

Quantum Error-Correcting Codes in QuantumClifford.jl

鄢语轩@IIIS, Tsinghua GSoC mentor: Stefan Krastanov

Outline

- Stabilizers
- Quantum error correction
- My contribution: random circuit codes, lifted product codes

Stabilizers

Quantum states are generally not classical simulatable (unless BPP=BQP).

Yet, not all of them.

Stabilizer states: +1 or -1 eigenstates of some Pauli strings, called its stabilizer group.

The stabilizer group of a pure state has n generators.

Describe the state by those generators.

Generators (Pauli strings) are Bool vectors.

```
using QuantumClifford
     st = ghz(3)
[6]:
[6]:
     + XXX
     + ZZ
     + ZZ
     stab to gf2(st)
[8]:
[8]: 3x6 Matrix{Bool}:
```

Clifford

Clifford operations transform a Pauli string always to another Pauli string.

-> Transform a stabilizer state to another stabilizer state.

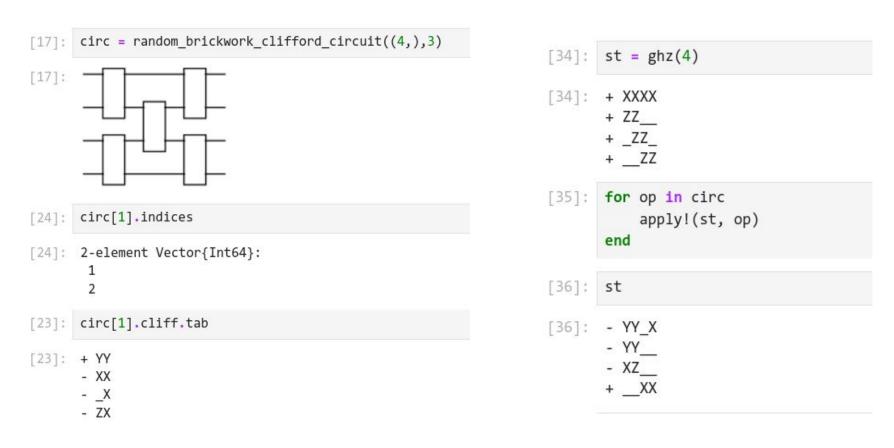
By transforming each generator.

- "Stabilizer states + Clifford" are classically simulatable.
- -> QuantumClifford.jl

```
[7]: cl = random_clifford(3)
[7]: X₁ → - Y
      X_2 \mapsto - YZ
      X_3 \mapsto + X
      Z_1 \mapsto - XX
      Z_2 \mapsto + X
      Z_3 \mapsto + Y
     cl * st
[8]: + ZX
      + XY
```

Clifford circuit

Sequence of sparse Clifford operations (gates).



Quantum error correction

Errors happen and destroy information.

Solution: quantum error correction.

Idea: encode information into a subspace.

Encode: k logical qubits -> subspace of n qubits

Decode: -> k logical qubits

(Similar to classical error correction, where qubits -> bits, space are binary.)

Stabilizer codes

Stabilizer codes: the subspace is stabilized by (n-k) Pauli generators.

Syndrome, parity checks: measure some of stabilizer group elements.

The left k generators are taken as logical operators.

Parameters [[n,k,d]]:

- n: size
- k: dimension
- d: distance (related to maximal correctable errors)

```
[38]: code = Steane7()

[38]: Steane7()

[39]: code_n(code)

[39]: 7

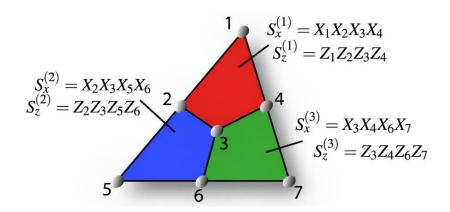
[40]: code_k(code)

[40]: 1
```

Parity checks

Choose (n-k) elements from stabilizer group as parity checks.

CSS code: if parity checks can be either "all X" or "all Z."



```
[43]:
      parity checks(code)
[43]:
      + XXXX
      + XX XX
      + X X X X
      + ZZZZ
      + ZZ ZZ
      + Z Z Z Z
[41]: parity checks x(code)
[41]: 3×7 Matrix{Bool}:
      parity_checks_z(code)
[42]: 3×7 Matrix{Bool}:
```

Simulating codes

Pipeline: encoding

- -> errors
- -> syndrome (parity checks)
- -> decoding
- -> correction

Everything in Clifford and QuantumClifford.jl!

```
@testset "belief prop decoders, good for sparse codes" begin
    codes = vcat(LP04, LP118, test gb codes, other lifted product codes)
    noise = 0.001
    setups = [
        CommutationCheckECCSetup(noise),
        NaiveSyndromeECCSetup(noise, 0),
        ShorSyndromeECCSetup(noise, 0),
    for c in codes
        for s in setups
            for d in [c -> PyBeliefPropOSDecoder(c, maxiter=2)]
                nsamples = 10000
                if true
                    @test broken false # TODO these are too slow to test in CI
                    continue
                end
                e = evaluate_decoder(d(c), s, nsamples)
                # @show c
                # @show s
                # @show e
                @test max(e...) <= noise</pre>
            end
```

Better (and larger) codes?

New code construction (GSoC 2024)

Skip it because too textbook

Concatenated quantum code #289

Krastanov merged 6 commits into QuantumSavory:master from royess:concat [on Jun 14 **}**→ Merged

add random Clifford circuit codes #298

}→ Merged

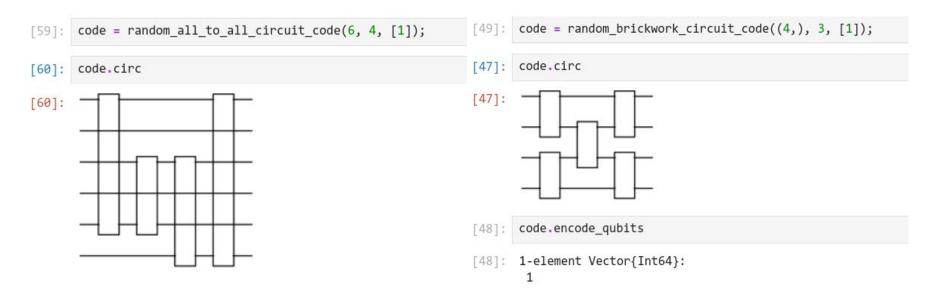
Lifted and lifted product codes via Hecke's GroupAlgebra #356

Krastanov merged 70 commits into QuantumSavory:master from royess:lift-dev-hecke r□ on Sep 27 **}**→ Merged

Random Clifford circuit codes

Idea: to encode a logical qubit, we need to spread it to the whole system.

Use random circuit as encoding circuit.



Lifted product codes

Quantum LDPC (QLDPC) code.

Idea: combine two classical codes into a quantum code.

Hypergraph product.

The hypergraph product yields two classical codes $C_{X,Z}$ with parity-check matrices

$$H_X = \begin{pmatrix} H_1 \otimes I_{n_2} & I_{r_1} \otimes H_2^T \end{pmatrix} \tag{1}$$

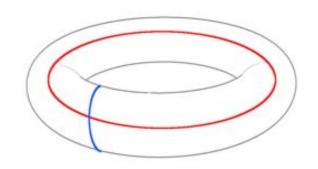
$$H_Z = \begin{pmatrix} I_{n_1} \otimes H_2 & H_1^T \otimes I_{r_2} \end{pmatrix}, \qquad (2)$$

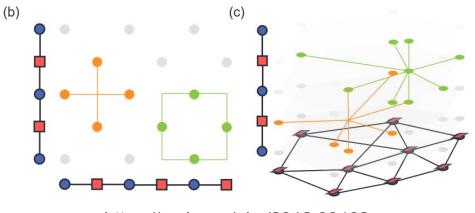
where I_m is the m-dimensional identity matrix. These two codes then yield a hypergraph product code via the CSS construction.

https://errorcorrectionzoo.org/ c/hypergraph_product

Geometrical view of hypergraph product

Canonical example: repetition code ⊗_{HG} repetition code -> toric code





https://arxiv.org/abs/2312.08462

Lifted product

Motivation: save qubits to make better codes. (Asymptotically good QLDPC codes.)

1. Use matrices of group algebra elements in hypergraph product.

$$egin{aligned} H_X &= ig(H_1 \otimes I_{n_2} & I_{r_1} \otimes H_2^T ig) \ H_Z &= ig(I_{n_1} \otimes H_2 & H_1^T \otimes I_{r_2} ig) \ , \end{aligned}$$

2. Then, lift each group algebra element to its binary matrix representation (e.g., permutation group representation).

Group algebra, group ring: group elements with coefficients, e.g., $g_1 + g_2$.

$$(g_1 + g_2) * g_3 = g_1 * g_3 + g_2 * g_3$$

Construction by group algebra (Hecke)

B1) [[882, 24, d]] code, $18 \le d \le 24$. The matri H_X and H_Z are (3,6)-regular ($\ell = 63$).

$$A = \begin{pmatrix} x_{54}^{27} & 0 & 0 & 0 & 0 & 1 & x^{54} \\ x_{54}^{54} & x_{54}^{27} & 0 & 0 & 0 & 0 & 1 \\ 1 & x_{54}^{54} & x_{54}^{27} & 0 & 0 & 0 & 0 \\ 0 & 1 & x_{54}^{54} & x_{54}^{27} & 0 & 0 & 0 \\ 0 & 0 & 1 & x_{54}^{54} & x_{54}^{27} & 0 & 0 \\ 0 & 0 & 0 & 1 & x_{54}^{54} & x_{54}^{27} & 0 \\ 0 & 0 & 0 & 0 & 1 & x_{54}^{54} & x_{54}^{27} \end{pmatrix},$$

$$B = (1 + x + x^{6})I_{7}.$$

https://doi.org/10.22331/q-2021-11-22-585

Cyclic group of order I=63

```
[69]: import Hecke: group_algebra, GF, abelian_group, gens
          import LinearAlgebra: diagind
   [70]: 1 = 63
         GA = group_algebra(GF(2), abelian_group(1))
         x = gens(GA)[]
          Element of Group algebra of group of order 63 over GF(2) with
[71]: A = zeros(GA, 7, 7);
        A[diagind(A)] = x^27;
        A[diagind(A, -1)] = x^54;
         A[diagind(A, 6)] = x^54;
         A[diagind(A, -2)] := GA(1);
         A[diagind(A, 5)] = GA(1);
          B = reshape([1 + x + x^6], (1, 1));
   [72]: c1 = LPCode(A, B);
         code n(c1), code k(c1)
   [73]: (882, 24)
```

Lesson: reinvent the wheel (group algebra)

Lifted and lifted product codes #312

10



Important but not covered in this talk

- Decoders, e.g., BP-OSD.
- Entanglement.

- Logical gates.
- T gates and magic.
- Mixed states, destabilizers.

More development experiences.

QuantumClifford.jl







