Nonlocal games and the Grothendieck-Tsirelson inequality

Introduction to quantum information theory

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   - A specific example
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   - Local correlation matrices
   - Quantum correlation matrices
   - The relations between quantum correlation and local correlation matrices
   - Motivation: The Grothendieck-Tsirelson Theorem
   - Grothendieck's Inequality
   - Tsirelson's Theorem
   - Gorthendieck-Tsirelson Theorem

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- We should be familiar with the *Dirac notation*
  - $|\psi\rangle$ is a vector in $\mathbb{C}^n$ and $\langle\psi|$ is its conjugate transpose

Quantum systems

- A quantum system is a portion of the whole universe. For example a set electrons.
- A quantum system X is associated with a copy of $\mathbb{C}^k$
- It may consist of subsystems $X_1, \dots, X_N$ each of which is associated with a copy of $\mathbb{C}^{n_i}$. In this case $k = n_1 \dots n_N$

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- This can be achieved:

Definition

Measurement

- We define a measurement by a set of psd matrices $\{F^a\}_{a\in\mathcal{A}} \subseteq \mathbb{C}^{n\times n}$ that sum up to the identity matrix, i.e. $\sum_{a\in\mathcal{A}} F^a = I$
- The outcome of a measurement is a random variable χ with probability distribution: $\mathbb{P}[\chi = a] = \mathrm{Tr}(\rho F^a)$
- A projective measurement is defined by psd matrices that satisfy $F^a F^b = \delta_{ab} F^a \ \forall a, b \in \mathcal{A}$

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- if we consider projective measurements we have

$$(F^+ - F^-)^2 = \underbrace{{F^+}^2}_{=F^+} - \underbrace{F^+F^-}_{\delta_{+-}=0} + \underbrace{{F^-}^2}_{F^-} = F^+ + F^- = I$$

- i.e. a $\{-1,1\}$-valued observable is both unitary an Hermitian

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The joint probability distribution of the N measurement outcomes $\chi_1, \dots, \chi_N$ is

$$\mathbb{P}\left[\chi_1 = a_1, \chi_2 = a_2, \dots, \chi_N = a_N\right] = \mathrm{Tr}(\rho F_1^{a_1} \otimes \cdots \otimes F_N^{a_N})$$

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Calculation

$$\begin{aligned}
\mathrm{Tr}(|\psi\rangle\langle\psi|F^a \otimes G^b) &= \langle\psi|F^a \otimes G^b|\psi\rangle \\
&= (\langle\psi_A| \otimes \langle\psi_B|)(F^a \otimes G^b)(|\psi_A\rangle \otimes |\psi_B\rangle) \\
&= ((\langle\psi_A|F^a) \otimes (\langle\psi_B|G^b))(|\psi_A\rangle \otimes |\psi_B\rangle) \\
&= \langle\psi_A|F^a|\psi_A\rangle \otimes \langle\psi_B|G^b|\psi_B\rangle \\
&= \langle\psi_A|F^a|\psi_A\rangle\langle\psi_B|G^b|\psi_B\rangle
\end{aligned}$$

This is equal to the product of the probabilities of Alice measuring a and Bob measuring b, i.e. the outcome do not correlate.

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- Alice sends answer a and Bob sends answer b back to the referee, who then decides whether both win or both lose

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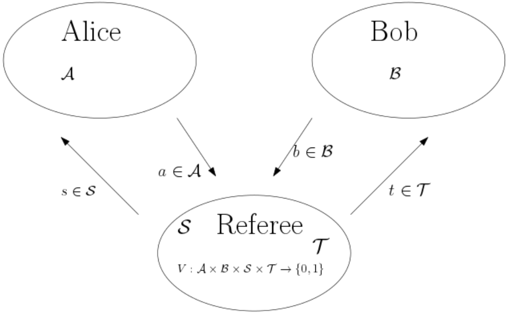
- A map $V : \mathcal{S} \times \mathcal{T} \times \mathcal{A} \times \mathcal{B} \to \{0, 1\}$
- They win if $V(s, t, a, b) = 1$ and lose otherwise



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- The winning probability then is:

$$\mathbb{E}_{(s,t)\sim\pi}\left[V(a(s), b(t), s, t)\right]$$

Quantum case

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- If the state is entangled measurements might give correlated measurement outcomes
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- Answering according to measurement outcomes could increase winning probability

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- They send their measurement outcome as their answer to the referee
- Their winning probability is:

$$\mathbb{E}_{(s,t)\sim\pi}\left[\sum_{a\in\mathcal{A}}\sum_{b\in\mathcal{B}} \mathrm{Tr}(\rho F_s^a \otimes G_t^b) V(a,b,s,t)\right]$$

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- A truth table for $a \oplus b$ looks like this

$\oplus$	0	1
0	0	1
1	1	0

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- Since $\frac{1}{2} + \gamma - (1 - \frac{1}{2} - \gamma) = 2\gamma$
- The violation ratio is defined as $\frac{\beta^*(G)}{\beta(G)}$

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- Let $a : \mathcal{S} \to \{0, 1\}$ and $b : \mathcal{T} \to \{0, 1\}$ be classical strategies and π the probability distribution the referee uses to pick s, t
- The bias is given by the probability under π that $a(s) \oplus b(t) = f(s, t)$ minus the probability under π that $a(s) \oplus b(t) \neq f(s, t)$

This means the bias can be written as:

$$\mathbb{E}_{(s,t)\sim\pi}\left[(-1)^{[a(s)\oplus b(t)=f(s,t)]}\right] =$$

$$= \mathbb{E}_{(s,t)\sim\pi}\left[(-1)^{a(s)\oplus b(t)+f(s,t)}\right] =$$

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Define the sign matrix $\Sigma_{s,t} = (-1)^{f(s,t)}$ and functions $\chi(s) = (-1)^{a(s)}$ and $\psi(t) = (-1)^{b(t)}$ The expected value becomes

$$\mathbb{E}_{(s,t)\sim\pi}\left[\chi(s)\psi(t)\Sigma_{st}\right]$$

Recap

- The outcomes in an XOR game are $\{0, 1\}$
- Alice and Bob have measurements $\{F_s^0, F_s^1\}$ and $\{G_t^0, G_t^1\}$ and share an entangled state
- The probability of Alice and Bob answering with a, b upon receiving s, t respectively is $\langle\psi|F_s^a \otimes G_t^b|\psi\rangle$

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Lets calculate the expected value of $(-1)^{a\oplus b}$

$$\begin{aligned}(1) \cdot \mathbb{P}[a = b] + (-1) \cdot \mathbb{P}[a \neq b] &= \\ = \langle\psi|F_s^0 \otimes G_t^0|\psi\rangle &+ \langle\psi|F_s^1 \otimes G_t^1|\psi\rangle \\ -\langle\psi|F_s^1 \otimes G_t^0|\psi\rangle &- \langle\psi|F_s^0 \otimes G_t^1|\psi\rangle \\ = \langle\psi|(F_s^0 - F_s^1) &\otimes (G_t^0 - G_t^1)|\psi\rangle\end{aligned}$$

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- Which is, if considering the $\{-1, 1\}$ basis, the expected value above
- We are looking to maximize this quantity

Classical strategies

When using classical strategies this is

$$\max\{\mathbb{E}_{(s,t)\sim\pi}\left[\Sigma_{st}\chi(s)\psi(t)\right] : \chi : \mathcal{S} \to \{-1,1\},\\ \psi : \mathcal{T} \to \{-1,1\}\}$$

Entangled strategies

When using entangled strategies the winning probability might increase indefinitely with the dimensions, so we use the $\sup_{n\in\mathbb{N}}$

$$\sup_{n\in\mathbb{N}}\{\mathbb{E}_{(s,t)\sim\pi}\left[\Sigma_{st}\langle\psi|F_s\otimes G_t|\psi\rangle\right] : |\psi\rangle\in\mathbb{C}^n\otimes\mathbb{C}^n,$$

$$F_s, G_t \in O(\mathbb{C}^n)\}$$

The CHSH game

- The CHSH game (Clauser, Horner, Shimony, Holt) is a two player XOR game with $\mathscr{A} = \mathscr{B} = \mathscr{S} = \mathscr{T} = \{0, 1\}$ and π being the uniform distribution
- $f(s, t) = s \wedge t$, i.e. $f(1, 1) = 1$ and $f(0, 0) = f(0, 1) = f(1, 0) = 0$
- Alice and Bob can win $\frac{3}{4}$ of the games by using deterministic strategies $(0, 0), (1, 0)$ or $(0, 1)$

Quantum strategy

- Let Alice and Bob share an EPR state
- Define

$$X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$$

- $XY + YX = 0$ and $X^2 = Y^2 = I$
- For Alice define the observable for question 0 by $F_0 = X$ and for question 1 by $F_1 = Y$
- Bobs observables are going to be $G_0 = (X - Y)/\sqrt{2}$ for question 0 and $G_1 = (X + Y)/\sqrt{2}$ for question 1

The following auxiliary calculations will be helpful later:

$$\langle \mathrm{EPR}|X \otimes X|\mathrm{EPR}\rangle = \frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}0 & 0 & 0 & 1\\ 0 & 0 & 1 & 0\\ 0 & 1 & 0 & 0\\ 1 & 0 & 0 & 0\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix}$$

$$= \frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix} = \frac{2}{2} = 1$$

$$\langle \mathrm{EPR}| Y \otimes Y |\mathrm{EPR}\rangle = \frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}0 & 0 & 0 & -1\\ 0 & 0 & 1 & 0\\ 0 & 1 & 0 & 0\\ -1 & 0 & 0 & 0\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix}$$

$$= \frac{1}{2}\begin{pmatrix}-1 & 0 & 0 & -1\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix} = -1$$

$$\langle \mathrm{EPR}|X \otimes Y|\mathrm{EPR}\rangle = \frac{1}{2}\begin{pmatrix}1 & 0 & 0 & 1\end{pmatrix}\begin{pmatrix}0 & 0 & 0 & -i\\ 0 & 0 & i & 0\\ 0 & -i & 0 & 0\\ i & 0 & 0 & 0\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix}$$

$$= \frac{1}{2}\begin{pmatrix}i & 0 & 0 & -i\end{pmatrix}\begin{pmatrix}1\\0\\0\\1\end{pmatrix} = 0$$

$$\langle \mathrm{EPR}|Y \otimes X|\mathrm{EPR}\rangle = 0$$

Lets calculate the expected values of the sign $a \oplus b$:

$$\begin{aligned}\langle \mathrm{EPR}|F_0 \otimes G_0|\mathrm{EPR}\rangle &= \langle \mathrm{EPR}|X \otimes \frac{1}{\sqrt{2}}(X - Y)|\mathrm{EPR}\rangle \\ &= \langle \mathrm{EPR}|X \otimes \frac{1}{\sqrt{2}}X|\mathrm{EPR}\rangle - \langle \mathrm{EPR}|X \otimes \frac{1}{\sqrt{2}}Y|\mathrm{EPR}\rangle \\ &= \frac{1}{\sqrt{2}} - 0 = \frac{1}{\sqrt{2}}\end{aligned}$$

$$
\begin{aligned}
\langle \mathrm{EPR} | F_1 \otimes G_1 | \mathrm{EPR} \rangle &= \langle \mathrm{EPR} | Y \otimes \frac{1}{\sqrt{2}}(X + Y) | \mathrm{EPR} \rangle \\
&= \langle \mathrm{EPR} | Y \otimes \frac{1}{\sqrt{2}} X | \mathrm{EPR} \rangle + \langle \mathrm{EPR} | Y \otimes \frac{1}{\sqrt{2}} Y | \mathrm{EPR} \rangle \\
&= 0 - \frac{1}{\sqrt{2}} = -\frac{1}{\sqrt{2}}
\end{aligned}
$$

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&= \frac{1}{\sqrt{2}} + 0 = \frac{1}{\sqrt{2}}
\end{aligned}
$$

$$\begin{aligned}\langle \mathrm{EPR}|F_1 \otimes G_0|\mathrm{EPR}\rangle &= \langle \mathrm{EPR}|Y \otimes \frac{1}{\sqrt{2}}(X - Y)|\mathrm{EPR}\rangle \\ &= \langle \mathrm{EPR}|Y \otimes \frac{1}{\sqrt{2}}X|\mathrm{EPR}\rangle - \langle \mathrm{EPR}|Y \otimes \frac{1}{\sqrt{2}}Y|\mathrm{EPR}\rangle \\ &= 0 - (-\frac{1}{\sqrt{2}}) = \frac{1}{\sqrt{2}}\end{aligned}$$

Thus, we have

$$\langle \mathrm{EPR}|F_s \otimes G_t|\mathrm{EPR}\rangle = \begin{cases} \frac{1}{\sqrt{2}}, (0,0), (1,0), (0,1) \\ -\frac{1}{\sqrt{2}}, (1.1) \end{cases}$$

which is equivalent to

$$\langle \mathrm{EPR}|F_s \otimes G_t|\mathrm{EPR}\rangle = \frac{(-1)^{s \wedge t}}{\sqrt{2}}, s, t \in \{0, 1\}$$

The bias of the entangled strategy equals

$$\mathbb{E}_{(s,t)\sim\pi}\left[\Sigma_{s,t}\langle\psi|F_s\otimes G_t|\psi\rangle\right] =$$

$$= \frac{1}{4}\sum_{s,t=0}^{1}(-1)^{s\wedge t}\langle \mathrm{EPR}|F_s\otimes G_t|\mathrm{EPR}\rangle$$

$$= \frac{1}{4}\cdot\frac{4}{\sqrt{2}} = \frac{1}{\sqrt{2}}$$

The bias is $\frac{1}{\sqrt{2}}$ from which follows that the winning probability is by definition:

$$\frac{1}{2}+\frac{1}{2}\cdot\frac{1}{\sqrt{2}} = \cos(\pi/8) \approx 0.85\ldots$$

Outline

1. Introduction
   - Basics
   - Nonlocal games
   - A special case of nonlocal games
   - A specific example

2. Local and quantum correlation matrices
   - Local correlation matrices
   - Quantum correlation matrices
   - The relations between quantum correlation and local correlation matrices
   - Motivation: The Grothendieck-Tsirelson Theorem
   - Grothendieck's Inequality
   - Tsirelson's Theorem
   - Gorthendieck-Tsirelson Theorem

Local correlation matrices

- Deterministic strategies of Alice and Bob correspond to vectors $\xi \in \{-1,1\}^{\mathcal{S}}$, respectively $\eta \in \{-1,1\}^{\mathcal{T}}$

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- Their common answer is $\mathbb{E}[X_i Y_j]$
- This information can be encoded in an $\mathcal{S} \times \mathcal{T}$ matrix ◂ Motivation Quantum

Definition

Let $(X_i)_{1\leq i\leq m}$ and $(Y_j)_{1\leq j\leq n}$ be families of random variables on a common probability space such that $|X_i|, |Y_j| \leq 1$ almost surely. Then $A = (a_{ij})$ is the corresponding *classical (or local) correlation matrix* if

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$$\mathrm{LC}_{m,n} = \mathrm{conv}\{\xi\eta^T \,|\, \xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n\}$$

- No matter which probabilistic strategy there is a deterministic one which as at least as good as the one one chooses

$\mathrm{LC}_{m,n} \supset \mathrm{conv}\{\xi\eta^T \mid \xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n\}$.

- $\xi\eta^T \in LC_{m,n}$ for all $\xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n$ (Choose $X_i \equiv \xi_i,\ Y_j \equiv \eta_j$)

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$$\beta a_{ij}^{(0)} + (1-\beta) a_{ij}^{(1)} = \mathbb{E}[X_i Y_j]$$

for $\beta \in [0,1]$

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  for $\beta \in [0,1]$
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- Then

$$\begin{aligned} \mathbb{E}[X_i Y_j] &= \mathbb{E}[X_i^{(\alpha)} Y_j^{(\alpha)} \mathbb{1}_{\{\alpha=0\}}] + \mathbb{E}[X_i^{(\alpha)} Y_j^{(1)}] \mathbb{1}_{\{\alpha=1\}}] \\ &= \beta \mathbb{E}[X_i^{(0)} Y_j^{(0)}] + (1-\beta) \mathbb{E}[X_i^{(1)} Y_j^{(1)}] \end{aligned}$$

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- Let $a_{ij} = \mathbb{E}[X_i Y_j]$ for $\mathbb{R}$-valued random variables $(X_i), (Y_j)$ defined on a common probability space Ω with $|X_i|, |Y_j| \leq 1$ almost surely.

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- Define random variables $\lambda_\xi^{(X)} : \Omega^m \to [0,1]$ such that

$$X(\omega) = \sum_{\xi \in \{-1,1\}^m} \lambda_\xi^{(X)}(\omega)\xi$$

almost surely and $\sum_{\xi \in \{-1,1\}^m} \lambda_\xi^{(X)}(\omega) = 1$

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- Using the same decomposition for Y we obtain

$$\begin{aligned} a_{ij} = \mathbb{E}[X_i Y_j] &= \mathbb{E}\Big[\Big(\sum_{\xi \in \{-1,1\}^m} \lambda_\xi^{(X)} \xi_i\Big)\Big(\sum_{\eta \in \{-1,1\}^n} \lambda_\eta^{(Y)} \eta_j\Big)\Big] \\ &= \sum_{\xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n} \mathbb{E}\big[\lambda_\xi^{(X)} \lambda_\eta^{(Y)}\big] \xi_i \eta_j \\ &= \Big(\sum_{\xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n} \mathbb{E}[\lambda_\xi^{(X)}]\mathbb{E}[\lambda_\eta^{(Y)}]\Big) \xi_i \eta_j \end{aligned}$$

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- Due to $\sum_{\xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n} \mathbb{E}[\lambda_\xi^{(X)}]\mathbb{E}[\lambda_\eta^{(Y)}] = 1$ the matrix (a_{ij}) is a convex combination of $\xi\eta^T$, $\xi \in \{-1,1\}^m, \eta \in \{-1,1\}^n$

□

Quantum correlation matrices

- Alice and Bob share a common state ρ and get inputs $i \in \mathcal{S}$, $j \in \mathcal{T}$

Motivation Locality

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- Alice and Bob share a common state ρ and get inputs $i \in \mathcal{S}$, $j \in \mathcal{T}$
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- Again we can encode this information in a matrix

Motivation Locality

Definition

Let $(X_i)_{1\leq i\leq m}$ and $(Y_j)_{1\leq j\leq n}$ be self-adjoint operators on $\mathbb{C}^{d_1}$, respectively $\mathbb{C}^{d_2}$ for some positive integers d_1, d_2, satisfying $\|X_i\|_\infty, \|Y_j\|_\infty \leq 1$. $A = (a_{ij})$ is called *quantum correlation matrix* if there exists a state on $\mathbb{C}^{d_1} \otimes \mathbb{C}^{d_2})$ such that

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Lemma

$$\mathrm{QC}_{m,n} = \{(\langle x_i, y_j\rangle)_{1\leq 1\leq m, 1\leq j\leq n} \,|\, x_i, y_j \in \mathbb{R}^{\min\{m,n\}}, |x_i| \leq 1, |y_j| \leq 1\},$$

$\mathrm{QC}_{m,n} \subset \{(\langle x_i, y_j \rangle)_{1\leq 1\leq m, 1\leq j\leq n} \mid x_i, y_j \in \mathbb{R}^{\min\{m,n\}}, |x_i| \leq 1, |y_j| \leq 1\}$

- $a_{ij} = \operatorname{Tr} \rho X_i \otimes Y_j$, sate ρ on a Hilbert space $\mathcal{H} = \mathbb{C}^{d_1} \otimes \mathbb{C}^{d_2}$ and Hermitian operators $(X_i)_{1\geq m}$, $(Y_j)_{1\geq n}$ on $\mathbb{C}^{d_1}$, respectively $\mathbb{C}^{d_2}$ satisfying $\|X_i\|_\infty, \|Y_j\|_\infty \leq 1$

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- Define a positive semidefinite symmetric bilinear form on the space of Hermitian operators on $\mathcal{H}$ by $\beta : \mathcal{H} \times \mathcal{H} \to \mathbb{R}$ where $\beta(S, T) = \mathrm{Re}(\operatorname{Tr} \rho ST)$.

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Proposition

For all $n \ge 1$ there is a subspace of the $2^n \times 2^n$ Hermitian matrices where every vector is the multiple of a unitary matrix.

- The proof is based on $n-$fold tensor products of the Pauli matrices which are the three matrices

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

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Proof.

- Define

$$U_i = I^{\otimes(i-1)} \otimes X \otimes Y^{\otimes(n-i)},$$
$$U_{n+i} = I^{\otimes(i-1)} \otimes Z \otimes Y^{\otimes(n-i)},\ i = 1, \dots n$$

- U_i's are anti-commuting traceless Hermitian unitaries, i.e. $U_iU_j = -U_jU_i$ for $i \neq j$ and $U_i^2 = I$
- For $X = \sum_{i=1}^{2n} \xi_i U_i$, $Y = \sum_{i=1}^{2n} = \eta_i U_i$ we can calculate

$$\begin{aligned} XY &= \sum_{i=1}^{2n} \xi_i \eta_i I + \sum_{1 \leq i,j \leq \leq 2n} \xi_i \eta_j U_i U_j \\ &= \sum_{i=1}^{2n} \xi_i \eta_i I + \sum_{1 \leq i < j \leq \leq 2n} \xi_i \eta_j U_i U_j - \sum_{1 \leq i < j \leq \leq 2n} U_i U_j = \sum_{i=1}^{2n} \xi_i \eta_i I \\ &= \langle \xi, \eta \rangle I. \end{aligned}$$

- The result follows by setting $X = Y$.

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- Let $|\phi\rangle = \frac{1}{\sqrt{d}} \sum_{i=1}^{d} |ii\rangle$ and $\rho = |\phi\rangle\langle\rho|$. Note that we can write ρ as

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- Then

$$\begin{aligned}\mathrm{Tr}\left(\rho X_i \otimes Y_j\right) &= \frac{1}{d} \sum_{1\leq k,l\leq d} \mathrm{Tr}\left(|k\rangle\langle l| X_i \otimes |k\rangle\langle l| Y_j\right) \\ &= \frac{1}{d} \sum_{1\leq k,l\leq d} \mathrm{Tr}\left(|k\rangle\langle l| X_i\right) \mathrm{Tr}\left(|k\rangle\langle l| Y_j\right) \\ &= \frac{1}{d} \mathrm{Tr}\, X_i Y_j^T = \langle x_i, y_j\rangle.\end{aligned}$$

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- Right dimension is obtained by projecting $(\tilde{x}_i)_{1 \leq i \leq m}$, $(\tilde{y}_j)_{1 \leq i \leq n}$ on $\mathrm{span}\{x_1, \ldots, x_m\}$ or $\mathrm{span}\{y_1, ..., y_n\}$, as in the proof before.

The relations between quantum correlation and local correlation matrices

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$$\mathrm{LC}_{2,2} = \mathrm{conv}\{\pm \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}, \pm \begin{pmatrix} -1 & -1 \\ 1 & 1 \end{pmatrix}, \pm \begin{pmatrix} -1 & 1 \\ -1 & 1 \end{pmatrix}, \pm \begin{pmatrix} -1 & 1 \\ 1 & -1 \end{pmatrix}\}.$$

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- W can also write it as in intersections of halfspaces:

$$\mathrm{LC}_{2,2} = \{A \in \mathbb{R}^{2\times 2} \mid -1 \le \operatorname{Tr} AM \le 1 \text{ for all } M \in \mathcal{K}\}, \tag{1}$$

where $\mathcal{K} = \{\frac{1}{2}\sigma(\begin{pmatrix}-1 & 1\\ 1 & 1\end{pmatrix}), \sigma(\begin{pmatrix}1 & 0\\ 0 & 0\end{pmatrix}) \mid \sigma \in \{\mathrm{id}, (1\ 2), (1\ 3), (1\ 4)\}\}$.

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- $(|x_1 + x_2| + |x_1 - x_2|)^2 \leq 4(|x_1|^2 + |x_2|^2)$

- $\max\{\mathrm{Tr}\,(AM) \mid A \in \mathrm{LC}_{2,2}\} = 2$
- Recall $QC_{2,2} = \{\langle x_i, y_j\rangle \mid x_1, x_2, y_1, y_2 \in \mathbb{R}^2, \; |x_i|, |y_j| le 1\}$
- $A \in \mathrm{QC}_{2,2}$ we obtain, by Cauchy-Schwarz and $|y_i| \leq 1$,

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- Bound is achieved by $A = \frac{1}{\sqrt{2}}\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$, induced by the vectors $x_1 = x_2 = \frac{1}{\sqrt{2}}(1,1)$ and $y_1 = y_2 = (1,0)$

Nonlocal games and the Grothendieck-Tsirelson inequality

Introduction to quantum information theory

Maxi Brandstetter, Felix Kirschner, Arne Heimendahl

University of Cologne

September 19, 2018

Theorem (Grothendieck-Tsirelson)

There exists an absolute constant $K \geq 1$ such that, for any positive integers m, n, the following three equivalent conditions hold:

(1) *We have the inclusion*

$$\mathrm{QC}_{m,n} \subset K\mathrm{LC}_{m,n}. \tag{2}$$

(2) *For any $M \in \mathbb{R}^{m\times n}$ and for any ρ, X_i, Y_j verifying the conditions of Definition 4.2.1 we have*

$$\sum_{i,j} M_{ij} \operatorname{Tr} \rho(X_i \otimes Y_j) \leq K \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \sum_{i,j} M_{ij}\xi_i\eta_j \tag{3}$$

$$\Leftrightarrow \qquad \operatorname{Tr} MA^\top \leq \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \operatorname{Tr} M(\xi\eta^\top)^\top. \tag{4}$$

(3) *For any $M \in \mathbb{R}^{m\times n}$ and for any (real) Hilbert space vectors x_i, y_j with $|x_i| \leq 1$, $|y_j| \leq 1$ we have*

$$\sum_{i,j} M_{i,j}\langle x_i, y_j\rangle \leq K \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \operatorname{Tr} \xi^\top M\eta. \tag{5}$$

Lemma (Grothendieck's identity)

Let $x, y \in \mathbb{R}^d$ be unit vectors. Let $r \in \mathbb{R}^d$ be a random unit vector chosen from $O(d)$-invariant probability distribution on the unit sphere. Then

i, $\mathbb{P}[\mathrm{sign}(\langle x, r\rangle) \neq \mathrm{sign}(\langle y, r\rangle)] = \frac{\arccos(\langle x,y\rangle)}{\pi}$

ii, $\mathbb{E}[\mathrm{sign}(\langle x, r\rangle)\,\mathrm{sign}(\langle y, r\rangle)] = \frac{2}{\pi}\arcsin(\langle x, y\rangle)$.

Proof.

- if x and y are linearly dependent, then

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- if x and y are linearly independent, then
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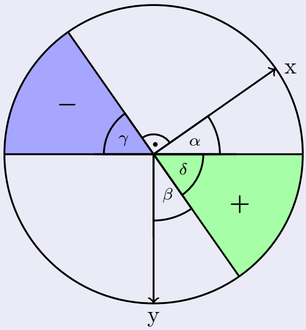
  - the normalized vector $n := s/\|s\|$ is uniformly distributed on the intersection of the unit sphere and $\text{span}\{x, y\}$ by the $O(d)$-invariance of the probability distribution

Proof (cont.).

Calculation of the probability that the signs of the scalar products $\langle x, n\rangle$ and $\langle y, n\rangle$ are unlike:



$$\mathbb{P}[\text{sign}(\langle x, n\rangle) \neq \text{sign}(\langle y, n\rangle)] = 2\,\frac{\frac{\pi}{2} + \alpha}{2\pi} = \frac{\arccos(\langle x, y\rangle)}{\pi}$$

Proof (cont.).

We conclude with the proof of the second part of Lemma 6:

$$
\begin{aligned}
&\mathbb{E}[\operatorname{sign}(\langle x, r\rangle)\operatorname{sign}(\langle y, r\rangle)] \\
&\qquad = 1 \cdot \mathbb{P}[\operatorname{sign}(\langle x, r\rangle) = \operatorname{sign}(\langle y, r\rangle)] - 1 \cdot \mathbb{P}[\operatorname{sign}(\langle x, r\rangle) \neq \operatorname{sign}(\langle y, r\rangle)] \\
&\qquad = 1 - 2\mathbb{P}[\operatorname{sign}(\langle x, r\rangle) \neq \operatorname{sign}(\langle y, r\rangle)] \\
&\qquad = 1 - 2\frac{\arccos(\langle x, y\rangle)}{\pi} \\
&\qquad = \frac{2}{\pi}\arcsin(\langle x, y\rangle),
\end{aligned}
$$

because $\arcsin(t) + \arccos(t) = \pi/2$. □

Lemma (Krivine's trick)

Let $x_1, \dots, x_m, y_1, \dots, y_n \in S^{m+n-1}$ *be given. Furthermore, let* $r \in \mathbb{R}^d$ *be a random unit vector chosen form the* $O(d)$*-invariant probability distribution on the unit sphere. Then there are* $x'_1, \dots, x'_m, y'_1, \dots, y'_n \in S^{m+n-1}$ *so that*

$$\mathbb{E}[\operatorname{sign}(\langle x'_i, r\rangle)\operatorname{sign}(\langle y'_j, r\rangle)] = \beta\langle x_i, y_j\rangle, \tag{6}$$

with $\beta = \frac{2}{\pi}\ln(1+\sqrt{2})$.

Definition (The k-th tensor product)

The *k-th tensor product* of $\mathbb{R}^n$ with orthonormal basis $e_1, \dots, e_n$ is denoted by $(\mathbb{R}^n)^{\otimes k}$ and it is a Euclidean vector space of dimension n^k with othonormal basis $e_{i_1} \otimes \cdots \otimes e_{i_k}$, $i_j \in \{1, \dots, n\}$. In particular

$$\langle e_{i_1} \otimes \cdots \otimes e_{i_k}, e_{j_1} \otimes \cdots \otimes e_{j_k} \rangle = \prod_{l=1}^{k} \langle e_{i_l}, e_{j_l} \rangle = \begin{cases} 1 & \text{, if } i_l = j_l \text{ for all } l = 1, \dots, n, \\ 0 & \text{, otherwise,} \end{cases} \tag{7}$$

and for $v \in \mathbb{R}^n$ with $v = v_1 e_1 + \cdots + v_n e_n$ we define $v^{\otimes k} \in (\mathbb{R}^n)^{\otimes k}$ by

$$\begin{aligned} v^{\otimes k} &:= (v_1 e_1 + \cdots + v_n e_n) \otimes \cdots \otimes (v_1 e_1 + \cdots + v_n e_n) \\ &= \sum_{i_1, \dots, i_k} v_{i_1} \cdots v_{i_k} e_{i_1} \otimes \cdots \otimes e_{i_k}. \end{aligned} \tag{8}$$

Thus, for $v, w \in \mathbb{R}^n$

$$\langle v^{\otimes k}, w^{\otimes k} \rangle = \langle v, w \rangle^k. \tag{9}$$

Proof of Krivine's trick.

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- idea: To find β, x_i', y_j' invert E:

$$E^{-1}(t) = \sin(\pi/2 \cdot t) = \sum_{k=0}^{\infty} \underbrace{\frac{(-1)^{2k+1}}{(2k+1)!}\left(\frac{\pi}{2}\right)^{2k+1}}_{=:g_{\mathbf{2}k+\mathbf{1}}} t^{2k+1}$$

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- define the infinite-dimensional Hilbert space

$$H = \bigoplus_{r=0}^{\infty} (\mathbb{R}^{m+n})^{\otimes 2k+1}. \tag{10}$$

Proof (cont.).

- define $\tilde{x}_i, \tilde{y}_j \in H$, $i = 1, \ldots, m, j = 1, \ldots, n$ componentwise:

$$(\tilde{x}_i)_k = \operatorname{sign}(g_{2k+1})\sqrt{|g_{2k+1}|\beta^{2k+1}}\, x_i^{\otimes 2k+1} \tag{11}$$

$$(\tilde{y}_j)_k = \sqrt{|g_{2k+1}|\beta^{2k+1}}\, y_j^{\otimes 2k+1} \tag{12}$$

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- then

$$\begin{aligned} \langle \tilde{x}_i, \tilde{y}_j \rangle &= \sum_{k=0}^{\infty} g_{2k+1}\beta^{2k+1}\langle x_i^{\otimes 2k+1}, y_j^{\otimes 2k+1}\rangle \\ &= \sum_{k=0}^{\infty} g_{2k+1}\beta^{2k+1}\langle x_i, y_j\rangle^{2k+1} \\ &= E^{-1}(\beta\langle x_i, y_j\rangle). \end{aligned}$$

Proof (cont.).

- hence, β is defined by the condition that the vectors $\tilde{x}_i, \ldots, \tilde{x}_m, \tilde{y}_1, \ldots, \tilde{y}_n$ are unit vectors:

$$1 = \langle \tilde{x}_i, \tilde{x}_i \rangle = \langle \tilde{y}_j, \tilde{y}_j \rangle = \sum_{k=0}^{\infty} \frac{1}{(2k+1)!} \left(\frac{\pi}{2}\right)^{2k+1} \beta^{2k+1} = \sinh(\frac{\pi}{2}\beta)$$

$$\Leftrightarrow \qquad \beta = \frac{2}{\pi} \operatorname{arcsinh}(1) = \frac{2}{\pi} \ln(1 + \sqrt(2))$$

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- problem: $\tilde{x}_1, \ldots, \tilde{x}_m, \tilde{y}_1, \ldots, \tilde{y}_n$ are infinite-dimensional
- solution: the positive definite and symmetric Gram matrix G

$$G = \begin{pmatrix} \langle \tilde{x}_1, \tilde{x}_1 \rangle & \cdots & \langle \tilde{x}_1, \tilde{x}_m \rangle & \langle \tilde{x}_1, \tilde{y}_1 \rangle & \cdots & \langle \tilde{x}_1, \tilde{y}_n \rangle \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ \langle \tilde{x}_m, \tilde{x}_1 \rangle & \cdots & \langle \tilde{x}_m, \tilde{x}_m \rangle & \langle \tilde{x}_m, \tilde{y}_1 \rangle & \cdots & \langle \tilde{x}_m, \tilde{y}_n \rangle \\ \langle \tilde{y}_1, \tilde{x}_1 \rangle & \cdots & \langle \tilde{y}_1, \tilde{x}_m \rangle & \langle \tilde{y}_1, \tilde{y}_1 \rangle & \cdots & \langle \tilde{y}_1, \tilde{y}_n \rangle \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ \langle \tilde{y}_n, \tilde{x}_1 \rangle & \cdots & \langle \tilde{y}_n, \tilde{x}_m \rangle & \langle \tilde{y}_n, \tilde{y}_1 \rangle & \cdots & \langle \tilde{y}_n, \tilde{y}_n \rangle \end{pmatrix} \tag{13}$$

Proof (cont.).

- due to the properties of G we can decompose G via a real orthogonal matrix Q with columns that are the eigenvectors of G and a real diagonal matrix Λ having the eigenvalues of G on the diagonal, thus

$$G = Q\Lambda Q^\top = (\underbrace{Q\Lambda^{1/2}}_{=:A})^\top (Q\Lambda^{1/2}) \tag{14}$$

□

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- the columns of A are the vectors $x'_1, \ldots, x'_m, y'_1, \ldots, y'_n \in S^{m+n-1}$ we are looking for

□

Definition

For $M \in \mathbb{R}^{m \times n}$ define the quadratic program

$$\|M\|_{\infty \to 1} = \max \left\{ \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} \xi_i \eta_j : \xi_i^2 = 1, i = 1, \ldots, m, \eta_j^2 = 1, j = 1, \ldots, n \right\}$$
$$= \max \left\{ \operatorname{Tr} M \eta \xi^\top : \xi \in \{-1, 1\}^m, \eta \in \{-1, 1\}^n \right\}. \quad (15)$$

Definition

The SDP relaxation of $\|M\|_{\infty \to 1}$ is given via:

$$\mathrm{sdp}_{\infty \to 1}(M) = \max \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} \langle x_i, y_j \rangle$$
$$x_i, y_j \in \mathbb{R}^{m+n}$$
$$\|x_i\| = 1, i = 1, \ldots, m$$
$$\|y_j\| = 1, j = 1, \ldots, n$$

Theorem (Grothendieck's inequality)

There exists a constant K such that for all $M \in \mathbb{R}^{m \times n}$:

$$\|M\|_{\infty\to 1} \le \mathrm{sdp}_{\infty\to 1}(M) \le K\|M\|_{\infty\to 1}. \tag{16}$$

Proof.

Use the following approximation algorithm with randomized rounding:

Algorithm 1: Approximation algorithm with randomized rounding for $\|M\|_{\infty\to 1}$

1. Solve $\mathrm{sdp}_{\infty\to 1}(M)$. Let $x_1, \ldots, x_m, y_1, \ldots, y_n \in S^{m+n-1}$ be the optimal unit vectors
2. Apply Krivine's trick (Lemma 7) and use vectors x_i, y_j to create new unit vectors $x'_1, \ldots, x'_m, y'_1, \ldots, y'_n \in S^{m+n-1}$.
3. Choose $r \in S^{m+n-1}$ randomly
4. Round: $u_i = \mathrm{sign}(\langle x'_i, r\rangle)$
   $v_j = \mathrm{sign}(\langle y'_j, r\rangle)$

Proof (cont.).

Expected quality of the outcome:

$$\begin{aligned}\|M\|_{\infty\to 1} &\geq \mathbb{E}\left[\sum_{i=1}^{m}\sum_{j=1}^{n} M_{ij} u_i v_j\right] \\ &= \sum_{i=1}^{m}\sum_{j=1}^{n} M_{ij}\mathbb{E}[\operatorname{sign}(\langle x_i', r\rangle)\operatorname{sign}(\langle y_j', r\rangle)] \\ &= \sum_{i=1}^{m}\sum_{j=1}^{n} M_{ij}\beta\langle x_i, y_j\rangle \\ &= \beta\,\mathrm{sdp}_{\infty\to 1}(M),\end{aligned}$$

where the last equality follows by Krivine's trick with $\beta = \frac{2\ln(1+\sqrt{2})}{\pi}$, thus $K \leq \beta^{-1}$.

Theorem (Tsirelson)

(Hard direction) For all positive integers n, r *and any* $x_1, \dots, x_n,\ y_1, \dots, y_n \in S^r$*, there exists a positive integer* $d := d(r)$*, a state* $|\psi\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$ *and* $\{-1,1\}$*-observables* $F_1, \dots, F_n, G_1, \dots, G_n \in O(\mathbb{C}^d)$*, such that for every* $i, j \in \{1, \dots, n\}$*, we have*

$$\langle \psi | F_i \otimes G_j | \psi \rangle = \langle x_i, y_j \rangle. \tag{17}$$

Moreover, $d \leq 2^{\lceil r/2 \rceil}$*.*
(Easy direction) Conversely, for all positive integers n, d*, state* $|\psi\rangle \in \mathbb{C}^d \otimes \mathbb{C}^d$ *and* $\{-1,1\}$*-observables* $F_1, \dots, F_n, G_1, \dots, G_n \in O(\mathbb{C}^d)$*, there exist a positive integer* $r := r(d)$ *and* $x_1, \dots, x_n, y_1, \dots, y_n \in S^r$ *such that for every* $i, j \in \{1, \dots, n\}$*, we have*

$$\langle x_i, y_j \rangle = \langle \psi | F_i \otimes G_j | \psi \rangle. \tag{18}$$

Moreover, $r \leq 2d^2$*.*

Since

Theorem (Grothendieck-Tsirelson)

There exists an absolute constant $K \geq 1$ such that, for any positive integers m, n, the following three equivalent conditions hold:

(1) *We have the inclusion*

$$\mathrm{QC}_{m,n} \subset K\mathrm{LC}_{m,n}. \tag{19}$$

(2) *For any $M \in \mathbb{R}^{m\times n}$ and for any ρ, X_i, Y_j verifying the conditions of Definition 4.2.1 we have*

$$\sum_{i,j} M_{ij} \operatorname{Tr} \rho(X_i \otimes Y_j) \leq K \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \sum_{i,j} M_{ij}\xi_i\eta_j \tag{20}$$

$$\Leftrightarrow \qquad \operatorname{Tr} MA^\top \leq \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \operatorname{Tr} M(\xi\eta^\top)^\top. \tag{21}$$

(3) *For any $M \in \mathbb{R}^{m\times n}$ and for any (real) Hilbert space vectors x_i, y_j with $|x_i| \leq 1$, $|y_j| \leq 1$ we have*

$$\sum_{i,j} M_{i,j}\langle x_i, y_j\rangle \leq K \max_{\xi\in\{-1,1\}^m, \eta\in\{-1,1\}^n} \operatorname{Tr} \xi^\top M\eta. \tag{22}$$

Proof.

Since (22) is a direct consequence of Grothendieck's inequality the only thing left to prove is the equivalence between (1)-(3). The equivalence of (3) and (2) (the Tsirelson's bound) is a consequence of either the proof of Lemma **??** or Tsirelsons Theorem (Theorem 1).

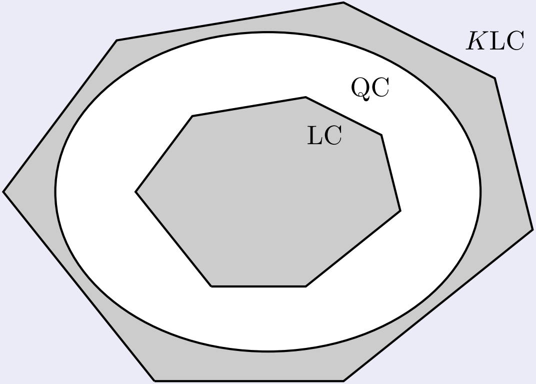

Figure: Visualization