Hash Tabel
A hash table is a data structure where data is stored in an
associative manner. The data is mapped to array positions by
a hash function that generates a unique value from each key.
Hash Table
Hash Table
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The value stored in a hash table can be searched in O(1) time,
by using the same hash function which generates an address
from the key. The process of mapping the keys to appropriate
locations (or indices) in a hash table is called hashing.
Advantages of Hashing
The main advantage of hash tables over other data structures
is speed. The access time of an element is on average O(1),
therefore lookup could be performed very fast. Hash tables are
particularly efficient when the maximum number of entries can
be predicted in advance.
Hash Functions

A hash function is a mathematical formula which, when
applied to a key, produces a value which can be used as an
index for the key in the hash table.
The main aim of a hash function is that elements should be
uniformly distributed. It produces a unique set of integers
within some suitable range in order to reduce the number of
collisions.
Properties of a Good Hash Function
Uniformity
A good hash function must map the keys as evenly as possible.
This means that the probability of generating every hash value
in the output range should roughly be the same. This also helps
in reducing collisions.
Deterministic
A hash function must always generate the same hash value
for a given input value.

Low Cost
The cost of executing a hash function must be small so that
using the hashing technique becomes preferable over other
traditional approaches.
Applications in Programming
Identification Databases: A hash function can make a unique
signature from never changing data like our Date of Birth.
This can then be used in combination with other variables to
uniquely identify a person.
Search Engines: As the number of pages to be crawled is huge,
a hash function can be used to determine if the page is unique
or it had already been crawled before, without comparing the
contents of the whole webpage.
Different Hash Functions
Division Method
This is the most simple method of hashing. Any integer, for
example, x is divided by a number M and the remainder
obtained is used as the hash.

Generally, M is chosen to be a prime number because	se a prime
number increases the likelihood that the keys are m	apped with
uniformity in the output range of values.	
This function could be represented as:	
$h(k) = k \mod M$	
Multiplication Method	
The Multiplication method has the following steps:	
A constant is chosen which is between 0 and 1, say	y A.
The key k is multiplied by A.	
The fractional part of kA is extracted.	
The result of Step 3 is multiplied by the size of the	hash table
(m).	
This can be represented as:	
h(k) = fractional_part[m(kA mod 1)]	
Mid-Square Method	
The Mid-Square method is as follows:	

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The value of the key is squared. That is, k^2 is found.
The middle r digits of the result are extracted.
The result r is the hash obtained.
The algorithm works well because most or all digits of the key-
value contribute to the resulting hash.
Collisions
Collisions occur when the hash function maps two different
keys to the same location. Two records cannot be stored in the
same location of a hash table normally.
The method used to solve the problem of collisions is called the
collision resolution technique.
There are two popular collision resolution techniques:
Open Addressing
Hash collision resolved by separate chaining
Hash collision resolved by open addressing
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Once a collision takes place, open addressing (also known as
closed hashing computes new positions using a probe sequence
and the next record is stored in that position. There are some
well-known probe sequences:
Linear Probing: The interval between the probes is fixed to 1.
This means that the very next available position in the table
would be tried.
Quadratic Probing: The interval between the probes increases
quadratically. This means that the next available position that
would be tried would increase quadratically.
Double Hashing: The interval between probes is fixed for each
record but the hash is computed again by double hashing.
Chaining
Chaining is another solution to the problem of collisions.
Hash collision resolved by separate chaining
Hash collision resolved by separate chaining
Jorge Stolfi [CC BY-SA 3.0]
In chaining, each location in a hash table stores a pointer to

a linked list that contains all the key values that were hashed
to that location. As new collisions occur, the linked list grows
to accommodate those collisions forming a chain.
This effectively means that each location in the hash table is
not limited to store one value. Searching for a value in a
chained hash table is as simple as scanning a linked list for an
entry with the given key.
Insertion operation appends the key to the end of the linked
list pointed by the hashed location.
Deleting a key requires searching the list and removing the
element.
This solution, however, presents a problem if the linked list
becomes large enough that it takes O(n) time to search one
position. This occurs if the hash table is too small and has to
accommodate many values.