

# BINARY SEARCH TREES

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Problem Solving with Computers-II

C++

```
#include <iostream>
using namespace std;

int main(){
    cout<<"Hola Facebook\n";
    return 0;
}
```

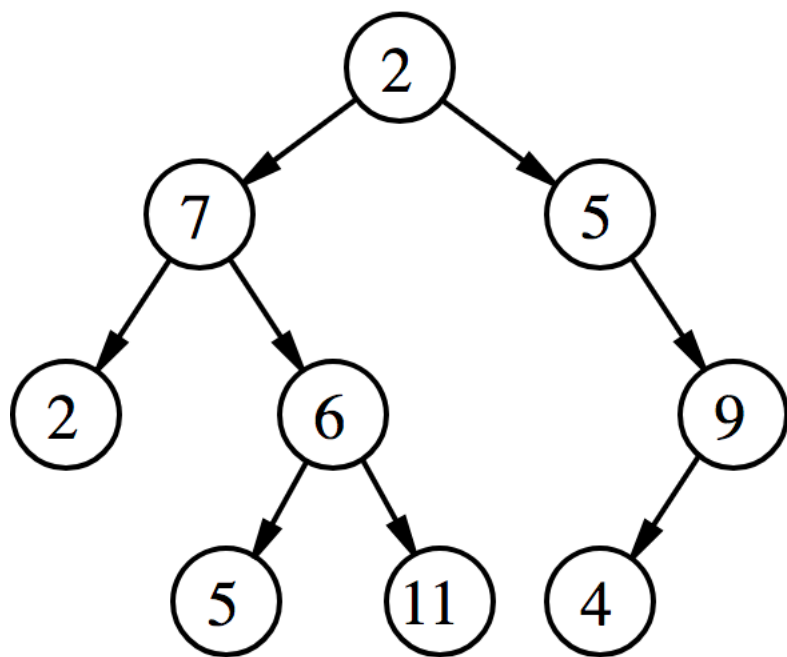
# Binary Search Trees

- What are the operations supported?
- What are the running times of these operations?
- How do you implement the BST i.e. operations supported by it?

# Operations supported by Sorted arrays and Binary Search Trees (BST)

Operations	Sorted Array	BST
Min		
Max		
Successor		
Predecessor		
Search		
Insert		
Delete		
Print elements in order		

# Trees



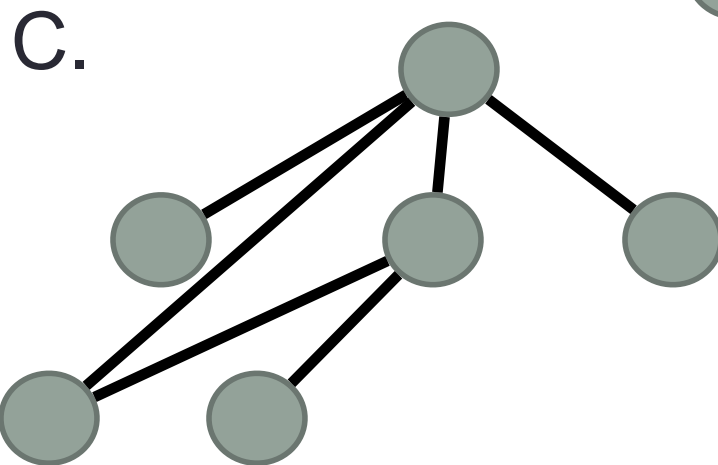
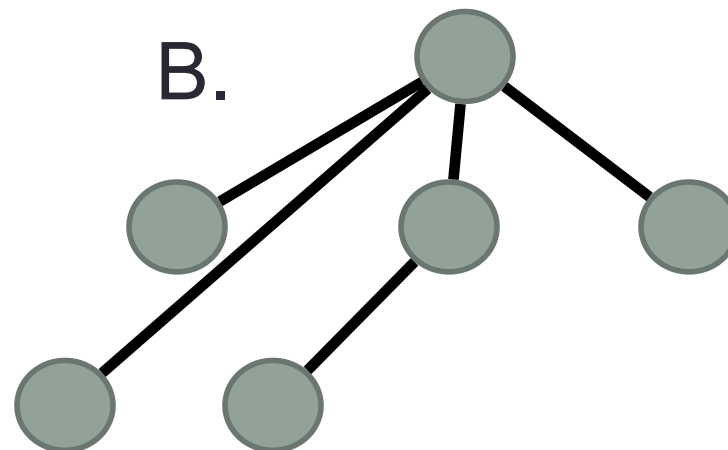
A tree has following general properties:

- One node is distinguished as a **root**;
- Every node (exclude a root) is connected by a directed edge *from* exactly one other node;

A direction is: *parent -> children*

- *Leaf node: Node that has no children*

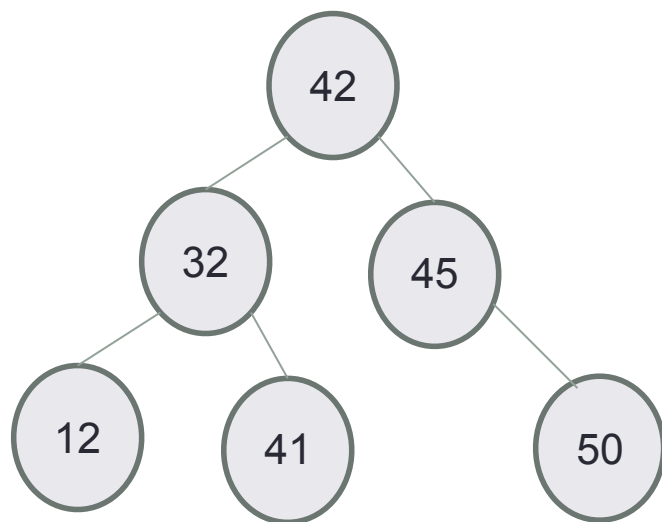
Which of the following is/are a tree?



D. A & B

E. All of A-C

# Binary Search Tree – What is it?

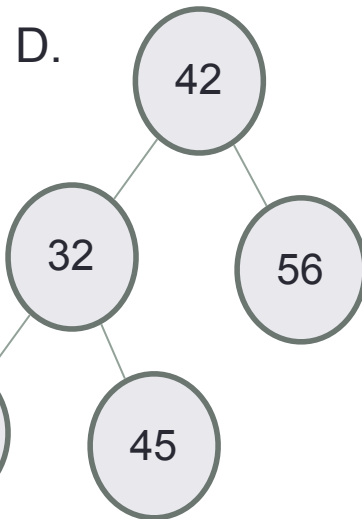
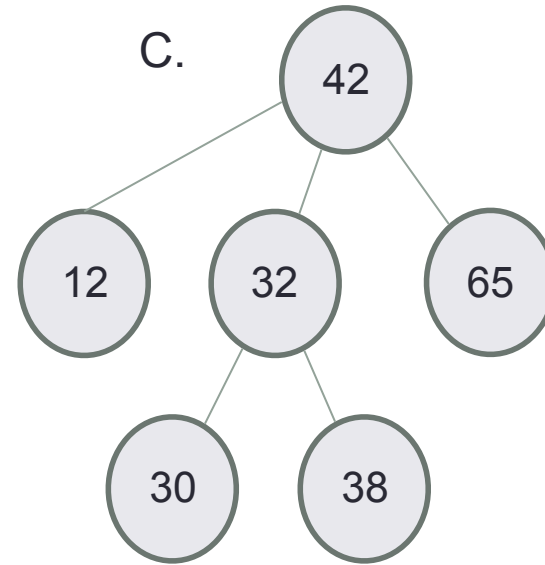
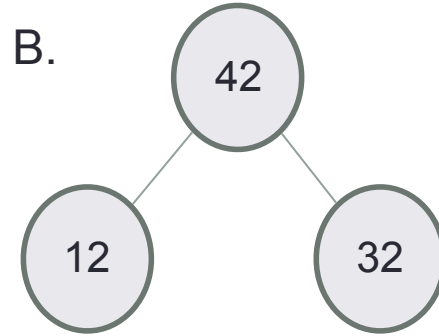
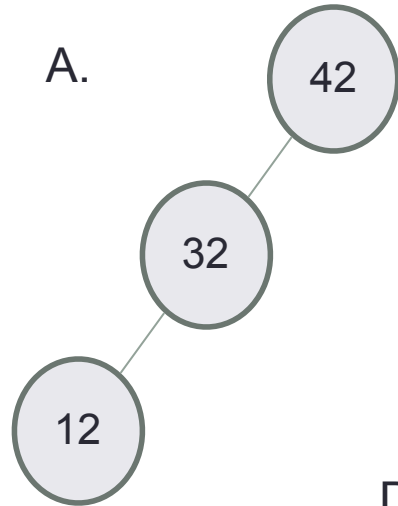


- Each node:
  - stores a key (k)
  - has a pointer to left child, right child and parent (optional)
- Satisfies the **Search Tree Property**

For any node,  
Keys in node's left subtree  $\leq$  Node's key  
Node's key  $<$  Keys in node's right subtree

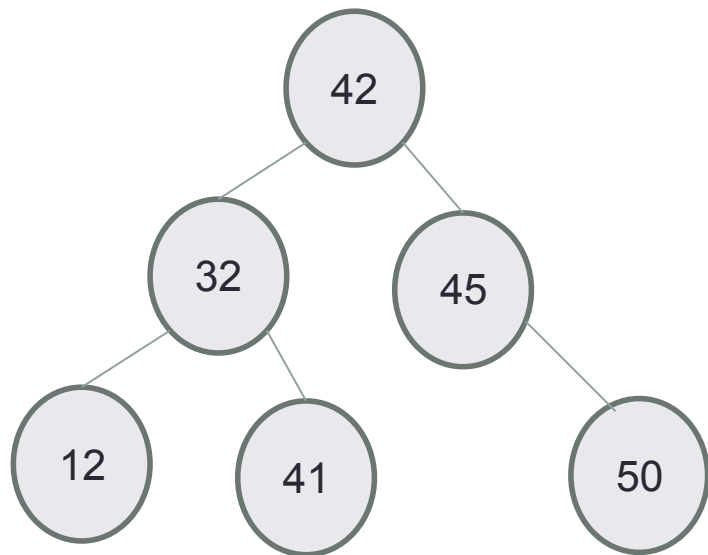
Do the keys have to be integers?

# Which of the following is/are a binary search tree?



E. More than one of these

# BSTs allow efficient search!



- Start at the root;
- Trace down a path by comparing  $k$  with the key of the current node  $x$ :
  - If the keys are equal: we have found the key
  - If  $k < \text{key}[x]$  search in the left subtree of  $x$
  - If  $k > \text{key}[x]$  search in the right subtree of  $x$



**Search for 41, then search for 53**



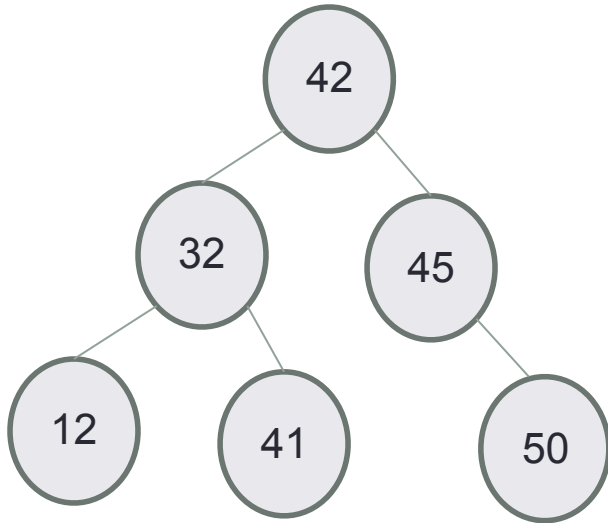
# A node in a BST

```
class BSTNode {  
  
public:  
    BSTNode* left;  
    BSTNode* right;  
    BSTNode* parent;  
    int const data;  
  
    BSTNode( const int & d ) : data(d) {  
        left = right = parent = 0;  
    }  
};
```

# Define the BST ADT

<b>Operations</b>	
Min	
Max	
Successor	
Predecessor	
Search	
Insert	
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Print elements in order	

# Insert



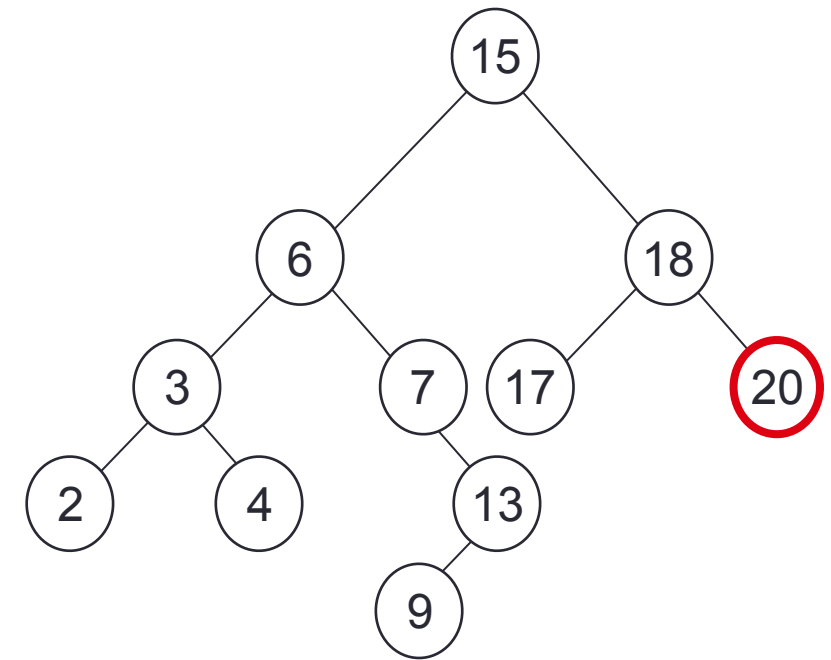
- Insert 40
- Search for the key
- Insert at the spot you expected to find it

# Max

**Goal:** find the maximum key value in a BST

Following right child pointers from the root, until a leaf node is encountered. The least node has the max value

**Alg:** `int BST::max()`



**Maximum = 20**

# Min

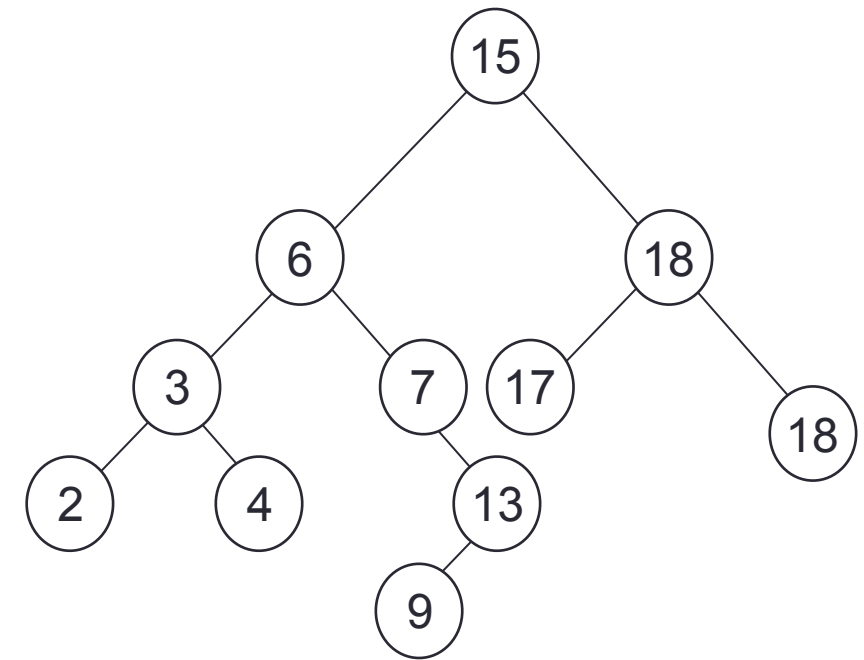
**Goal:** find the minimum key value in a BST

Start at the root.

Follow \_\_\_\_\_ child pointers from the root, until a leaf node is encountered

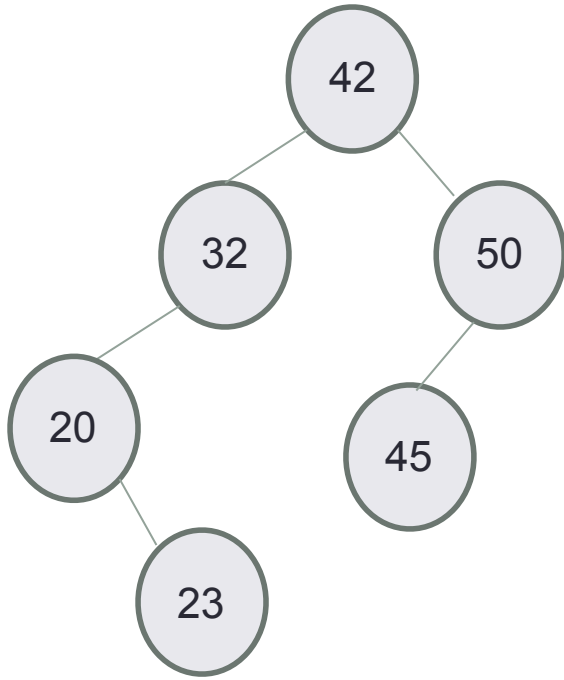
Leaf node has the min key value

**Alg:** `int BST::min()`



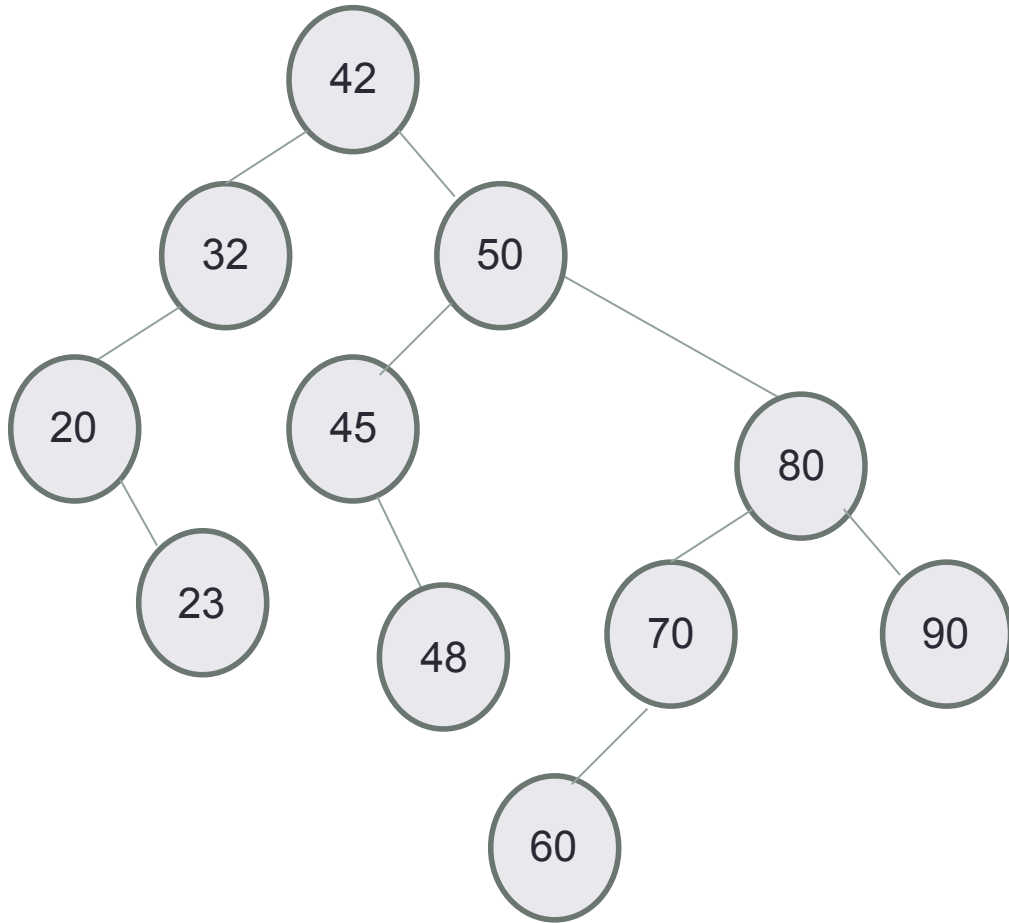
Min = ?

# Predecessor: Next smallest element



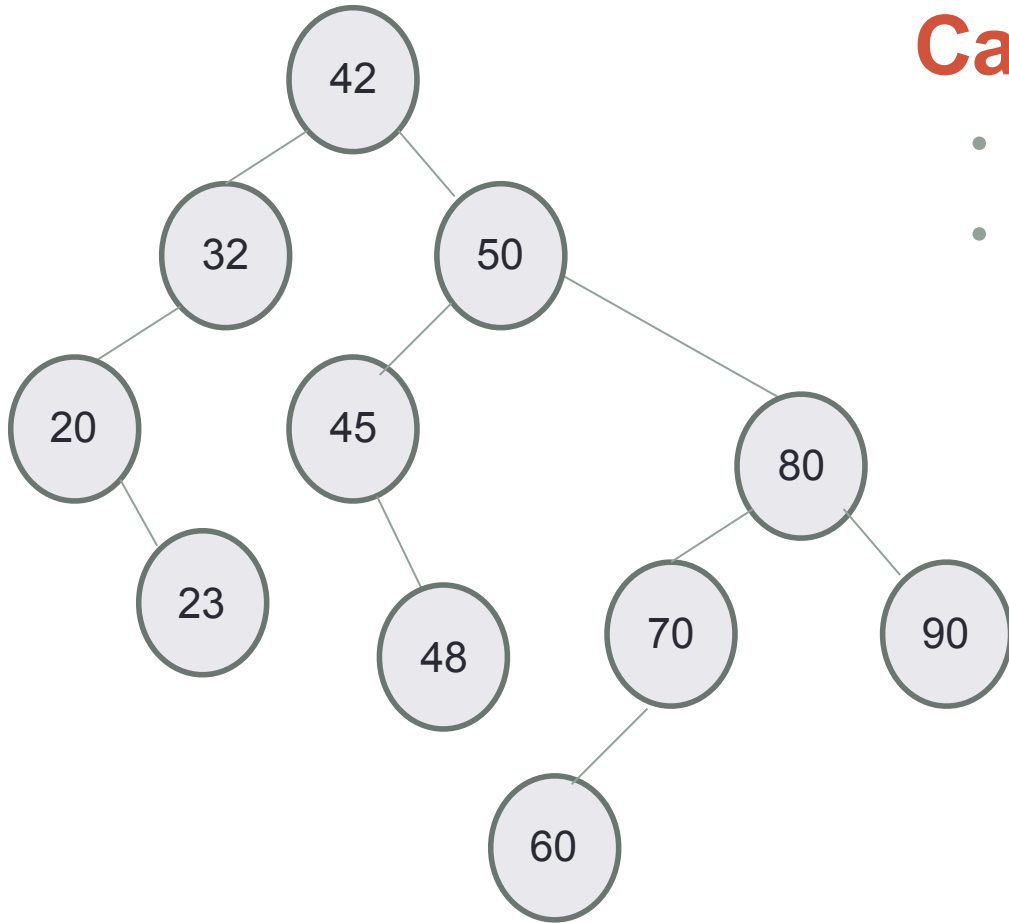
- What is the predecessor of 32?
- What is the predecessor of 45?

# Successor: Next largest element



- What is the successor of 45?
- What is the successor of 50?
- What is the successor of 60?

# Delete: Case 1

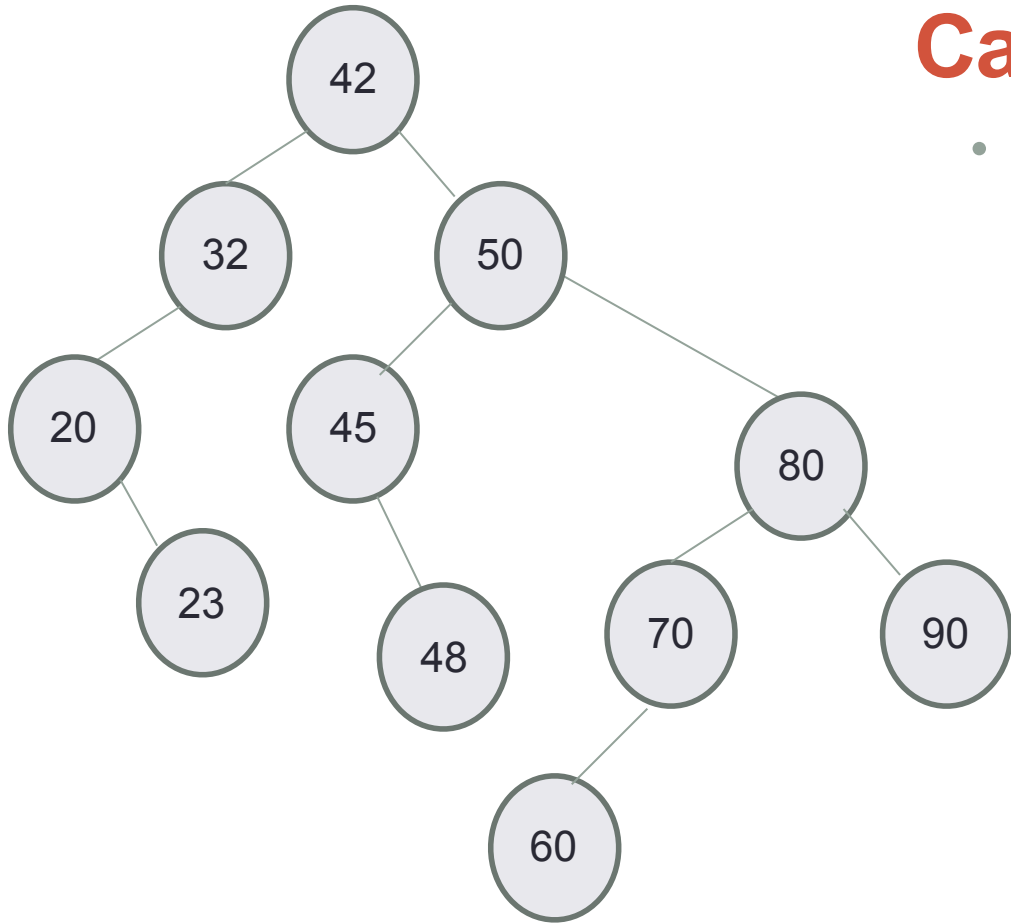


## Case 1: Node is a leaf node

- Set parent's (left/right) child pointer to null
- Delete the node



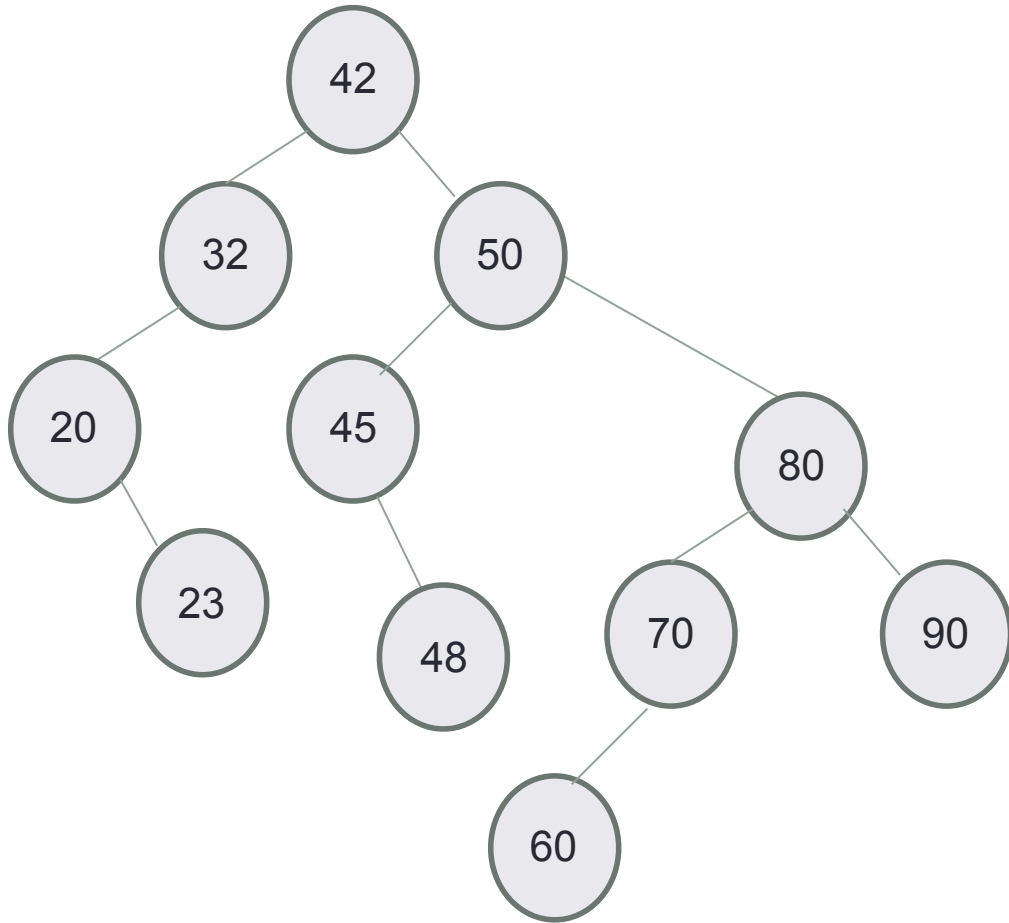
# Delete: Case 2



## Case 2 Node has only one child

- Replace the node by its only child

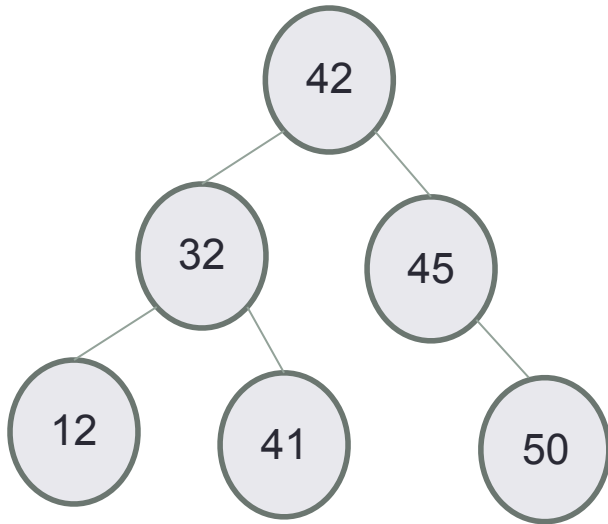
# Delete: Case 3



## Case 3 Node has two children

- Can we still replace the node by one of its children? Why or Why not?

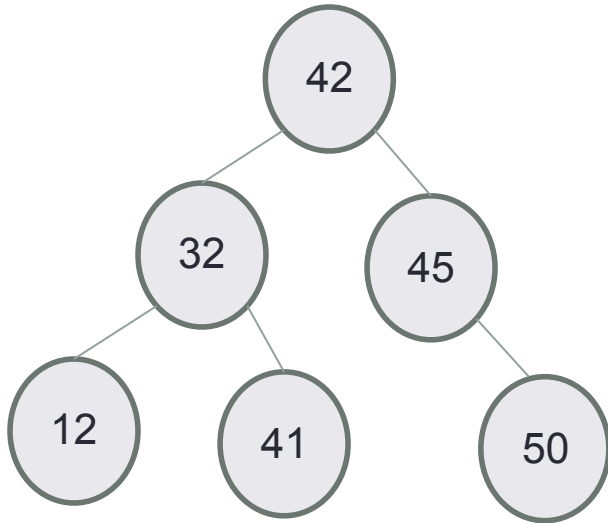
# In order traversal: print elements in sorted order



Algorithm Inorder(tree)

1. Traverse the left subtree, i.e., call Inorder(left-subtree)
2. Visit the root.
3. Traverse the right subtree, i.e., call Inorder(right-subtree)

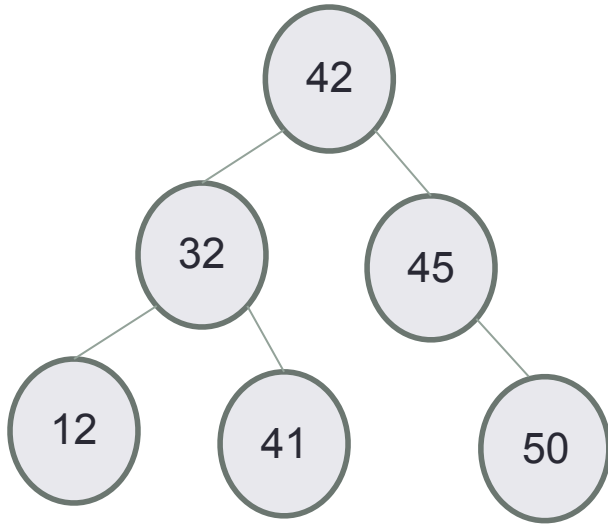
# Pre-order traversal: nice way to linearize your tree!



Algorithm Preorder(tree)

1. Visit the root.
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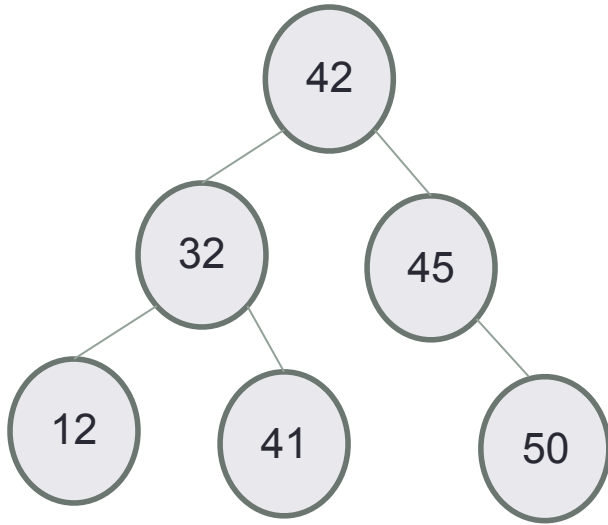
# Post-order traversal: use in recursive destructors!



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# Post-order traversal: use in recursive destructors!



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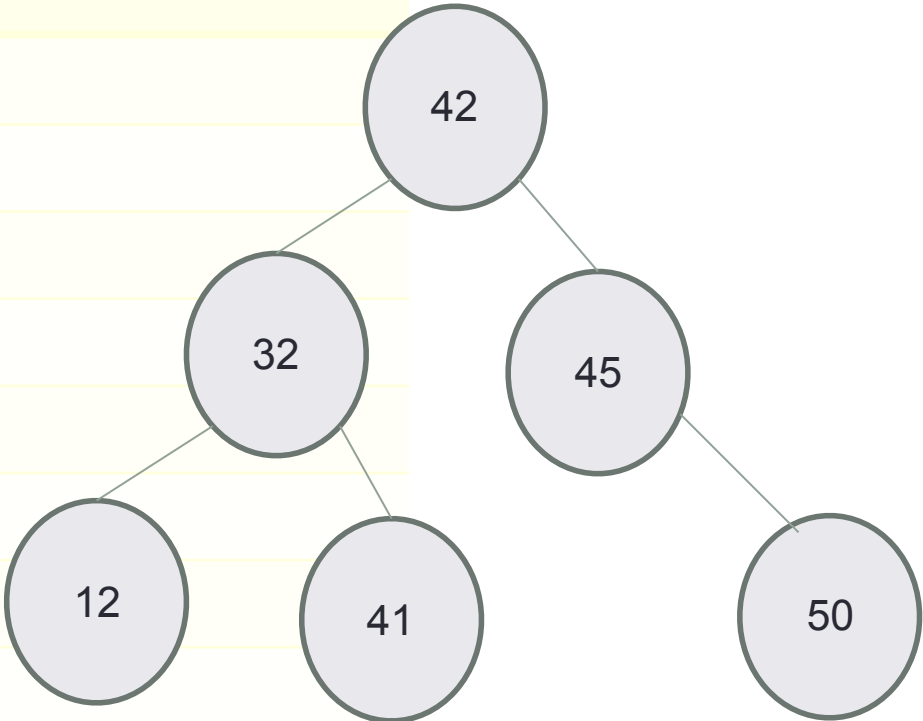
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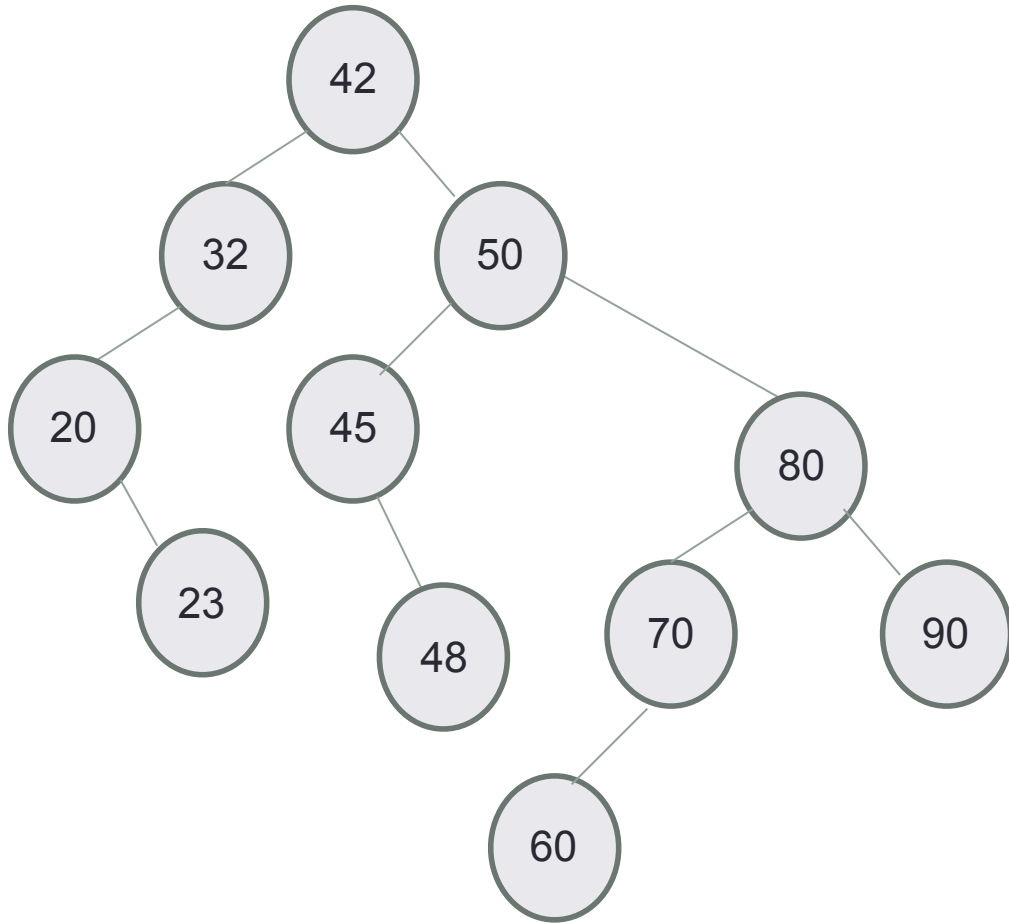
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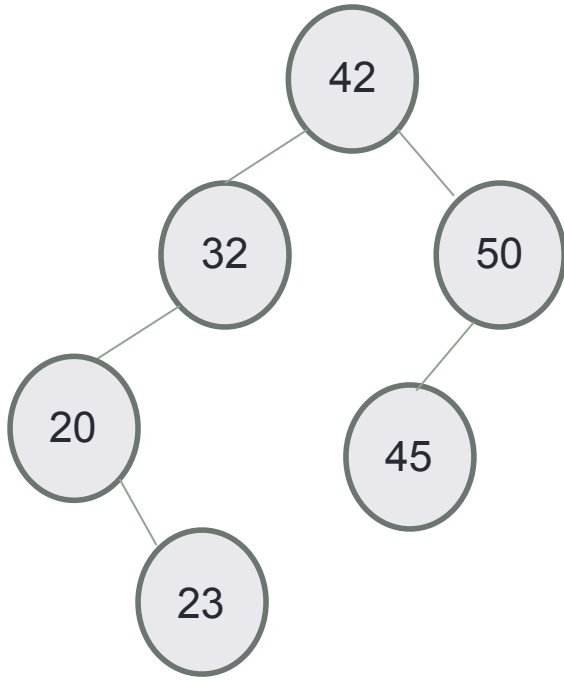


# Successor: Next largest element



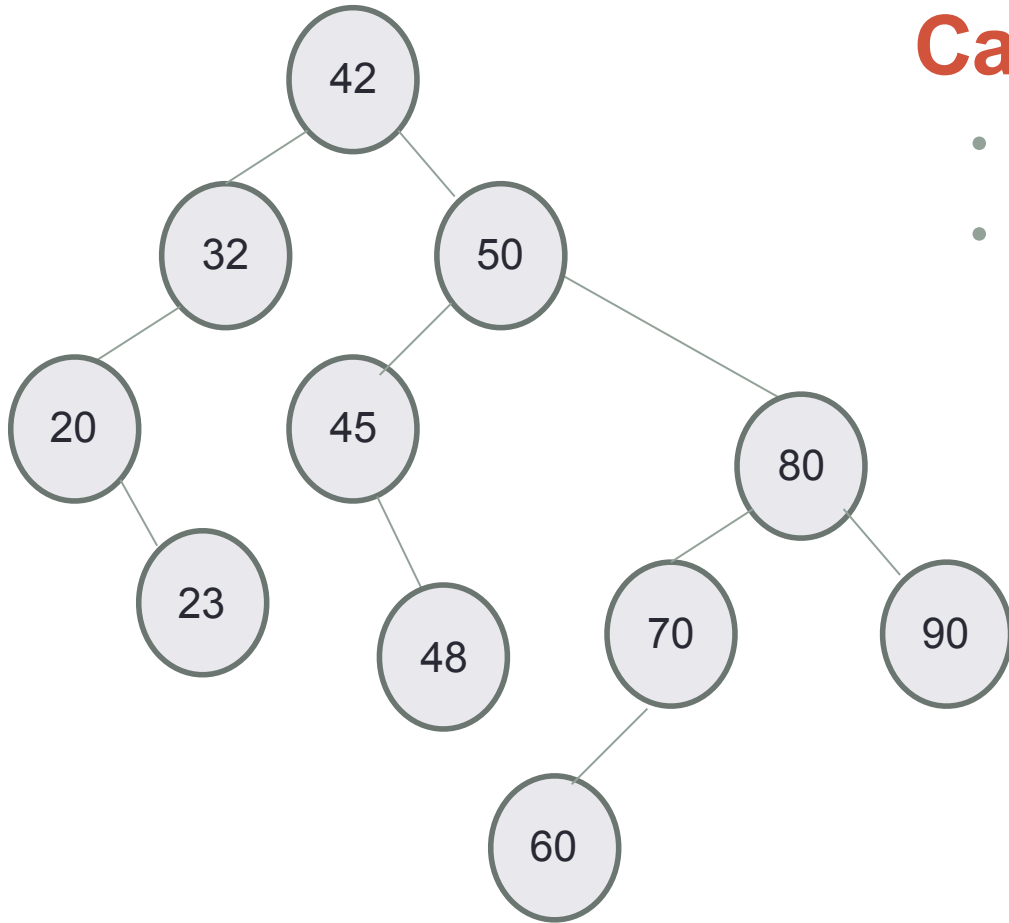
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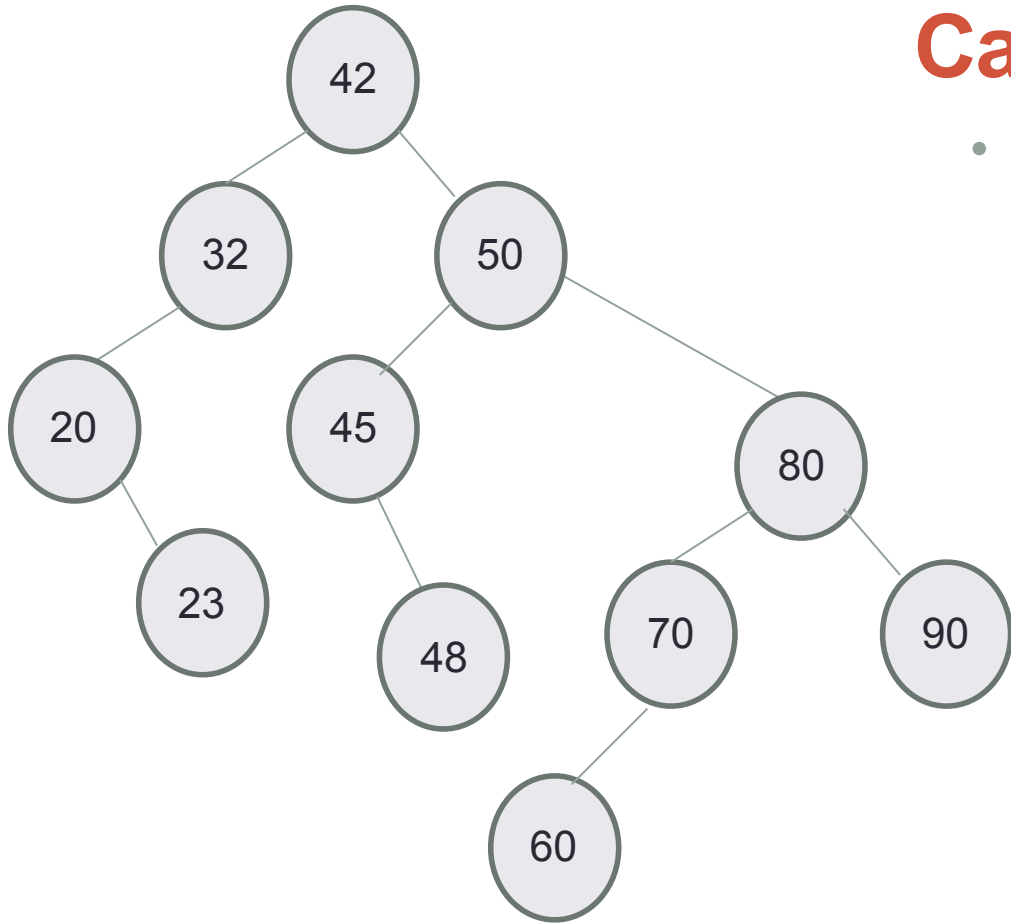
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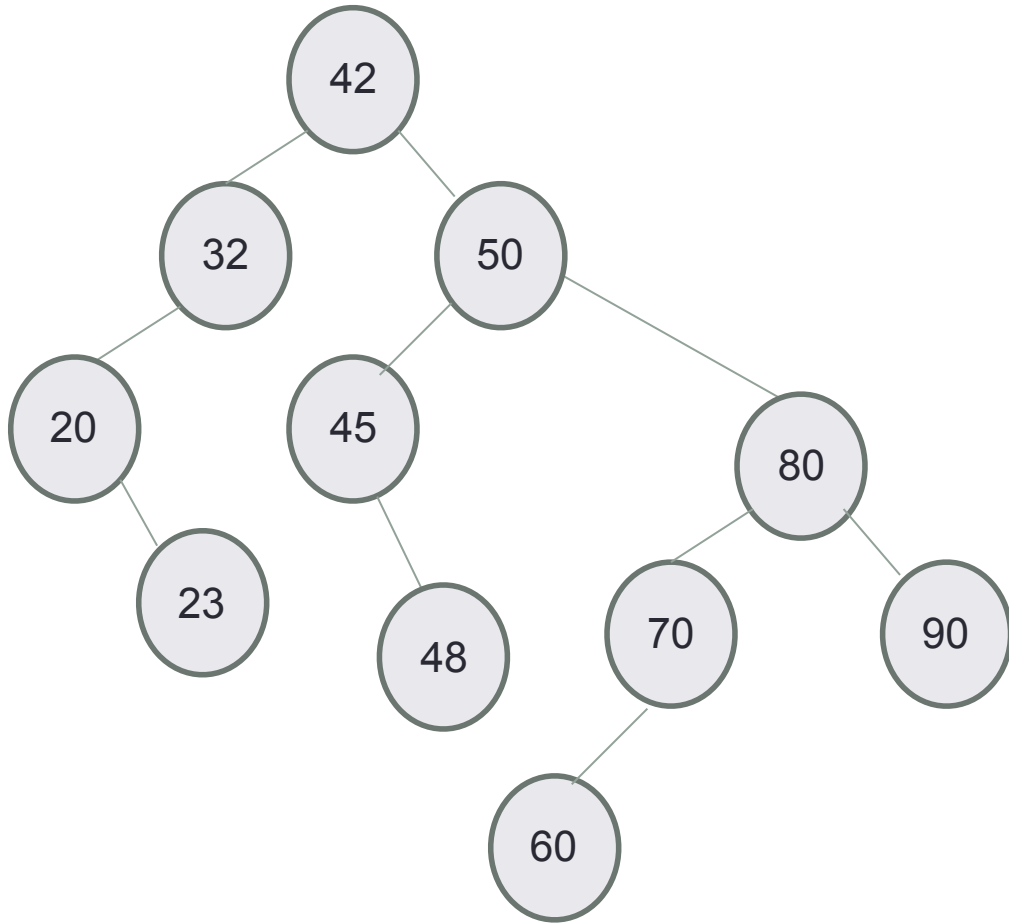
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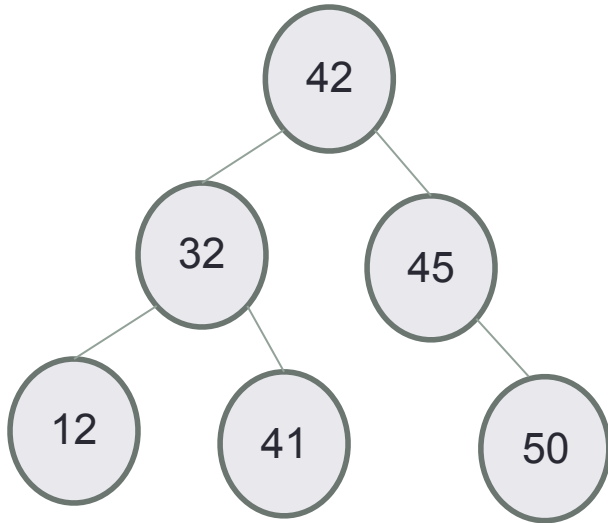
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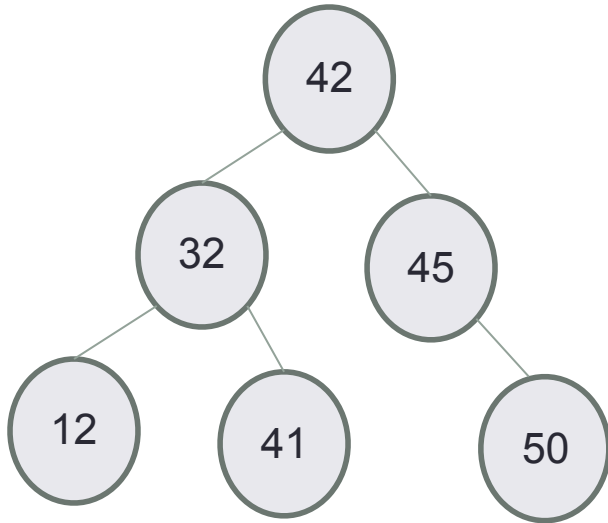
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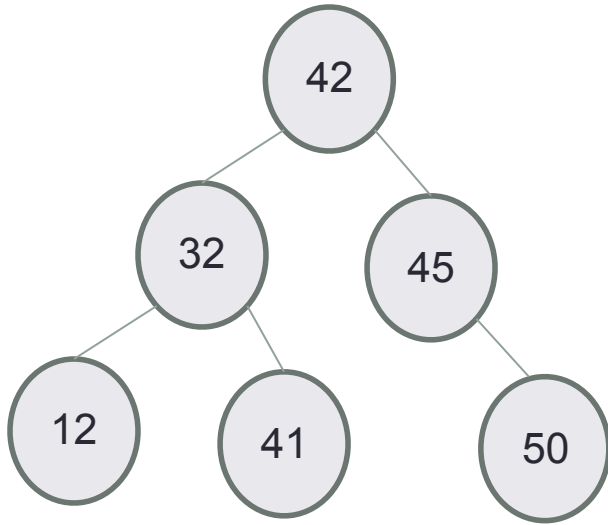
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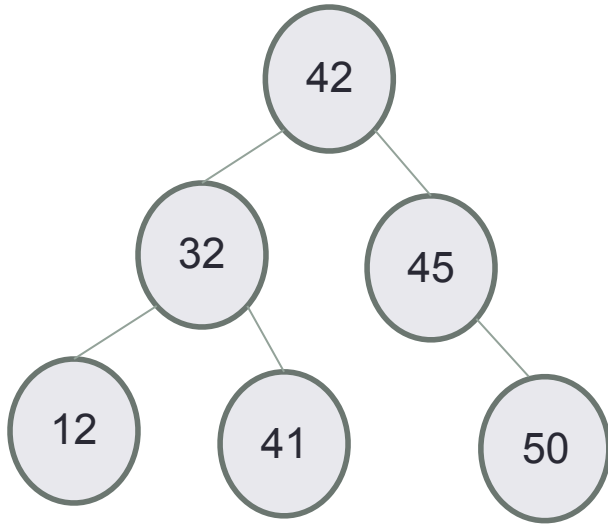


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