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5/5/20

Assignment: Evaluate the effectiveness of four different hash functions and find which works best for the corresponding data sets provided.

Collaboration Statement: I worked on this alone.

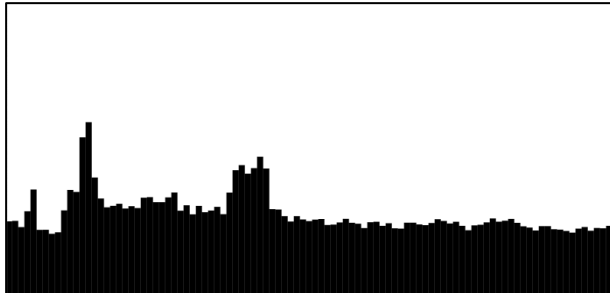
Note: I made no changes to HashEvaluation.java and therefore have not included it in this file

## **Words. Txt – case where all inputs are distinct**

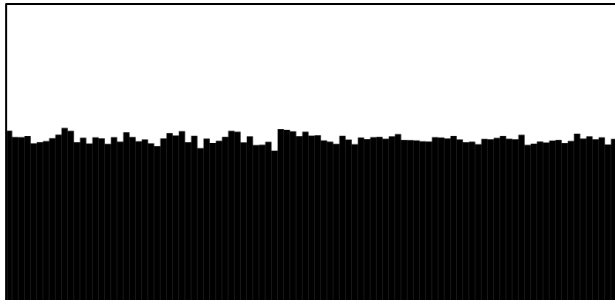
### ***M-Hash***

#### **Spread of Hash values (100 bins)**

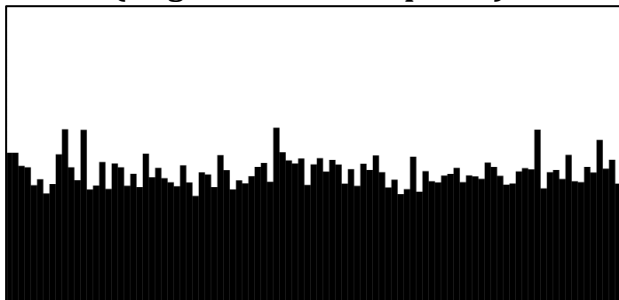
**R = 3 (small number, prime)**



**R = 31 (middle range number, prime)**



**R = 140 (large number, non prime)**



Amongst the 3 multiplier values for R, R = 31 performs the best with the most even distribution of hash values.

### "Time to collision" for 100 experiments

**R = 3**

N	min	mean	stdev	max
100	210.000	2764.720	1323.659	6860.000

**R = 31 (Time to collision is vastly higher)**

N	min	mean	stdev	max
100	10071.000	<b>77987.880</b>	37797.191	168031.000

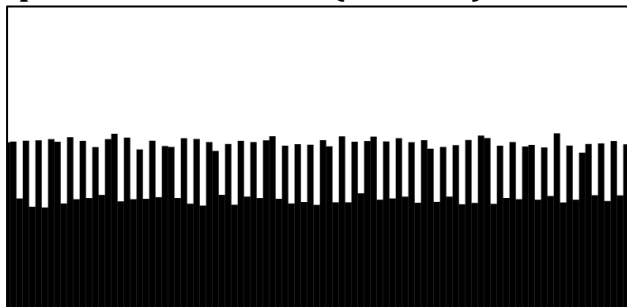
**R = 140**

N	min	mean	stdev	max
100	1544.000	8771.030	4632.829	21870.000

Based on the above data, M hash using a prime number within the approximate range of  $20 < x < 80$ , should be optimal in delivering both a high time to collision as well.

### ***P-Hash***

#### **Spread of Hash values (100 bins)**



P-hash has an uneven distribution of hash values even where every input string is unique, making it unsuitable to be used. There is a pattern observed of certain equally spaced ranges where hash values occur more frequently.

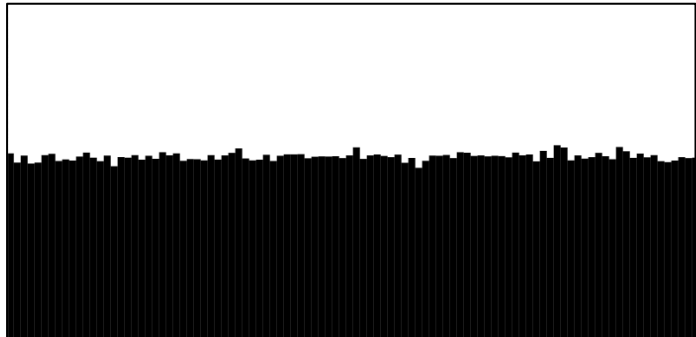
### "Time to collision" for 100 experiments

N	min	mean	stdev	max
100	2.000	21.300	9.485	45.000

Compare to M hash, the time to collision is far smaller, implying that collisions appear very frequently even when the input is distinct. As a result, this version of P hash is not viable in this situation

***Q-Hash (P-Hash Variant)***

**Spread of Hash values (100 bins)**



For Q-Hash, the distribution is far more uniform than its P-Hash counterpart, comparable to that of M-hash with an optimal choosing of the multiplier R.

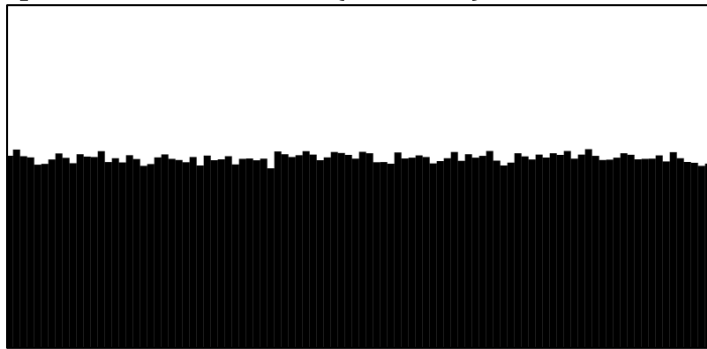
**“Time to collision” for 100 experiments**

N	min	mean	stdev	max
<hr/>				
100	8756.000	94533.560	45049.104	201674.000

Furthermore, the time to collision for this version of P-Hash is better than the earlier data from M-Hash.

***J-Hash***

**Spread of Hash values (100 bins)**



J-Hash also appears to have a fairly uniform distribution of hash values comparable to that of M-Hash and Q-Hash.

**“Time to collision” for 100 experiments**

N	min	mean	stdev	max
<hr/>				
100	2604.000	17594.270	9455.062	46689.000

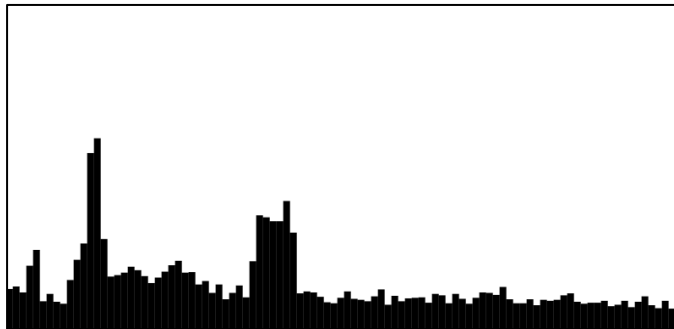
Despite its visible uniformity, J-Hash still falls short compared to Q-Hash with respect to its time to collision Statistic. **As a result, Q-hash seems to be the best option for use in words.txt due to an equal distribution of hash values as well as the highest time to collision in a set of fully distinct inputs.**

### **Tale.txt – case for distinct inputs but also some frequently repeated inputs**

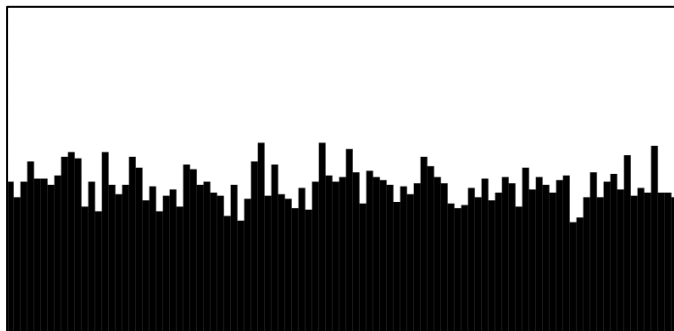
*M-Hash*

**Spread of Hash values (100 bins)**

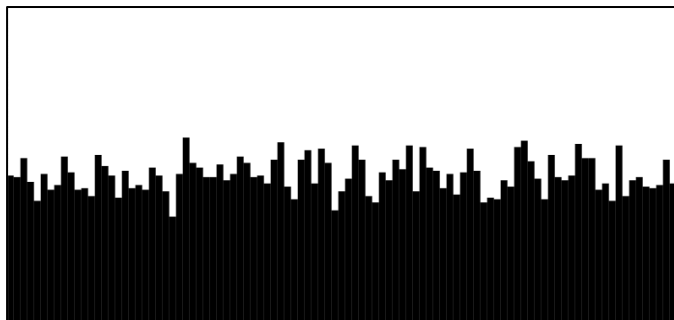
**R = 3**



**R = 31**



**R = 140**



*Note: Apart from the images above, I also checked R values of 7,19,67,91. All led to similar distributed hash values as the R = 31 and R = 140 case.*

Reviewing the 3 cases above, M-hash seems to work best when R = 31 as the hash values are the most uniform of the three. However, with the introduction of

the file "tale.txt" which has regular sentences compared to a dictionary of distinct words, the values are no longer as well distributed as before.

### **"Time to collision" for 100 experiments**

**R = 3**

N	min	mean	stdev	max
100	24.000	111.480	56.987	298.000

**R = 31**

N	min	mean	stdev	max
100	93.000	2549.440	1550.258	8312.000

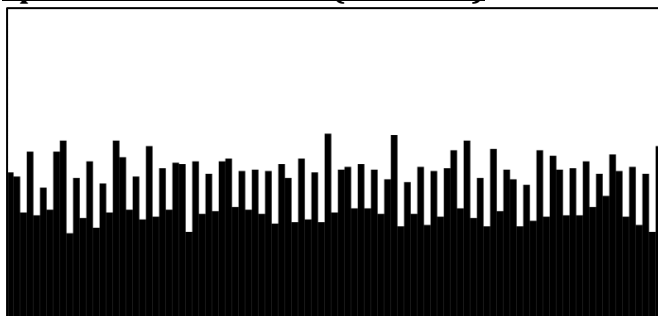
**R = 140**

N	min	mean	stdev	max
100	3833.000	12719.180	3645.146	19191.000

In this scenario, our arbitrary case of R = 140 has the best time till collision. It is possible that in this case, words that are used in regular sentences might be smaller and simpler than those in the dictionary and hash to similar values . Thus using a large value of R may spread them out.

### ***P-Hash***

#### **Spread of Hash values (100 bins)**



Again, this version of P hash has an uneven distribution of hash values. With a pattern observed of certain equally spaced ranges of hash values occurring more frequently.

### **"Time to collision" for 100 experiments**

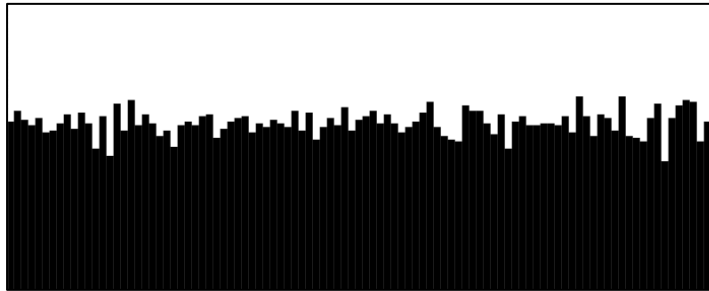
N

min	mean	stdev	max	
100	2.000	16.780	7.618	39.000

The time to collision for this hash function is still very small even though there are still a significant number of distinct inputs.

### ***Q-Hash***

#### **Spread of Hash values (100 bins)**



For Q-hash, the hash values seem relatively well distributed, more so than any of the M-hash histograms.

#### **"Time to collision" for 100 experiments**

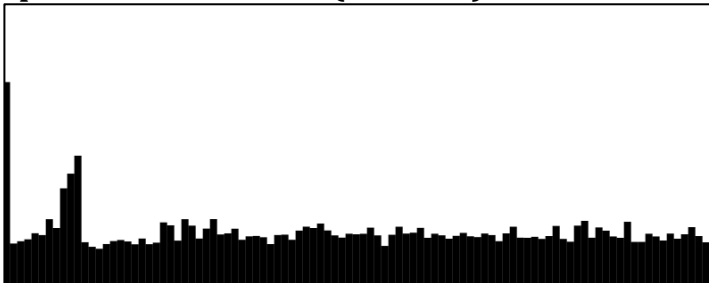
N	min	mean	stdev	max
100	19696.000	19696.000	0.000	19696.000

The results for time to collision in the case of Q-hash for "tale.txt" suggests that there are no collisions whatsoever.

Using the Zipf and frequency counter classes from Lab 8&9, I verified that the file tale.txt has 19695 distinct words. Q-hash spreads these values over the  $2^{31}$  range for any positive integer value (primitive type). Therefore it is definitely viable that Q-hash in this case prevents the case of collisions completely.

### ***J-Hash***

#### **Spread of Hash values (100 bins)**



For J-hash, the hash values seem to be very well distributed amongst the larger hash values values, but poorly distributed for smaller hash values.

#### **"Time to collision" for 100 experiments**

N	min	mean	stdev	max
100	1045.000	6495.850	3381.027	15475.000

Compared to the case of  $R = 140$ , J hash a similar standard deviation, but it's mean time to collision value is almost half as small, suggesting that both M-hash and Q-hash are superior.

As a result, Q-hash seems to be the best option as it not only distributes the hashes regularly but also ensures no collisions whatsoever in a case where the number of distinct entries is much smaller than  $2^{31}$ .

### Dna3. Txt – case of very large strings composed of repeating characters

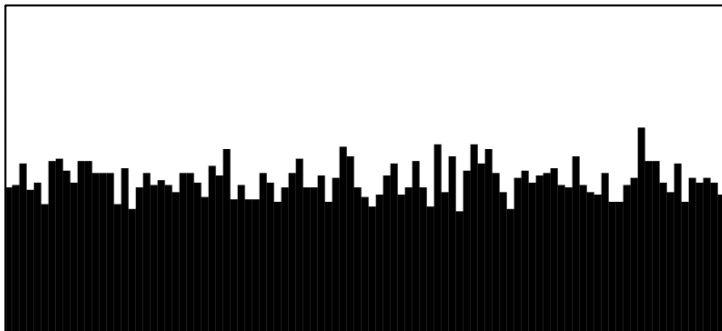
*M-Hash*

Spread of Hash values (100 bins)

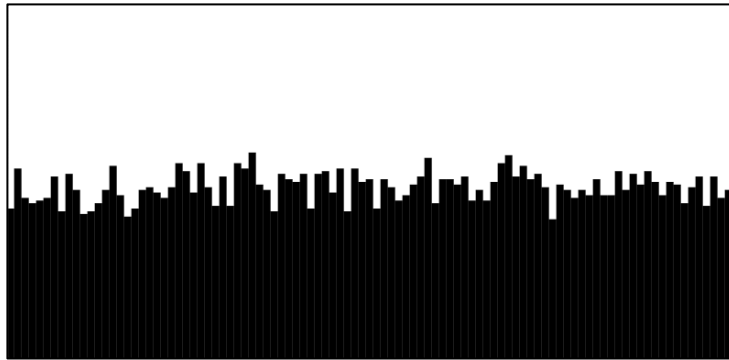
$R = 3$



$R = 31$



$R = 140$



Reviewing the 3 cases above, M-hash seems to work well for  $R = 31$  and  $R = 140$  but very poorly for  $R = 3$ , with a large number of majority of hash values within a fairly small range.

### **“Time to collision” for 100 experiments**

**$R = 3$**

N	min	mean	stdev	max
100	310.000	3056.370	1563.117	8349.000

**$R = 31$  (No Collision)**

N	min	mean	stdev	max
100	13086.000	13086.000	0.000	13086.000

**$R = 140$**

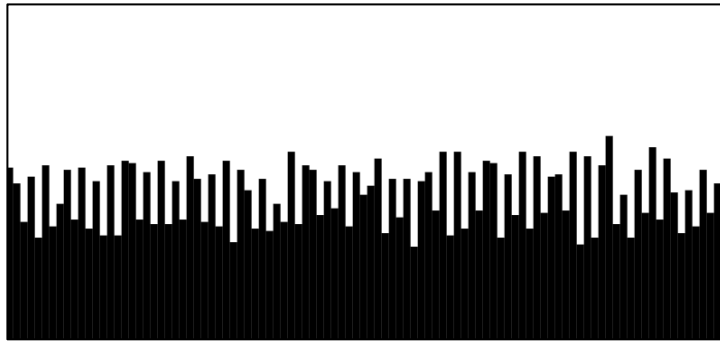
N	min	mean	stdev	max
100	244.000	2440.960	1274.995	6951.000

Reviewing the 3 cases above, M-hash seems to work best for  $R = 31$  as there seems to be a case of no collisions across 100 experiments. Upon verifying with the frequency counter class from Lab 8 & 9 indeed has 13085 distinct DNA sequences present in the text file. Furthermore,  $R = 3$  and  $R = 140$  have similar time to collision values. Given the huge lengths of each DNA sequence, it is possible that the hash values of with  $R = 140$  exceed the range of  $2^{31}$  and as a result of the “& 0x7FFFFFFF” operation, we might have skewed results in this situation.

### ***P-Hash***

#### **Spread of Hash values (100 bins)**





Again, this version of P hash has an uneven distribution of hash values. With a pattern observed of certain equally spaced ranges of hash values occurring more frequently.

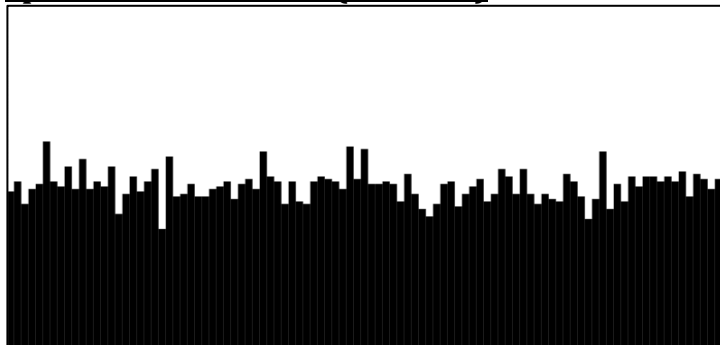
#### **"Time to collision" for 100 experiments**

N	min	mean	stdev	max
100	3.000	19.130	9.056	46.000

The time to collision for this hash function is still very small compared to those from M-hash, making it ineffective.

#### ***Q-Hash***

##### **Spread of Hash values (100 bins)**



Like the case for  $R = 31$  in M-Hash, Q-hash has a well distributed graph. However, on closer visual inspection it seems (to me) like the M-Hash function still has a slightly more even distribution.

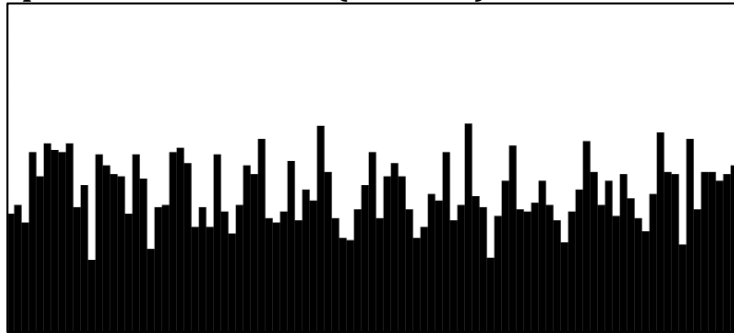
#### **"Time to collision" for 100 experiments**

N	min	mean	stdev	max
100	13086.000	13086.000	0.000	13086.000

Similar to the  $R=31$  case, Q-Hash exhibits no collisions.

### ***J-Hash***

#### **Spread of Hash values (100 bins)**



Compared to M-Hash and Q-Hash, the distribution of hash values for J-hash is far more uneven, making it less effective than the other two hashes.

#### **“Time to collision” for 100 experiments**

N	min	mean	stdev	max
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100	125.000	1082.840	591.200	2729.000

J hash has a much smaller value for time to collision as well, even less so than the M-hash case for  $R = 3$  which had a very uneven distribution in the centre. This suggests it is less effective for this case.

**Reviewing all the cases above, both Q-hash and M-Hash for  $R = 31$  seem to work very well, with a good distribution of Hash values as well as no collisions. However, the M-hash seems to have a slightly more even distribution based on a visual comparison and is also less computationally intensive than Q-Hash, performing noticeably faster upon multiple attempts. Therefore for this data set, I would recommend using M-hash with  $R = 31$ .**