1. Since the [optimize_Oplus] transformation doesn't change the value of [aexp]s, we should be able to apply it to all the [aexp]s that appear in a [bexp] without changing the [bexp]'s value. Write a function that performs this transformation on [bexp]s and prove it is sound. Use the tacticals we've just seen to make the proof as short and elegant as possible.

代码

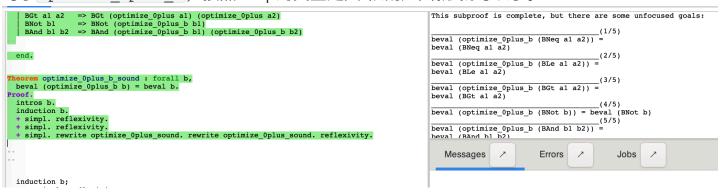
```
Fixpoint optimize Oplus b (b : bexp) : bexp :=
 2
     match b with
 3
      BTrue
                   => b
      BFalse
                  => b
 4
                   => BEq (optimize Oplus a1) (optimize Oplus a2)
 5
     BEq a1 a2
     BNeq a1 a2 => BNeq (optimize_0plus a1) (optimize_0plus a2)
 6
      BLe a1 a2 => BLe (optimize Oplus a1) (optimize Oplus a2)
 7
     BGt a1 a2 => BGt (optimize Oplus a1) (optimize Oplus a2)
 8
 9
      BNot b1 => BNot (optimize Oplus b b1)
      BAnd b1 b2 => BAnd (optimize Oplus b b1) (optimize Oplus b b2)
10
     end.
11
12
   Theorem optimize Oplus b sound : forall b,
13
     beval (optimize Oplus b b) = beval b.
14
   Proof.
15
     intros b.
16
17
     induction b;
    try (simpl; reflexivity);
18
     try (simpl; repeat rewrite optimize Oplus sound; reflexivity).
19
     + simpl. rewrite IHb. reflexivity.
20
     + simpl. rewrite IHb1. rewrite IHb2. reflexivity.
21
22
   Qed.
```

分析

提示的意思就是,由于optimize_Oplus转换不会改变aexps的值,因此应该能够将其应用于bexp中出现的所有aexps,而无需更改bexp的值。编写一个在bexps上执行此转换的函数,并证明它是正确的。使用我们刚刚看到的tacitics,使证明尽可能简洁优雅。

```
Inductive bexp : Type :=
2
      BTrue
3
       BFalse
      BEq (a1 a2 : aexp)
4
5
     BNeq (a1 a2 : aexp)
     BLe (al a2 : aexp)
6
      BGt (a1 a2 : aexp)
      BNot (b : bexp)
8
      BAnd (b1 b2 : bexp).
9
```

对于 optimize_0plus_b, 按照bexp的类型定义"依葫芦画瓢"就可以了。



对于 optimize_0plus_b_sound , intros 和 induction 之后,根据观察可以把需要证明的8个分支分为3类。1-2直接 simpl. reflexivity.即可,第二种,根据试验,只需要 simpl. 之后不同 rewrite optimize_0plus_sound.,最后再 reflexivity。最后的 BNot 和 BAnd 则需要用一些假设去 rewrite。因为要求证明要尽可能的"elegant",所以我们使用 try 和 repeat。具体如上面的代码。

运行结果

```
Fixpoint optimize_Oplus_b (b : bexp) : bexp :=
   match b with
       BTrue
      BFalse
                         => b
   BFalse => D

BEq al a2 => BEq (optimize_Oplus a1) (optimize_Oplus a2)

BNeq al a2 => BNeq (optimize_Oplus a1) (optimize_Oplus a2)

BLe al a2 => BLe (optimize_Oplus a1) (optimize_Oplus a2)

BGt al a2 => BGt (optimize_Oplus a1) (optimize_Oplus a2)

BNot b1 => BNot (optimize_Oplus_b b1)

BANd b1 b2 => BANd (optimize_Oplus_b b1) (optimize_Oplus_b b2)
end.
Theorem optimize_Oplus_b_sound : forall b,
  beval (optimize_Oplus_b b) = beval b.
Proof.
  intros b.
   induction b;
  try (simpl; reflexivity);
try (simpl; repeat rewrite optimize_Oplus_sound; reflexivity).
  + simpl. rewrite IHb. reflexivity.
  + simpl. rewrite IHb1. rewrite IHb2. reflexivity.
Qed.
```