

# Advanced Programming in the UNIX Environment — *Using UNIX*

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# UNIX Architecture I

- ▶ In a strict sense, an **operating system** can be defined as the software that controls the hardware resources of the computer and provides an environment under which programs can run. Generally, we call this software the **kernel**. Figure 1 shows a diagram of the UNIX System architecture.



## UNIX Architecture II

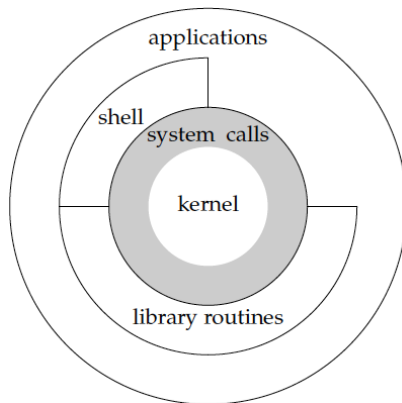


Figure: Architecture of the UNIX operating system



## UNIX Architecture III

- ▶ The interface to the kernel is a layer of software called the **system calls** (the shaded portion in Figure). **Libraries** of common functions are built on top of the system call interface, but **applications** are free to use both. The **shell** is a special application that provides an interface for running other applications.
- ▶ In a broad sense, an operating system consists of the kernel and all the other software that makes a computer useful and gives the computer its personality: system utilities, applications, shells, libraries of common functions, and so on.



# Login

- ▶ Before login into a UNIX system, you should have a valid account in that host.
- ▶ Install an Ubuntu/FreeBSD system on your own PC is a nice choice.
- ▶ For detail about using a UNIX (Linux) operating system please refer another lecture slides at room 2-301 13:30-15:20 each Friday in this semester.
- ▶ If you choose to installing a UNIX os on your own PC, you can operate under the console.
- ▶ If you choose to using the account that someone else provided to you, you must login the host through ethernet by an **ssh** client with the login account and the password provided by the owner of the host.



# Shell

- ▶ A *shell* is a command interpreter and a big UNIX program too that reads user input and executes commands.
- ▶ There are several popular shells like **bsh**, **csh**, **ksh**, etc.
- ▶ The last field of file **/etc/passwd** determine which kind of shell should be executed after login.
- ▶ The most frequently used shell under Linux and MacOS is **bash**. FreeBSD and Solaris often use **csh**.
- ▶ The shell was standardized in the POSIX 1003.2 standard. The specification was based on features from the Korn shell and Bourne shell.
- ▶ All examples in this book were tested under the Bourne shell, the Korn shell, and the Bourne-again shell environments.



# Files and Directories I

- ▶ There must be only one logical filesystem in an UNIX OS. The UNIX filesystem is a hierarchical arrangement of files.
- ▶ Files are placed in various directories. A directory is a special file that contains directory entries.
- ▶ All directories are arranged like a upset tree. The top most directory named **root**.
- ▶ The entries' name in a directory are called filenames. The only two characters that cannot appear in a filename are the slash character(/) and the null character.
- ▶ Two filenames are automatically created whenever a new directory is created: `.` and `..`.
- ▶ Dot refers to the current directory and dot-dot refers to the parent directory.





## Files and Directories II

- ▶ A sequence of zero or more filenames, separated by slashes, and optionally starting with a slash, forms a **pathname**. A pathname leading by a slash is named an **absolute pathname**, otherwise it's called a **relative pathname**.
- ▶ Example (Figure 1.3, `intro/ls1.c`).
- ▶ **Working Directory** is the directory from which all relative pathnames are interpreted.
- ▶ The second last field of file `/etc/passwd` denotes the **Home Directory** of each user. When we login successfully, Home Directory is our Working Directory.



# Input and Output

- ▶ UNIX kernel uses small nonnegative integers to identify the files being accessed by a particular process, named **file descriptors**.
- ▶ By default, all shells open three descriptors wherever a new program is run: **standard input**, **standard output** and **standard error**.
- ▶ All these three file are connected to your terminal by default.
- ▶ Example (Figure 1.4, `intro/mycat.c`).
- ▶ Example (Figure 1.5, `intro/getcputc.c`).



# Programs and Processes I

- ▶ A **program** is an executable file residing in a disk file. A program is read into memory and executed by the kernel as a result of one of the six **exec** functions.
- ▶ An executing instance of a program is called a **process**.
- ▶ Every UNIX process is guaranteed to have a unique nonnegative integer as numeric identifier called **process ID**.
- ▶ Example (Figure 1.6, `intro/hello.c`).
- ▶ Process ID 0 is usually scheduler process and is often know as the **swapper**. Process ID 1 is usually the **init** process and is invoked by the kernel at the end of the bootstrap procedure.
- ▶ There are three primary functions used for process control: **fork**, **exec** and **waitpid**.
- ▶ Example (Figure 1.7, `intro/shell1.c`).



## Error Handling I

- ▶ When an error occurs in one of the UNIX System functions, a negative value is often returned, and the integer `errno` is usually set to a value that gives additional information.
- ▶ Most functions that return a pointer to an object return a null pointer to indicate an error.
- ▶ The file `errno.h` defines the symbol `errno` and constants for each value that `errno` can assume.
- ▶ The first page of Section 2 of the UNIX system manuals, named `intro`, usually lists all these error constants. On Linux, the error constants are listed in the `errno` manual page.
- ▶ The traditional definition of `errno` is:

```
1 extern int errno;
```



## Error Handling II

Modern UNIX supports multithreaded access to `errno` by defining it as

```
1 extern int *_ _errno_location(void);  
2 #define errno (*_ _errno_location())
```

- ▶ There are two rules to be aware of with respect to `errno`.
  1. First, its value is never cleared by a routine if an error does not occur. Therefore, we should examine its value only when the return value from a function indicates that an error occurred.
  2. Second, the value of `errno` is never set to 0 by any of the functions, and none of the constants defined in `errno.h` has a value of 0.
- ▶ Two functions are defined by the C standard to help with printing error messages.



## Error Handling III

```
1 #include <string.h>
2 char *strerror(int errnum);
3 #include <stdio.h>
4 void perror(const char *msg);
```

- ▶ Figure 1.8 shows the usage of these two error functions.
- ▶ Instead of calling either `strerror` or `perror` directly, all the examples in this text use the error functions shown in Appendix B of *APUE, 3rd Edition*.
- ▶ Example (Figure 1.8, `intro/testerror.c`)



## User Identification I

- ▶ The **user ID** from our entry in the password file is a numeric value that identifies us to the system.
- ▶ This user ID is assigned by the system administrator when our login name is assigned, and we cannot change it.
- ▶ We call the user whose user ID is 0 either **root** or the **superuser**. The **superuser** has free rein over the system.
- ▶ Our entry in the password file also specifies our numeric **group ID**.
- ▶ Groups are normally used to collect users together into projects or departments. This allows the sharing of resources among members of the same group.
- ▶ There is also a group file that maps group names into numeric group IDs, usually `/etc/group`.



## User Identification II

- ▶ In addition to the group ID specified in the password file for a login name, most versions of the UNIX System allow a user to belong to additional groups.
- ▶ These **supplementary group IDs** are obtained at login time by reading the file `/etc/group` and finding the entries that list the user as a member.
- ▶ Example (Figure 1.9, `intro/uidgid.c`).





# Signals

- ▶ Signals are a technique used to notify a process that some condition has occurred.
- ▶ The process has three choices for dealing with the signal: ignore, default or custom.
- ▶ Many conditions generate signals: some special key stroke, `kill(1)` and `killall(1)` commands, `kill(2)`, etc.
- ▶ Example (Figure 1.10, `intro/shell2.c`).



## Time Values

- ▶ Historically, UNIX systems have maintained two different time values:
  1. Calendar time, the number of seconds since the Epoch.
  2. Process time, the CPU resource used by process, measured in clock ticks.
- ▶ Command `time(1)` will give three time values of a running process: clock time, user CPU time, system CPU time.



# System Calls and Library Functions I

- ▶ All operating systems provide service points through which programs request services from the kernel.
- ▶ All implementations of the UNIX System provide a well-defined, limited number of entry points directly into the kernel called **system calls**.
- ▶ The technique used on UNIX systems is for each system call to have a function of the same name in the standard C library.
- ▶ For our purposes, we can consider the system calls as being C functions.
- ▶ We call the general-purpose functions available to programmers **library function**.
- ▶ Normally, we can replace the library functions, if desired, whereas the system calls usually cannot be replaced.

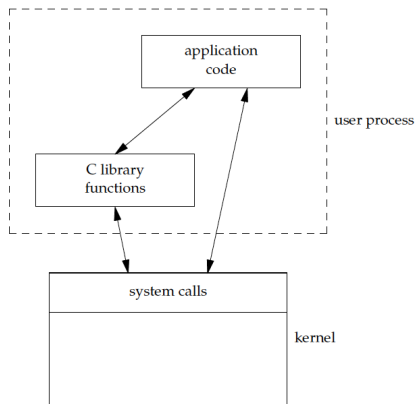


## System Calls and Library Functions II

- ▶ An application can call either a system call or a library routine. Also realize that many library routines invoke a system call.
- ▶ System calls usually provide a minimal interface, whereas library functions often provide more elaborate functionality.
- ▶ In this text, we'll use the term **function** to refer to both system calls and library functions, except when the distinction is necessary.



# System Calls and Library Functions III



**Figure:** Difference between C library functions and system calls



# ISO C Features

- ▶ All the examples in these course are written in ISO C.



## The End

The End of Chapter 1.

