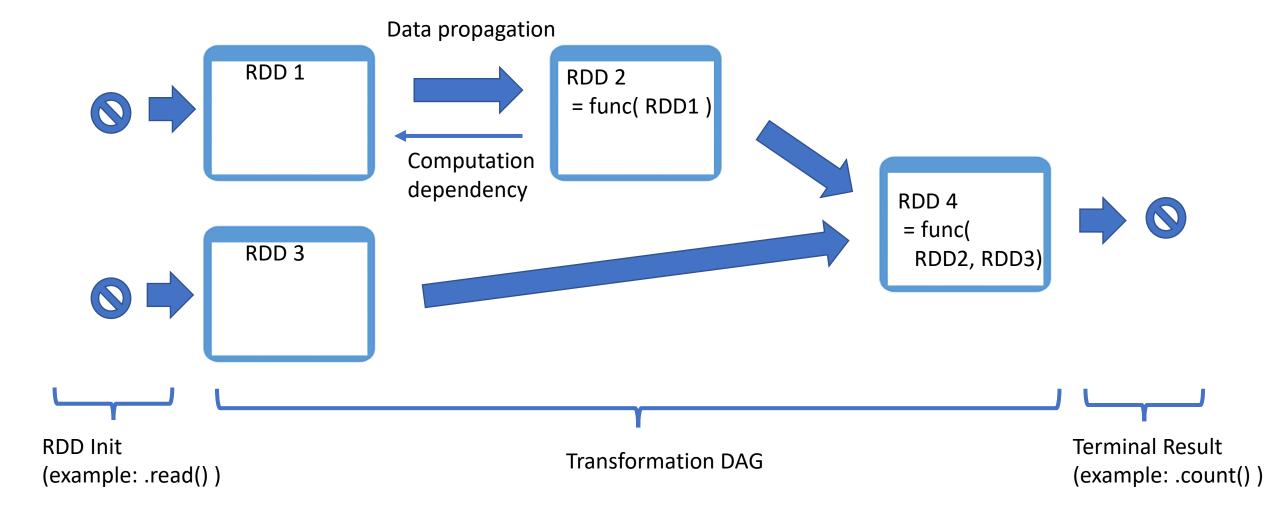
# Call transform function => Create new RDD (linked)



# Call compute() => ... dependency.compute()



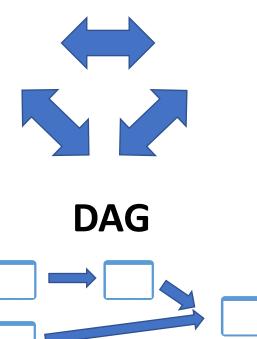
# RDD Dependencies: DAG (Directed Acyclic Graph)



# 3 equivalent formalisms: SSA create Api, Expression Algebra, DAG

# **SSA = Single State Assignments RDD API**

```
RDD rdd1 = ... read
RDD rdd2 = rdd1 .map( .. )
RDD rdd3 = .. read
RDD rdd4 = rdd2 .join( rdd3 )
```



#### **Expression Algebra, Sql**

```
SELECT map(t1) FROM Table1 t1
JOIN Table2 t2 on ..

new JoinRDD(
new MapRDD( readRDD(table1) ),
 readRDD(table2)
)
```

## RDD Doc (3/3)

All (..) in Spark is done based on these methods, allowing each RDD to implement its own way of computing itself.

Indeed, users can implement custom RDDs (e.g. for reading data from a new storage system) by overriding these functions.

Please refer to the

<a href="http://people.csail.mit.edu/matei/papers/2012/nsdi\_spark.pdf">Spark paper</a>for more details on RDD internals.

#### RDD Paper

A Fault-Tolerant
Abstraction
For In-Memory
Cluster
computing

To achieve fault tolerance efficiently,
RDDs provide a restricted form of shared memory,
based on coarse-grained transformations

1 / 14 | - 67% + | 🕏 🔕

#### Resilient Distributed Datasets: A Fault-Tolerant Abstraction for In-Memory Cluster Computing

Matei Zaharia, Mosharaf Chowdhury, Tathagata Das, Ankur Dave, Justin Ma, Murphy McCauley, Michael J. Franklin, Scott Shenker, Ion Stoica University of California, Berkeley

#### Abstract

We present Resilient Distributed Datasets (RDDs), a distributed memory abstraction that lets programmers perform in-memory computations on large clusters in a fault-tolerant manner. RDDs are motivated by two types of applications that current computing frameworks handle inefficiently: iterative algorithms and interactive data mining tools. In both cases, keeping data in memory can improve performance by an order of magnitude. To achieve fault tolerance efficiently, RDDs provide a restricted form of shared memory, based on coarsegrained transformations rather than fine-grained updates to shared state. However, we show that RDDs are expressive enough to capture a wide class of computations, including recent specialized programming models for iterative jobs, such as Pregel, and new applications that these models do not capture. We have implemented RDDs in a system called Spark, which we evaluate through a variety of user applications and benchmarks.

#### 1 Introduction

Cluster computing frameworks like MapReduce [10] and Dryad [19] have been widely adopted for large-scale data analytics. These systems let users write parallel computations using a set of high-level operators, without having to worry about work distribution and fault tolerance.

Although current frameworks provide numerous abstractions for accessing a cluster's computational resources, they lack abstractions for leveraging distributed memory. This makes them inefficient for an important class of emerging applications: those that reuse intermediate results across multiple computations. Data reuse is common in many iterative machine learning and graph algorithms, including PageRank, K-means clustering, and logistic regression. Another compelling use case is interactive data mining, where a user runs multiple adhoc queries on the same subset of the data. Unfortunately, in most current frameworks, the only way to reuse data between computations (e.g., between two MapReduce jobs) is to write it to an external stable storage system, e.g., a distributed file system. This incurs substantial overheads due to data replication, disk I/O, and serialization, which can dominate application execution times.

Recognizing this problem, researchers have developed specialized frameworks for some applications that require data reuse. For example, Pregel [22] is a system for iterative graph computations that keeps intermediate data in memory, while HaLoop [7] offers an iterative MapReduce interface. However, these frameworks only support specific computation patterns (e.g., looping a series of MapReduce steps), and perform data sharing implicitly for these patterns. They do not provide abstractions for more general reuse, e.g., to let a user load several datasets into memory and run ad-hoc queries across them.

In this paper, we propose a new abstraction called resilient distributed datasets (RDDs) that enables efficient data reuse in a broad range of applications. RDDs are fault-tolerant, parallel data structures that let users explicitly persist intermediate results in memory, control their partitioning to optimize data placement, and manipulate them using a rich set of operators.

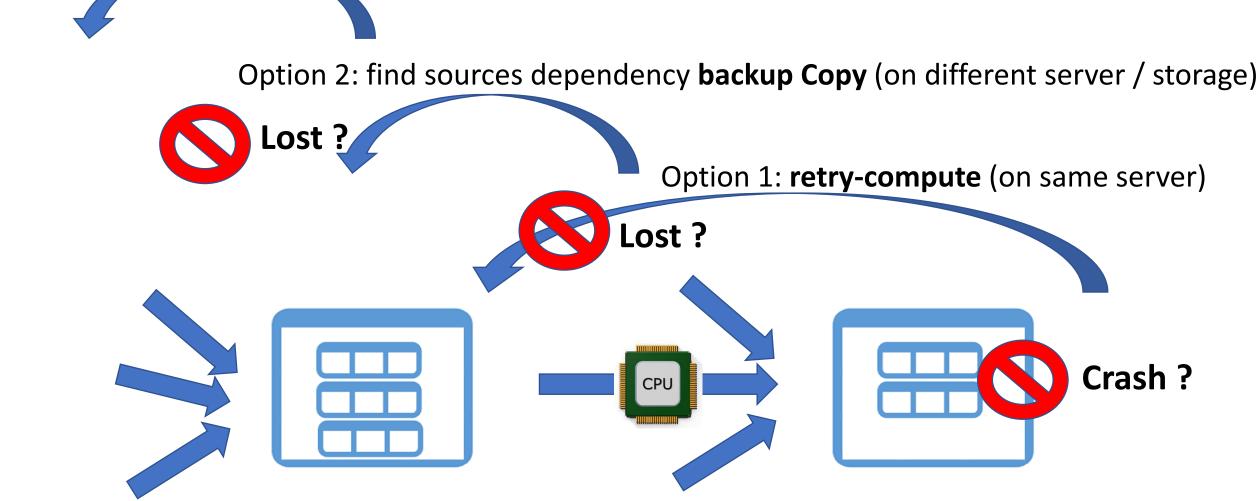
The main challenge in designing RDDs is defining a programming interface that can provide fault tolerance efficiently. Existing abstractions for in-memory storage on clusters, such as distributed shared memory [24], keyvalue stores [25], databases, and Piccolo [27], offer an interface based on fine-grained updates to mutable state (e.g., cells in a table). With this interface, the only ways to provide fault tolerance are to replicate the data across machines or to log updates across machines. Both approaches are expensive for data-intensive workloads, as they require copying large amounts of data over the cluster network, whose bandwidth is far lower than that of RAM, and they incur substantial storage overhead.

In contrast to these systems, RDDs provide an interface based on coarse-grained transformations (e.g., map, filter and join) that apply the same operation to many data items. This allows them to efficiently provide fault tolerance by logging the transformations used to build a dataset (its lineage) rather than the actual data. If a partition of an RDD is lost, the RDD has enough information about how it was derived from other RDDs to recompute

 $<sup>^1 \</sup>mbox{Checkpointing}$  the data in some RDDs may be useful when a lineage chain grows large, however, and we discuss how to do it in §5.4.

# Fault Tolerant - Computation

Option 3: recompute dependency sources



#### CoarseGrain ... Scheduler/Executer

#### **Spark-driver**

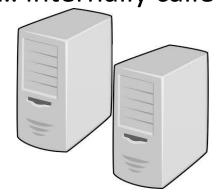
Implements Fault Tolerance+Distribution
... internally called « CoarseGrainScheduler »



#### **Spark-executor**

Implements task main loop

... internally called « CoarseGrainExecutor »

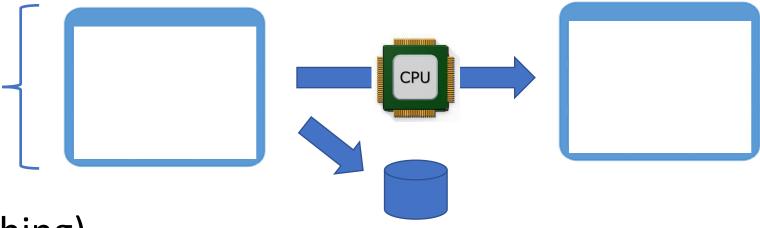


#### CoarseGrain

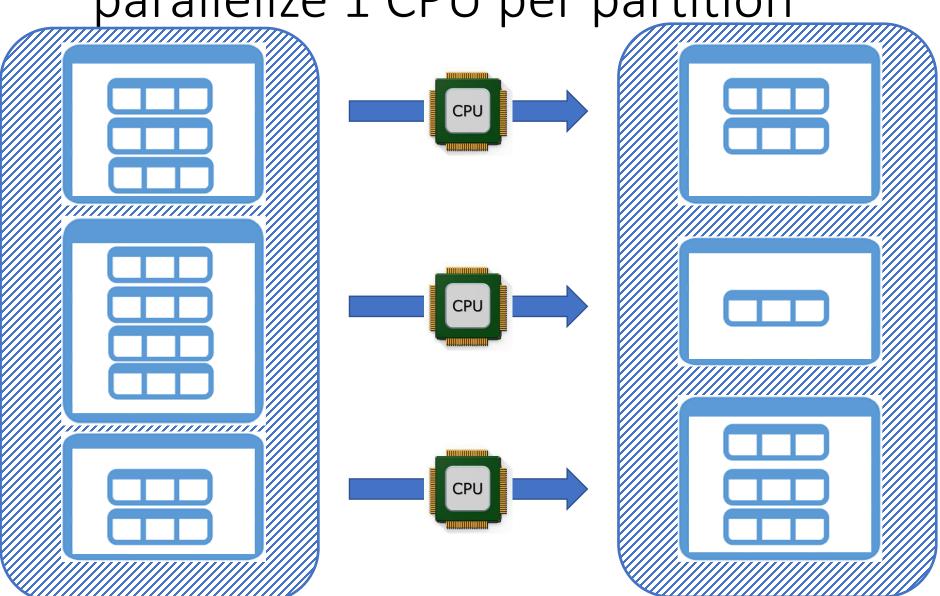
partition

= unit of caching/
recomputation

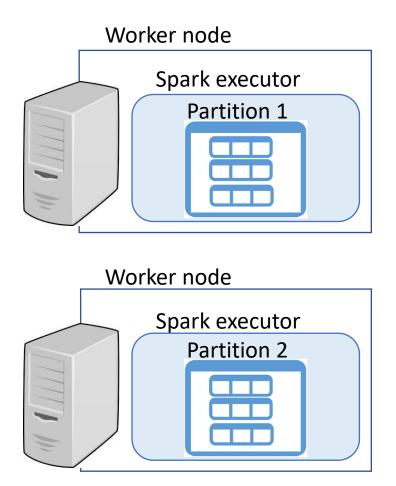
(all elements or nothing)



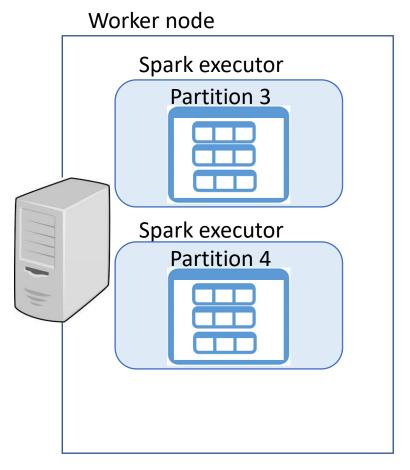
DataSet: Collection of Objects, parallelize 1 CPU per partition



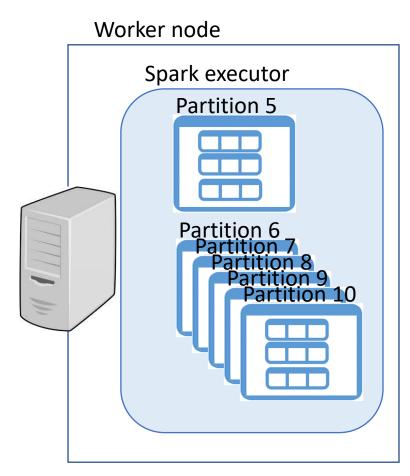
#### Distribution: Partition < Executor < Node







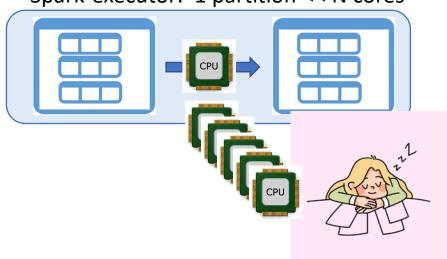
Several spark-executor processes per nodes



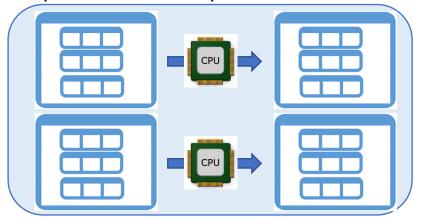
Several partitions Per spark-executor

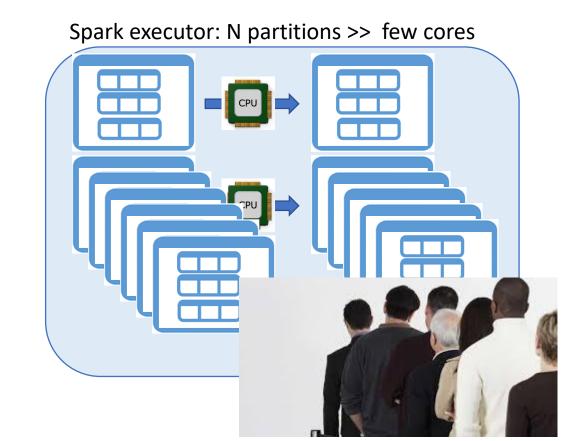
# Optimize parralelism: Adapt partitions to number of Cores

Spark-executor: 1 partition << N cores

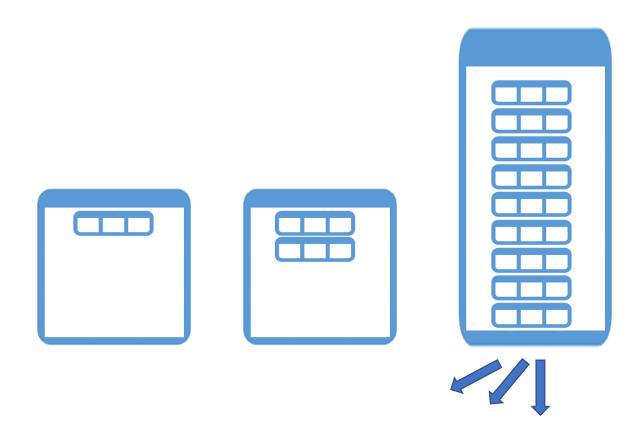


Spark executor: N partitions ~ N cores



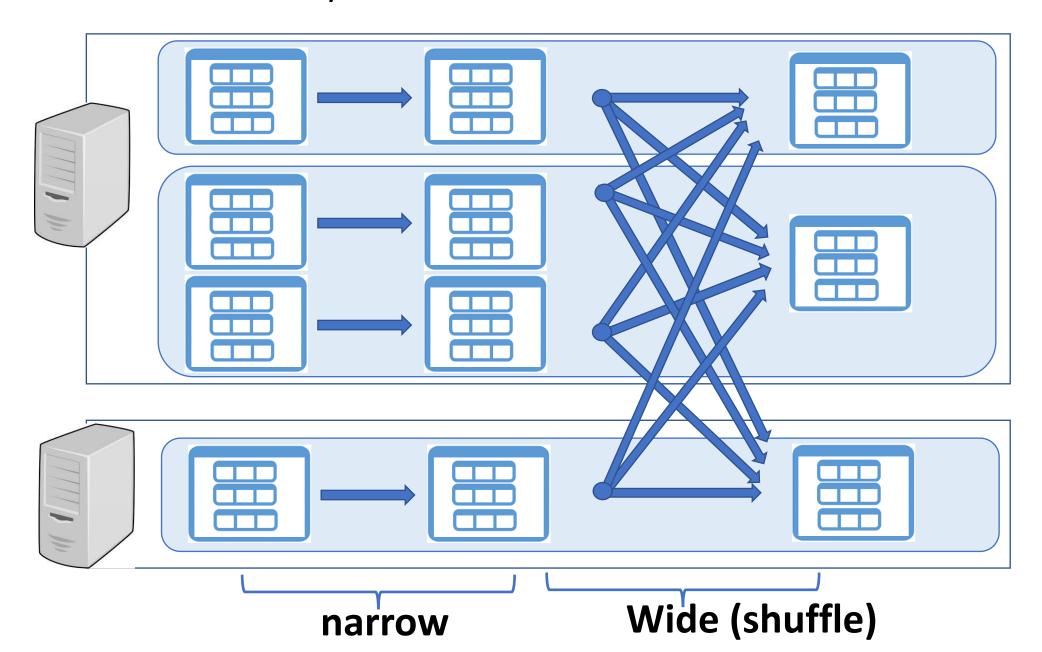


# Skewed Data ... need Repartitioned equally



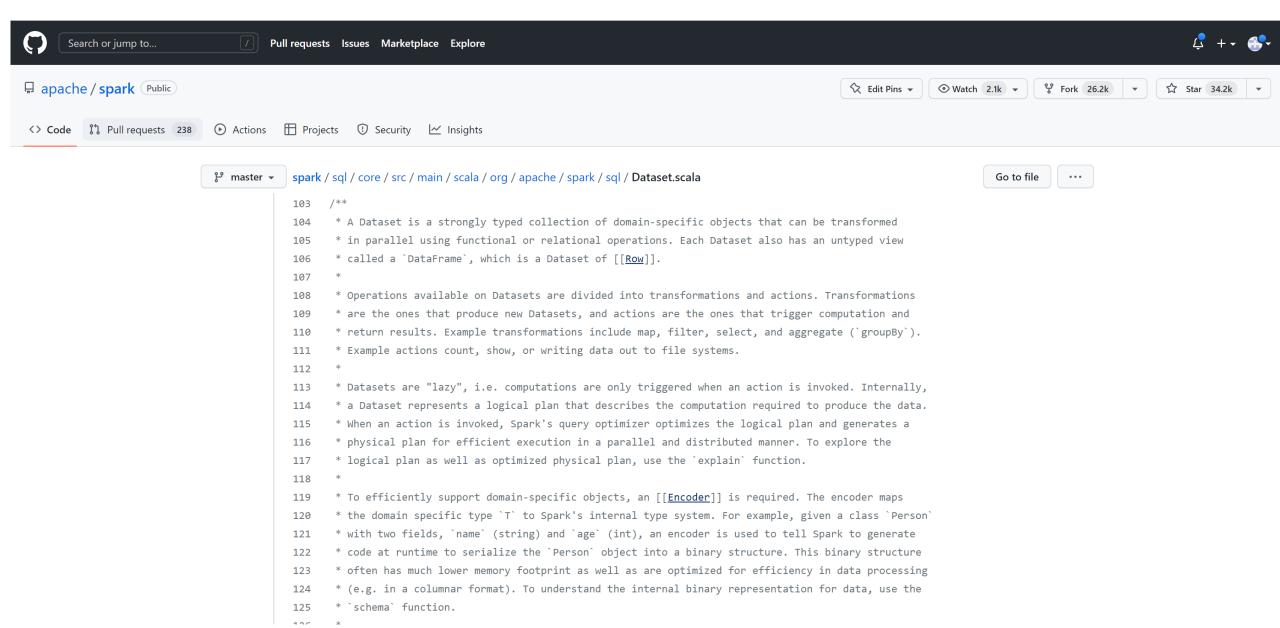
Target: move each row[i] to node[j] j = « rowId modulo N »

## « Narrow » / « Wide » Transformations



RDD ... OK so What is a Dataset ?

#### Dataset source doc



## Dataset Doc (1/4)

A strongly typed collection of domain-specific objects

that can be transformed in parallel

using functional or relational operations

Associated untyped view: DataFrame = Dataset<Row>

#### Dataset Doc (2/4)

Operations on Datasets:

**Transformations = produce new Datasets** 

**Actions = trigger computation and return results.** 

Example transformations: map, filter, select, groupBy...

Example actions: count, show, write ...

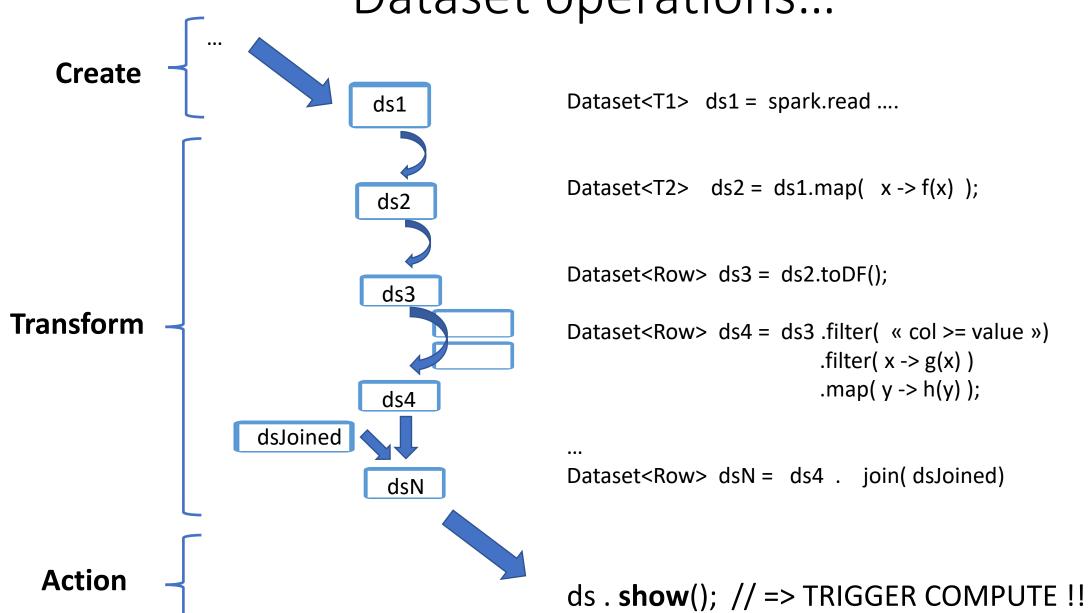
#### Datasets are "lazy",

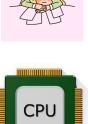
i.e. computations are only triggered when an action is invoked.

# Transformation... « produce » new Dataset NO UPDATE method Dataset are computable / « immutable »

```
dataset. <noSetter> ();
Transform Dataset newDataset = dataset . <createNewWithTransform> ( ... )
Example: chain method
dataset = ds .filter(..) .filter(..) .map(..) .groupBy(..),
```

#### Dataset operations...





#### Dataset Doc (3/4)

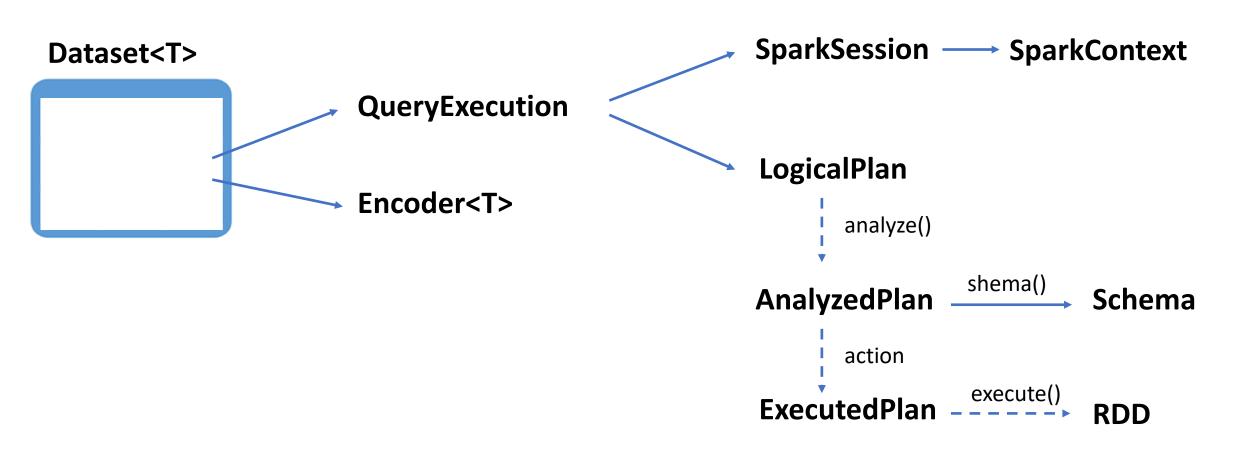
Internally, a Dataset represents a logical plan that **describes** the computation required **to produce the data**.

When an action is invoked,

Spark's query optimizer optimizes the logical plan
and generates a physical plan
for efficient execution in a parallel and distributed manner.

To explore.. Use 'explain plan'

#### Internally...



#### Dataset Doc 4/4

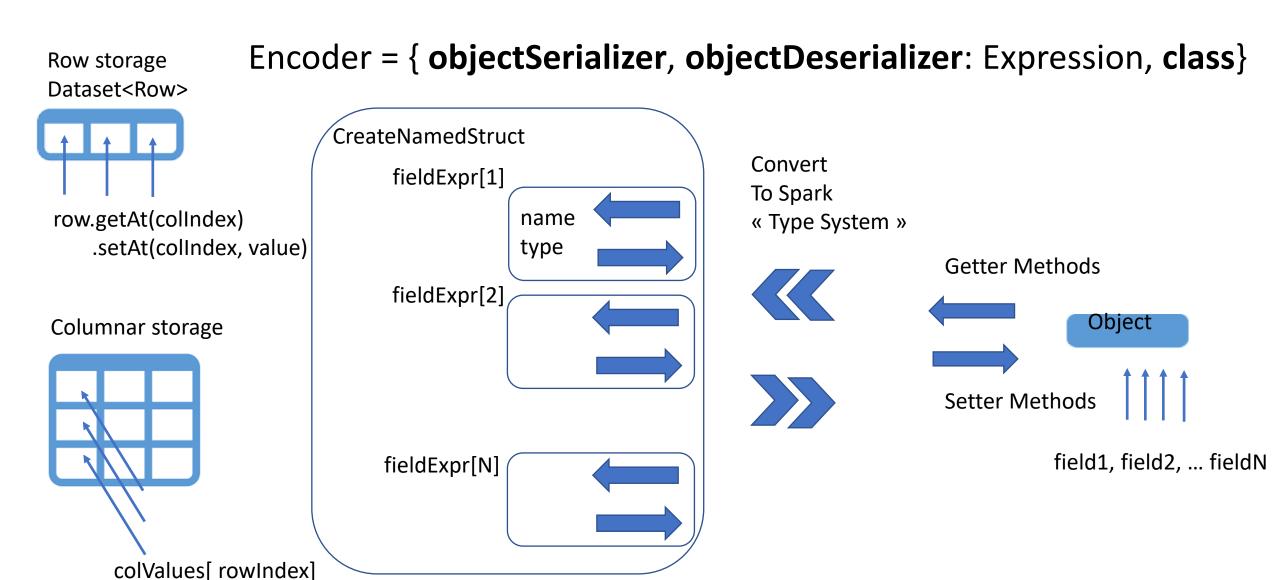
To efficiently support domain-specific objects, an 'Encoder' is required.

The encoder maps the domain specific type `T` to Spark's internal type system.

- (...) to generate code at runtime to serialize object into a much efficient binary structure
- (..) with lower memory footprint
- (..) optimized for processing (e.g. in a columnar format).

To explore representation of data, use `schema` function.

#### Encoder<T>



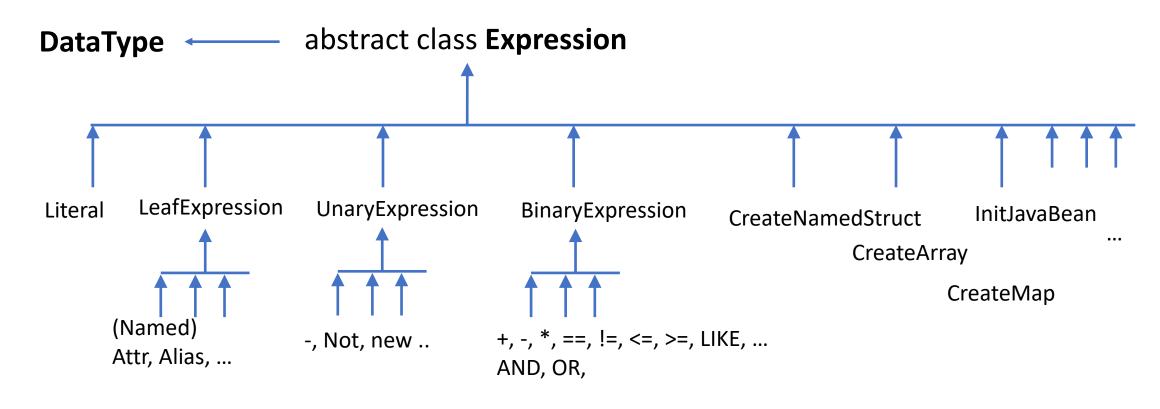


## Spark « Type System » DataTypes ...



#### Encoder .. Internal Expression with DataType

Expression abstract class AST (Abstract Syntaxic Tree) for Sql / CodeGenerator / Java Getter-Setter





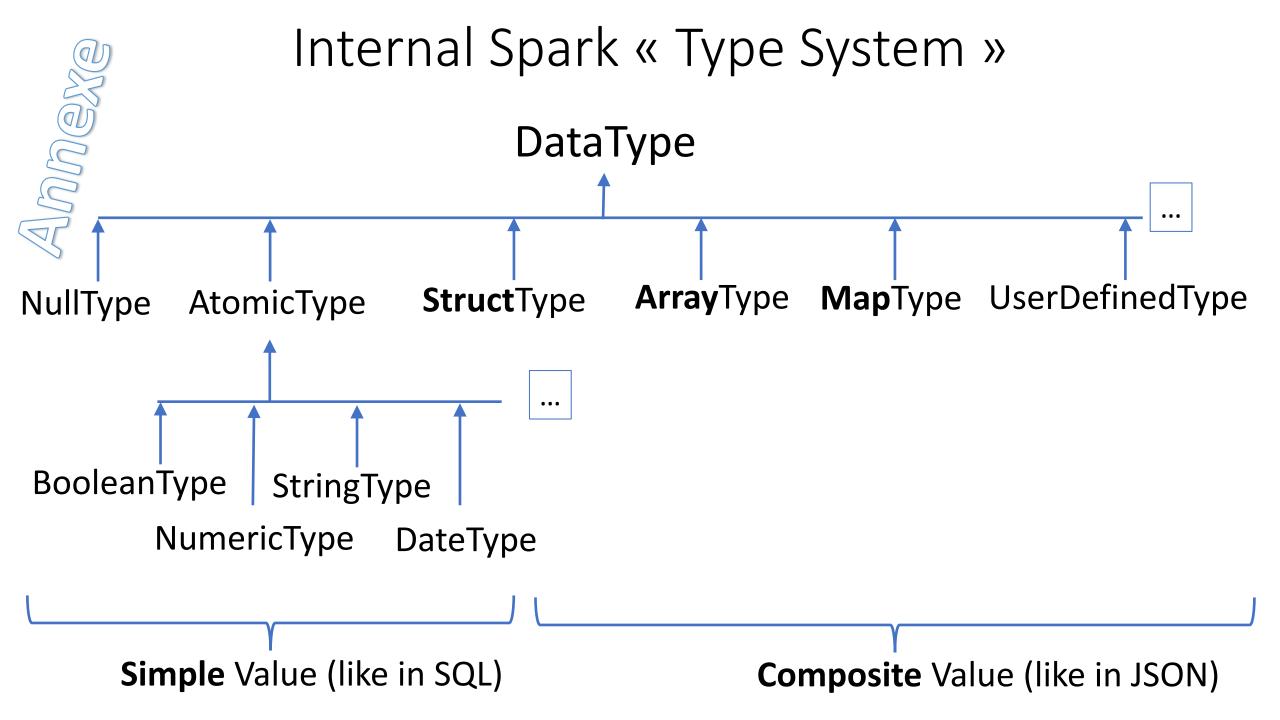
#### Expression ...

Complex ... INTERNAL ... ( not need to understand, to use Spark)

Looks like a « **SQL Compiler** »

... support **NOT ONLY Sql literal values**, but also **NamedStruct**, **Array**, **Map**, **Java Objects** !!!

Used internally to compile + generate code





#### Spark DataType

```
♠ AbstractDataType (org.apache.spark.sql.types)
* * PataType (org.apache.spark.sql.types)
      ObjectType (org.apache.spark.sql.types)
   > ( UserDefinedType (org.apache.spark.sql.types)
   > CalendarIntervalType (org.apache.spark.sql.types)
      😭 🗀 MapType (org.apache.spark.sql.types)
      G ■ NullType (org.apache.spark.sql.types)

✓ (☐ ? AtomicType (org.apache.spark.sql.types),
      > G BinaryType (org.apache.spark.sql.types)
      > G BooleanType (org.apache.spark.sql.types)
      > ( NumericType (org.apache.spark.sql.types)
      > ( DatetimeType (org.apache.spark.sql.types)
      > ( AnsiIntervalType (org.apache.spark.sql.types)
         🖕 🖆 VarcharType (org.apache.spark.sql.types)
         CharType (org.apache.spark.sql.types)
      > G StringType (org.apache.spark.sql.types)
      🖕 🖿 StructType (org.apache.spark.sql.types)
      🖕 🖿 ArrayType (org.apache.spark.sql.types),
```

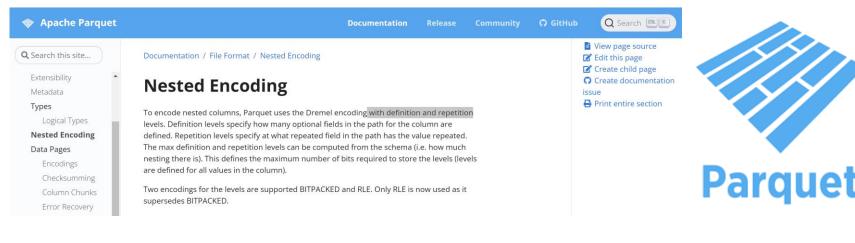
#### Hive - Spark SQL supports Struct, List, Map ...

```
Example:

CREATE EXTERNAL TABLE `student` (
    firstName string, lastName string,

practicedSports array< named_struct< name: string, numberYear: int > >,
    diploma map<string, named_struct<mention: string, obtentionDate: Date > >)
```

#### Nested fields in File Format: Parquet / Orc / Json



Parquet DataType ~ Spark DataType
Nested Encoding with « definition » + « repetition »







JSON DataType (map with string only)

Spark DataType

#### Nested Fields in Spark SQL UDF

SELECT ename, dept\_list FROM employee

```
+-----+
| ename | dept_list |
+-----+
| Tom | [20] |
| Jerry | [10,20] |
| Riley | [20,30,40] |
+-----+
```

SELECT ename,
 exists(dept\_list, x -> x = 10) as found10
FROM employee

```
+-----+
| ename | found10 |
+-----+
| Tom | false |
| Jerry | true |
| Riley | false |
+-----+
```

#### SQL Grammar Extension: « lateral view »

### SELECT ename, dept\_list FROM employee

```
+-----+
| ename | dept_list |
+-----+
| Tom | [20] |
| Jerry | [10,20] |
| Riley | [20,30,40] |
```

SELECT ename, dept\_id FROM employee LATERAL VIEW explode(dept\_list) depts AS dept\_id;

#### More SQL: collect\_list(row) -> List

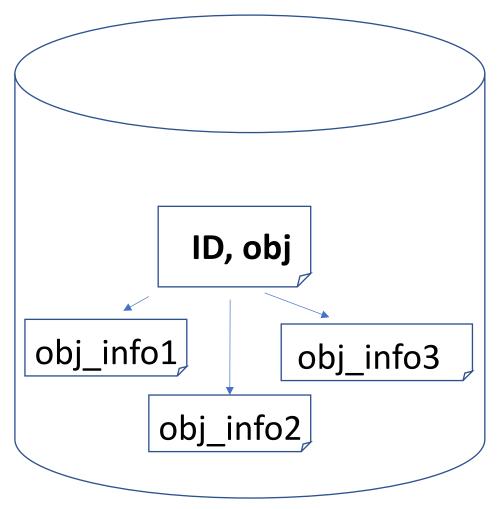
SELECT ename, dept\_list FROM employee SELECT ename, collect\_list(dept\_id + 1) as ls FROM (SELECT employee LATERAL VIEW explode(dept\_list) depts AS dept\_id ) GROUP BY ename

```
+-----+
| ename | dept_list |
+-----+
| Tom | [20] |
| Jerry | [10,20] |
| Riley | [20,30,40] |
+-----+
```

#### UDF / UDAF (User Defined Aggregate Function)

```
Example UDF: f(x, y) { return x + y }
       Function like in Math: idempotent, side-effect less, ...
! = UDAF : Aggregate / Accumulator
      like in « SELECT count(..), sum(..), average(..) FROM .. GROUP BY .. »
      Object instance, Class with 3 methods:
        init()
        add(value)
        Result getResult()
```

#### List, Map, Struct ... denormalize data, avoid Joins



[ID, obj, obj\_info1, obj\_info2, ...obj\_infoN

Normalized relationnal database

Efficient DE-normalized analytics system

Enough for Today!

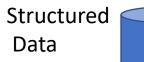
... but more to come

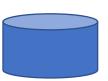
Questions are Welcome

Take Away

What Did you learn?

#### Spark-Core + ...







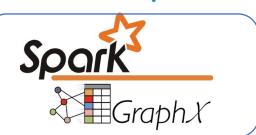














Modules











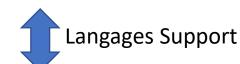










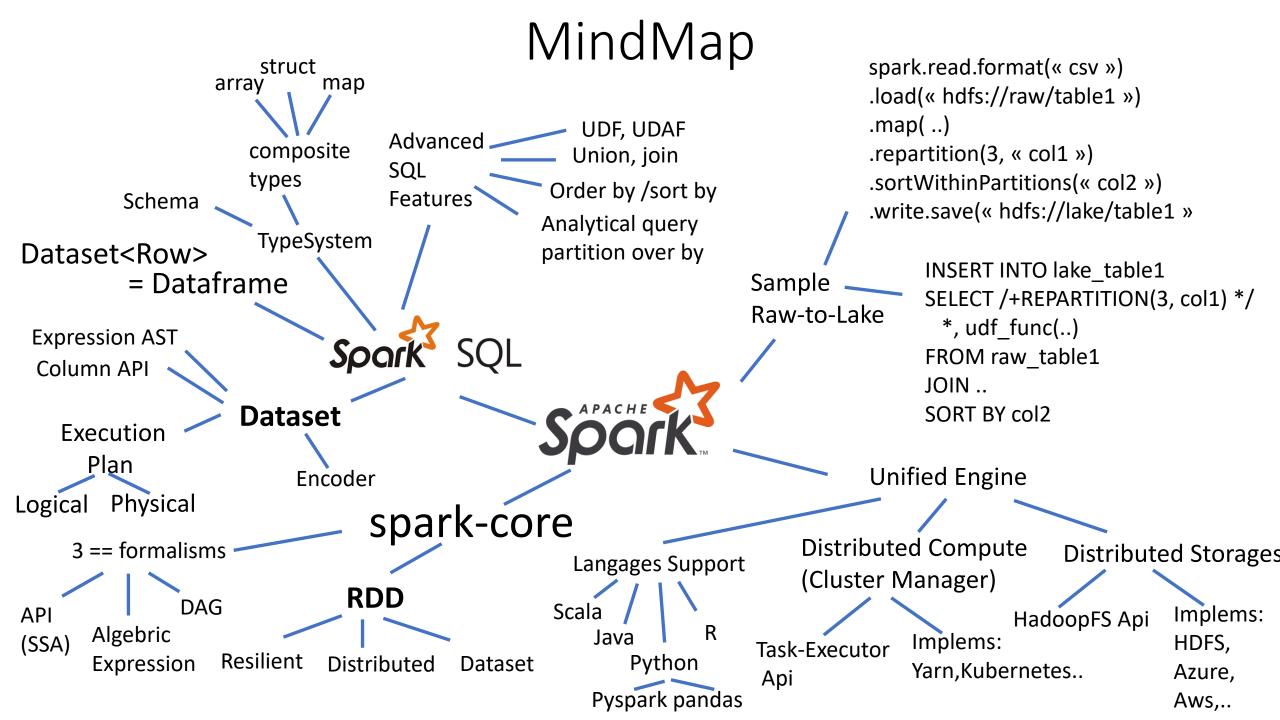












#### Next..

Lesson 1: introduction to BigData / Hadoop cluster / Spark



Lesson 2: this document: spark-core / spark-sql

Lesson 3: cluster management, tuning args, deploy-mode

Lesson 4: Advanced Spark: UI, parquet PPD, java api, spark-streaming, ...