Assignment 0x03

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Task 1: Number Theory and Algebra 1

1.1 Explain why the integers Z form a ring, and the rationals Q form a field.

A ring is a set R with two operations "+" and "*". (R,+) is an abelian group, that means it is associative, commutative and has a neutral and inverse element. (R,*) is a monoid, that means it is associative and has a neutral element. Lastly multiplication is distributive with respect to addition. If all axioms hold, a set is a ring. And they hold for Z.

A field is a ring but with the special case that $(R\setminus\{0\},*)$ is an abelian group. Meaning that associativity and commutativity hold for all elements except 0 and that there also exists a neutral and inverse element for every element of the field.

Find the inverses of all elements in \mathbb{Z}_7^* . Why do all 1.2numbers between 1 and 6 have an inverse?

```
Z_7^* = 1, 2, 3, 4, 5, 6
A number is an inverse, if number mod inverse = 1 (the neutral element)
    Number:
               1
                   2 3
                          4
                               5
                                  6
    Inverse:
               1
                       5
                           2
                               3
```

If n is a prime, then all numbers between 1 and n-1 have an inverse with the modulo-Operation. This is due to n not being a multiple of any of the elements in the set.

As an exercise, calculate the inverse of 331 in \mathbb{Z}_{1234} using the extended euclidean algorithm.

```
GCD = 1, if Rest is 0 in the end
1234 == 3 * 331 + 241
331 == 1 * 241 + 90
241 == 2 * 90 + 61
90 == 1 * 61 + 29
61 == 2 * 29 + 3
29 == 9 * 3 + 2
3 == 1 * 2 + 1
2 == 2 * 1 + 0
   Applying the algorithm:
```

```
1 == 3 - 1 * 2
1 == 3 - (29 - 9 * 3)
1 == 10 * 3 - 1 * 29
1 == 10 * (61 - 2 * 29) - 1 * 29
1 == 10 * 61 - 21 * 29
1 == 10 * 61 - 21 * (90 - 1 * 61)
1 == 31 * 61 - 21 * 90
```

```
1 == 31 * (241 - 2 * 90) - 21 * 90
1 == 31 * 241 - 83 * 90
1 == 31 * 241 - 83 * (331 - 241 * 1)
1 == 114 * 241 - 83 * 331
1 == 114 * (1234 - 3 * 331) - 83 * 331
1 == 114 * 1234 - 425 * 331
```

Inverse of 331 is -425. Since it is not in Z we add 1234 to it until it is: -425 + 1234 = 809 = 809 is the inverse of 331 in canonical form.

Test: 331 * 809 = 267779

267779%1234 = 1

So the calculations are correct.

1.4 Find all generators of \mathbf{Z}_{11}

	1	2	3	4	5	6	7	8	9	10	
1	1	1	1	1	1	1	1	1	1	1	
2	2	4	8	5	10	9	7	3	6	1	TRUE
3	3	9	5	4	1	3	9	5	4	1	FALSE
4	4	5	9	3	1	4	5	9	3	1	FALSE
5	5	3	4	9	1	5	3	4	9	1	FALSE
6	6	3	7	9	10	5	8	4	2	1	TRUE
7	7	5	2	3	10	4	6	9	8	1	TRUE
8	8	9	6	4	10	3	2	5	7	1	TRUE
9	9	4	3	5	1	9	4	3	5	1	FALSE
10	10	1	10	1	10	1	10	1	10	1	TRUE

ExcelFormula = MOD((linke Spalte Obere Zeile);11)

The marked lines (2, 6, 7, 8) are generators.

1.5 Calculate 42^{497} in Z_{1361} using fast exponentiation and a handheld calculator (not the one running on your computer or a programming language). Document a, e and n for each step.

1st Iteration: e = 1n is odd: $e = (42 * 1) \mod 1361 = 42$ $a = 42 \mod 1361 = 1764 \mod 1361 = 403$ n = 497 / 2 = 2482): e = 42n is even: $a = 403 \mod 1361 = 450$

Fastexp(a=42, n=497)

```
n = n/2 = 124
   3)
   e = 42
n is even:
a=450 \bmod 1361=1072
n = 124 / 2 = 62
   4)
   e = 42
n is even:
a = 1072 \bmod 1361 = 500
n = 62 / 2 = 31
   5)
   e = 42
n is odd:
e = (500 * 42) \mod 1361 = 585
a = 500 \mod 1361 = 937
n = 31 / 2 = 15
   6)
   e = 1072
n is odd:
e = (937 * 585) \mod 1361 = 1023
a = 937 \mod 1361 = 124
n = 15 / 2 = 7
   7)
   e = 46
n is odd:
e = (124 * 1023) \mod 1361 = 279
a = 124 \bmod 1361 = 405
n = 7 / 2 = 3
   8)
   e = 260
n is odd:
e = (405*279) \mod 1361 = 32
a = 405 \mod 1361 = 705
```

$$n=3\ /\ 2=1$$

9)

$$e = 503$$

n is odd:

 $e = (705*32) \mod 1361 = 784$

 $a = 705 \mod 1361 = 260$

n = 1 / 2 = 0

10)

n = 0 = > return e = 784

 $=>42497 \mod 1361=784$

2 Task 2: The RSA cryptosystem

2.1 Using that private key, decrypt c

$$PHI(n) = (p-1)(q-1) = (17-1) * (23-1) = 16 * 22 = 352$$

$$e = 17; Z_{352}^*$$

(e * d) mod $352 = 1$

Using Excel we found the inverse using the formula above: d = 145

$$m = c^d \mod n = 282^{145} \mod 391 = 197$$

2.2 Did Donald choose reasonable parameters? Explain!

When e = 1, the message itself is encrypted just by applying the modulo with the public Key part N. Resulting in $c = m \mod N$, which means that if m is smaller than N, the encrypted message is just the plain message => c = m

2.3 Find a value of e for which RSA encryption is the identity function

I don'T know:(

2.4 How could the attacker Eve get the AES key without factoring?

Since $2^{4096} > 256^3$ the encrypted key will never have the modulo applied. So the resulting cipher will just be a 1234 bit long key (len(2^{4096})) with the first 256 bit being the AES-Key encrypted with potency e=3. I'm not sure whether the rest are zeroes without padding or garbage bytes or non-existent at all, but the actual key should be well distinguishable. Now you just have to do c^{-e} to get the message.

2.5 How could an adversary Eve get the combination for the lock if she intercepts the message?

Bob fills the entire message, up to the mod length, with 9s, plus a 5 digit number that is the actual message. As it's always just nines plus the number, they is insufficient randomness, which padding usually provides in RSA, allowing Eve to just repeatedly guess messages and encrypt them until one comes out similar. If she knows the subject of the message, she can just try tenthousand numbers with added nines