

Explanation Catalog explained

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1 Disclaimer

The explanation generator is a prototype. The decompositions are chosen to experiment its capabilities. One should not use these decompositions nor these explanations without care.

2 AllDifferent($\{X_1, \dots, X_n\}$)

2.1 Decomposition

$$\begin{aligned} X_i = t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\ \sum_{i \in \llbracket 1, n \rrbracket} b_{it} &\leq 1 & \forall t \in \llbracket 1, m \rrbracket \end{aligned}$$

2.2 Explanation

$$\langle X_i \neq t \rangle \rightarrow \langle X_{i'} = t \rangle_{\forall i', i' \neq i, i' \in \llbracket 1, n \rrbracket, i \in \llbracket 1, n \rrbracket}$$

A value is removed of a domain because an other variable is instantiated to it.

3 AllEqual($\{X_1, \dots, X_n\}$)

3.1 Decomposition

$$\begin{aligned} X_i = t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\ \bigwedge_{i \in \llbracket 1, n \rrbracket} b_{it} & & \forall t \in \llbracket 1, m \rrbracket \end{aligned}$$

3.2 Explanation

$$\langle X_i = t \rangle \rightarrow \langle X_i = t \rangle_{\exists i, i \in \llbracket 1, n \rrbracket}$$

$$\langle X_i \neq t \rangle \rightarrow \langle X_i \neq t \rangle_{\exists i, i \in \llbracket 1, n \rrbracket}$$

A value instantiation or removal is due to an other variable same event.

3.3 Decomposition

$$\begin{array}{ll}
X_i \geq t \iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
\bigwedge_{i \in \llbracket 1, n \rrbracket} b_{it} \iff b1_t & \forall t \in \llbracket 1, m \rrbracket \\
\bigwedge_{i \in \llbracket 1, n \rrbracket} \neg b_{it} \iff b2_t & \forall t \in \llbracket 1, m \rrbracket \\
(b1_t \vee b2_t) & \forall t \in \llbracket 1, m \rrbracket
\end{array}$$

3.4 Explanation

$$\begin{array}{l}
\langle X_i \geq t \rangle \rightarrow \langle X_i \geq t \rangle_{\exists i, i \in \llbracket 1, n \rrbracket} \\
\langle X_i < t \rangle \rightarrow \langle X_i < t \rangle_{\exists i, i \in \llbracket 1, n \rrbracket}
\end{array}$$

A modified bound is due to an other variable having the same bound.

4 NValue($N, \{X_1, \dots, X_n\}$)

4.1 Decomposition

$$\begin{array}{ll}
X_i = t \iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
\bigvee_{i \in \llbracket 1, n \rrbracket} b_{it} \iff b2_t & \forall t \in \llbracket 1, m \rrbracket \\
\sum_{t \in \llbracket 1, m \rrbracket} b2_t = p \iff b3_p & \forall t \in \llbracket 1, m \rrbracket \forall p \in \llbracket 1, n \rrbracket \\
N = p \iff b3_p & \forall p \in \llbracket 1, n \rrbracket
\end{array}$$

4.2 Explanation

$$\langle X_i = t \rangle \rightarrow \langle X_i \neq t \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t, t \in \llbracket 1, m \rrbracket} \langle N = p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

An instantiation is explained by N instantiated and all the variables that cannot take the value t .

$$\langle X_i \neq t \rangle \rightarrow \langle X_i = t \rangle_{\exists i, i \in \llbracket 1, n \rrbracket, \forall t, t \in \llbracket 1, m \rrbracket} \langle N = p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

An value removal is explained by N instantiated and all the variables that are instantiated to the value t .

$$\langle N = p \rangle \rightarrow \langle X_i \neq t \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t, t \in \llbracket 1, m \rrbracket} \langle X_i = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

A N instantiation is explained by all the variables that takes the value t or not.

$$\langle N \neq p \rangle \rightarrow \langle X_i = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

$$\langle N \neq p \rangle \rightarrow \langle X_i \neq t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

A N value removal can be explained either by all the variables that are instantiated to t or all the variables that does not have it in their domain.

5 AtLeastNValue($N, \{X_1, \dots, X_n\}$)

5.1 Decomposition

$$\begin{aligned}
X_i = t &\iff b_{it} && \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
\bigvee_{i \in \llbracket 1, n \rrbracket} b_{it} &\iff b_{2t} && \forall t \in \llbracket 1, m \rrbracket \\
\sum_{t \in \llbracket 1, m \rrbracket} b_{2t} \geq p &\iff b_{3p} && \forall t \in \llbracket 1, m \rrbracket \\
N \geq p &\iff b_{3p} && \forall t \in \llbracket 1, m \rrbracket
\end{aligned}$$

5.2 Explanation

$$\langle X_i = t \rangle \rightarrow \langle X_i \neq t' \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t', t' \neq t, t' \in \llbracket 1, m \rrbracket, t \in \llbracket 1, m \rrbracket} \langle N \geq p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

An instantiation is explained by the upper bound of N and all the variables that cannot take the value t .

$$\langle X_i \neq t \rangle \rightarrow \langle X_i = t' \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t', t' \neq t, t' \in \llbracket 1, m \rrbracket, t \in \llbracket 1, m \rrbracket} \langle N < p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

A value removal is explained by the lower bound of N and all the variables that are instantiated to the value t .

$$\langle N \geq p \rangle \rightarrow \langle X_i = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

The lower bound modification of N is explained by all the variables instantiated to t

$$\langle N < p \rangle \rightarrow \langle X_i \neq t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

The upper bound modification of N is explained by all the variables that cannot be affected to the value t .

6 AtMostNValue($N, \{X_1, \dots, X_n\}$)

6.1 Decomposition

$$\begin{aligned}
X_i = t &\iff b_{it} && \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
\bigvee_{i \in \llbracket 1, n \rrbracket} b_{it} &\iff b_{2t} && \forall t \in \llbracket 1, m \rrbracket \\
\sum_{t \in \llbracket 1, m \rrbracket} b_{2t} < p &\iff b_{3p} && \forall t \in \llbracket 1, m \rrbracket \\
N \geq p &\iff b_{3p} && \forall t \in \llbracket 1, m \rrbracket
\end{aligned}$$

6.2 Explanation

$$\langle X_i = t \rangle \rightarrow \langle X_i \neq t' \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t', t' \neq t, t' \in \llbracket 1, m \rrbracket, t \in \llbracket 1, m \rrbracket} \langle N < p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

$$\langle X_i \neq t \rangle \rightarrow \langle X_i = t' \rangle_{\forall i, i \in \llbracket 1, n \rrbracket, \forall t', t' \neq t, t' \in \llbracket 1, m \rrbracket, t \in \llbracket 1, m \rrbracket} \langle N \geq p \rangle_{\forall p, p \in \llbracket 1, n \rrbracket}$$

$$\langle N < p \rangle \rightarrow \langle X_i = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

$$\langle N \geq p \rangle \rightarrow \langle X_i \neq t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket, \forall i, i \in \llbracket 1, n \rrbracket}$$

Same explanations as AtLeastNValues but inverted.

7 Cumulative($\{X_1, \dots, X_n\}, \{d_1, \dots, d_n\}, c$)

7.1 Decomposition

$$\begin{aligned}
X_i \geq t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
(b_{i(t-d_i)} \wedge \neg b_{it}) &\iff b_{2it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
\sum_{i \in \llbracket 1, n \rrbracket} b_{2it} &\leq c & \forall t \in \llbracket 1, m \rrbracket
\end{aligned}$$

7.2 Explanation

$$\langle X_i \geq t \rangle \rightarrow \langle X_i \geq t' \rangle_{t'=t-d_i} \langle X_{i'} < t \rangle_{\forall i', i' \neq i, i' \in \llbracket 1, n \rrbracket, i \in \llbracket 1, n \rrbracket} \langle X_{i'} \geq t' \rangle_{\forall i', i' \neq i, i' \in \llbracket 1, n \rrbracket, i \in \llbracket 1, n \rrbracket, t'=t-d_{i'}}$$

A lower bound modification is explained by all the tasks that saturate the time t and the previous lower bound because the filtering cannot make holes in the domain

$$\langle X_i < t \rangle \rightarrow \langle X_i < t' \rangle_{t'=t+d_i} \langle X_{i'} < t' \rangle_{\forall i', i' \neq i, i' \in \llbracket 1, n \rrbracket, i \in \llbracket 1, n \rrbracket, t'=t+d_i} \langle X_{i'} \geq t'' \rangle_{\forall i', i' \neq i, i' \in \llbracket 1, n \rrbracket, i \in \llbracket 1, n \rrbracket, t''=t'-d_{i'}, t'=t+d_i}$$

Same idea for the upper bound

Except for the heights of the tasks it is the similar to the explanation in: <http://arxiv.org/abs/1208.3015>

8 Element($I, \{X_1, \dots, X_n\}, V$)

8.1 Decomposition

$$\begin{aligned}
X_i = t &\iff b_{it}^X & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
I = i &\iff b_i^I & \forall i \in \llbracket 1, n \rrbracket \\
V = t &\iff b_t^V & \forall t \in \llbracket 1, m \rrbracket \\
\neg b_t^V \wedge \neg b_i^I \wedge b_{it}^X & & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
b_t^V \wedge \neg b_i^I \wedge \neg b_{it}^X & & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket
\end{aligned}$$

8.2 Explanation

$$\begin{aligned}
\langle X_i = t \rangle &\rightarrow \langle I = i \rangle_{V=t} \\
\langle X_i \neq t \rangle &\rightarrow \langle I = i \rangle_{V \neq t} \\
\langle I \neq i \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket} \langle V = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket} \\
\langle I \neq i \rangle &\rightarrow \langle X_i = t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket} \langle V \neq t \rangle_{\forall t, t \in \llbracket 1, m \rrbracket} \\
\langle V = t \rangle &\rightarrow \langle X_i = t \rangle_{\forall i, i \in \llbracket 1, n \rrbracket} \langle I = i \rangle_{\forall i, i \in \llbracket 1, n \rrbracket} \\
\langle V \neq t \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall i, i \in \llbracket 1, n \rrbracket} \langle I = i \rangle_{\forall i, i \in \llbracket 1, n \rrbracket}
\end{aligned}$$

9 Increasing($\{X_1, \dots, X_n\}$)

9.1 Decomposition

$$\begin{aligned}
X_i \geq t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\
(\neg b_{(i-1)t} \vee b_{it}) & & \forall i \in \llbracket 2, n \rrbracket \forall t \in \llbracket 1, m \rrbracket
\end{aligned}$$

9.2 Explanation

$$\begin{aligned}
\langle X_i \geq t \rangle &\rightarrow \langle X_{i'} \geq t \rangle_{i'=i-1} \\
\langle X_i < t \rangle &\rightarrow \langle X_{i'} < t \rangle_{i'=i+1}
\end{aligned}$$

10 Decreasing($\{X_1, \dots, X_n\}$)

10.1 Decomposition

$$\begin{aligned} X_i \geq t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\ (b_{(i-1)t} \vee \neg b_{it}) & & \forall i \in \llbracket 2, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \end{aligned}$$

10.2 Explanation

$$\begin{aligned} \langle X_i \geq t \rangle &\rightarrow \langle X_{i'} \geq t \rangle_{i'=i+1} \\ \langle X_i < t \rangle &\rightarrow \langle X_{i'} < t \rangle_{i'=i-1} \end{aligned}$$

11 Among($c, \{X_1, \dots, X_n\}, D_4$)

11.1 Decomposition

$$\begin{aligned} X_i = t &\iff b_{it} & \forall i \in \llbracket 1, n \rrbracket \forall t \in \llbracket 1, m \rrbracket \\ \bigvee_{t \in D_4} b_{it} &\iff b_{2i} & \forall i \in \llbracket 1, n \rrbracket \\ \sum_{i \in \llbracket 1, n \rrbracket} b_{2i} &= c \end{aligned}$$

11.2 Explanation

$$\begin{aligned} \langle X_i = t \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall t, t \in D_4, \forall i, i \in \llbracket 1, n \rrbracket} \\ \langle X_i = t \rangle &\rightarrow \langle X_i \neq t' \rangle_{\forall t', t' \neq t, t' \in D_4, t \in D_4} \\ \langle X_i \neq t \rangle &\rightarrow \langle X_i = t \rangle_{\exists t, t \in D_4 \forall i, i \in \llbracket 1, n \rrbracket} \end{aligned}$$

12 Roots($\{X_1, \dots, X_n\}, I, V$)

12.1 Decomposition

$$\begin{aligned} X_i = t &\iff b_{it} & \forall i, t \\ \sum_{t \in V} b_{it} &= 1 & \forall i \in I \\ \sum_{t \in D \setminus V} b_{it} &= 1 & \forall i \in \llbracket 1, n \rrbracket \setminus I \end{aligned}$$

12.2 Explanation

$$V = D_5 \text{ and } D \setminus V = D_6$$

$$\begin{aligned} \langle X_i = t \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall t, t \in D_5} \\ \langle X_i = t \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall t, t \in D_6} \\ \langle X_i \neq t \rangle &\rightarrow \langle X_i = t \rangle_{\forall t, t \in D_5} \\ \langle X_i \neq t \rangle &\rightarrow \langle X_i = t \rangle_{\forall t, t \in D_6} \end{aligned}$$

13 Range($\{X_1, \dots, X_n\}, I, V$)

13.1 Decomposition

$$\begin{aligned}
 X_i = t &\iff b_{it} && \forall it \\
 \sum_{i \in I} b_{it} &\geq 1 && \forall t \in V \\
 \sum_{t \in V} b_{it} &= 1 && \forall i \in I
 \end{aligned}$$

13.2 Explanation

$I = D_5$ and $V = D_6$

$$\begin{aligned}
 \langle X_i = t \rangle &\rightarrow \langle X_i \neq t' \rangle_{\forall t', t' \neq t, t' \in D_5, t \in D_5} \\
 \langle X_i = t \rangle &\rightarrow \langle X_i \neq t \rangle_{\forall t, t \in D_6} \\
 \langle X_i \neq t \rangle &\rightarrow \langle X_i = t \rangle_{\forall t, t \in D_6}
 \end{aligned}$$

14 Xor($b1, b2, b3$)

14.1 Decomposition

$$b1 = (b2 \neq b3)$$

14.2 Explanation

$$\begin{aligned}
 \langle b1 \rangle &\rightarrow \langle b2 \rangle \langle \neg b3 \rangle \\
 \langle b1 \rangle &\rightarrow \langle b3 \rangle \langle \neg b2 \rangle \\
 \langle \neg b1 \rangle &\rightarrow \langle b2 \rangle \langle b3 \rangle \\
 \langle \neg b1 \rangle &\rightarrow \langle \neg b2 \rangle \langle \neg b3 \rangle
 \end{aligned}$$

And conversly for $b2$ and $b3$