Ben Bitdiddle's home network connection can upload at 125,000 bytes/second. His router has a 100,000 byte first in first out buffer for packets awaiting transmission.

If the buffer is completely full, how long will it take for the buffer to clear?

- A. 0.4 seconds
- B. 0.6 seconds
- C. 0.8 seconds
- D. 1 second
- E. 1.25 seconds

Ben Bitdiddle's home network connection can upload at 125,000 bytes/second. His router has a 100,000 byte first in first out buffer for packets awaiting transmission.

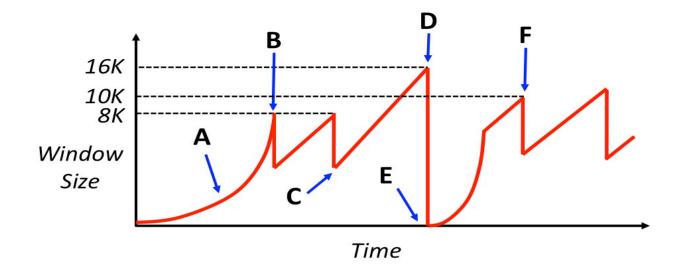
At time 0, Ben's client starts sending 1,000 byte packets at 150 packets/s. When will the first packet be dropped by the router?

- A. 2 seconds
- B. 3 seconds
- C. 4 seconds
- D. Buffer will never discard a packet in this case

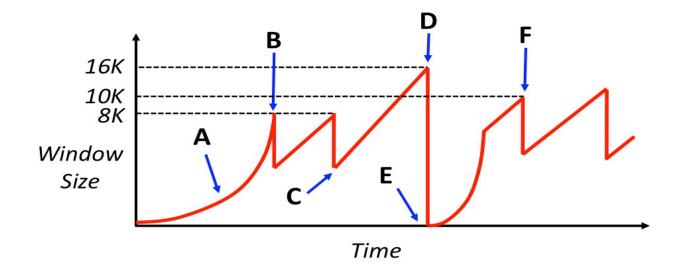
Alyssa P. Hacker and Ben Bitdiddle communicate over a link with capacity of 100 pkts / sec. The latency (RTT) on this link is 100 ms.

If a sliding window protocol with acknowledgement packets is used, and there is a *FIXED* window size of 4 packets, what is the maximum rate of traffic on the link?

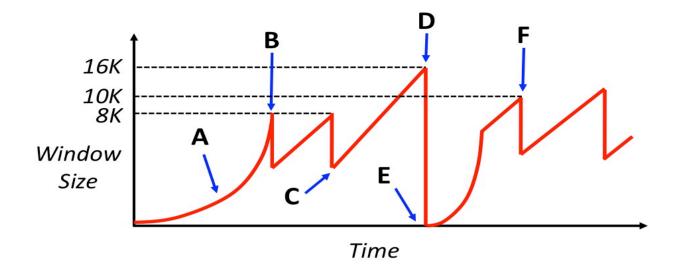
- A. 20 pkts / s
- B. 40 pkts / s
- C. 80 pkts / s
- D. 100 pkts /s



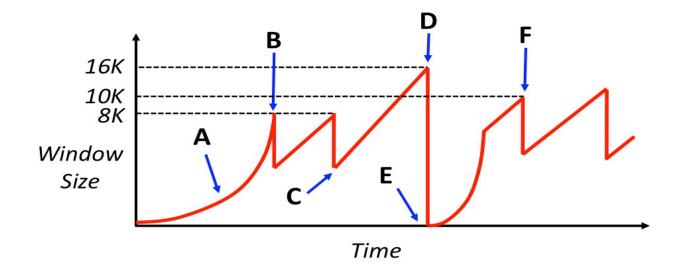
- 1. Name the event at B which occurs that causes the sender to decrease its window
 - (a) Triple Duplicate Ack
 - (b) Slow Start
 - (c) Packet loss
 - (d) Time out



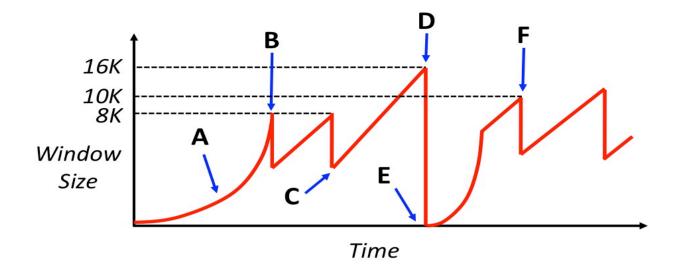
- 2. Does the event at B necessitate that the network discarded a packet ?
 - (a) Yes
 - (b)No
 - (c) Don't know



- 3. Name the event at D which occurs that causes the sender to decrease its window.
 - (a) Triple Duplicate Ack
 - (b) Slow Start
 - (c) Packet loss
 - (d) Time out

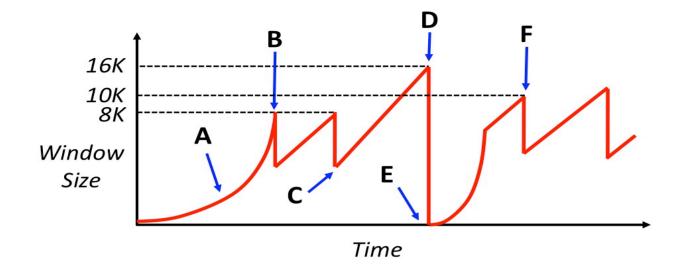


- 4. Does the event at D necessitate that the network discarded a packet
 - (a) Yes
 - (b)No
 - (c) Don't know

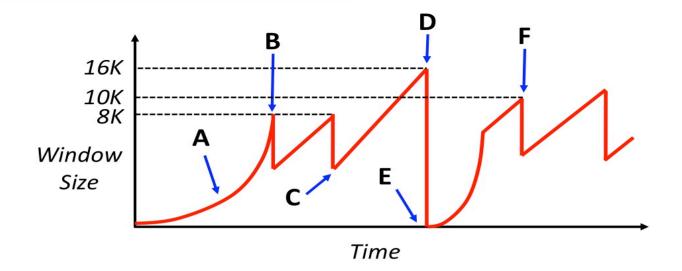


5. For a lightly-loaded network, is the event at D MORE likely or LESS likely to occur when the sender has multiple TCP segments outstanding

- (a) MORE
- (b) LESS
- (c) ALMOST SAME



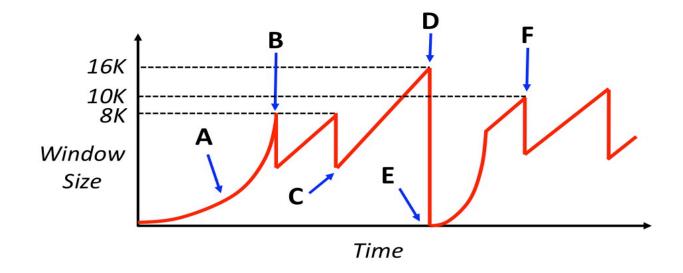
6. Consider the curved slope labeled by point A. Why does the TCP window behave in such a manner, rather than have a linear slope? (Put another way, why would it be bad if region A had a linear slope?)



Assume that the network has an MSS of 1000 bytes and the round-trip-time between sender and receiver of 100 milliseconds.

Assume at time 0 the sender attempts to open the connection.

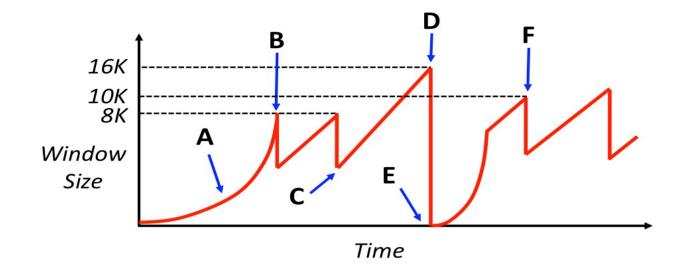
Also assume that the sender can "write" a full window's worth of data instantaneously, so the only latency you need to worry about is the actual propagation delay of the network.



RTT = 100ms, MSS = 1000 bytes

7. How much time has progressed by point B?

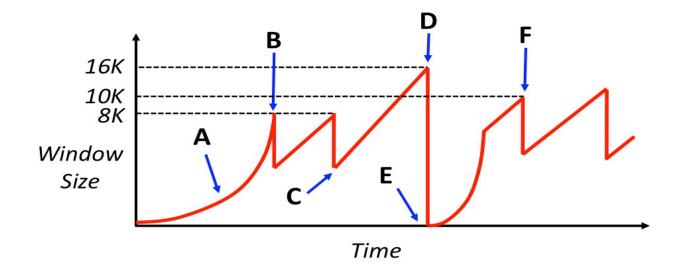
- (a)200ms
- (b)300ms
- (c) 400ms
- (d)600ms
- (e)700ms



RTT = 100ms, MSS = 1000 bytes

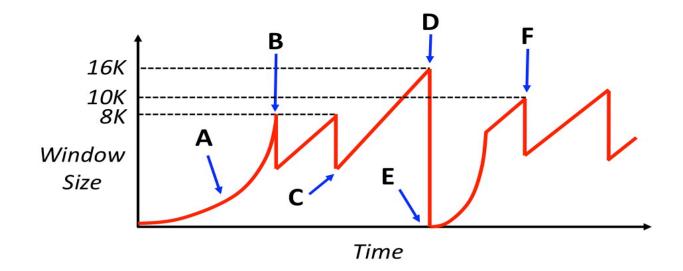
8. How much time has progressed between points C and D?

- (a)800ms
- (b) 1000ms
- (c) 1200ms
- (d) 1400ms



RTT = 100ms, MSS = 1000 bytes

- 9. How much time has progressed between points E and F?
- (a) 400ms
- (b)500ms
- (c)800ms
- (d) 900ms



10. If the sender shares its network with other clients whose traffic traverses the same IP routers, give one explanation for why point D is higher than point B?

1 In slow start, a sender doubles its window size every RTT if all sent packets were acknowledged

T/F