Lesson 12c: Backtracking Introduction

Backtracking

Backtracking is an algorithmic technique for solving problems recursively by trying to build a solution incrementally. It incrementally builds candidates to the solutions and abandons a candidate ("backtracks") as soon as it determines that the candidate cannot lead to a valid solution (the candidate does not satisfy the constraints of the problem). Each candidate solution is only evaluated once.

Ex. A Sudoku being solved by backtracking. Each cell is tested for a valid number, moving "back" when there is a violation, and moving forward again until the puzzle is solved.

Difference between Recursion and Backtracking:

- In recursion, the function calls itself until it reaches a base case.
- In backtracking, we use recursion to explore all the possibilities until we get the best result for the problem.

Backtracking

Backtracking works incrementally, and it is an improvement to the naïve or brute force approach (where all configurations are generated and tried).

The backtracking technique searches for a solution to a problem among all the available options.

PseudoCode for Backtracking:

```
boolean findSolutions(n, other params) :
   if (found a solution):
        displaySolution()
        return True

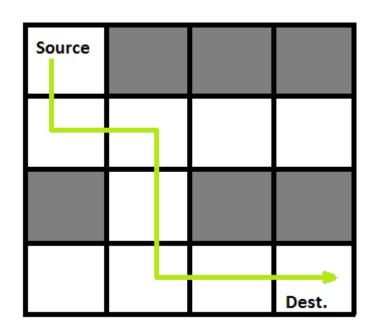
for (val = first to last) :
        if (isValid(val, n)):
            applyValue(val, n)
        if (findSolutions(n+1, other params))
            return True
        removeValue(val, n)
        return False
```

A Maze is given as N*N binary matrix of blocks where source block is maze[0][0] and destination block is maze[N-1][N-1]. A rat starts from source and has to reach the destination. The rat can move only in vertical and horizontal directions

In the maze matrix, 0 means the block is a dead end and 1 means the block can be used in the path from source to destination.

maze =
$$[[1, 0, 0, 0],$$

 $[1, 1, 1, 1],$
 $[0, 1, 0, 0],$
 $[1, 1, 1, 1]]$



Approach: Implement a recursive function, which will follow a path and check if the path reaches the destination or not. If the path cannot reach the destination then **backtrack** and try other paths.

Steps:

- 1. Create a solution matrix, initially filled with 0's.
- 2. Create a recursive function, which takes initial matrix, output matrix and position of rat (i, j).
- 3. If the position is out of the matrix or position not valid then discard it and return.
- 4. Mark the position output[i][j] as 1 and check if the current position is the destination or not. If destination is reached print the output matrix and return.
- 5. Recursively search for next position (movement)
- 6. Unmark position (i, j), i.e output[i][j] = 0 (backtracking)

```
def solveMaze(maze):
    # Creating a n x n matrix
    res = [[0 for i in range(n)] for i in range(n)]
    res[0][0] = 1
    move_x = [-1, 1, 0, 0] # x matrix for each direction
    move y = [0, 0, -1, 1] # y matrix for each direction
    if RatMaze(n, maze, move x, move y, 0, 0, res):
        # print solution
        for i in range(n):
            for j in range(n):
                print(res[i][j], end=' ')
            print()
    else:
        print('Solution does not exist')
# MAIN
n = 4 # Maze size
maze = [[1, 0, 0, 0],
        [1, 1, 0, 1],
        [0, 1, 0, 0],
        [1, 1, 1, 1]
solveMaze(maze)
```

res = []
for i in range(n):
 res.append([0]*n)

```
# To check if x, y is a valid index for N * N Maze
def isValid(n, maze, x, y, res):
    if 0 \le x \le n and 0 \le y \le n and maze[x][y] == 1 and res[x][y] == 0:
        return True
    return False
# A recursive function to solve Maze problem
def RatMaze(n, maze, move x, move y, x, y, res):
    if x == n-1 and y == n-1: # if (position (x,y) is goal/destiny)
        return True
    for i in range(4):
        x \text{ new} = x + \text{move } x[i] \# \text{ Generate new value of } x
        y new = y + move y[i] # Generate new value of y
        # Check if maze[x][y] is valid
        if isValid(n, maze, x_new, y_new, res):
            res[x new][y new] = 1 # mark x, y as part of solution path
            if RatMaze(n, maze, move_x, move_y, x_new, y_new, res):
                 return True
            res[x new][y new] = 0 # backtracking
    return False
```

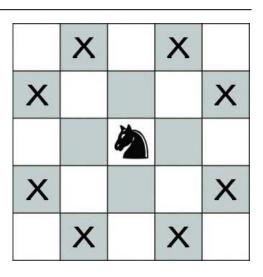
Exercise - Knight's Tour

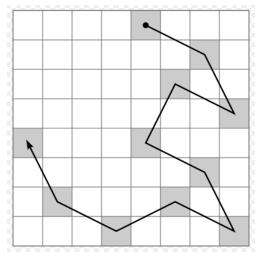
The "Knight's Tour" is an ancient puzzle in which the objective is to move a knight, starting from one square on a chessboard, to every other position, landing on each square only once.

Problem:

N*N board and the Knight placed on the first position of an empty board. Moving according to the rules of chess, the knight must visit each square exactly once. Print the order of each cell in which they are visited.

<pre>Input : N = 8 Output:</pre>							
0			33	30	17	8	63
37	34	31	60	9	62	29	16
58	1	36	39	32	27	18	7
35	48	41	26	61	10	15	28
42	57	2	49	40	23	6	19
47	50	45	54	25	20	11	14
56	43	52	3	22	13	24	5
51	46	55	44	53	4	21	12





Exercise - Knight's Tour

Naive Algorithm for Knight's tour

The Naive Algorithm is to generate all tours one by one and check if the generated tour satisfies the constraints.

```
while there are untried tours
{
   generate the next tour
   if this tour covers all squares
   {
      print this path;
   }
}
```

Exercise - Knight's Tour

We start from an empty solution vector and add items one by one (an item is a Knight's move). When we add an item, we check if this item violates the problem constraint. If it does then we remove it and try other alternatives.

If none of the alternatives works out then we go back (backtracking) and remove the item added in the previous stage. If we reach the initial stage, then no solution exists. If adding an item doesn't violate constraints then we recursively add items one by one. If the solution vector becomes complete, then we print the solution.

If all squares are visited print the solution

Else

- a) Add one of the next moves to solution vector and recursively check if this move leads to a solution. (A Knight can make 8 moves. Choose one of them).
- b) If the move chosen in the above step doesn't lead to a solution, then remove this move from the solution vector and try other alternative moves.
- c) If none of the alternatives work then return false (Returning false will remove the previously added item in recursion. If false is returned by the initial call of recursion then "no solution exists")