Routing in The Dark Scalable Searches in Dark P2P Networks

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The Freenet Project

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- But when individual users come under attack, decentralisation is not enough.
- Future networks may need to limit connections to trusted friends.
- The big question is: Can such networks be useful?

Overview of "Peer to Peer" networks

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- Information is spread across many interconnected computers
- Users want to find information
- Some are centralised (eg. Napster), some are semi-centralised (eg. Kazaa), others are distributed (eg. Freenet)

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- Advantage: Globally scalable with the right routing algorithm
- Disadvantage: Vulnerable to "harvesting", ie. people you don't know can easily discover whether you are part of the network

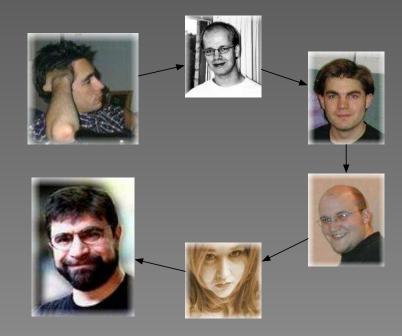
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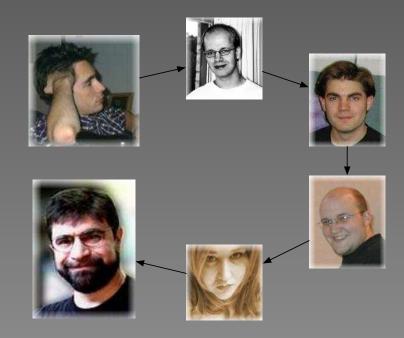
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- Examples: Waste
- Advantage: Only your trusted friends know you are part of the network
- Disadvantage: Networks are disconnected and small, they typically don't scale well

The Small World Phenomenon



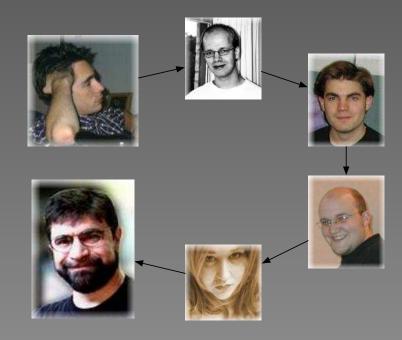
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- Short paths may exist but they may not be easy to find

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- This is called "Greedy Routing"
- Freenet and "Distributed Hash Tables" rely on this principal to find data in a scalable decentralised manner

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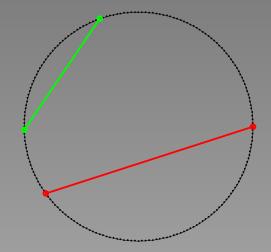
- A Darknet is, essentially, a social network of peoples trusted relationships.
- If people can route in a social network, then it should be possible for computers.
- Jon Kleinberg explained in 2000 how small world networks can be navigable.

Kleinberg's Result

The possibility of routing efficiently depends on the proportion of connections that have different lengths with respect to the "position" of the nodes.

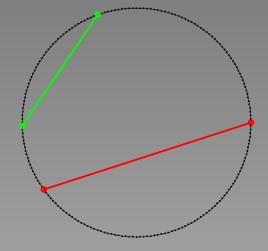
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- If the positions are in a ring, the proportion of connections with a certain length should be inverse to the length:



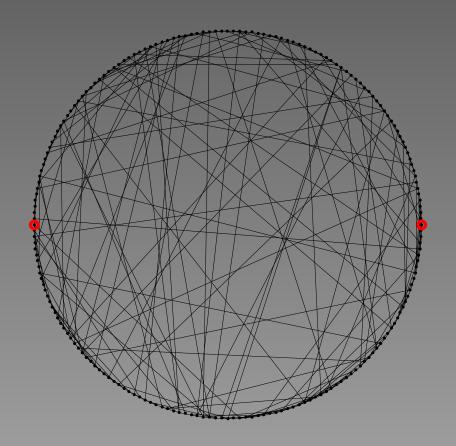
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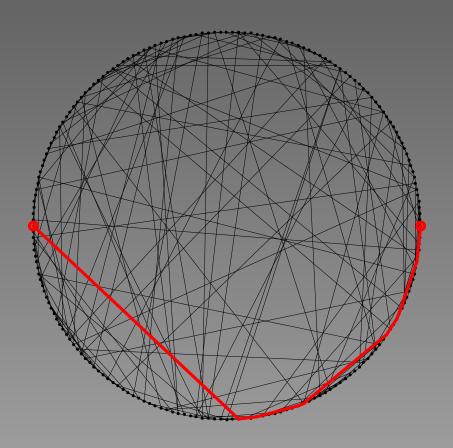


• In this case a simple greedy routing algorithm performs in $O(\log^2 n)$ steps.

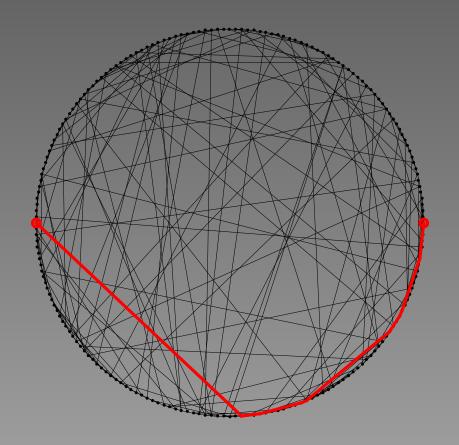
Kleinbergs Result, cont.



Kleinbergs Result, cont.



Kleinbergs Result, cont.



But in a social network, how do we see if one person is closer to the destination than another?

Application, cont.

Is Alice closer to Harry than Bob?

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- One cannot, in practice, expect a computer to route based on such things.
- Instead, we let the network tell us!

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- Then greedy route with respect to these numerical identities.

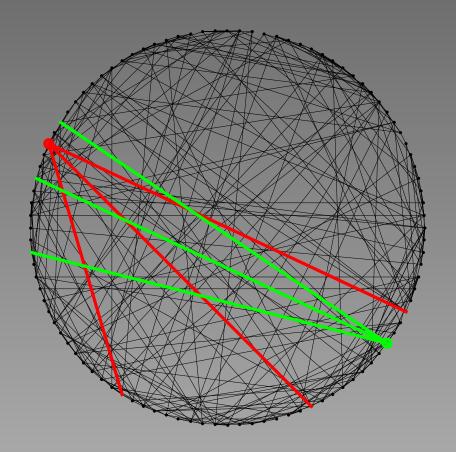
The Method

When nodes join the network, they choose a position on the circle randomly.

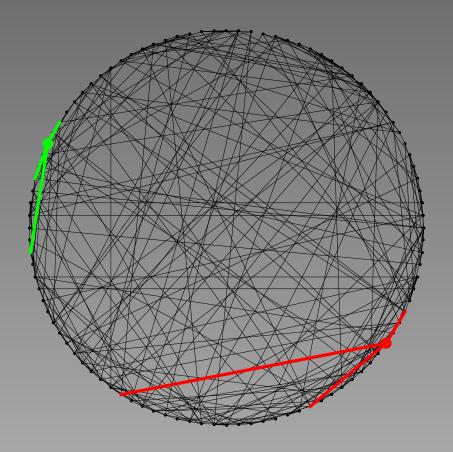
The Method

- When nodes join the network, they choose a position on the circle randomly.
- They then switch positions with other nodes, so as to minimize the product of the edge distances.

An advantageous switch of position:



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- Because this is an ongoing process as the network grows (and shrinks) it will be difficult to keep permanent positions.

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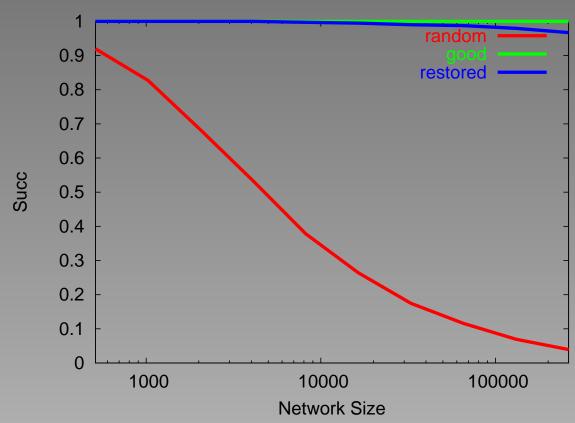
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- Random walk search: "random".
- Greedy routing in Kleinberg's model with identities as when it was constructed: "good".
- Greedy routing in Kleinberg's model with identities assigned according to our algorithm (2000 iterations per node): "restored".

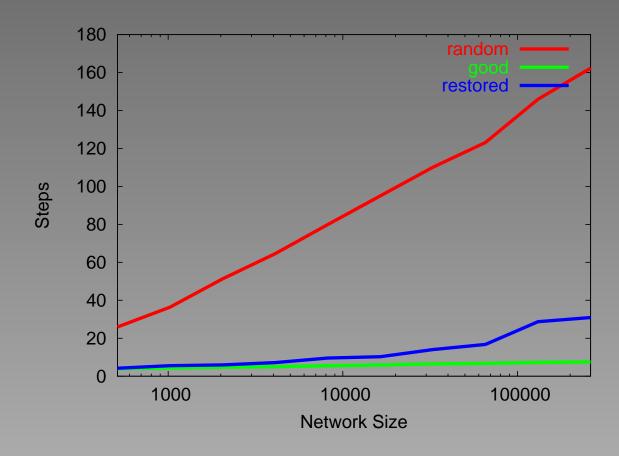
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- We borrowed some data from orkut.com. 2196 people were spidered, starting with Ian.

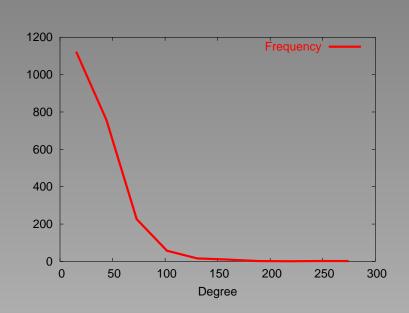


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The degree distribution is approximately Power-Law:



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	Success Rate	Mean Steps
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Our Algorithm		

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Random Search Our Algorithm

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	Success Rate	Mean Steps
Random Search	0.72	43.85
Our Algorithm	0.97	7.714

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Our algorithm takes advantage of there being people who have many connections, but it does not depend on them.

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Practical Concerns

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- Key concerns:
 - Preventing malicious behaviour
 - Ensuring ease of use
 - Storing data

Preventing Malicious Behaviour

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- Selection of identity to attract certain data
- Manipulation of other node's identities

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- What about NATs and firewalls
 - Could use UDP hole- punching (as used by Dijjer, Skype)
 - Would require third- party for negotiation

We believe very strongly that building a navigable, scalable Darknet is possible. *And we intend to do it!*

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 - Can other models work better?
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 - It needs to be tested on more data.

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People who are interested can join the discussion at http://freenetproject.org/.

Long Live the Darknet!

