15-122 Bootcamp: Point-to-Proofs

Spring 2024

get materials here!

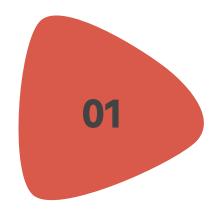


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Today's Agenda



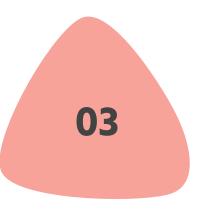
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Quick Proof Recap



TA Example: Increasing A



Exercises: Find Min

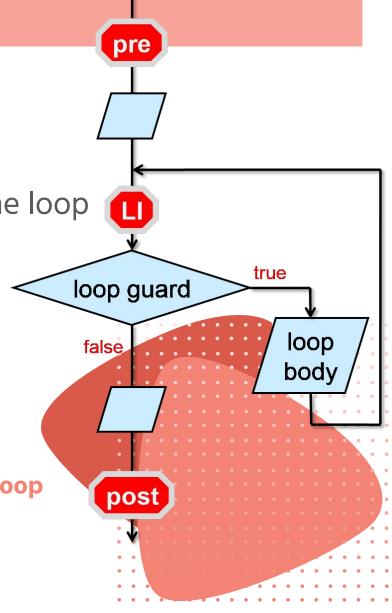


Safety (Arrays)

- We prove something is <u>safe</u> by showing that its preconditions are satisfied
 - function calls we won't practice these today
 - primitive operations ←
 - with array accesses A[i], we need to show that the index i is in bounds
 - In other words, show that $0 \le i \le \text{length}(A)$

Correctness Proofs

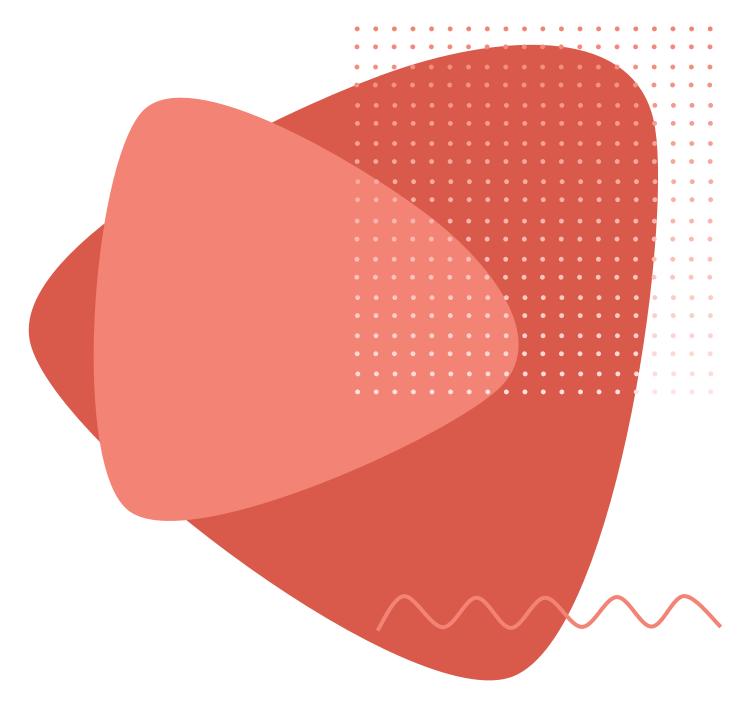
- INIT: Proving LI is true before the loop
 - Use code <u>before</u> the loop and preconditions
- PRES: Proving the LI is true in an arbitrary iteration of the loop
 - Use code <u>within</u> the loop and the loop guard
- **TERM:** Proving that the loop terminates
 - This follows a template we won't be practicing these today
- EXIT: proving that the loop is correct (and/or function returns correct value)
 - Use loop invariants, <u>negation</u> of the loop guard, code after the loop





TA Example:





Safety Example

To Show: 0 <= i - 1 && i - 1 < \length(A)

```
To Show: 0 <= i - 1
```

```
A. 0 \le i by line 7
```

B. i!= 0 by short-circuiting line 8

C. 0 < i by math on A and B

D. $0 \le i - 1$ by math on C

To Show: i - 1 < \length(A)

```
E. i \le n by line 7
```

F. $n == \left(A\right)$ by line 3, n unchanged

G. $i - 1 < \operatorname{length}(A)$ by math on **A**, E, F

```
1 // returns array containing 1, 2, 3, 4... n
 2 int[] increasing A(int[] A, int n)
 3 //@requires n >= 1 && \length(A) == n;
4 //@ensures A[n-1] == n
5 {
     for (int i = 0; i < n; i++)
     //@loop_invariant 0 <= i && i <= n;</pre>
     //@loop invariant i == 0 || A[i-1] == i;
      A[i] = i + 1;
10
11
12
13
     return A;
14 }
15
```

INIT Example

```
To Show: i == 0 || A[i - 1] == i;
```

```
A. i = 0 by line 6
B. \underline{i == 0 \parallel A[i-1] == i} by logic on A (short-circuiting)
```

```
1 // returns array containing 1, 2, 3, 4... n
 2 int[] increasing_A(int[] A, int n)
 3 //@requires n >= 1 && \length(A) == n;
 4 //@ensures A[n-1] == n
5 {
     for (int i = 0; i < n; i++)
    //@loop_invariant 0 <= i && i <= n;</pre>
     //@loop_invariant i == 0 || A[i-1] == i;
      A[i] = i + 1;
11
12
     return A;
14 }
15
```

Pres Example

```
Assuming: i == 0 || A[i - 1] == i

To Show: i' == 0 || A[i' - 1] == i'
```

```
A. i' = i + 1 by line 6

B. A[i] = i + 1 by line 10

C. i' - 1 = i by math on A

D. i == 0 || A[i'-1] = i' by math on B and C
```

```
1 // returns array containing 1, 2, 3, 4... n
 2 int[] increasing A(int[] A, int n)
 3 //@requires n >= 1 && \length(A) == n;
 4 //@ensures A[n-1] == n
5 {
    for (int i = 0; i < n; i++)
    //@loop_invariant 0 <= i && i <= n;</pre>
    //@loop_invariant i == 0 | A[i-1] == i;
     A[i] = i + 1;
10
11
12
13
     return A;
14 }
15
```

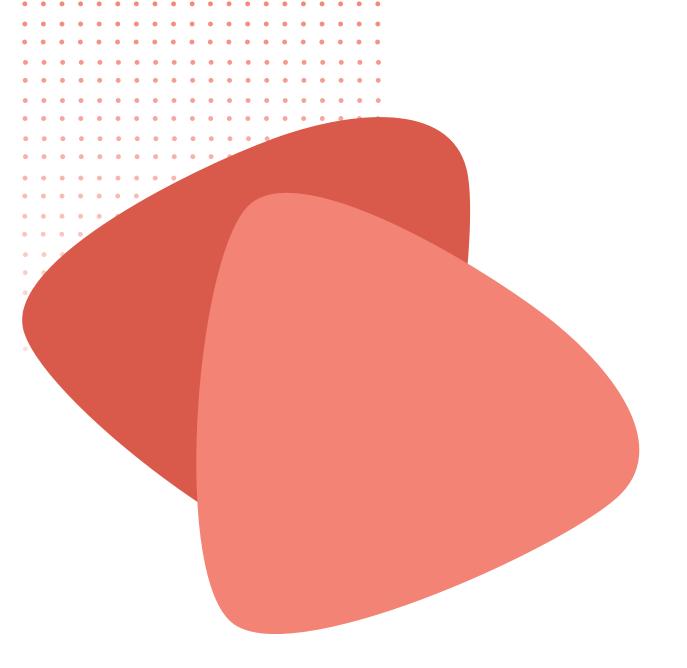
Exit Example

```
To Show: A[n - 1] == n
```

```
A. i \ge n by negation of LG, line 6
B. n == \lambda(A) by line 3, n unchanged
C. i \le n by line 7, B
D. i == n by A and C
E. n \ge 0 by line 3, n unchanged
F. i == 0 \mid A[i-1] == i by logic, D, E and F
```

```
LI must be true upon exiting loop!
```

```
1 // returns array containing 1, 2, 3, 4... n
 2 int[] increasing A(int[] A, int n)
 3 //@requires n >= 1 && \length(A) == n;
 4 //@ensures A[n-1] == n
5 {
     for (int i = 0; i < n; i++)
    //@loop_invariant 0 <= i && i <= n;</pre>
    //@loop_invariant i == 0 || A[i-1] == i;
     A[i] = i + 1;
11
12
13
     return A;
14 }
15
```



Exercises:

Find Min

Safety Example

To Show: 0 <= j && j < \length(A);

```
To Show: 0 <= j

A. lo < j by line 8
```

B. 0 <= lo by line 2, lo unchanged

C. $0 \le j$ by math on A, B

To Show: j < \length(A)

```
D. j < hi by line 7
```

E. hi <= \length(A) by line 2, hi unchanged

F. $\underline{i < \operatorname{length}(A)}$ by math on D, E

```
1 int find min (int[] A, int lo, int hi)
 2 //@requires 0 <= lo && lo < hi && hi <= \length(A);</pre>
   //@ensures lo <= \result && \result < hi;</pre>
 4 //@ensures A[\result] <= A[lo];
 5 {
      int m = lo;
      for (int j = lo + 1; j < hi; j++)
      //@loop invariant lo < j && j <= hi;</pre>
      //@loop_invariant lo <= m && m < hi;</pre>
      //@loop invariant A[m] <= A[lo];</pre>
          if (A[j] < A[m])
12
13
             m = j;
14
15
      return m;
16 }
```

INIT Example

To Show: A[m] <= A[lo]

```
A. m = lo
```

 $B. \quad A[lo] \le A[lo]$

C. $A[m] \leq A[lo]$

by line 6

by math

by math on A and B

```
1 int find_min (int[] A, int lo, int hi)
 2 //@requires 0 <= lo && lo < hi && hi <= \length(A);</pre>
   //@ensures lo <= \result && \result < hi;</pre>
 4 //@ensures A[\result] <= A[lo];</pre>
 5 {
      int m = lo;
      for (int j = lo + 1; j < hi; j++)
      //@loop invariant lo < j && j <= hi;</pre>
      //@loop_invariant lo <= m && m < hi;</pre>
      //@loop_invariant A[m] <= A[lo];</pre>
11
12
          if (A[j] < A[m])
13
             m = j;
14
15
      return m;
16 }
```

PRES Example

Assuming: A[m] <= A[lo] **To Show:** A[m'] <= A[lo']

A. A' = A

B. lo' = lo

by A unchanged by lo unchanged

Case 1: A[j] < A[m]

C. m' = j

D. A[j] < A[m] <= A[lo]

E. $A[j] \leq A[lo]$

F. $A[m'] \leq A[lo']$

by line 13

by case, assumption

by math on D

by A, B, math on C and E

Case 2: A[j] >= A[m]

G. m' = m

H. $A[m'] \leq A[lo']$

by m unchanged

by math on B, G & assum.

```
1 int find min (int[] A, int lo, int hi)
 2 //@requires 0 <= lo && lo < hi && hi <= \length(A);</pre>
  //@ensures lo <= \result && \result < hi;</pre>
 4 //@ensures A[\result] <= A[lo];
 5 {
      int m = lo;
      for (int j = lo + 1; j < hi; j++)
      //@loop invariant lo < j && j <= hi;</pre>
      //@loop_invariant lo <= m && m < hi;</pre>
      //@loop invariant A[m] <= A[lo];</pre>
          if (A[j] < A[m])
13
             m = j;
14
15
      return m;
16 }
```

EXIT Example

To Show: A[\result] <= A[lo];

```
A. \result = m by line 15

B. A[m] <= A[lo] by line 10

C. <u>A[\result] <= A[lo]</u> by math on A and B
```

```
1 int find_min (int[] A, int lo, int hi)
 2 //@requires 0 <= lo && lo < hi && hi <= \length(A);</pre>
   //@ensures lo <= \result && \result < hi;</pre>
 4 //@ensures A[\result] <= A[lo];
 5 {
      int m = lo;
      for (int j = lo + 1; j < hi; j++)
      //@loop invariant lo < j && j <= hi;</pre>
      //@loop_invariant lo <= m && m < hi;</pre>
      //@loop_invariant A[m] <= A[lo];</pre>
11
12
          if (A[j] < A[m])
13
             m = j;
14
15
      return m;
16 }
```

Thanks for Coming!



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