


⇒ Agenda → continue the discussion of funcⁿ,
passes by value & references

Recursion

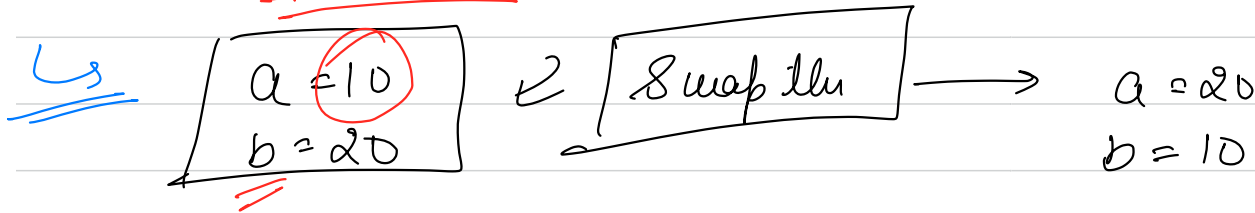
functions

def fun(a, b):
 a = a * 2, b = b - 1
 return sonucu
 # return sonucu parametre

x = 10
y = 20
fun(—, —)

Qⁿ Given 2 numbers m, write a program to

Swap them.



temp = a
→ a = b
b = temp

temp ← 10
a ← 20

def swap(x, y):

Swap (a, b)

→ Does this work for lists also ??

→ When we pass a list, inside a funcⁿ, then it is passed by reference (means actual original list is passed) So whatever change you will do persists.

But in case of int, str, bool, float they are passed by value (means a copy of them is passed) So whatever changes you will do, they won't persist outside the funcⁿ

local copy of a is passed

df

dist
then E

fun

(

$$a_1 = 20$$

$Q = 5$

g to baf(a)

inf

```
print(a)
```

S

$$f_{\text{un}}(a)$$

```
print(a)
```

diff mem
loc

def gum(l):

then
cars come

p.append(20)

li = [1, 2, 3, 4]

print(li) → [1, 2, 3, 4]

gum(li)

print(li) → [1, 2, 3, 7, 20]

[1, 2, 3, 4, 20]

→ in memory

→ lists are passed by reference

Q3

def update(x):
x = "i"

y = "10"
print(y)
update(y)
print(y)

pass by value \rightarrow int, str, bool, float...

pass by reference \rightarrow lists

Recursion

$$y = f(x)$$

o/p
input

$$y' = g(x)$$

o/p
i/p

$$f(g(x))$$

composite funcⁿ → calling another funcⁿ inside a fu

$$y = f(x)$$

$$x = f(x')$$

$$f(f(x'))$$

↓

$$f(f(f(x'')))$$

⋮

$$x' = f(x'')$$

special composite funcⁿ which calls itself inside it
only but with a different parameters.

→ Mathematically, Recursion is the process where we have special composite funcⁿ, where the composition is made such that funcⁿ calls itself inside it only with or without the same parameters.

when it will be infinite?

$$x = f(x)$$
$$\underline{f(f(x))}$$

In programming, recursion is defined as a tool where a function solving a bigger problem, calls itself inside it to solve another problem, till the time we reach to a stage where we have the smallest already solved problem with an extra memory buffer/space

→

Principal of mathematical induction

→ Prove that sum of first n natural no. is $\frac{n(n+1)}{2}$

using PMI

→ for $n=1$, we already know that sum is $\frac{1 \times (1+1)}{2} = 1$

so formula works: → Base Case → Smallest

subproblem for which we already know ans.

→ assume that it works for $n = \underline{k}$ → assumption
→ prove that if it works for $n = k$, then it also works
for $n = \underline{k+1}$ → self works for any k

if → $(k+1)$
↳ first get (k)
if → (k)
get $(k-1)$ if $(k-1)$ $\hookrightarrow k-2$
... if (2)
↳ get (1)

factorial \rightarrow Let's say $f(n)$ is a funcⁿ $5! = 5 \times 4!$
which returns $n!$ $f(s)$

For $n=1$, we know, $f(1) = 1$ or $f(0) = 1$ Base Case
assume $f(n)$ works for $n=k$ $f(k)$ $\xrightarrow{\text{assume}}$ correct value

$$\overset{n=k+1}{f(k+1)} = \underbrace{(k+1)}_1 \times \overbrace{f(k)}^{\text{correct}}$$

$$\begin{aligned} f(5) &= 5 \times \underbrace{f(4)}_{\rightarrow 4 \times \underbrace{f(3)}_{\rightarrow 3 \times \underbrace{f(2)}_{\rightarrow 2 \times f(1)}}} \end{aligned}$$

↳ fibonacci?

0th 1st 2nd 3rd 4th 5th 6th 7th 8th ...
0, 1, 1, 2, 3, 5, 8, 13, 21, ...

any i^{th} term is the sum of previous 2 terms

Let say $f(n)$ → is a funcⁿ that takes n^{th} fibonacci

→ Base case → $f(0) = 0$, $f(1) = 1$

assume
→ for $n = k$ $f(k)$ correctly takes k^{th} fibonacci
for $n = k-1$ $f(k-1)$ correctly takes $(k-1)^{\text{th}}$ fib } correct

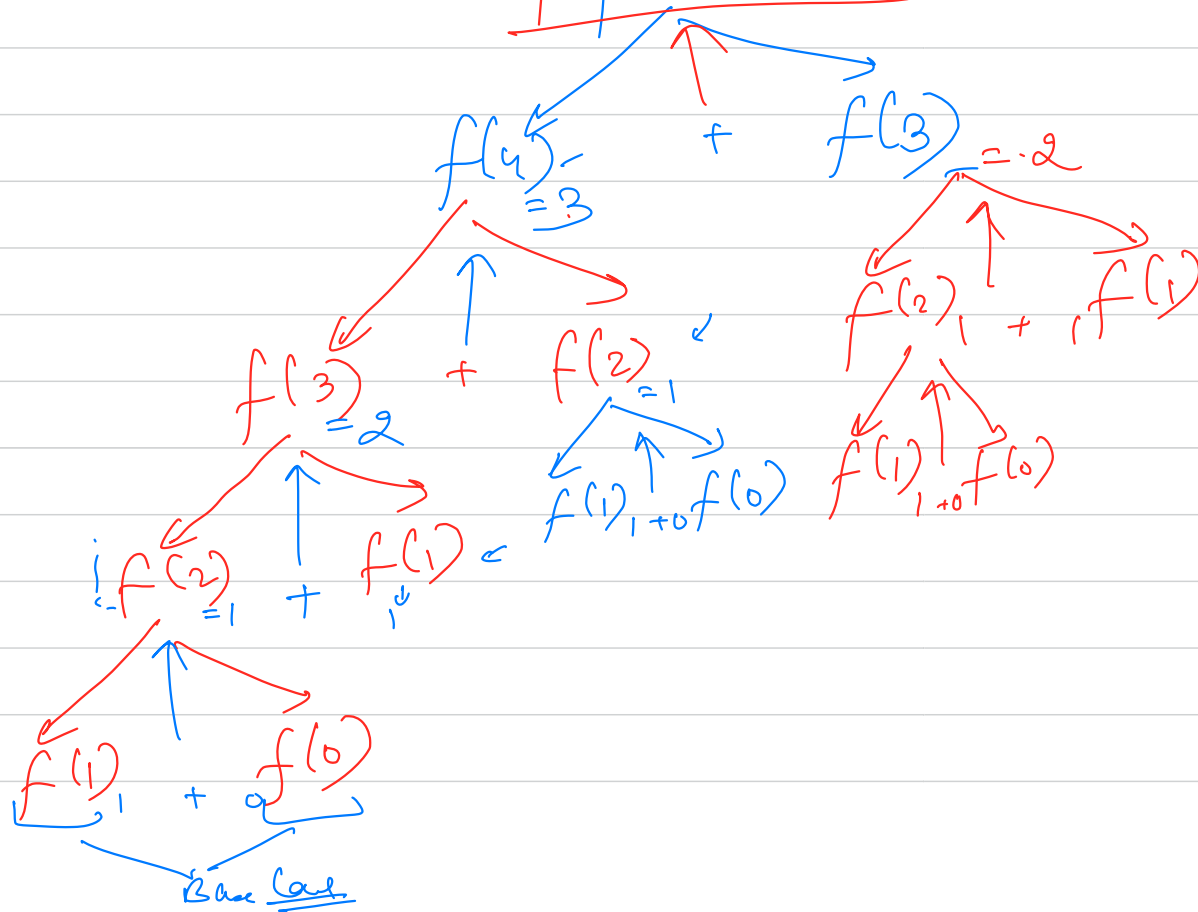
→ self work proves for $n = k+1$

$$f(k+1) = f(k) + f(k-1)$$

so

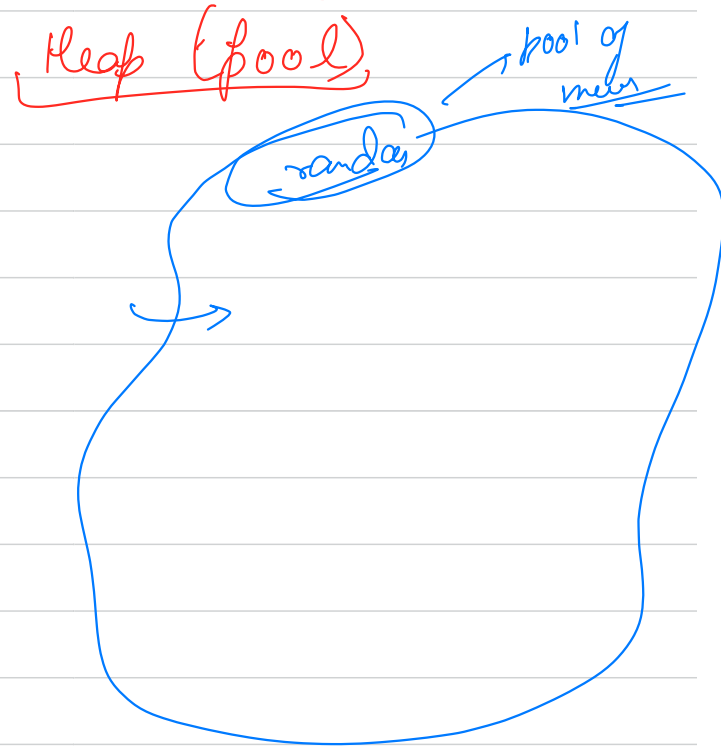
$$f(n) = \underbrace{f(n-1)}_{\text{correctly}} + \underbrace{f(n-2)}_{\text{correctly}}$$

$$f(s) = s$$



Memory has 2 parts (major) There are some minerals

Call stack → linear memory storage



Call stack \rightarrow whenever you call a function, then we add an entry to the top of stack. This entry is called as stack frame. (frames of funcⁿ calls)

Qⁿ what is more in a frame??

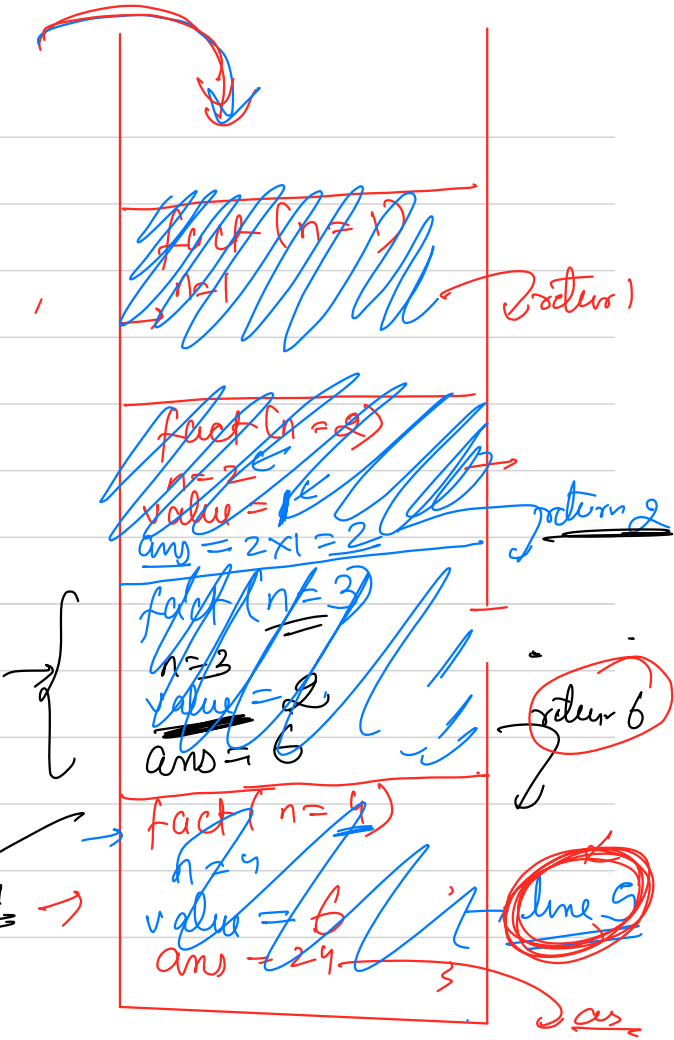
\hookrightarrow all the local variables of a funcⁿ are there in the frames.

```
1 def fact(n):
2     if n == 1: # base case
3         return 1
4
5     value = fact(n - 1) # assume this works correctly
6     ans = n * value # self work ans -> f(n)
7     return ans
8
9
10 print(fact(4))
11
```

27

return → removes the entry of the frame from call stack & returns the value.

call stack →



loops are space optimized, Recursion is not

some algorithms are really easy to implement via recursion!