# 1.

### 分析:

根据提示,Fixpoint中的定义必须都是递归的,所以这里需要同时对l1和l2进行模式匹配。基线条件就是l1或l2中有一个是空,那么就返回另外一个。若都非空则将l1和l2的头元素重新组合成列表再在后面进行递归。

## 代码:

```
1
    (** **** Exercise: 3 stars, advanced (alternate)
 2
3
        Complete the following definition of [alternate], which
        interleaves two lists into one, alternating between elements taken
 4
5
        from the first list and elements from the second. See the tests
        below for more specific examples.
 6
 7
        Hint: there is an elegant way of writing [alternate] that fails to
8
9
        satisfy Coq's requirement that all [Fixpoint] definitions be
        _structurally recursing_, as mentioned in [Basics]. If you
10
        encounter that difficulty, consider pattern matching against both
11
12
        lists at the same time with the "multiple pattern" syntax we've
        seen before. *)
13
14
    Fixpoint alternate (11 12 : natlist) : natlist :=
15
      match 11, 12 with
16
17
      nil, => 12
      _, nil => 11
18
      | h1 :: t1, h2 :: t2 => h1 :: h2 :: (alternate t1 t2)
19
20
      end.
2.1
22
23
   Example test alternate1:
24
     alternate [1;2;3] [4;5;6] = [1;4;2;5;3;6].
25
    Proof. reflexivity.
     (* FILL IN HERE *) Admitted.
26
2.7
28
   Example test_alternate2:
29
     alternate [1] [4;5;6] = [1;4;5;6].
30
   Proof. reflexivity.
     (* FILL IN HERE *) Admitted.
31
32
   Example test alternate3:
33
34
     alternate [1;2;3] [4] = [1;4;2;3].
35
   Proof. reflexivity.
     (* FILL IN HERE *) Admitted.
36
```

## 运行结果:

```
Fixpoint alternate (I1 I2 : natlist) : natlist :=
match I1, I2 with
I nil, _ => 12
I _, nil => I1
I h1 :: t1, h2 :: t2 => h1 :: h2 :: (alternate t1 t2)
end.
Example test_alternate1:
alternate [1;2;3] [4;5;6] = [1;4;2;5;3;6].
Proof. reflexivity.
(* FILL IN HERE *) Admitted.
Example test_alternate2:
alternate [1] [4;5;6] = [1;4;5;6].
Proof. reflexivity.
(* FILL IN HERE *) Admitted.
Example test_alternate3:
alternate [1;2;3] [4] = [1;4;2;3].
Proof. reflexivity.
(* FILL IN HERE *) Admitted.
```

# 2.

## 分析:

由于不知道题目中所说的集合中是否包含重复的元素,所以两种理解都进行了实现。实现的方法都是递归。若不能出现相同元素,则通过已经实现的count函数,判断其返回值是否为0,就能够判断该元素是否在交集之中。若能出现相同元素,则需要在递归函数中用已经实现的remove\_one函数进行处理。

## 代码:

```
1
   Fixpoint inter (s1: bag) (s2: bag) : bag :=
2
      match s1 with
      nil => nil
 3
      h :: t => if ((count h s2) =? 0) then inter t s2
 4
 5
                     else h :: inter t s2
 6
      end.
7
    Example test_inter1: inter [1;3;2;4;5] [1;2;3] = [1;3;2].
8
9
    Proof. simpl. reflexivity. Qed.
10
    Example test_inter2: inter [1;3;2;4;5] [ ] = [ ].
11
    Proof. simpl. reflexivity. Qed.
12
13
14
    Example test_inter3: inter [] [1;2;3] = [].
    Proof. simpl. reflexivity. Qed.
15
16
```

```
Example test_inter4: inter [1;3;5] [1;2;3;5;7;9] = [1;3;5].

Proof. simpl. reflexivity. Qed.
```

```
1
    Fixpoint inter' (s1: bag) (s2: bag) : bag :=
 2
      match s1, s2 with
      nil, _ => nil
 3
 4
      _, nil => nil
      h1 :: t1, _ => if ((count h1 s2) =? 0) then inter t1 s2
 5
                             else h1 :: inter' t1 (remove_one h1 s2)
 6
 7
      end.
 8
9
    Example test_inter'1: inter' [1;3;1;1;2;4;5] [1;2;1;3] = [1;3;1;2].
10
    Proof. simpl. reflexivity. Qed.
11
12
   Example test_inter'2: inter' [1;3;3;2;4;5] [ ] = [ ].
13
    Proof. simpl. reflexivity. Qed.
14
15
16
    Example test_inter'3: inter' [] [1;2;3;1] = [].
17
    Proof. simpl. reflexivity. Qed.
18
19
   Example test_inter'4: inter' [1;3;5;8;6;1] [1;2;3;5;7;1;9] = [1;3;5;1].
20
   Proof. simpl. reflexivity. Qed.
```

#### 运行结果:

```
Fixpoint inter (s1: bag) (s2: bag) : bag :=
match s1 with
I nil => nil
I h :: t => if ((count h s2) =? 0) then inter t s2
else h :: inter t s2
end.

Example test_inter1: inter [1;3;2;4;5] [1;2;3] = [1;3;2].

Proof. simpl. reflexivity. Qed.

Example test_inter2: inter [1;3;2;4;5] [] = [].

Proof. simpl. reflexivity. Qed.

Example test_inter3: inter [] [1;2;3] = [].

Proof. simpl. reflexivity. Qed.

Example test_inter4: inter [1;3;5] [1;2;3;5;7;9] = [1;3;5].

Proof. simpl. reflexivity. Qed.
```