

Agenda → Code Bubble Sort

↳ Insertion sort

↳ Counting Sort

5, 4, 3, 2, 1

4 5 3 2 1

4 3 5 2 1

4 3 2 5 1

4 3 2 1 5

3 4 2 1 5

3 2 4 1 5

3 2 1 4 5

5, 4, 3, 2, 1

4 5 3 2 1

4 3 5 2 1

4 3 2 5 1

4 3 2 1 5

3 4 2 1 5

3 2 4 1 5

3 2 1 4 5

not sorted

~~1 = 0~~ (1)

5 4 3 2 1

$j \rightarrow [0, \text{len}(arr) - 1]$

$i \rightarrow$ iterator no
for Greed

$S - 0 - 1$

4 5 3 2 1

9

4 3 5 2 1

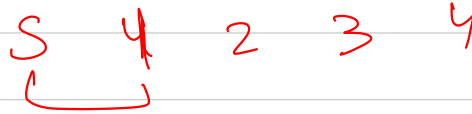
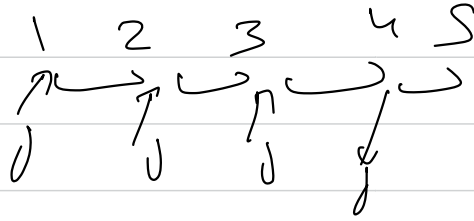
$S - 1 - 1$

3

4 3 2 5 1

$(j, j+1)$ Coff

4 3 2 1



→ No swap we call
stop

```

1 def bubble_sort(arr):
2     for i in range(0, len(arr)):
3         swapped = False
4         for j in range(0, len(arr) - i - 1):
5             if arr[j] > arr[j+1]:
6                 arr[j], arr[j+1] = arr[j+1], arr[j]
7                 swapped = True
8
9         if not swapped:
10             return # why ?? because now the list will be
11                 sorted
12
13     return

```

[5, 4, 3]

$\rightarrow i \rightarrow$ start from 0 & ends at $\frac{\text{len(arr)}-1}{2}$

$i \rightarrow [0, 4]$

when $i=0$ $j [0, 3]$

0 1 2 3 4
1 2 3 4 5

when $i=1 \rightarrow j [0, 2]$

$i=2 \rightarrow j [0, 1]$

5 4 3
4 3 2

1 2 3 4 5

4 3 2
3 2 1

i	j	swapped
0	0	True
	1	True
	2	True
	3	True
1	0	True
	1	True
	2	True
2	0	False
	1	False

$$n-1 + n-1 + n-1 + \dots + n-1$$

$$\underline{\underline{n(n-1)}}$$

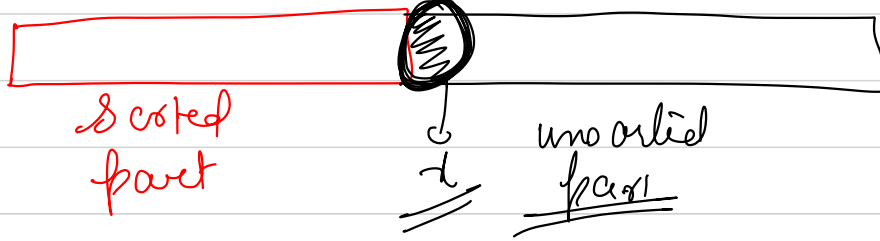
$$n-1 + n-2 + n-3 + \dots + 1$$



Insertion Sort :

Cards

insert it to
the right pos.



1, 2, 3, 4, ...

→ search first for the right position of x^{th} card

→ Insert that x^{th} card to the right pos.

25, 17, 31, 13, 2

17, 25, 31, 13, 2

sorted
unsorted

17, 25, 31, 13, 2

(1) → pick 17

17, 25, 13, 31, 2

(2) → pick 31

17, 13, 25, 31, 2

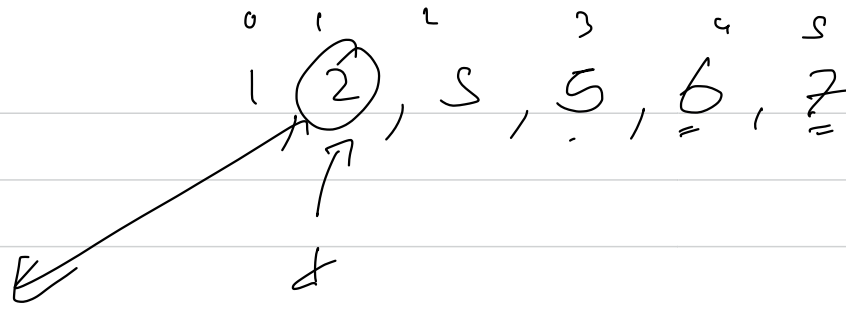
(3) pick 13

13, 17, 25, 31, 2

13, 17, 25, 2, 31

13, 17, 2, 25, 31

13, 17, 25, 31 → 2, 13, 17, 25, 31



$i = 5$
 ~~$j = 5$~~ ~~22~~

$\text{key} = 4$

2 was the first element less than key

So right place of key will be to insert

right of 2 ($j+1$)

```

1 def insertion_sort(arr):
2     for i in range(1, len(arr)):
3         key = arr[i] # this element will be inserted
4         j = i - 1 # because the right element is already sorted
5         while j >= 0 and key < arr[j]:
6             arr[j+1] = arr[j]
7             j -= 1
8         arr[j+1] = key
9
10    return arr
11
12    li = [5,4,3,2,1]

```

$i \rightarrow [1, 4]$

1 2 4 3 0

1 2 4 4 0

1 2 3 4 0

1 2 3 4 4

1 2 3 3 4

1 2 2 3 4 \rightarrow 1 1 2 3 4

sorted array
 0 1 2 3 4
 1 1 2 3 4

i		key
1	0	2
2	1	4
3	2	3
3	1	3
4	3	0
4	2	0
4	1	0
4	0	0
	-1	

TC \rightarrow $O(n^2)$, ^{$O(n^2)$} $\Omega(n)$

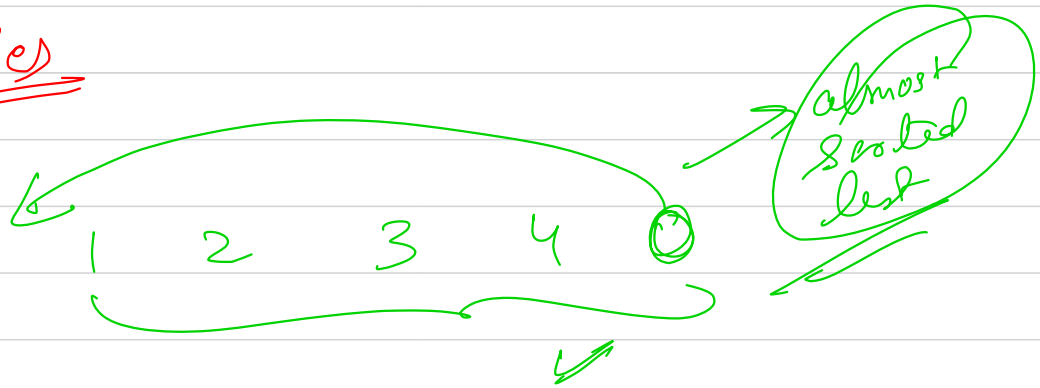
Space $\rightarrow O(1)$

No of swaps \rightarrow $O(n^2)$

Inplace \rightarrow Yes

No of comparisons \rightarrow $O(n^2)$

Stable \rightarrow Yes



5 4 3 2 1

1 5 4 3 2

$$\frac{n}{n-1}$$

1 2 ~~5~~ ~~4~~ 3

$$n-2$$

5 2 3 ~~5~~ 4

$$n-3$$

⋮

$$1 + n-1 + n-2 + \dots + 1$$

$$\underline{\underline{O(n^2)}}$$

5 4 3 2 1 1

Counting Sort

4 * 1 * 6

O(n)

1 - ~~3~~ ~~2~~ ~~0~~

8 - 2

4 - 2

3 - 1

2 - 1

size n

nd

O(n)

1, 1, 3, 4, 5, 3, 4, 1, 2

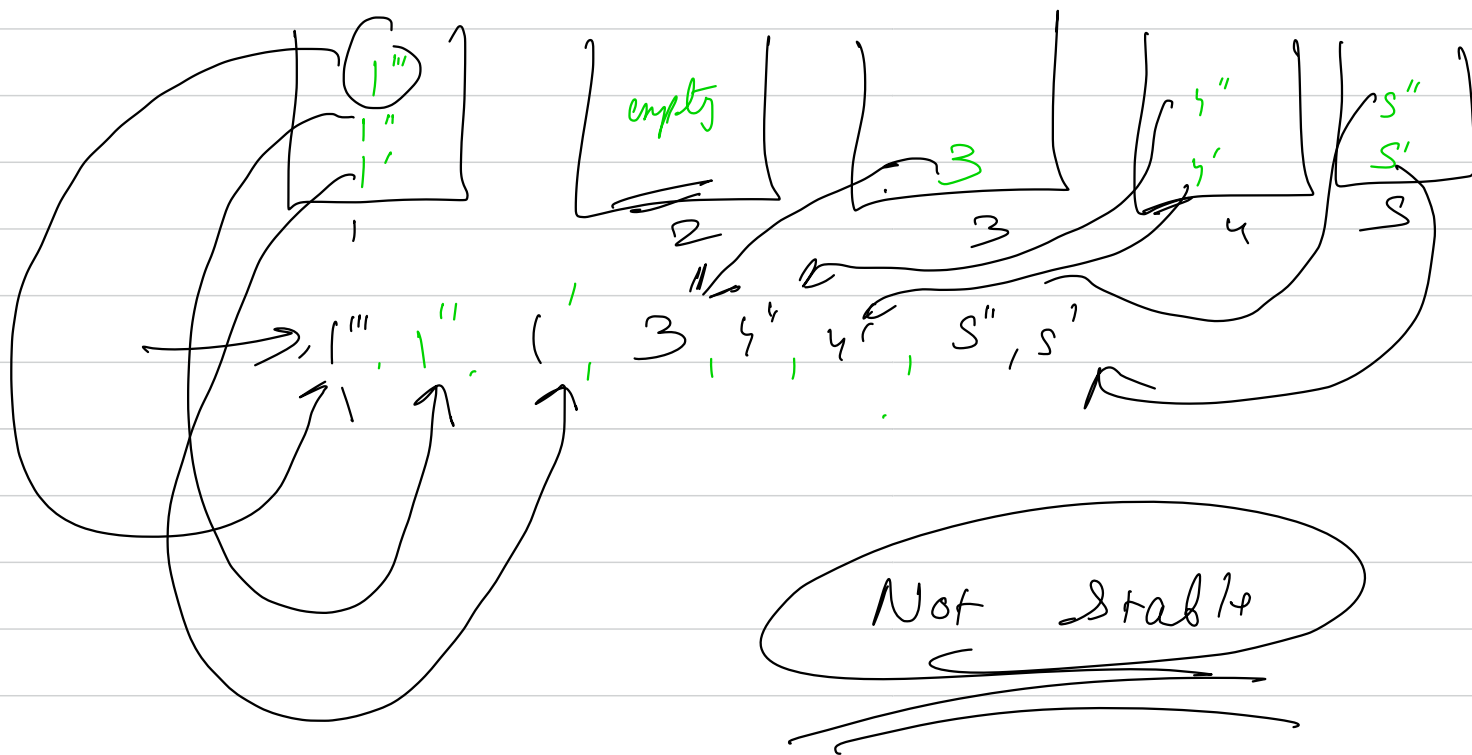
1, 1, 1, 2, 3, 4, 4, 5, 5

min elem — max elem

1, 2, 3, 4

while (mp.get(i) != 0) {

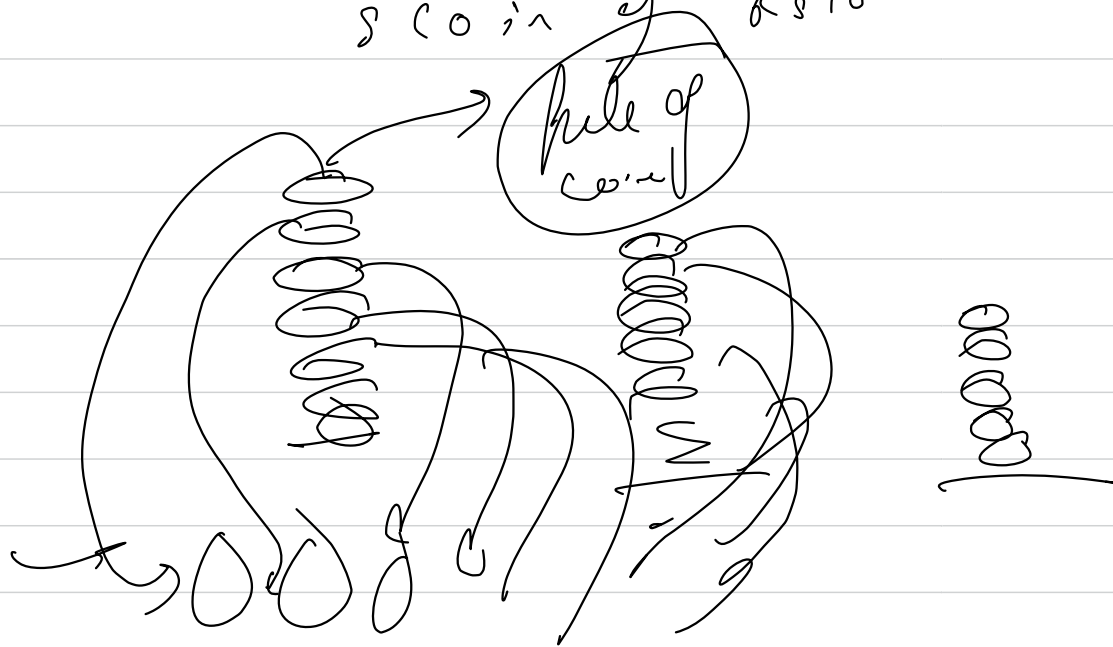
}



Coins → 10 coin of Rs 1

10 coin of Rs 2

8 coin of Rs 10



for $\underline{\underline{n}}$

for $\underline{\underline{max - min}}$ $\rightarrow \underline{\underline{k}}$

$$\underline{\underline{O(n + (max - min))}}$$

$$\underline{\underline{O(n + k)}}$$

→ 1, 4, 1, 2, 2, 5, 2 → arr

0	2	2	0	1	1	0	1	0	0
0	1	2	3	4	5	6	7	8	9

count (last)

max-min

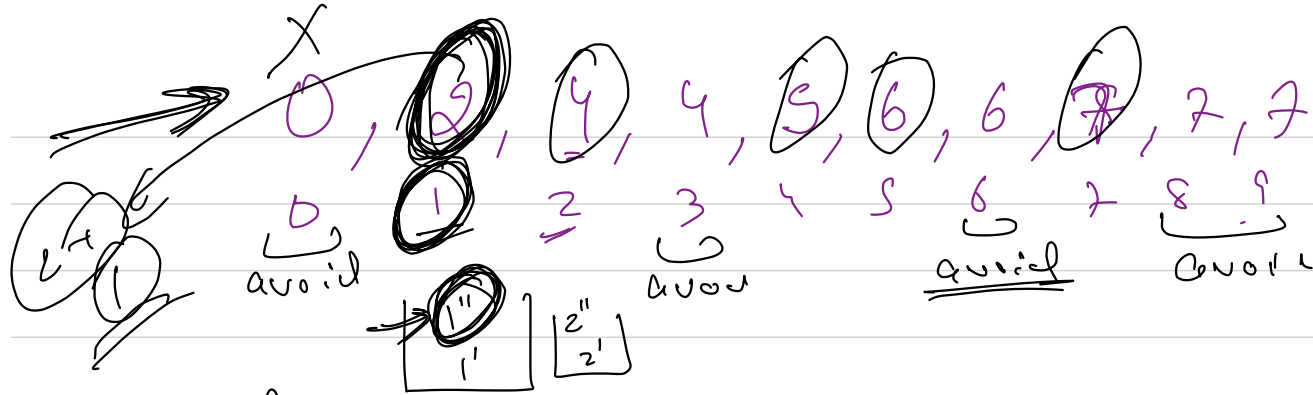
prefer
sum

0, 2, 4, 4, 5, 6, 6, 7, 7 → prefer sum

under 0 1 2 3 4 5 6 7 8 9

Consider 2 → current value is 4 which denotes the position of the last 2 in the sorted list

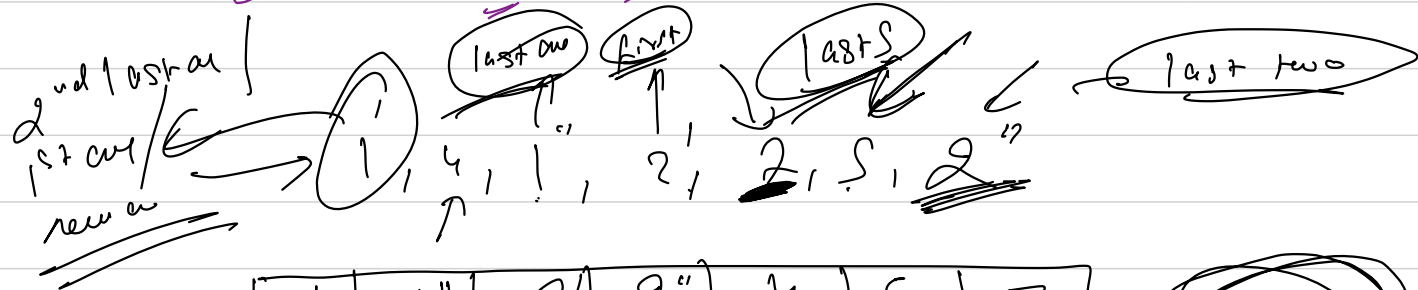
→ value of prefer[i] → the position of the last arr[i] in the sorted form.



→ final sorted list →

1	1	1	2	2	4	5	7	
<u>posim</u>	1	2	3	4	5	6	7	✓✓

0, 1, 3, 4, 5, 6, 7, 7



	1	1	2	2	4	5	7
index →	0	1	2	3	4	5	6

Output

pos → 4

idx 4 - 1 → 3

6 - 1 → 5

2 - 1 → 6

1 - 1 → 0

5 - 1 → 4

3 - 1 → 2

2 - 1 → 1



all positive

stable \rightarrow yes

Inplace \rightarrow NO

Space complexity

$O(n + k)$

$k \rightarrow$ size of elements value

$TC \rightarrow O(n + k)$

Swaps? \rightarrow NO

Comparison \rightarrow no

