Arpita Singh

231071005

Experiment Task 1

1 — Algorithm (Divide & conquer)

```
merge and Count (arr, temp_arr, left, mid, right)
Initialise i,j,k and inversion_count to 0
 If  arr[i] <= arr[j]
     temp_arr[k] = arr[i]
     i++
 Else
  temp_arr[k] = arr[j]
  Increment inv_count by (mid - i + 1)
  j++
k++
  copy remaining element of left subarray, if any
  copy remaining element of right subarray, if any
Return inversion count.

 merge_sort_and_count(arr, temp_arr, left, right)
If left < right
  Find mid = (left + right)/2
 Count left inv = merge_sort_and_count(arr, tem_arr, left, mid)
 count right inv = merge_sort_and_count(arr, temp_arr, mid+1, right)
 count split inv = merge and Count (arr, temp_arr, left, mid, right)
 Return (left + right + split)
```

— Algorithm (Brute force)

```
Initialize inv_count to 0
 For i from 0 to n-1:
     For j from i+1 to n:
     If  arr[i] > arr[j], increament inv_count
Return inv_count
```

Test Case

2] Positive test Cases

① Input : 165370
Expected output : 8

② Input : 373708
Expected output : 5

③ Input : 592290
Expected output : 9

④ Input : 686641
Expected output : 11

⑤ Input : 827581
Expected output : 9

Negative test Cases

① Input : 123abc
Expected Output : Error (invalid input)

② Input : 5
Expected output : Error (single digit input)

③ Input : 111111
Expected output : 0 as all digits are same

④ Input : -123567
Expected output : Error (as negative input)

⑤ Input : 1234
Expected output : Error (source code should be six digit only)

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3] Time Complexity

A] Brute force Algorithm

Outer loop, loop runs n times, for
where n is no. of elements
Inner loop runs n-i times

when i=0, inner loop runs n-1 times
when i=1, inner loop runs n-2 times
⋮
when i=n-1, inner loop runs 0 times

Total $(n-1)+(n-2)+(n-3)+\dots+1+0 = \frac{n(n-1)}{2} \approx \frac{n^2}{2} = O(n^2)$

∴ Brute force time complexity $O(n^2)$

B] Divide and Conquer Algorithm

merge step takes linear time O(n) for merging two halves
— each level of recursion processes n element
— height of recursion tree is log n

Total time complexity

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

From masters theorem
$a=2, b=2, f(n)=O(n)$ $\quad n^{\log_2 2} = n$
∴ $T(n) = O(n \log n)$
and also it is piggy back on merge sort $O(n \log n)$

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Experiment Task - 2

1] - Brute force Algorithm

// Input : Two integers x and y
Initialization : Initialize variable result to 0
Loop:
For each digit in y :
- Multiply x by current digit of y (from least significant to most significant)
- Shift the result acc to position of the current digit in y (using power of 10)
- Add the shifted result to result

~~Outp~~ return result
//Output : result of multiplication

- Karatsuba Algorithm (Divide and conquer)

//Input : Two integer x and y
Base Case : If either number has only one digit, return product
Split : For two n digit numbers, split in 2 halves
a and b are high and low part of x
c and d are high and low part of y
$x = a \cdot 10^m + b$, $y = c \cdot 10^m + d$ $(m = \lceil n/2 \rceil)$

Recursive calls :
- Compute a.c (Z_0 = karatsuba(a,c))
- compute b.d (z_1 = karatsuba(b,d))
- compute (a+b).(c+d) (z_2 = karatsuba(a+b, c+d))

Combine results :
- Use formula

result = $z_0 10^{2m} + (z_2 - z_1 - z_0) \cdot 10^m + z_1$

return result

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Test Cases

2]

Positive test cases

① Input : (123, 456)
Expected output : 56088

② Input : (12, 34)
Expected output : 408

③ Input : (0, 1237)
Expected output : 0

④ Input : (721, 314)2
Expected output : 226394

⑤ Input : (+3, 67)
Expected output : 2881

Negative Test Cases

① Input : (None, 100)
Expected output : Error (None is not valid input type)

② Input : (12, 'abc')
Expected output : Error (Invalid input type)

③ Input : (10.5, 2)
Expected output : Error (Floating-point numbers are invalid)

④ Input : (-123, 456)
Expected output : -56088 (or error if neg multiplication not allowed)

⑤ Input : (0, 0)
Expected output : 0 (not error but multiplication with 0)

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Time Complexity - Task 2

3]

A) Brute force algorithm

outer loop: iterates over each digit of y

inner loop: $O(n)$ for multiplying with single digit

$O(n)$ for shifting

Total cost is $O(n) + O(n) = O(n)$

outer loop runs n times

Each iteration takes $O(n)$

Total time complexity =

$$T(n) = n \times O(n) = O(n^2)$$

so brute force time complexity is $\underline{O(n^2)}$

B) Divide and conquer - karatsuba

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$3T(n/2)$ as three recursive multiplication of half-sized number

$O(n)$ for combining of result

Masters theorem -

$a = 3 \quad b = 2$

so $n^{\log_b a} = n^{\log_2 3} = n^{1.585}$

∴ time complexity of karatsuba algorithm is $O(n^{\log_2 3}) \approx \underline{O(n^{1.585})}$