

Prolog use cases *other than* genealogy (Part I/II)

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<http://developers.svn.sourceforge.net/viewvc/developers/repository/prolog-crash-course/>

What's Prolog?

- A language based on **logic** (say, Hoare clauses).
- A full-blown **declarative** programming language.
- A super-weapon of a computer scientist.

To be continued.

Simple examples

main :-

hello.pro

write('Hello, world!'),
nl.

\$ swipl

Welcome to SWI-Prolog (Multi-threaded, 64 bits, Version 5.10.4)

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?- ['hello.pro'].

% hello.pro compiled 0.00 sec, 992 bytes

true.

?- main.

Hello, world!

true.

?- halt.

\$

... or use CTRL-D

```
main :-
```

```
    write('Hello, world!'),  
    nl.
```

```
auto.pro
```

```
:- main, halt.
```

```
$ swipl -f auto.pro
```

```
Hello, world!
```

```
$
```

The genealogy use case

% Steve's adopted parents

```
sex(steve,male).  
father(paul,steve).  
mother(clara,steve).
```

% Steve's biological parents

```
father(abdul,steve).  
mother(joanne,steve).
```

% Sister of Steve

```
sex(mona,female).  
father(abdul,mona).  
mother(joanne,mona).
```

% Steve's daughter back from his sterile period

```
sex(lisa,female).  
father(steve,lisa).  
mother(anne,lisa).
```

...

<http://www.applegazette.com/feature/the-family-tree-of-steve-jobs/>

Genealogy relations

```
grandfather(X,Y) :-  
    father(X,Z),  
    father(Z,Y).
```

```
sibling(X,Y) :-  
    father(F,X),  
    father(F,Y),  
    mother(M,X),  
    mother(M,Y),  
    X \== Y.
```

```
sister(X,Y) :-  
    sibling(X,Y),  
    sex(X,female).
```

Prolog queries

% Do we know who Steve's grandfather is?

?- grandfather(X,steve).
false.

% Do we know who Reed's grandfather is?

?- grandfather(X,reed).
X = paul ;
X = abdul ;
false.

Genealogy relations cont'd

```
halfsister(X,Y) :-  
    sex(X,female),  
    father(FX,X),  
    mother(MX,X),  
    father(FY,Y),  
    mother(MY,Y),  
    overlap(FX,FY,MX,MY).
```

```
overlap(F,F,MX,MY) :- MX \== MY.  
overlap(FX,FY,M,M) :- FX \== FY.
```

Use of “disjunction”

```
halfsister(X,Y) :-  
    sex(X,female),  
    father(FX,X),  
    mother(MX,X),  
    father(FY,Y),  
    mother(MY,Y),  
    ( FX == FY, MX \== MY; FX \== FY, MX == MY ).
```

List processing

```
member(X,[X|T]).  
member(X,[_|T]) :- member(X,T).
```

```
append([H|T],L2,[H|R]) :- append(T,L2,R2).  
append([],R,R).
```

```
?- member(X,[a,b,c]).
```

```
X = a ;
```

```
X = b ;
```

```
X = c.
```

```
?- append([1,2,3],[4,5,6],X).
```

```
X = [1, 2, 3, 4, 5, 6].
```

Directed graphs

```
node(1).  
node(2).  
node(3).
```

```
edge(1,2).  
edge(2,3).
```

```
connected(X,Y) :-  
    edge(X,Y).
```

```
connected(X,Y) :-  
    edge(X,Z),  
    connected(Z,Y).
```

```
?- connected(1,2).  
true
```

```
?- connected(1,3).  
true
```

```
?- connected(2,1).  
false
```

Implementing Peano axioms

```
add(zero,X,X).  
add(succ(X),Y,succ(Z)) :- add(X,Y,Z).
```

```
?- add(succ(succ(zero)),succ(zero),X).  
X = succ(succ(succ(zero))).
```

```
?- eval(add(add(num(1),num(2)),num(3)),X).  
X = 6.
```

A simple expression interpreter

```
eval(num(N),N) :-  
    number(N).
```

```
eval(add(E1,E2),N) :-  
    eval(E1,N1),  
    eval(E2,N2),  
    N is N1 + N2.
```

```
?- eval(add(add(num(1),num(2)),num(3)),X).  
X = 6.
```

Totaling salaries

<http://10lcompanies.org/index.php/10limplementation:prolog>

```
total(company(_,Ds),R) :-  
    total(Ds,R).
```

```
total([],0).
```

```
total([H|T],R) :-  
    total(H,R1),  
    total(T,R2),  
    R is R1 + R2.
```

```
total(dept(_,M,Units),R) :-  
    total(M,R1),  
    total(Units,R2),  
    R is R1 + R2.
```

```
total(employee(_,_,S),S).
```

```
?- total(company(me,[dept(leadership,employee(ralf,b127,42),[])]),X).  
X = 42.
```

Cutting salaries

<http://l0lcompanies.org/index.php/l0limplementation:prolog>

```
cut( company(N,Ds1),  
      company(N,Ds2)) :-  
    cut(Ds1,Ds2).
```

```
cut(N1,N2) :-  
    number(N1), N2 is N1 / 2.
```

```
cut([],[]).  
cut([H1|T1],[H2|T2]) :-  
    cut(H1,H2), cut(T1,T2).
```

```
cut( dept(X,M1,Units1),  
      dept(X,M2,Units2)) :-  
    cut(M1,M2),  
    cut(Units1,Units2).
```

```
cut( employee(X,Y,S1),  
      employee(X,Y,S2)) :-  
    cut(S1,S2).
```

```
?- cut(company(me,[dept(leadership,employee(ralf,b127,42),[])]),X).  
X = company(me, [dept(leadership, employee(ralf, b127, 21), [])])
```


Prolog — why?

- Highly declarative.
- Highly operational.
- Highly scripted.
- Highly untyped.
- Highly typeable.
- Highly debuggable.
- **Highly under-appreciated.**
- ...

*A super-weapon
of a
computer scientist*

Prerequisites

- Propositional logic
- Predicate logic
- Herbrand universe
- Unification
- SLD resolution

I/O

File I/O Edinburgh style

```
test :-  
    see('eval.sample'),  
    read(E),  
    seen,  
    eval(E,V),  
    write(V),  
    nl.
```

```
?- test.  
6  
true.
```

File I/O ISO style

```
test :-  
    open('eval.sample',read,In),  
    read(In,E),  
    close(In),  
    eval(E,V),  
    write(V),  
    nl.
```

```
?- test.  
6  
true.
```

I/O predicates

- `see/1`: open file for input, set it as current input
- `seen/0`: close current input, return to previous one
- `read/1`: read a term from the input
- `tell/1`: open file for output, set it as current output
- `told/0`: close current output, return to previous one
- `write/1`: write a term to the output
- `nl/0`: start a new line in the output
- `format/2`: formatted output
- `open/3`: open a stream for input or output
- `close/1`: close a stream
- `write/2`: write a term to a stream

Debugging

Debugging with traces

```
?- trace, expr(add(num(1),num(2))).  
  Call: (7) expr(add(num(1), num(2))) ? creep  
  Call: (8) expr(num(1)) ? creep  
  Call: (9) number(1) ? creep  
  Exit: (9) number(1) ? creep  
  Exit: (8) expr(num(1)) ? creep  
  Call: (8) expr(num(2)) ? creep  
  Call: (9) number(2) ? creep  
  Exit: (9) number(2) ? creep  
  Exit: (8) expr(num(2)) ? creep  
  Exit: (7) expr(add(num(1), num(2))) ? creep  
true.
```


Breakpoints

```
?- spy(number/1).
```

```
% Spy point on number/1  
true.
```

```
[debug] ?- expr(add(num(1),num(2))).
```

```
* Call: (8) number(1) ? creep
```

```
* Exit: (8) number(1) ? creep
```

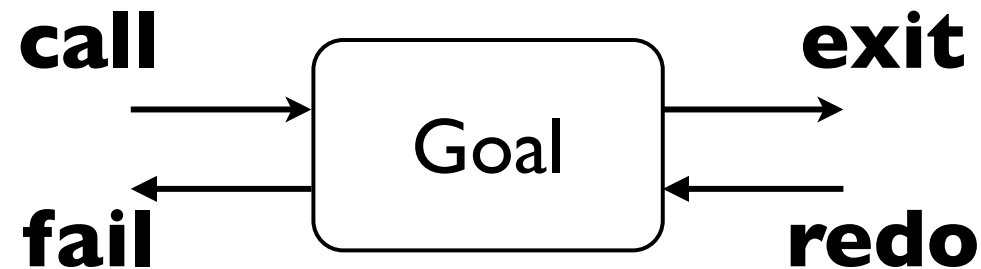
```
Exit: (7) expr(num(1)) ? leap
```

```
* Call: (8) number(2) ? leap
```

```
* Exit: (8) number(2) ? leap
```

```
true.
```

The *Box Model* of goal execution



- **call**: enter the goal when first attempting proof
- **exit**: leave the goal when completing proof
- **redo**: re-entering goal upon backtracking
- **fail**: ultimately finishing goal when without (further) proof

```
?- spy(number/1).  
% Spy point on number/1  
true.
```

```
[debug] ?- expr(add(num(1),num(2))).
```

```
* Call: (8) number(1) ? creep
```

```
* Exit: (8) number(1) ? creep
```

```
Exit: (7) expr(num(1)) ? creep
```

```
Call: (7) expr(num(2)) ? Options:
```

+:	spy	-:	no spy
/c e r f u a goal:	find	.:	repeat find
a:	abort	A:	alternatives
b:	break	c (ret, space):	creep
[depth] d:	depth	e:	exit
f:	fail	[ndepth] g:	goals (backtrace)
h (?):	help	i:	ignore
l:	leap	L:	listing
n:	no debug	p:	print
r:	retry	s:	skip
u:	up	w:	write
m:	exception details		
C:	toggle show context		

```
Call: (7) expr(num(2)) ? skip
```

```
Exit: (7) expr(num(2)) ?
```

“skip” can be used to
go from “call” to “exit”
port right away.

Types and modes

“Types are programs.”

```
expr(num(N)) :- number(N).  
expr(add(E1,E2)) :- expr(E1), expr(E2).
```

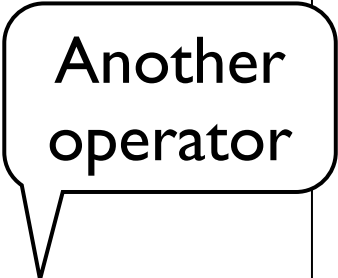
```
?- expr(add(num(1),num(2))).  
true.
```

```
?- expr(foo).  
false.
```

Another example: finding the max leaf in a tree

```
tree(leaf(X)) :- integer(X).  
tree(fork(T1,T2)) :- tree(T1), tree(T2).
```

```
max(leaf(X),X).  
max(fork(T1,T2),X) :- max(T1,Y), max(T2,Z), X is max(Y,Z).
```



Another
operator

```
?- max(fork(leaf(1),fork(leaf(42),leaf(88))),X).  
X = 88.
```

Built-in type tests

- number/1
- integer/1
- atom/1
- is_list/1
- ...

```
?- number(1.1).  
true.  
?- number(foo).  
false.  
?- integer(1.1).  
false.  
?- integer(42).  
true.  
?- atom(42).  
false.  
?- atom(foo).  
true.  
?- is_list(foo).  
false.  
?- is_list([foo]).  
true.
```

Flexible modes

?- add(X,Y,Z).

X = zero,

Y = Z ;

X = succ(zero),

Z = succ(Y) ;

X = succ(succ(zero)),

Z = succ(succ(Y)) .

?- add(X,Y,succ(succ(zero))).

X = zero,

Y = succ(succ(zero)) ;

X = Y, Y = succ(zero) ;

X = succ(succ(zero)),

Y = zero ;

false.

Inflexible modes

?- X is 1 + 1.
X = 2.

?- 2 is 1 + 1.
true.

?- 2 is X + 1.
ERROR: is/2: Arguments are not sufficiently instantiated

Documentation of modes

- Modes
 - `+`: needs to be instantiated upon call
 - `-`: will be instantiated upon exit
 - `?`: neither of the two above
- Application to example
 - `add(?X,?Y,?Z)`
 - `is(-X,+Y)`

Mode not sufficient here.
We need groundness.

Examples of modes in the list library

- `member(?Elem, ?List)`
- `append(?List1, ?List2, ?List1AndList2)`
- `append(+ListOfLists, ?List)`
- `selectchk(+Elem, +List, -Rest)`
- `permutation(?Xs, ?Ys)`
- `subset(+SubSet, +Set)`

Basic modularization

```
:- ['Company.pro'].
:- ['Total.pro'].
:- ['Cut.pro'].
:- ['Depth.pro'].

:-
    see('sampleCompany.trm'),
    read(C1),
    seen,
    isCompany(C1),
    total(C1,R1),
    format('total = ~w~n',[R1]),
    cut(C1,C2),
    total(C2,R2),
    format('cut = ~w~n',[R2]),
    depth(C1,R3),
    format('depth = ~w~n',[R3]).

:- halt.
```

Loading all 101 companies
modules and running tests.

% That's a term to be "read".

```
company(  
  'meganalysis',  
  [ dept(  
    'research',  
    employee('Craig','Redmond',123456),  
    [ employee('Erik','Utrecht',12345),  
      employee('Ralf','Koblenz',1234)  
    ]  
  ),  
  dept(  
    'dev',  
    employee('Ray','Redmond',234567),  
    [ dept(  
      'dev1',  
      employee('Klaus','Boston',23456),  
      [ dept(  
        'dev1.1',  
        employee('Karl','Riga',2345),  
        [ employee('Joe','Wifi City',2344)  
        ]  
      )  
    ]  
  )  
  ]  
)  
].
```

Basic modularization

% Basic form of input

`:- consult('MyPrologFile.pro').`

% Concise notation

`:- ['MyPrologFile.pro'].`

% Ensure import (avoid repeated import)

`:- ensure_loaded('MyPrologFile.pro').`

Related predicates

% Predicate may be defined in more than file.

:- multifile father/2.

% Clauses may appear discontinuously in file.

:- discontinuous father/2.

% Re-load all files (typically after edits).

:- make.

Declarative vs. operational

Lists versus sets of answers

```
max(X,Y,X) :- X >= Y.  
max(X,Y,Y) :- X <= Y.
```

A single answer is
preferred.

```
?- max(42,88,X).  
X = 88.
```

```
?- max(42,42,X).  
X = 42 ;  
X = 42.
```

Efficiency

```
max(X,Y,X) :- X >= Y.  
max(X,Y,Y) :- X < Y.
```

Backtracking
ultimately fails.

```
?- max(42,88,X).  
X = 88.
```

```
?- max(42,42,X).  
X = 42 ;  
false.
```

Operational reasoning

$\text{max}(X,Y,X) \text{ :- } X \geq Y, !.$
 $\text{max}(X,Y,Y) \text{ :- } X < Y.$

A green cut

No more superfluous
backtracking

$?- \text{max}(42,88,X).$
 $X = 88.$

$?- \text{max}(42,42,X).$
 $X = 42.$

Destroyed declarative semantics

```
max(X,Y,X) :- X >= Y, !.  
max(X,Y,Y).
```

A red cut

No problem?

```
?- max(42,88,X).  
X = 88.
```

```
?- max(42,42,X).  
X = 42.
```

A red cut

$\text{max}(X,Y,X) :- X \geq Y, !.$
 $\text{max}(X,Y,Y).$

A green cut

$\text{max}(X,Y,X) :- X \geq Y, !.$
 $\text{max}(X,Y,Y) :- X < Y.$

?- $\text{max}(88,42,42).$
true.

?- $\text{max}(88,42,42).$
false.

Structured cut

```
(If -> Then); _Else :- If, !, Then.  
(If -> _Then); Else :- !, Else.  
If -> Then :- If, !, Then.
```

```
max(X,Y,Z) :-  
  X >= Y -> Z = X; X = Y.
```

Looks all good!

```
?- max(42,88,X).  
X = 88.
```

```
?- max(42,42,X).  
X = 42.
```

```
?- max(42,88,42).  
false.
```

Graph example

```
connected(X,Y) :-  
    edge(X,Y).
```

```
connected(X,Y) :-  
    edge(X,Z),  
    connected(Z,Y).
```

```
connected(X,Y) :-  
    edge(X,Y) ->  
    true;  
    edge(X,Z),  
    connected(Z,Y).
```


Free and bound variables

Terms with variables

- **var/!**: test a term to be a variable
- **ground/!**: test a term to be ground

?- var(42).
false.

?- var(X).
true.

?- X=42, var(X).
false.

?- var(foo(X)).
false.

?- ground(42).
true.

?- ground(X).
false.

?- X=42, ground(X).
X = 42.

?- ground(foo(X)).
false.

Use of non-ground terms

```
?- member(Y,[X,Z]).  
Y = X ;  
Y = Z.
```

```
member(X,[X|T]).  
member(X,[_|T]) :- member(X,T).
```

```
?- varmember(Y,[X,Z]).  
false  
?- varmember(X,[X,Z]).  
true
```

```
varmember(V,[H|_]) :- V==H.  
varmember(V,[H|T]) :- V\==H, varmember(V,T).
```

Term de-/composition

Inspection of terms

- **functor/3**: observe functor symbol and arity
- **=../2**: take apart compound terms

```
?- functor(foo(bar),X,A).  
X = foo,  
A = 1.
```

```
?- foo(bar) =.. X.  
X = [foo, bar].
```

Print terms with indentation

```
print_term(T) :-  
    print_term(T,0).
```

```
print_term(T,N) :-  
    spaces(N),  
    ( var(T) ->  
        format('~w~n',[T])  
    ; T =.. [F|Ts],  
        format('~w~n',[F]),  
        M is N + 1,  
        print_terms(Ts,M) ).
```

```
print_terms([],_).
```

```
print_terms([H|T],N) :-  
    print_term(H,N),  
    print_terms(T,N).
```

```
spaces(N) :-  
    N > 0 -> write(' '), M is N - 1, spaces(M); true.
```

```
?- print_term(add(num  
(1),add(num(2),num(3)))).  
add  
  num  
    1  
  add  
    num  
      2  
    num  
      3  
true.
```

Application to *Programming Language Theory*

A trivial imperative language: ***syntax***

```
program(Es) :- exprs(Es).
```

```
exprs([]).
```

```
exprs([E|Es]) :- expr(E), exprs(Es).
```

```
expr(N) :- number(N).
```

```
expr(E1+E2) :- expr(E1), expr(E2).
```

```
expr(V) :- atom(V).
```

```
expr(V=E) :- atom(V), expr(E).
```

```
?- program([x=1,y=x+41]).  
true
```


A trivial imperative language: ***interpretation***

`eval(Es,V) :- eval(Es,V,[],_).`

`eval([E],N,M1,M2) :-
 eval(E,N,M1,M2).`

`eval([E|Es],N,M1,M2) :-
 Es \== [], eval(E,_,M1,M0), eval(Es,N,M0,M2).`

`eval(N,N,M,M) :-
 number(N).`

`eval(E1+E2,N,M1,M2) :-
 eval(E1,N1,M1,M0), eval(E2,N2,M0,M2), N is N1+N2.`

`eval(V,N,M,M) :-
 atom(V), lookup(V,M,N).`

`eval(V=E,N,M1,M2) :-
 atom(V), eval(E,N,M1,M0), update(V,N,M0,M2).`

`?- eval([x=1,y=x+41],N).
N = 42`

List-processing convenience

```
lookup(V,[(V,N)|_],N).
```

```
lookup(V,[(W,_)|R],N) :- V \== W, lookup(V,R,N).
```

```
update(V,N,[],[(V,N)]).
```

```
update(V,N,[(V,_)|R],[(V,N)|R]).
```

```
update(V,N,[(W,M)|R],[(W,M)|S]) :- V \== W, update(V,N,R,S).
```

End of Lecture (EOL)