## **Scaling Definitions**

# Proposed Scaling Definitions (July 2025 update)

## 1) Consensus (Old)

## sus (Old) 1) Consensus (New)

Define

 $T_R$  = 3000 bytes; Reference Transaction weight. Note:  $T_R$  must be greater than  $T_2$ .  $T_2$  equals the weight in bytes of a 2 input and 2 output transaction.

 $Z_M = 300000$  bytes; Minimum penalty free zone.

 $M_B$  = Block weight in bytes.

 $M_L$  = The median over the last 100000 blocks of max((min (M\_B , 1.7M\_L), Z\_M , M\_L/1.7) ; recursive calculation for M\_L with M\_L starting at M\_L of previous 100001 block (currently =  $Z_{\mbox{\scriptsize M}}$ ); Long term median

 $M_L$ ; Dynamic penalty free zone.

 $M_{\text{S}}$  = the median over the last 100 blocks of max( $M_{\text{B}}$ ,  $M_{\text{L}}$ ); Effective short term median.

 $M_N = min(M_S, 50M_L)$ ; Median for Penalty calculation.

R<sub>Base</sub> = Block Reward.

 $0 < M_B \le 2M_N$ ; Requirement for valid block.

 $B = M_B / M_N - 1$  where  $-1 < B \le 1$ 

 $P_B = R_{Base}B^2$  for B > 0; Monero applies a penalty  $P_B$ , to increase the block weight by B;  $P_B = 0$  for  $B \le 0$ 

 $T_R$  = 10000 bytes; Reference Transaction weight. Note:  $T_R$  must be greater than  $T_2$ .  $T_2$  equals the weight in bytes of a 2 input and 2 output transaction.

 $Z_M = 1000000$  bytes; Minimum penalty free zone.

 $M_B$  = Block weight in bytes.

Define

 $M_L$  = The median over the last 100000 blocks of max((min  $(M_B\,,\,2M_L),\,Z_M\,,\,M_L/2)$ ; recursive calculation for  $M_L$  with  $M_L$  starting at  $M_L$  of previous 100001 block (currently =  $Z_M$ ); Long term median

 $M_{L}$ ; Dynamic penalty free zone.

 $M_S$  = the median over the last 100 blocks of max(min( $M_B$ , 16 $M_L$ ),  $M_L$ ); Effective short term median.

 $M_N = M_S$ ; Median for Penalty calculation.

 $R_{Base}$  = Block Reward.

 $0 < M_B \le 2M_N$ ; Requirement for valid block.

 $B = M_B / M_N - 1$  where  $-1 < B \le 1$ 

 $P_B = R_{Base}B^2$  for B > 0; Monero applies a penalty  $P_B$ , to increase the block weight by B;  $P_B = 0$  for  $B \le 0$ 

#### Changes

- 1) Maximum growth of  $M_S$  is reduced from  $50M_L$  to  $16M_L$
- 2) Rate of growth and decline of  $M_L$  is increased from 1.7x to 2x
- 3)  $T_R$  is increased from 3000 bytes to 10000 bytes
- 4)  $Z_M$  is increased from 300000 bytes to 1000000 bytes

## 2a) Minimum Fee For Node Relay (Old)

to a block of weight MB

Define

$$B_T = T_T / M_N$$

 $P_{BT} = R_{Base}(B+B_T)^2 = R_{base}(B^2 + 2BB_T + B_T^2)$ ; The new penalty, where B +  $B_T > 0$ 

 $P_T = P_{BT} - P_B = R_{Base}(2BB_T + B_T^2)$ ; Increase in penalty from adding transaction T

 $F_T = R_{Base}(2BB_T + B_T^2)$ ; The additional fee required to overcome the increase in penalty  $P_{T}$ 

For the case B = 0 this reduces to  $F_T = R_{Base}B_T^2$ 

 $M_F = M_L$ ; Median for minimum fee calculation

To calculate the minimum fee we consider a transaction of weight  $T_R$  at the start of the penalty, B = 0 with  $M_N = M_F$ 95% of the fee required to pay the penalty incurred is the minimum fee.

$$B_{RL} = T_R / M_F$$
;

 $F_R = R_{Base} B_{RL}^2$ ; Fee required to pay the penalty incurred

 $f_R = R_{Base}B_{RL}/M_F$ ; Fee required to pay the penalty incurred per byte for a given M<sub>F</sub>

 $f_1 = 0.95 f_R$ ; Minimum fee per byte

## 2a) Minimum Fee For Node Relay (New)

We add a, penalty attracting, transaction T with a size of  $T_T$  We add a, penalty attracting, transaction T with a size of  $T_T$ to a block of weight M<sub>B</sub>

Define

$$B_T = T_T / M_N$$

 $P_{BT} = R_{Base}(B+B_T)^2 = R_{base}(B^2 + 2BB_T + B_T^2)$ ; The new penalty, where B +  $B_T > 0$ 

 $P_T = P_{BT} - P_B = R_{Base}(2BB_T + B_T^2)$ ; Increase in penalty from adding transaction T

 $F_T = R_{Base}(2BB_T + B_T^2)$ ; The additional fee required to overcome the increase in penalty  $P_{T}$ 

For the case B = 0 this reduces to  $F_T = R_{Base}B_T^2$ 

 $M_F = M_L$ ; Median for minimum fee calculation

To calculate the minimum fee we consider a transaction of weight  $T_R$  at the start of the penalty, B = 0 with  $M_N = M_E$ 95% of the fee required to pay the penalty incurred is the minimum fee.

$$B_{RL} = T_R / M_F$$
;

 $F_{RL} = R_{Base}B_{RL}^2$ ; Fee required to pay the penalty incurred

 $f_{RL} = R_{Base}B_{RL}/M_F$ ; Fee required to pay the penalty incurred per byte for a given M<sub>F</sub>

 $f_{IL} = 0.95f_{R}$ ; Minimum fee per byte

Additional, minimum fee, node relay requirements for large transactions, if such transactions are permitted by consensus

If a transaction has a weight  $T_T > 10000$  bytes and / or more than 8 inputs and / or more than 8 outputs

$$f_{IN} = 4f_{IL}$$

If a transaction has a weight  $T_T > 20000$  bytes and / or more than 16 inputs and / or more than 16 outputs

$$f_{\text{IM}}=16(M_{\text{F}}/Z_{\text{M}})f_{\text{IL}}$$

If a transaction has a weight  $T_T > 40000$  bytes and / or more than 32 inputs and / or more than 32 outputs

$$f_{IH} = 200(M_F/Z_M)f_{IL}$$

## 2b) Wallet Fees (Old)

For the calculation of wallet fees we assume that the next 10 blocks have no transactions, other than the coinbase transaction, the empty blocks, We then calculate M<sub>LW</sub> and M<sub>SW</sub> by following the calculation of M<sub>L</sub> and M<sub>S</sub> at this future point. We use the previous 99990 blocks and the future 10 empty blocks (100000 blocks) for M<sub>1</sub> and the previous 90 blocks and future 10 blocks (100) blocks for Ms.

## 2b) Wallet Fees (New)

For the calculation of wallet fees we assume that the next 1000 blocks have no transactions, other than the coinbase transaction, the empty blocks, We then calculate M<sub>LW</sub> by following the calculation of  $M_L$  at this future point. We use the previous 99000 blocks and the future 1000 empty blocks (100000 blocks) for M<sub>1</sub>.

#### Define

 $M_{BW} = M_B$  for the last 99990 blocks

 $M_{BW}$  = 0 for the future 10 blocks; A value of 0 bytes can be used for the empty blocks for the purposes of calculating  $M_{LW}$ .

M<sub>LW</sub> = The median over the last 99990 blocks and future 10 M<sub>LW</sub> = The median over the last 99000 blocks and future blocks (100000 blocks) of max((min (MBW, 2ML), ZM,  $M_1/2$ ): The current value of  $M_1$  from consensus is used: Effective long term median for wallet fees

M<sub>LW</sub>; Penalty free zone for wallet fees

M<sub>SW</sub> = the median over the last 90 blocks and future 10 blocks (100 blocks) of max(M<sub>BW</sub>, M<sub>LW</sub>); Effective short term median for wallet fees

 $M_{NW} = min(M_{SW}, 50M_{LW})$ 

 $M_{FW} = M_{LW}$ ; Median for wallet fee calculation

 $B_{RLW} = T_R / M_{EW}$ ; Used for the low and normal fees

 $B_R = T_R / Z_M$ ; Used for the medium and high fees

 $F_L = R_{Base} B_{RLW}^2$ ; Low transaction fee for reference transaction

 $f_L = R_{Base}B_{RLW}/M_{FW}$ ; Low transaction fee per byte for a given  $f_L = R_{Base}B_{RLW}/M_{FW}$ ; Low transaction fee per byte for a given  $M_{FW}$ 

 $f_N = 4f_L$ ; Normal transaction fee per byte for a given  $M_{EW}$ 

 $f_M = 16 R_{Base} B_R / M_{FW}$ ; Medium Transaction fee per byte for a given M<sub>FW</sub>

Transaction fee per byte for a given M<sub>NW</sub>

 $f_H = 4f_M max (1, M_{FW}/(32B_R M_{NW}))$ ; High Transaction fee per byte

#### Define

 $M_{BW} = M_B$  for the last 99000 blocks

 $M_{BW} = 0$  for the future 1000 blocks; A value of 0 bytes can be used for the empty blocks for the purposes of calculating MLW.

1000 blocks (100000 blocks) of max((min ( $M_{BW}$ ,  $2M_L$ ),  $Z_M$ ,  $M_1/2$ ): The current value of  $M_1$  from consensus is used: Effective long term median for wallet fees

M<sub>LW</sub>; Penalty free zone for wallet fees

 $M_{EW} = M_{LW}$ ; Median for wallet fee calculation

 $B_{RLW} = T_R / M_{FW}$ ; Used for the low and normal fees

 $B_R = T_R / Z_M$ ; Used for the medium and high fees

 $F_L = R_{Base} B_{RLW}^2$ ; Low transaction fee for reference transaction

 $f_N = 4f_L$ ; Normal transaction fee per byte for a given  $M_{EW}$ 

 $f_M = 16 R_{Base} B_R / M_{FW}$ ; Medium Transaction fee per byte for a given M<sub>FW</sub>

f<sub>P</sub> = 2R<sub>Base</sub>/M<sub>NW</sub> = f<sub>M</sub>M<sub>FW</sub>/(8B<sub>R</sub>M<sub>NW</sub>); Maximum Penalty (B =1) f<sub>P</sub> = 2R<sub>Base</sub>/M<sub>FW</sub> Maximum Penalty (B =1) Transaction fee per byte assuming  $M_S = M_{FW}$ 

 $f_H = f_P$ ; High Transaction fee per byte

#### Changes

- 1) All fees including the high fee are now based upon M<sub>LW</sub> with the ratio between fees constant for a given M<sub>IW</sub>. 2) The grace period is increased to 1000 blocks.

3) Transitional considerations for Reference transaction,  $T_R$ , Minimum penalty free zone,  $Z_M$ , and Median calculations after the fork.

3) Transitional considerations for Reference transaction,  $T_R$ , Minimum penalty free zone,  $Z_M$ , and Median calculations after the fork.

Calculation of  $M_L$ ,  $M_S$ ,  $M_{LW}$  and  $M_{SW}$  where blocks from a previous the Monero version are included in a calculation after the fork.  $M_B$  is modified as follows:

Define:

 $Z_{MOId}$  = 300000 bytes ( $Z_{M}$  before hard fork)

M<sub>BOld</sub> = Block Weight in bytes (before hard fork)

 $T_R = 10000$  bytes

 $Z_{M} = 1000000 \text{ bytes}$ 

For blocks before the hard fork  $M_B = M_{BOld} (Z_M / Z_{MOld})$ 

The medians are then calculated normally.

## 4) Wallet Fee Rounding

Wallet fees,  $f_N$ ,  $f_L$ ,  $f_M$ , and  $f_H$  are rounded up to the desired number of significant digits in the significant

## **Wallet Fee Rounding Examples**

## Two significant digits

27810 Rounded to : 28000 37.94 Rounded to : 38 0.5555 Rounded to : 0.56 0.002342 Rounded to : 0.0024

## 4) Wallet Fee Rounding

Wallet fees,  $f_N,\,f_L,\,f_M,$  and  $f_H$  are rounded to 2 significant digits in the significant

## **Wallet Fee Rounding Examples**

## Two significant digits

27810 Rounded to : 28000 37.94 Rounded to : 38 0.5555 Rounded to : 0.56 0.002342 Rounded to : 0.0023

## 5) Transaction Weights

Transaction weighs are used to account for the different growth rate the output proof verification time with the number of outputs. This is done at consensus level and can lead to double charging if the fee per byte rate is increased because of an increase in the transaction weight.

## 5) Transaction Weights (Proposed)

Breakdown the transaction weight as follows:

1) Base transaction weight:

Use a standard weights based upon the current size in bytes.

2) Input proofs weight:

Use a standard weights based upon the current FCMP++ proof size in bytes and the number of inputs.

1 input: Current FCMP++ proof size for 1 input

2 inputs: Current FCMP++ proof size for 2 inputs

3 or 4 inputs: Current FCMP++ proof size for 4 inputs

5, 6, 7 or 8 inputs: Current FCMP++ proof size for 8 inputs

More than 8 inputs if applicable: Round up to next power of 2; namely 16 inputs, 32 inputs, etc. and use the current FCMP++ proof size for the corresponding power of 2 number of inputs.

3) Output proofs weight:

Use a standard weights based upon the current proof size in bytes and the number of outputs.

2 outputs: Current proof size for 2 outputs

3 or 4 outputs: Current proof size for 4 outputs

5, 6, 7 or 8 outputs: Current proof size for 8 outputs

More than 8 outputs if applicable: Round up to next power of 2; namely 16 outputs, 32 outputs, etc. and use the current proof size for the corresponding power of 2 number of outputs.

4) TX extra weight:

Use actual size.

The 1), 2) and 3) standard weights do not change if there is a change in the proof size due to significant adoption (this can happen with FCMP++) or a future hard fork. Verification time is then priced at node relay using the multiple node relay fees.