

**This homework is due at 11 PM on November 15, 2023.**

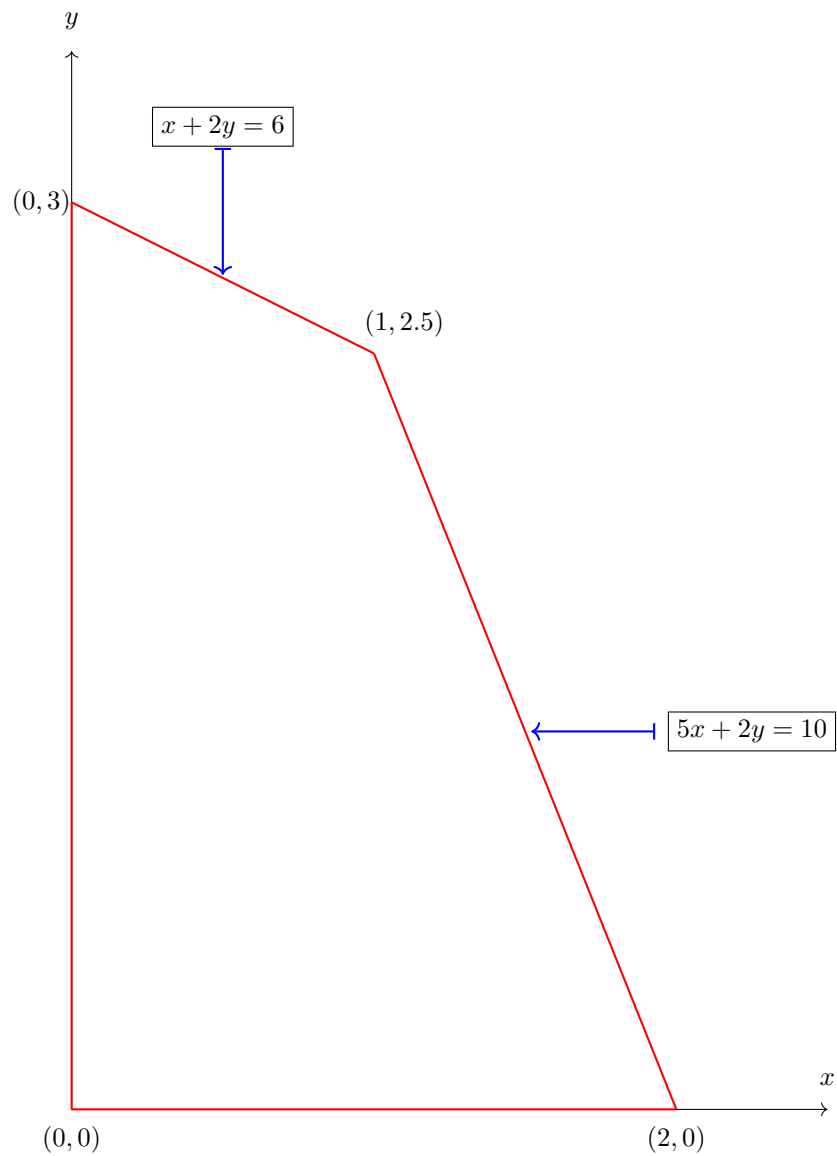
**Submission Format:** Your homework submission should consist of a single PDF file that contains all of your answers (any handwritten answers should be scanned).

### 1. Optimization over a Polytope

Consider the optimization problem:

$$\begin{aligned} \min_{x,y \in \mathbb{R}} \quad & ax + by + c \\ \text{s.t.} \quad & x \geq 0, \\ & y \geq 0, \\ & x + 2y \leq 6, \\ & 5x + 2y \leq 10. \end{aligned}$$

The feasible set is the closed bounded set enclosed by the red lines as given in Fig. 1. Find the optimal solutions for objectives (i)  $-2x + 3y + 5$ , (ii)  $-x - 2y + 5$  and (iii) 5 by finding the values at the vertices of the feasible set and then using the `cvxpy` package in the `optimization_over_polytope.ipynb` file. Please refer to <https://www.cvxpy.org/tutorial/index.html> to gain some familiarity with `cvxpy`, a widely used package for convex optimization.



**Figure 1:** The feasible set for Problem 2 is the polytope bounded by the red lines.

## 2. LP Duality, Part 1

This problem explores basic features of linear programming duality. Consider the following linear programming problem

$$\begin{aligned}
 p^* := \min_{x_1, x_2, x_3 \in \mathbb{R}} \quad & x_1 + x_3 \\
 \text{s.t.} \quad & x_1 + 2x_2 \leq 5, \\
 & x_1 + 2x_3 = 6, \\
 & x_1 \geq 0, \\
 & x_2 \geq 0, \\
 & x_3 \geq 0.
 \end{aligned}$$

- Show that the problem is feasible and that the feasible set is a polytope. (Recall that a polytope is a bounded polyhedron, so what this problem is asking you to show is that the feasible set is nonempty, is a polyhedron, and is bounded.)
- Determine the vertices of the feasible polytope, show that the optimal primal value is  $p^* = 3$ , and find all the optimal feasible points.
- Write the Lagrangian for traditional Lagrange duality for this problem. There will be five dual variables, one for each of the four inequality constraints, and one for the equality constraint.
- Find the dual objective function and show that the dual problem can be simplified to read

$$\begin{aligned}
 \max_{u, v} \quad & -5u - 6v \\
 \text{s.t.} \quad & u \geq 0, \\
 & v \geq -\frac{1}{2}.
 \end{aligned}$$

Note that this is also a linear program.

- Verify that the optimal dual value is  $d^* = 3$  and so strong duality holds.
- Is Slater's condition satisfied in this problem?

### 3. LP Duality, Part 2

This problem continues to explore basic features of linear programming duality with a variation on the preceding problem. Consider the following linear programming problem

$$\begin{aligned} p^* &:= \min_{\vec{x} \in \mathbb{R}^3} && x_1 + x_3 \\ \text{s.t.} &&& x_1 + 2x_2 \leq -5, \\ &&& x_1 + 2x_3 = 6, \\ &&& x_1 \geq 0, \\ &&& x_2 \geq 0, \\ &&& x_3 \geq 0. \end{aligned}$$

- (a) Show that the problem is infeasible and therefore the optimal primal value is  $p^* = \infty$ .
- (b) Write the Lagrangian for traditional Lagrange duality for this problem. There will be five dual variables, one for each of the four inequality constraints, and one for the equality constraint.
- (c) Find the dual objective function and show that the dual problem can be simplified to a linear program involving two variables.
- (d) Show that the optimal dual value is  $d^* = \infty$  and so strong duality once again holds.
- (e) Is Slater's condition satisfied in this problem?

#### 4. Dual Norms and SOCP

Consider the problem

$$p^* = \min_{\vec{x} \in \mathbb{R}^n} \|A\vec{x} - \vec{y}\|_1 + \mu \|\vec{x}\|_2, \quad (1)$$

where  $A \in \mathbb{R}^{m \times n}$ ,  $\vec{y} \in \mathbb{R}^m$ , and  $\mu > 0$ .

- Express this (primal) problem in standard SOCP form.
- Find a dual to the problem and express it in standard SOCP form.

*HINT: Recall that for every vector  $\vec{z}$ , the following dual norm equalities hold:*

$$\|\vec{z}\|_2 = \max_{\vec{u}: \|\vec{u}\|_2 \leq 1} \vec{u}^\top \vec{z}, \quad \|\vec{z}\|_1 = \max_{\vec{u}: \|\vec{u}\|_\infty \leq 1} \vec{u}^\top \vec{z}. \quad (2)$$

*Thus, we can rewrite the objective function of the original problem as*

$$\|A\vec{x} - \vec{y}\|_1 + \mu \|\vec{x}\|_2 = \max_{\vec{u}: \|\vec{u}\|_\infty \leq 1} \vec{u}^\top (A\vec{x} - \vec{y}) + \mu \max_{\vec{v}: \|\vec{v}\|_2 \leq 1} \vec{v}^\top \vec{x}. \quad (3)$$

*We can then express the original (primal) problem as*

$$p^* = \min_{\vec{x}} \max_{\vec{u}, \vec{v}: \|\vec{u}\|_\infty \leq 1, \|\vec{v}\|_2 \leq 1} \vec{u}^\top (A\vec{x} - \vec{y}) + \mu \vec{v}^\top \vec{x}. \quad (4)$$

*To form the dual, we reverse the order of min and max.*

- Assume strong duality holds<sup>1</sup> and that  $m = 100$  and  $n = 10^6$ , i.e.,  $A$  is  $100 \times 10^6$ . Which problem would you choose to solve using a numerical solver: the primal or the dual? Justify your answer.

<sup>1</sup>In fact, you can show that strong duality holds using Sion's theorem, a generalization of the minimax theorem that is beyond the scope of this class.

### 5. Magic with constraints

In this question, we will represent a problem in two different ways and show that strong duality holds in one case but doesn't hold in the other.

Let

$$f_0(x) \doteq \begin{cases} x^3 - 3x^2 + 4, & x \geq 0 \\ -x^3 - 3x^2 + 4, & x < 0 \end{cases}.$$

1) Consider the minimization problem

$$\begin{aligned} p^* = \min_{x \in \mathbb{R}} \quad & f_0(x) \\ \text{s.t.} \quad & -1 \leq x, \\ & x \leq 1. \end{aligned} \tag{5}$$

(a) Show that  $f_0(x)$  is differentiable everywhere and compute its derivative.

(b) Show that  $p^* = 2$  and that the set of optimizers is  $\mathcal{X}^* = \{-1, 1\}$ .

(c) Show that the dual problem can be represented as

$$d^* = \max_{\lambda_1, \lambda_2 \geq 0} g(\vec{\lambda}),$$

where

$$g(\vec{\lambda}) = \min \left\{ g_1(\vec{\lambda}), g_2(\vec{\lambda}) \right\},$$

with

$$g_1(\vec{\lambda}) = \min_{x \geq 0} x^3 - 3x^2 + 4 - \lambda_1(x + 1) + \lambda_2(x - 1),$$

$$g_2(\vec{\lambda}) = \min_{x < 0} -x^3 - 3x^2 + 4 - \lambda_1(x + 1) + \lambda_2(x - 1).$$

(d) Next, show that

$$g_1(\vec{\lambda}) \leq -3\lambda_1 + \lambda_2,$$

$$g_2(\vec{\lambda}) \leq \lambda_1 - 3\lambda_2.$$

Use this to show that  $g(\vec{\lambda}) \leq 0$  for all  $\lambda_1, \lambda_2 \geq 0$ .

(e) Show that  $g(\vec{0}) = 0$  and conclude that  $d^* = 0$ .

(f) Does strong duality hold?

2) Now, consider a problem equivalent to the minimization in (5):

$$p^* = \min_{x \in \mathbb{R}} f_0(x) \tag{6}$$

$$\text{s.t. } x^2 \leq 1. \tag{7}$$

Observe that  $p^* = 2$  and the set of optimizers is  $\mathcal{X}^* = \{-1, 1\}$ , since this problem is equivalent to the one in part 1).

(a) Show that the dual problem can be represented as

$$d^* = \max_{\lambda \geq 0} g(\lambda),$$

where

$$g(\lambda) = \min(g_1(\lambda), g_2(\lambda)),$$

with

$$g_1(\lambda) = \min_{x \geq 0} x^3 - 3x^2 + 4 + \lambda(x^2 - 1),$$

$$g_2(\lambda) = \min_{x < 0} -x^3 - 3x^2 + 4 + \lambda(x^2 - 1).$$

- (b) Show that  $g_1(\lambda) = g_2(\lambda) = \begin{cases} 4 - \lambda, & \lambda \geq 3, \\ -\frac{4}{27}(3 - \lambda)^3 + 4 - \lambda, & 0 \leq \lambda < 3. \end{cases}$
- (c) Conclude that  $d^* = 2$  and  $\lambda^* = \frac{3}{2}$ .
- (d) Does strong duality hold?