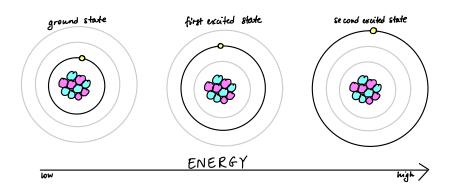
# CPEN 400Q Lecture 15 Solving combinatorial optimization problems with the Quantum Approximate Optimization Algorithm (part 1)

Monday 6 March 2023

#### Announcements

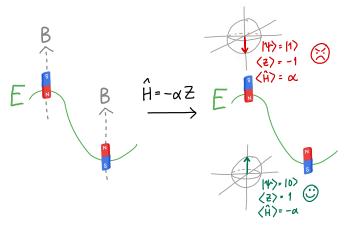
- Quiz 6 end of class today
- Assignment 2 due 13 March at 23:59
- Updated class schedule:
  - Friday March 10: Zoom (same link)
  - Monday March 13: in person (quiz 7)
  - Friday March 17: pre-recorded "infotainment" lecture about compilation
- Meetings this week for project prototypes on Zoom

We tackled the problem of computing the *ground state* and *ground state energy* of a physical system.

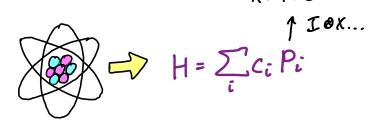


A **Hamiltonian** is an operator that describes the energy of a quantum system.

To run on a quantum computer, we *map* the system to an equivalent qubit-based Hamiltonian.

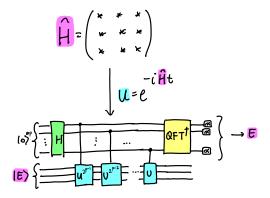


Any qubit Hamiltonian can be expressed as a linear combination of multi-qubit Pauli observables.



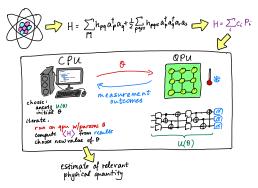
The Hamiltonian is a matrix. Its **spectrum** (**eigenvalues** and **eigenstates**) represents the possible measurement outcomes, just like single Paulis.

We can run *quantum phase estimation* to get the energy spectrum.



But this is not practical on near-term quantum devices.

The variational quantum eigensolver (VQE) is more practical.



The optimization works because of the variational principle:

for any other state  $|\psi\rangle$ .

You were given the Hamiltonian for a *deuteron* represented by a 4-qubit system.

I told you its ground state had the form

$$|\psi_{\rm g}\rangle = {\it a}|1000\rangle + {\it b}|0100\rangle + {\it c}|0010\rangle + {\it d}|0001\rangle$$

#### Challenge:

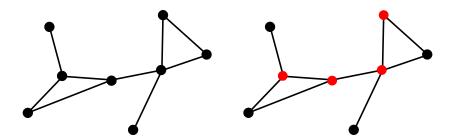
- Design a circuit ansatz that produces states that look like this
- Run VQE to determine the ground state energy

# Learning outcomes

- Describe the underlying ideas of adiabatic quantum computation and QAOA
- Convert an optimization problem over binary variables into an equivalent problem over qubits.

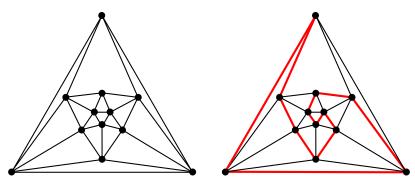
## Combinatorial optimization

Example: Given a graph G = (V, E), what is the *smallest number* of vertices you can colour such that every edge in the graph is attached to at least one coloured vertex?



# Combinatorial optimization

Example: Given a graph, can we find a path through it that visits every node *exactly once* and returns to the starting point?



(In graph theory terms, can we find a Hamiltonian cycle?)

#### Combinatorial optimization

Example: You have the opportunity to purchase 100 units of stocks from a fixed list of assets. You know the average returns of each stock, and their covariances.

Stock	Avg. return
AAA	3.44 %
BBB	2.21 %
CCC	-0.28 %
:	:

Cov.	AAA	BBB	
AAA	0.0038	0.002	
BBB	0.002	-0.006	
CCC	0.014	-0.0008	
:	:	:	٠

Suppose you're restricted to buying no more than 5 of any stock.

Which stocks, and how many of each, should you purchase, to maximize your profits?

The structure of a classical optimization problem is something like:

$$\min_{\vec{x}} \ \, \mathsf{cost}(\vec{x}) \quad \mathsf{subject to constraints}(\vec{x})$$

where  $\vec{x}$  is a multi-dimensional vector of variables of the problem.

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where  $\vec{x}$  is a multi-dimensional vector of variables of the problem.

In a physical context, optimization can be interpreted as an **energy minimization problem**.

Optimization	Physical system
$\vec{x}$	State of the system
$cost(ec{x})$	Hamiltonian
Optimum $\vec{x^*}$	Ground state
$cost(\vec{x^*})$	Ground state energy

Recall that every unitary U is associated with a Hamiltonian  $\hat{H}$  under the correspondence

We know that we can use gate model QC to *simulate* the evolution of a Hamiltonian.

Instead of simulating the Hamiltonians, adiabatic quantum computing works with them directly to perform computations. It is generally used to solve optimization problems.

#### The adiabatic theorem

#### Theorem:

"A physical system remains in its instantaneous eigenstate if a given perturbation is acting on it slowly enough and if there is a gap between the eigenvalue and the rest of the Hamiltonian's spectrum."

#### What we can take from this:

If we initialize a system in the lowest energy state and perturb it slowly enough, it will remain in the lowest energy state (with respect to the changed system)

Let  $H_c$  be a **cost Hamiltonian** whose ground state represents the solution to a problem of interest.

Let  $H_m$  be a **mixer Hamiltonian** whose ground state can be easily prepared.

Adiabatic evolution is expressed mathematically as the function

The parameter s is representative of time; s goes from 0 to 1; A(s) decreases to 0 with time and B(s) increases from 0.

# Quantum annealing

D-Wave makes quantum annealers: these are a physical implementation of AQC for a limited set of Hamiltonians.



Image credit:
www.dwavesys.com/tutorials/background-reading-series/introduction-d-wave-quantum-hardware

# Quantum approximate optimization algorithm (QAOA)

QAOA is a gate-model algorithm that can obtain approximate solutions to combinatorial optimization problems.

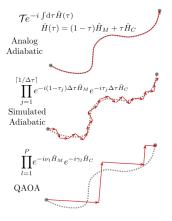
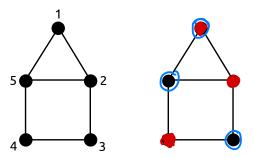


Image credit: G. Verdon, M. Broughton, J. Biamonte. A quantum algorithm to train neural networks using

low-depth circuits. https://arxiv.org/abs/1712.05304

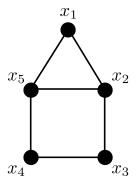
**Vertex cover**: given a graph G = (V, E) find the *minimum* number of coloured vertices such that every edge is adjacent to a coloured vertex. This problem is **NP-hard**.



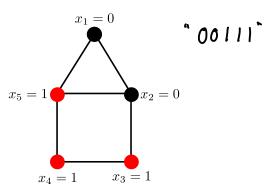
#### We will:

- define a cost function that is minimized for a valid colouring
- map it to a Hamiltonian
- run a quantum algorithm to find the ground state

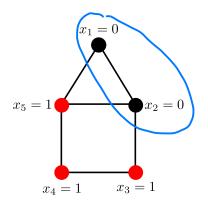
Whether or not a vertex is coloured is a binary variable.



Let's assign coloured vertices to have value 1, and un-coloured 0.

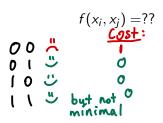


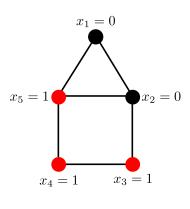
Now that we have our variables, how do we come up with a minimizable cost function that represents the problem?



We need every edge to be next to a coloured vertex.

Design a function on two vertices that is minimized when colouring is valid, and maximized if it's not.

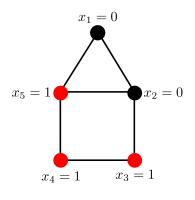




Consider for each edge *ij* the function

$$f(x_i, x_j) = (1 - x_i)(1 - x_j)$$

The possible values are:



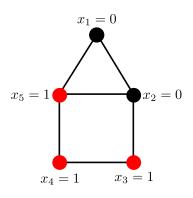
Then in an optimal colouring,

$$f(x_i,x_j)=(1-x_i)(1-x_j)=0$$

for all edges  $ij \in E$ .

So we can write

$$\min_{\vec{X}} \sum_{ij \in E} (1-x_i)(1-x_j)$$



However recall that we also want to colour the fewest vertices. The cost should also depend on the number of coloured vertices.

Solution: add to our cost

The full cost function is then

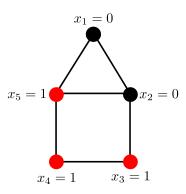
$$\min_{\vec{X}} \left( \sum_{ij \in E} (1-x_i)(1-x_j) + \sum_{i \in V} x_i \right)$$

#### Next steps:

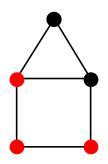
- 1. Turn this into a Hamiltonian over qubits
- 2. Find its minimum energy

# We will start here next time!

How should we do a mapping from this to qubits?



What should we use to represent cost?



Mathematically, we can make the mapping

$$x_i \to \frac{1}{2}(1-z_i), \quad z_i \in \{-1, 1\}$$

This associates

- $x_i = 0$  to  $z_i = 1$  (corresponds to  $|0\rangle$ )
- $x_i = 1$  to  $z_i = -1$  (corresponds to  $|1\rangle$ )

A complete derivation of the cost function is provided as an appendix at the end of the lecture slides.

We will take the result:

$$\sum_{ij \in E} (1 - x_i)(1 - x_j) + \sum_{i \in V} x_i \to \sum_{ij \in E} (z_i + z_j + z_i z_j) - 2 \sum_{i \in V} z_i$$

Remember what the  $z_i$  represent; how can we express this cost function as a Hamiltonian?

$$\hat{H} = \sum_{ij \in E} (Z_i + Z_j + Z_i Z_j) - 2 \sum_{i \in V} Z_i$$

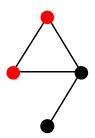
$$\hat{H} = \sum_{ij \in E} (Z_i + Z_j + Z_i Z_j) - 2 \sum_{i \in V} Z_i$$

This makes sense:

$$\langle \hat{H} \rangle = \sum_{ij \in E} (\langle Z_i \rangle + \langle Z_j \rangle + \langle Z_i Z_j \rangle) - 2 \sum_{i \in V} \langle Z_i \rangle$$

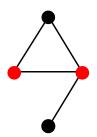
$$\hat{H} = \sum_{ij \in E} (Z_i + Z_j + Z_i Z_j) - 2 \sum_{i \in V} Z_i$$

Try it: what is the energy of this *invalid* colouring?



$$\hat{H} = \sum_{ij \in E} (Z_i + Z_j + Z_i Z_j) - 2 \sum_{i \in V} Z_i$$

Try it: what is the energy of this valid colouring?



#### Next time

#### Content:

 Solving combinatorial optimization problems with QAOA in PennyLane

#### Action items:

- 1. Technical assignment 2
- 2. Work on prototype implementation for project

#### Recommended reading (PennyLane demos):

- https://pennylane.ai/qml/demos/tutorial\_qaoa\_intro.html
- https://pennylane.ai/qml/demos/tutorial\_qaoa\_maxcut.html

In order

We will make the mapping

$$x_i \to \frac{1}{2}(1-z_i), \quad , z_i \in \{-1,1\}$$

This associates  $x_i=0$  to  $z_i=1$  (corresponds to  $|0\rangle$ ), and  $x_i=1$  to  $z_i=-1$  (corresponds to  $|1\rangle$ ).

Let's expand our cost function and make this substitution.

$$\sum_{ij\in E}(1-x_i)(1-x_j)+\sum_{i\in V}x_i$$

$$\sum_{ij\in E}(1-x_i-x_j+x_ix_j)+\sum_{i\in V}x_i$$

$$\sum_{ij\in E} (1-x_i-x_j+x_ix_j) + \sum_{i\in V} x_i$$

Substitute:

$$\sum_{ij\in E} \left(1 - \frac{1}{2}(1 - z_i) - \frac{1}{2}(1 - z_j) + \frac{1}{4}(1 - z_i)(1 - z_j)\right) + \sum_{i\in V} \frac{1}{2}(1 - z_i)$$

Expand:

$$\sum_{ij\in E} \left(1 - \frac{1}{2} + \frac{1}{2}z_i - \frac{1}{2} + \frac{1}{2}z_j + \frac{1}{4} - \frac{1}{4}z_i - \frac{1}{4}z_j + \frac{1}{4}z_iz_j\right) + \sum_{i\in V} \frac{1}{2}(1 - z_i)$$

Collect:

$$\sum_{i \in F} \left( \frac{1}{4} + \frac{1}{4}z_i + \frac{1}{4}z_j + \frac{1}{4}z_i z_j \right) + \sum_{i \in V} \frac{1}{2} (1 - z_i)$$

$$\sum_{ij\in E} \left(\frac{1}{4} + \frac{1}{4}z_i + \frac{1}{4}z_j + \frac{1}{4}z_iz_j\right) + \sum_{i\in V} \frac{1}{2}(1-z_i)$$

Consider now that: the total number of edges and vertices are constant - they will provide only an "offset" to the cost, and the values of the variables don't matter.

$$\sum_{ij \in E} \left( \frac{1}{4} z_i + \frac{1}{4} z_j + \frac{1}{4} z_i z_j \right) - \sum_{i \in V} \frac{1}{2} z_i$$

And finally, the absolute value doesn't matter, so we can rescale:

$$\sum_{ij\in E} (z_i + z_j + z_i z_j) - 2 \sum_{i\in V} z_i$$

Can also weight the terms differently depending on which constraint is more important (i.e., if you care more about just getting a valid colouring, weight the first one more).

$$\gamma \sum_{ij \in E} (z_i + z_j + z_i z_j) - 2\lambda \sum_{i \in V} z_i$$

To turn this into a Hamiltonian, recall that

- Each  $z_i$  represents an expectation value of  $Z_i$
- Computing expectation values is linear

$$\gamma \sum_{ij \in E} (z_i + z_j + z_i z_j) - 2\lambda \sum_{i \in V} z_i$$

$$\hat{H} = \gamma \sum_{ij \in E} (Z_i + Z_j + Z_i Z_j) - 2\lambda \sum_{i \in V} Z_i$$

Next time: we will look at the actual QAOA that can find the optimal configuration / minimum energy.