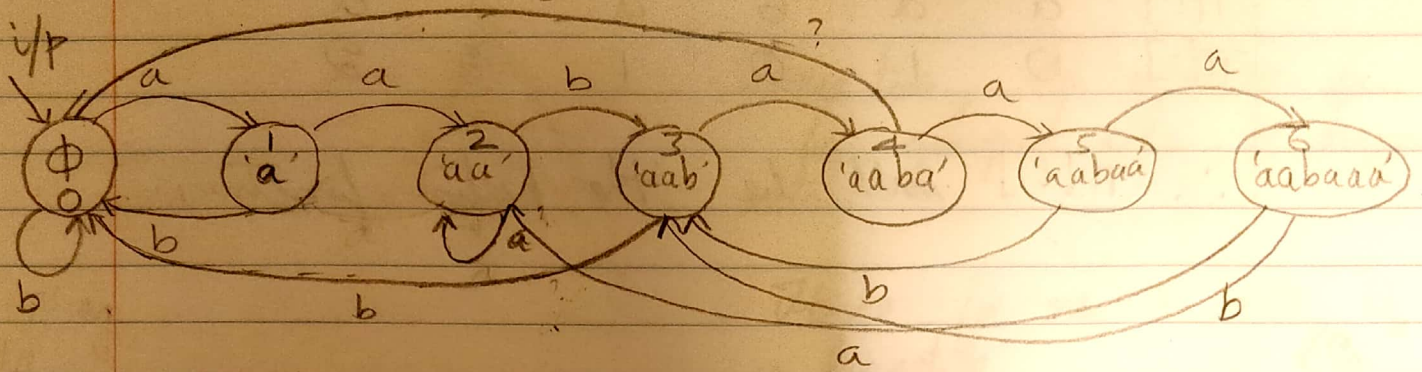
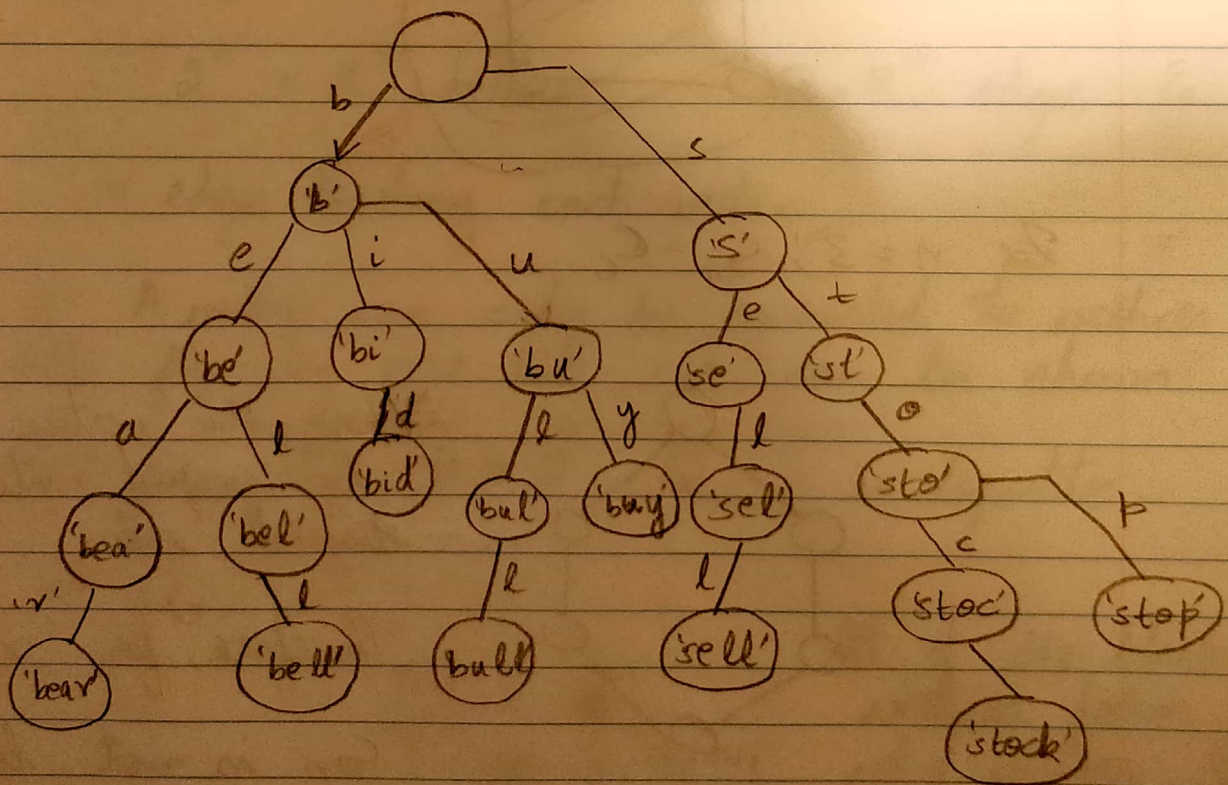


⇒ ARSHITHA BASAVARAJ

17a.  $= "aabaaa"$ ,  $\Sigma = \{a, b\}$



Anything other than  $a, b$  goes to the initial state.



c.

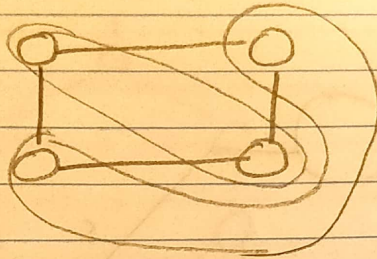
	1	1	2	3	4	5	6
$P[i]$	a	a	b	a	a	a	a
$\pi[i]$	0	1	0	1	2	2	2

$\therefore \pi[i]$  will be the prefix function.



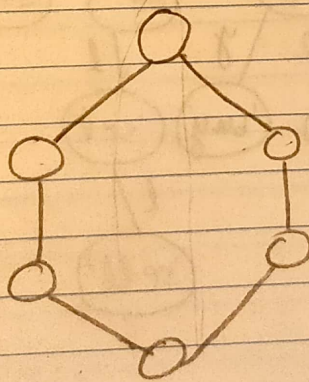
2/  
a)

For  $n=2$ ,  $C_{2 \times 2} = C_4$



$C_4$  is bipartite.

For  $n=3$ ,  $C_{2n} = C_6$

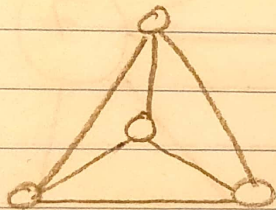


Here the vertices can't be grouped into 2 disjoint sets.

Therefore, for  $n \geq 3$ ,

$C_{2n}$  is not bipartite.

b)  $W_{2n}$  for  $n=2$ ,  $W_{2n} = W_4$



Here too, the vertices can't be grouped into 2 disjoint sets.

Hence, for no  $n > 0$ ,  $W_{2n}$  is bipartite



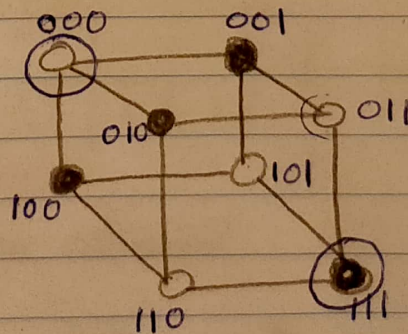
$$\Rightarrow Q_{n+7} = Q_8$$

$2^n = 2^8$  vertices  $8 \cdot 2^7$  edges  $E$

$n$  edges touching each vertex.

$\therefore$  A vertex can only be connected to vertices whose bits flip by one-bit, it can be shown that for all  $n > 0$ ,  $Q_{n+7}$  is bipartite.

For  $Q_3$ ,



All of the blackened vertices form one set of vertices  $E$

the other set of vertices are the white ones.

$\Rightarrow Q_3$  is bipartite

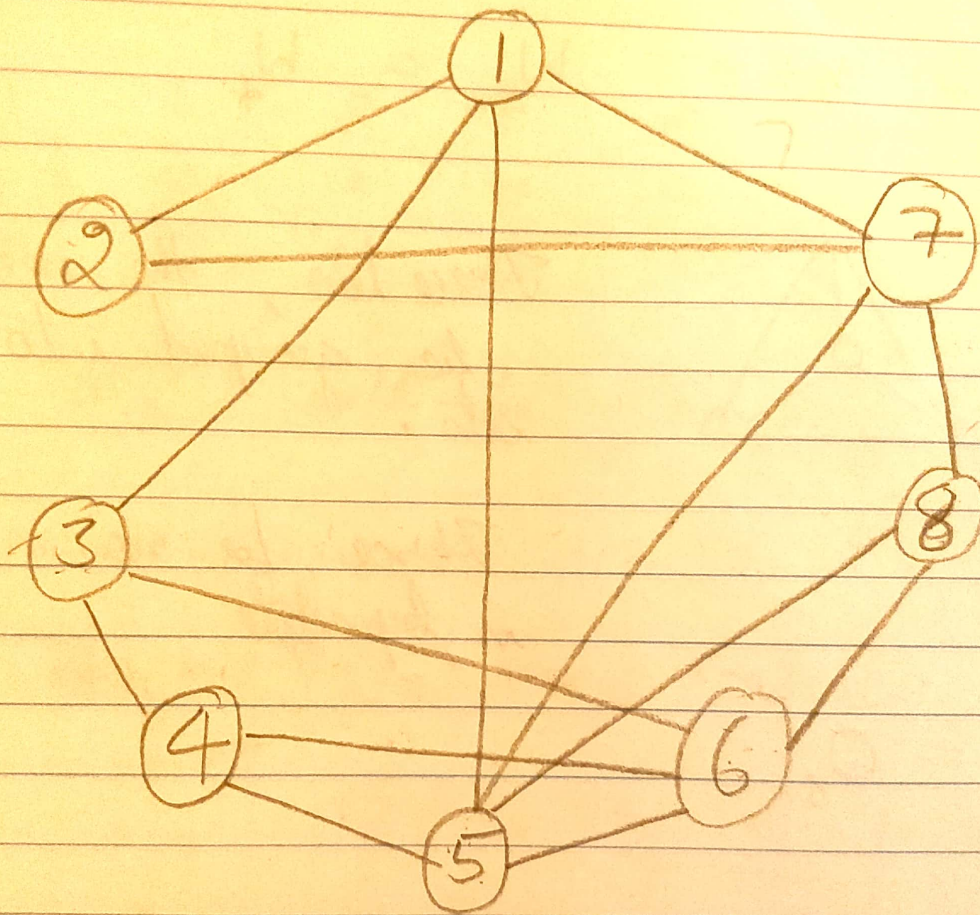
These sets are disjoint.

This theory can be extended to all  $n > 0$   
of  $Q_{n+7}$  hypercube graphs.



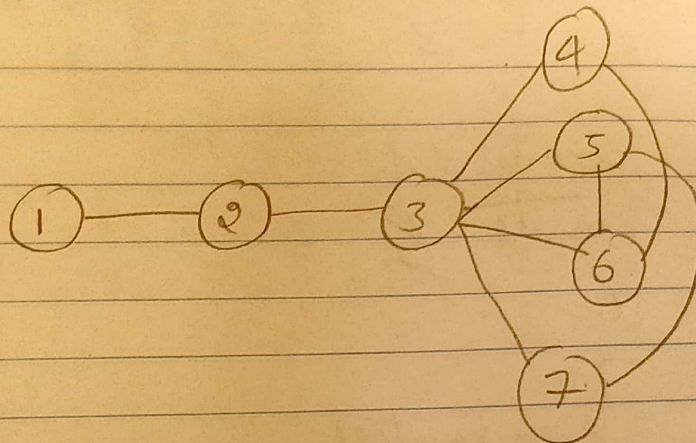
3.

a.



This would be the resulting graph  $G_1$ .

b.



Adjacency list

<u>Vertex</u>	<u>Edges</u>
1	2
2	1, 3
3	2, 4, 5, 6, 7
4	3, 6
5	3, 6, 7
6	3, 4, 5
7	3, 5



## Algorithm

Step 1: Go to every vertex & check if there exists a one and only vertex which has a single edge. -  $\Theta(V)$

Step 2: Vertex found in Step 1 would be the tail. Go to the vertex pointed by the tail.  
 $\Theta(1)$

Step 3: Check if this vertex has edges as the tail and another vertex. This would be the center vertex.  
 $\Theta(1)$

Step 4: go to the vertex other than tail of the center vertex.  $\Theta(1)$

Step 5: If the vertex found in step 4 has  $|V|-2$  edges, then it's the body vertex.  $\Theta(1)$

Step 6: If the above conditions are satisfied then it's a kite graph.

Asymptotic runtime of the algorithm is  $\Theta(V)$

Adjacency list is the most efficient graph representation of all the others.