# **Segment Trees**

## ARTHUR CONMY\*†

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Think about how to solve one of the below problems (perhaps for smaller bounds). We will discuss an approach able to solve both problems.

### **Problem** (Dynamic Range Sum Queries)

Given an array of n integers  $x_1x_2...x_n$ , process q queries of the following types:

- 1. Update the value at position k to u.
- 2. Output the sum of values in the range [a, b].

The bounds are  $n, q \le 2 \times 10^5$  and  $x_i, u \le 10^9$ .

#### **Problem** (LCAs in a Tree)

You are given a tree of n vertices, rooted at vertex 1. Given q queries u v, for each query output the lowest common ancestor. Again,  $n, q \leq 2 \times 10^5$ 

The lowest common ancestor of two vertices u and v is the lowest vertex that's on the path from 1 to u and from 1 to v:

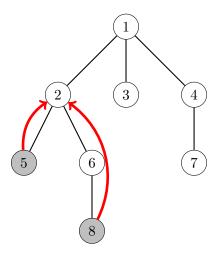


Figure 1: The LCA of vertices 5 and 8 is vertex 2.

<sup>\*</sup>asc70@cam.ac.uk for feedback or help or anything.

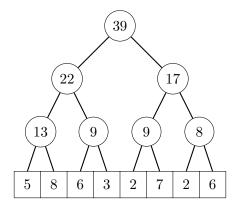
<sup>&</sup>lt;sup>†</sup>This is based on Competitive Programmer's Handbook, https://cses.fi/book/book.pdf, Antti Laaksonen.

## **Segment Trees**

A segment tree is a binary tree such that the nodes on the bottom level of the tree correspond to the array elements, and the other nodes contain information needed for processing range queries<sup>1</sup>.

0	1	2	3	4	5	6	7
5	8	6	3	2	7	2	6

The corresponding segment tree is as follows:



i.e, each vertex has value the sum of the two vertices below it in this case.

**Question.** Can you see how to perform the range and the sum queries given the data structure of this form?

You should be able to see a way to perform both range queries and updates in  $O(\log n)$ . For the Dynamic Range Queries problem this leads to an  $O(n \log n)$  solution! Segment trees often allow us to support both updating and querying efficiently, rather than just one of them.

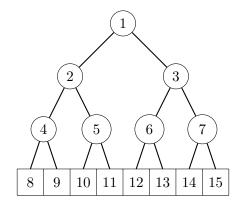
## **Implementation**

I'd recommend sticking quite close to the following segment tree implementation, perhaps even memorising the query function<sup>2</sup>.

We store the segment tree as an array tree of length 2n, where the array is kept in the array tree[n], tree[n+1], ..., tree[2\*n - 1], and the indices correspond to the binary tree vertices as shown:

<sup>&</sup>lt;sup>1</sup>In this section, we assume that the size of the array is a power of two, because it is convenient to build a segment tree for such an array. You might want to think about how we could use a power of two segment tree when our array doesn't have length that's a power of two.

<sup>&</sup>lt;sup>2</sup>It turns out a more complicated implementation is needed to allow for both range and point updates and queries to be supported; see 'lazy propagation' in ?. I don't memorise this implementation, but rewrite it, and it takes a fair bit longer for me.



This has the very nice property that where defined, vertex i has children 2i (which is even) to the left and 2i + 1 (which is odd) to the right, and parent  $\lfloor i/2 \rfloor$ .

Question. How can we implement the function void update(int index, int new\_value) that updates a value in our sum segment tree?

It's slightly trickier to implement the sum function<sup>5</sup>:

```
int sum(int a, int b) {
    a += n; b += n;
    int s = 0;
    while (a <= b) {
        if (a%2 == 1) s += tree[a++];
        if (b%2 == 0) s += tree[b--];
        a /= 2; b /= 2;
    }
    return s;
}</pre>
```

We can implement the range maximum query for the LCA problem by replacing the s += tree[a--] line with s = max(s, tree[a--]) and similar for the line below.

#### **Problems**

Implement Dynamic Range Sum Queries first: https://cses.fi/problemset/task/1648. If you get stuck implementing any part of the segment tree, you should be able to find all the parts needed (and more explanation) in https://cses.fi/book/book.pdf.

```
Problem. 1: https://cses.fi/problemset/task/1651
```

Problem. 2: https://cses.fi/problemset/task/1749

Problem. 3: https://cses.fi/problemset/task/1143

Problem. 4: https://cses.fi/problemset/task/1144

Problem. 5: https://codeforces.com/contest/1195/problem/F

Problem. 6: https://codeforces.com/contest/1467/problem/E

**Problem.** 7: https://oj.uz/problem/view/IOI13\_game http://www.ioi2013.org/wp-content/uploads/tasks/day2/game/game.pdf

 $<sup>^5{</sup>m I}$  committed this to memory at first, but so long as you've got the basic structure the checks modulo 2 and whether to increment or decrement should be rederivable, as I can go through.