# Assignment 1: Getting to Know Network Traffic

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## 1 Measurement Tools

### 1.1 ping

```
• (myenv) arpit@arpit-linux:-/Desktop/iitd/sem_7/COL334/projects/Getting-To-Know-Network-Traffic$ ping -c 10 -4 google.com PING google.com (172.217.26.110) is 6(84) bytes of data.
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=1 ttl=116 time=36.6 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=2 ttl=116 time=30.5 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=3 ttl=116 time=32.8 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=5 ttl=116 time=27.8 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=5 ttl=116 time=28.2 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=5 ttl=116 time=29.6 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=8 ttl=116 time=29.6 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=8 ttl=116 time=29.5 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=9 ttl=116 time=29.5 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=9 ttl=116 time=29.5 ms
64 bytes from kix05501-in-f14.le100.net (172.217.26.110): icmp_seq=1 ttl=116 time=36.5 ms

--- google.com ping statistics ---
10 packets transmitted, 10 received, 0% packet loss, time 9016ms
rtt min/avg/max/mdev = 27.837/30.543/36.586/3.149 ms
```

Figure 1: Ten Pings to google.com

```
arpit@arpit-linux:~/Desktop/iitd/sem_//COL334/projects/Getting-To-Know-Network-Traffic$ ping -c 10 craigslist.com
PING craigslist.com (208.82.238.135) 56(84) bytes of data.
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=1 ttl=48 time=264 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=2 ttl=48 time=262 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=3 ttl=48 time=265 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=4 ttl=48 time=264 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=5 ttl=48 time=262 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=6 ttl=48 time=262 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=6 ttl=48 time=264 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=8 ttl=48 time=264 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=8 ttl=48 time=266 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=9 ttl=48 time=263 ms
64 bytes from nonorg.craigslist.org (208.82.238.135): icmp_seq=9 ttl=48 time=261 ms
--- craigslist.com ping statistics ---
10 packets transmitted, 10 received, 0% packet loss, time 9013ms
rtt min/avg/max/mdev = 261.003/263.328/266.148/1.586 ms
```

Figure 2: Ten Pings to craigslist.com

- 1. **Protocols Used**: Ping uses ICMP protocol. It sends the Echo Request packet to the destination host and recieves Echo Reply packet from the same. It sits on the third layer of the protocol stack, which is the network layer.
- 2. Latency:

- (a) Avg Latency of Craigslist: 263.328 ms
- (b) Avg Latency of Google: 30.543 ms
- (c) Google's host has **smaller RTT** than Craigslist
- (d) Reason for different latencies of websites:
  - i. Google has more number of hosts than Craigslist, which splits newtwork traffic
  - ii. Google's traffic might directly peer with most of the ISPs in the same tier of internet hierarchy (Regional IPSs).

#### (e) Reason for differnt latencies across pings for same website:

- i. Length of Queue for service at destination host is not constant in time and varies according to the number of user who requested for service before we place any request (queueing of packets, at the host)
- ii. Network congestion is not constant, hence each switch in the network may not have same number of packets it has to route across differnt time
- iii. Different routes may be taken for different pings leading to different paths and hence latencies

#### 3. Using IPv6 for both websites

```
(myenv) arpit@arpit-linux:-/Desktop/iitd/sem_7/COL334/projects/Getting-To-Know-Network-Traffic$ ping -c 10 -6 google.com
PING google.com (2404:6800:4002:831::200e) 56 data bytes
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=1 ttl=116 time=15.7 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=2 ttl=116 time=8.70 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=3 ttl=116 time=5.84 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=3 ttl=116 time=7.43 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=5 ttl=116 time=8.47 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=5 ttl=116 time=8.16 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=7 ttl=116 time=8.16 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=8 ttl=116 time=6.01 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=8 ttl=116 time=5.75 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=8 ttl=116 time=5.06 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=8 ttl=116 time=5.06 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=9 ttl=116 time=5.06 ms
64 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=9 ttl=116 time=5.06 ms
65 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=10 ttl=116 time=5.06 ms
66 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=10 ttl=116 time=5.06 ms
67 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=10 ttl=116 time=5.06 ms
68 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=10 ttl=116 time=5.06 ms
69 bytes from tzdelb-bj-in-x0e.le100.net (2404:6800:4002:831::200e): icmp_seq=10 ttl=116 time=5.06 ms
60 bytes from tzdelb-bj-in-x0e.le100.net
```

Figure 3: Ten Pings to www.google.com using IPv6

```
oarpit@arpit-linux:-/Desktop/iitd/sem_7/COL334/projects/Getting-To-Know-Network-Traffic$ ping -c 10 -6 craigslist.com
ping: craigslist.com: Address family for hostname not supported
```

Figure 4: Error when pinging craigslist.com

- (a) **How to force IPv6**: pass a flag -6 to force ping to follow IPv6
- (b) **Result**: IPv6 was supported by Google's host but not by Craigslist's host
- (c) Why ping 6 Failed for Craigslist's host:
  - i. When checking the IPv6 address for craigslist.com using dig AAAA craigslist.com, my computer does not find any IPv6 address as can be seen from Fig. 5, hence does not know which address to resolve to. Therefore, forcing ping to follow IPv6 Addresing cannot be executed.

ii. Also, since IPv6 worked for Google's host, implies that my computer and Google's host both support IPv6. If a failure of Address Family support has occured it must have occured on Craigslist's server. This implies Craigslist's server does not support IPv6 addresses.

#### 4. Max size of the packets

- (a) Max Size = 65535 bits
- (b) The length of the field "total length" in the packet structure which indicates the size of the payeload, 16 for both IPv4 and IPv6 Addressing in ping packets. Hence, the total payload size =  $2^{16} 1$  (the 1 excludes all bits zero) = 65535 bits. Theoretically, the protocol allows these many number of bits to be sent as data in the payload.

#### 1.2 traceroute

- 1. IPv4 Address for www.google.com : 172.217.26.110
- 2. IPv4 Address for craigslist.com: 208.82.238.135

```
• (myenv) arpit@arpit-linux:-/Desktop/iitd/sem_7/COL334/projects/Getting-To-Know-Network-Traffic$ traceroute -n google.com traceroute to google.com (172.217.26.110), 30 hops max, 60 byte packets
1 10.184.0.13 21.698 ms 21.656 ms 34.829 ms
2 10.255.109.100 21.612 ms 34.777 ms 34.765 ms
3 10.255.109.100 21.612 ms 34.777 ms 34.756 ms
5 10.1.207.607 38.159 ms 39.761 ms 39.751 ms
6 10.1.207.607 38.159 ms 39.761 ms 39.751 ms
6 10.1.200.137 42.984 ms 40.973 ms 40.910 ms
7 10.255.237.94 40.822 ms 10.255.238.122 34.004 ms 10.255.238.254 39.950 ms
8 10.152.7.214 35.379 ms 39.884 ms 40.722 ms
9 72.14.204.62 39.860 ms * *
10 ***
10 ***
11 92.178.86.240 44.019 ms 142.250.235.10 27.200 ms 142.250.227.74 37.167 ms
12 192.178.110.108 37.081 ms 192.178.111.60 37.114 ms 192.178.110.206 28.826 ms
13 142.251.198.3 28.801 ms *
14 192.178.86.224 36.758 ms 209.85.143.186 28.157 ms 192.178.252.124 35.947 ms
15 192.178.86.225 27.800 ms 216.239.62.181 26.212 ms 216.239.54.93 27.922 ms
16 142.251.52.215 33.708 ms 142.251.52.217 29.069 ms 23.618 ms
17 172.217.26.110 32.026 ms 31.378 ms 39.884 ms
```

Figure 5: Trace Route of sending packet to www.google.com

```
** arpit@arpit-linux:-/Desktop/iitd/sem 7/COL334/projects/Getting-To-Know-Network-Traffic$ traceroute 208.82.238.135 traceroute to 208.82.238.135 (208.82.238.135) ab hops max, 60 byte packets

1 10.184 0.13 (10.184 0.13) 3.480 ms 3.487 ms 3.460 ms

2 10.255.109.100 (10.255.109.100) 2.358 ms 2.343 ms 2.394 ms

3 10.255.109.3 (10.255.107.3) 3.406 ms 3.383 ms 3.352 ms

4 10.119.233.65 (10.119.233.65) (10.119.233.65) 3.358 ms 3.319 ms 3.331 ms

5 ***

6 10.119.234.162 (10.119.234.162) 4.165 ms 8.388 ms 8.327 ms

7 59.144.234.93 (59.144.234.93) 8.292 ms 6.017 ms 5.993 ms

8 116.119.57.82 (116.119.57.82) 135.534 ms 116.119.57.86 (116.119.57.86) 134.643 ms 116.119.112.88 (116.119.112.88) 135.513 ms

9 mel-b5-link.ip.twelve99.net (62.115.42.118) 192.451 ms 191.107 ms 194.825 ms

10 ***

1 ael.6.bar2.sanfranciscol.net.lumen.tech (4.69.140.153) 266.800 ms 266.769 ms 267.847 ms

12 craigslist.bar2.sanfranciscol.level3.net (4.53.134.6) 284.747 ms 282.738 ms 283.102 ms

13 nonorg.craigslist.org (208.82.238.135) 266.215 ms 266.202 ms 266.177 ms
```

Figure 6: Trace Route of sending packet to craigslist.com

#### A Number of Hops:

Table 1: Number of Hops for Websites

Host	Number of Hops
www.google.com	17
craigslist.com	13

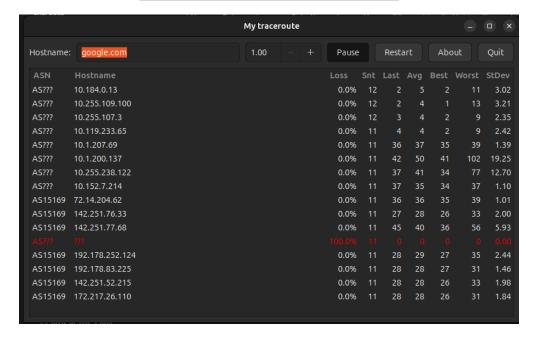


Figure 7: Autonomous System Number (denoted as ASN here) for each IP Address in the trace route of www.google.com (produces using the tool: mtr)

- B Explanation for "\*": Some servers do not cater to traceroute packets, for security reasons and traffic control, hence do not send the Time Exceeded packet back to the source and hence, we do not have information about this node. This is represented in the traceroute by "\*"
- C Multiple IP Addresses for the same Hop Count: traceroute sends three packets for each hop. The three packets may opt for independent routes, depending on the congestion of the network. traceroute lists all the unique ip addresses of the nodes encountered by the three packets
- D The following tables (Table 2 for google.com and Table 3 for craigslist.com) lists the IP Addresses and their correspoding RTTs and Geolocations
  - (a) **For google.com**: most of the network devices that relay the packet are private and hence no information is obtained about them. However we observe a delta in 9th hop, therefore we can assume that the packet has crossed the country.
  - (b) For craigslist.com: The bigger deltas are observed when there is a change in country. For eg., from Table 4, we observe that upto Bangalore the RTT was 5.77 ms but as the relay proceeded to the US, the RTT becomes signficantly higher, 197.05 ms.

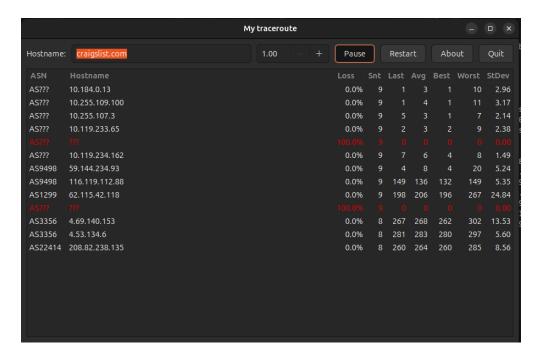


Figure 8: Autonomous System Number (denoted as ASN here) for each IP Address in the trace route of craigslist.com (produces using the tool: mtr)

Table 2: IP Addresses and their GeoLocations for google.com. Note: NA means Not Avail-

able from the respective method listed in the column

Sl No	IP Address	DNS	DNS Geolocation	Maxmind Geolocation	RTT (ms)
1	10.184.0.13	NA	NA	NA	8.457
2	10.255.109.100	NA	NA	NA	9.288
3	10.255.107.3	NA	NA	NA	9.235
4	10.119.233.65	NA	NA	NA	9.184
5	10.1.207.69	NA	NA	NA	37.041
6	10.1.200.137	NA	NA	NA	42.454
7	10.255.238.122	NA	NA	NA	34.988
8	10.152.7.214	NA	NA	NA	37.466
9	*	NA	NA	NA	*
10	*	NA	NA	NA	*
11	72.14.233.58	NA	NA	United States (US), North America	39.810
12	192.178.110.204	NA	NA	United States (US), North America	27.622
13	*	NA	NA	NA	142.251.198.3
14	192.178.252.110	NA	NA	United States (US), North America	31.866
15	216.239.54.93	NA	NA	United States (US), North America	34.751
16	142.251.52.215	NA	NA	United States (US), North America	26.822
		kix05s01-in-f14.1e100.net or		1	
17	172.217.26.110	kix05s01-in-f110.1e100.net or	Osaka Japan or Tanzania	United States (US), North America	27.744
		tzdelb-bj-in-f14.1e100.net			

Hence, from the above explanation, the data intuitively makes sense.

#### E Three Tier Architecture:

- (a) **craigslist.com**: Here we observe the three tier architecture clearly. Since first the packet travels from Delhi to Bangalore (which is a 2nd tier ISP transfer), then the exchange is observed from Bangalore to France and France to San Fransico which are a 1st Tier ISP Transfer. Finally 2nd and 3rd tier transfers occur for the packet to reach craigslist.com server
- (b) **google.com**: Here, we do not observe the three tier architecture clearly. Most of the transfers are through private ip addresses. Google peers its packets in the

Table 3: IP Addresses and their GeoLocations for craigslist.com. Note: NA means Not

Available from the respective method listed in the column

Sl No	IP Address	DNS	DNS Geolocation	Maxmind Geolocation	RTT (ms)
1	10.184.0.13	NA	NA	NA	685.733
2	10.255.109.100	NA	NA	NA	685.597
3	10.255.107.3	NA	NA	NA	685.551
4	10.119.233.65	NA	NA	NA	685.507
5	*	NA	NA	NA	*
6	10.119.234.162	NA	NA	NA	685.376
7	59.144.234.93	NA	NA	Bengaluru, Karnataka, India	5.715
8	116.119.112.88	NA	NA	India	138.551
9	62.115.42.118	mei-b5-link.ip.twelve99.net	Meridian, Mississippi, USA	France	197.050
10	*	NA	NA	NA	*
11	4.69.140.153	ae1.6.bar2.SanFrancisco1.net.lumen.tech	San Francisco, California, United States (US), North America	United States, North America	268.401
12	4.53.134.6	CRAIGSLIST.bar2.SanFrancisco1.Level3.net	San Francisco, California, United States (US), North America	San Francisco, California, United States (US), North America	270.387
13	208.82.238.135	nonorg.craigslist.org	San Francisco, California, United States (US), North America	San Francisco, California, United States (US), North America	259.620

same level of hierarchy (the Regional ISPs). This is the reason why we observe so many regional ISPs (Unites States (US), North America).

## 2 Network Traffic Collection and Analysis

### 2.1 Traffic Capture

A Median Time taken for the DNS request-response to comlete: 42.438s. Note Median was taken on query response times observed in the pcap file when applied with the filter of DNS. (This was done using python script, code present in traffic\_analysis.py)

#### B HTTP

- (a) Number of HTTP Requests = 363 (fitler used: http.request)
- (b) How webpages are structured: The webpage is structured like a tree. The root of the tree is the webpage itself. Root's children are the first level abstraction in the webpage, such as the title, the box that contains green ticks (as present in the website httpvshttps.com) etc.. The next level with each abstraction, contains the green ticks itself, in the case of the abstraction mentioned previously.
- (c) How browsers render complex pages with multiple images and files: The browser recursively calls on the nodes of the tree (mentioned above) and fetches the required information, and displays them according to their HTML Code. This way it is able to render complex images and texts.

#### C TCP Connection

- (a) Number of TCP Connections = 31 (filter used: tcp.flags.syn == 1 and tcp.flags.ack == 0)
- (b) Are Number of HTTP Conbnections==Number of TCP Connections: No, they are generally not equal, however they are related. In one TCP Connection multiple HTTP requests can be made.

(c) Content Object Fetch over same TCP Connection: Yes, some contents gets fetched over the same TCP Connection. This can be supported from the fact that HTTP transfer are made with HTTP/1.1, which implies sustained TCP Connection for multiple HTTP transfers. This was verified using the filter tcp.stream == 0 where we checked the upstream flow. This showed multiple to and fro packet flows.



#### D HTTPs packet trace

- (a) Is HTTP traffic there? No, there is no HTTP Traffic (filter used: http)
- (b) Content transfer of HTTP and JavaScript files? and why?: There is no content transfer of HTTP and JavaScript. HTTPs is secured data transfer, therefor the file content is encrypted and hence not visible without the key to the file. This is the reason why wireshark does not show this content.
- (c) dns traffic: Yes, DNS Traffic is present. DNS is required for look ups to convert web addresses to their correspoding IP Addresses. Once the ip addresses for the websites are resolved no more DNS Traffic is present for that particular domain name. Note: I was able to see DNS traffic because DoH, i.e., DNS over HTTPs is not enabled on my browser. Hence, DNS Traffic could not tunnel through HTTPs. (filter used: dns)
- (d) Number of tcp connections logged = 7

(e) Is Number of TCP Conbnections in HTTPS case == Number of TCP Connections in HTTP case?: They are not equal. However the HTTPS case has smaller number of TCP Connections. A potential reason is multiplexing, allowing many HTTPS request to be sent simultaneously. Also, apart from TCP Handshake, https use TLS Handshake, which is expensive and hence tend to perform this Handshake lesser number of times. Due to this lesser number of TCP Connections are established.

### 2.2 Performance Analysis

#### E HTTP vs HTTPs

- (a) Comparision of time taken for downloads: HTPP::17.132s and HTTPs::1.614 (93% faster than HTTP)
- (b) Observation from the plots:
  - i. The download plot for https.pcap has significant throughput for download in comparision to that for http.pcap. Since the amount of data downloaded is the same, hence if the time of download reduces (as inferred from the time on the website and also the x axis of the plots), this implies throughput increases. The download throughput are upto 100 times more in https than http
  - ii. Time of downloads can also be justified from the RTT plot. RTT is much smaller in https.pcap case than in comparision to http.pcap. This implies smaller time for data transfers from https.pcap than http.pcap
- (c) Assumption for RTT Plots: We assume no retransmission of signal. Hence, we confirm ACK from the receiver by matching it to the estimated Acknowledgment Number, that is, we find if A == S + L. Here A is the Acknowledgment number received from the receiver, S is the sequence number when the message in consideration was sent to the receiver from the sender and L is the number of bytes of the data.
- (d) **Note**: RTT values are recorded at the time of ACK of the sent message
- (e) **Note**: The images for throughputs have been found for each second 1, 2, 3, ... on the wall clock

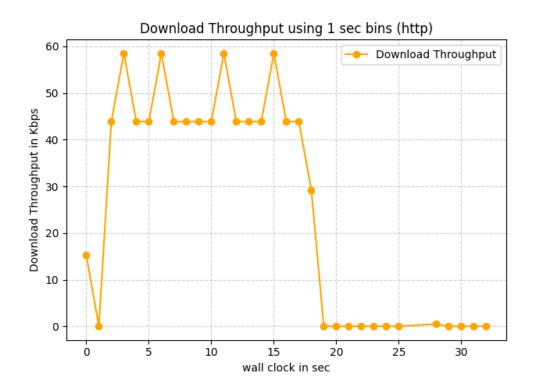


Figure 9: Download Throughput using http.pcap

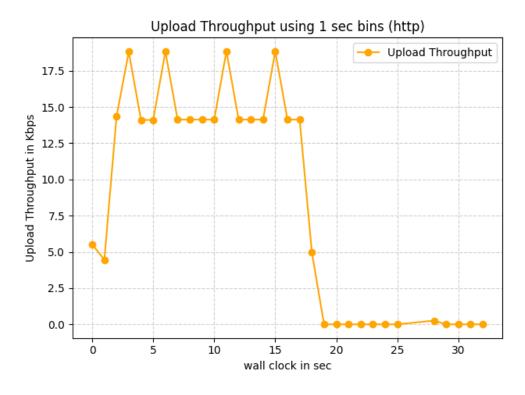


Figure 10: Upload Throughput using http.pcap

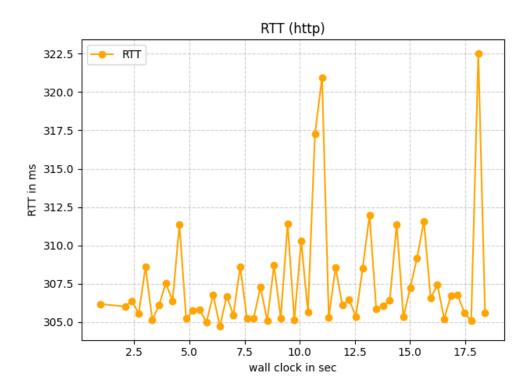


Figure 11: RTT using http.pcap

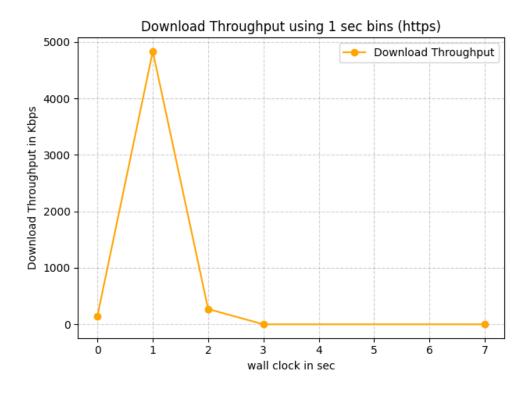


Figure 12: Download Throughput using https.pcap

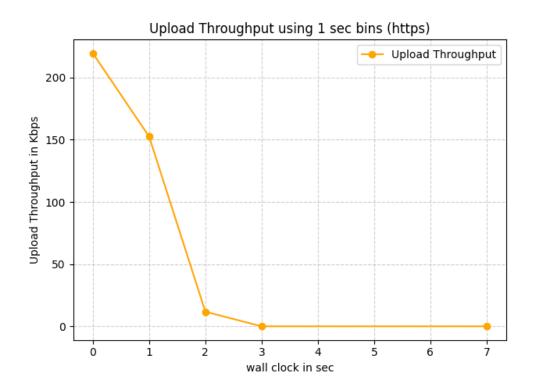


Figure 13: Upload Throughput using https.pcap

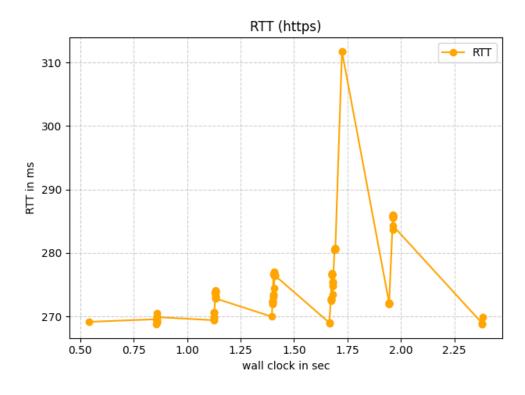


Figure 14: RTT using https.pcap