

CS118 Homework 4

1. (Cyclic Redundancy Check) Consider the 7-bit generator, $G = 10011$, and suppose that D has the value 1010101010 . What is the value of R ?

→ For this problem and the next, we are taking G to have 5 bits and $r = 4$. If instead $G = 0010011$ with 7 bits, then $r = 6$ and the answers will slightly vary.

Since $r = 4$, we pad D with four 0s, then “divide” this value by $G = 10011$. The intermediate steps involve XORing the values rather than subtracting. We stop once all $|D| = 10$ bits are found and the remainder R has length $r = |R| = 4$.

$$1010101010(0000) / 10011 = 1011011100 \text{ with } R = 0100$$

*For more details, see last attached page showing handwritten work.

2. Consider the previous problem, but suppose that D has the value:

→ We follow the same method as before with a different D value.

- a. **1001010101**
→ 1000110000 with $R = 0000$
- b. **0101101010**
→ 0101010101 with $R = 1111$
- c. **1010100000**
→ 1011010111 with $R = 1001$

3. Slotted ALOHA

- a. Recall that when there are N active nodes, the efficiency of slotted ALOHA is $Np(1-p)^{N-1}$. Find the value of p that maximizes this expression.

→ In order to find a maximal p , we need to find the roots of the equation above using the derivative and determine if, at these values of p , the expression is at its maximum. We begin by using the product rule and chain rule.

$$f = Np(1-p)^{N-1}$$

$$f = gh$$

$$g = Np$$

$$g' = N$$

$$h = (1-p)^{N-1}$$

$$h' = -(N-1)(1-p)^{N-2}$$

$$f' = g'h + gh'$$

$$f' = N(1-p)^{N-1} + -NP(N-1)(1-p)^{N-2}$$

$$f' = N(1-p)^{N-2}[(1-p) - p(N-1)]$$

$$f' = N(1-p)^{N-2}(1-Np)$$

Thus the roots are $p=1$ and $p=1/N$. However, at $p=1$, we know that $f=0$, unlike when $p=1/N$. Therefore, $p=1/N$ maximizes this expression.

- b. Using the value of p found in (a), find the efficiency of slotted ALOHA by letting N approach infinity.

→ We simply let $p=1/N$ and evaluate f where N approaches infinity. Let N approach infinity for the equations below.

$$f = Np(1-p)^{N-1}$$

$$f = N(1/N)(1-1/N)^{N-1}$$

$$f = (1-1/N)^{N-1}$$

$$f = (1-1/N)^N$$

$$f = 1/e$$

Since N approaches infinity, $N-1$ approaches N .

Therefore, maximum efficiency of slotted ALOHA with infinite N is $1/e$, or about 0.368.

4. Suppose nodes A and B are on the same 10Mbps broadcast channel, and the propagation delay between the two nodes is 325 bit times. Suppose CSMA/CD and Ethernet packets are used for this broadcast channel. Suppose node A begins transmitting a frame and, before it finishes, node B begins transmitting a frame. Can A finish transmitting before it detects that B has transmitted? Why or why not?

→ Assuming that A transmits at time $t=0$ some frame with size of at least $512+64$ bit times. Then in the worst case, B begins transmitting at $t=324$.

With a propagation delay of 325 times, the first packet of B is received at $t=324+325=649$ bit times. Since 649 bit times exceeds 576 bit times, then A may be able to finish transmitting before B is transmitted. (On the other hand, if the number of bit times taken to receive the first packet of B is less than 576 bit times, then A detects transmission of B – in this case, A would abort without finishing its own transmission.)

1.)

$$\begin{array}{r}
 1011011100 \\
 10011 \overline{) 10101010100000} \\
 \underline{10011} \\
 1100 \\
 \underline{10011} \\
 10100 \\
 \underline{10011} \\
 11110 \\
 \underline{10011} \\
 11010 \\
 \underline{10011} \\
 10010 \\
 \underline{10011} \\
 0100 \quad R
 \end{array}$$

2a.)

$$\begin{array}{r}
 1000110000 \\
 10011 \overline{) 10010101010000} \\
 \underline{10011} \\
 11010 \\
 \underline{10011} \\
 10011 \\
 \underline{10011} \\
 00000 \\
 \underline{00000} \\
 00000 \quad R
 \end{array}$$

2b.)

$$\begin{array}{r}
 0101010101 \\
 10011 \overline{) 01011010100000} \\
 \underline{10011} \\
 10110 \\
 \underline{10011} \\
 10110 \\
 \underline{10011} \\
 10110 \\
 \underline{10011} \\
 11100 \\
 \underline{10011} \\
 11111 \quad R
 \end{array}$$

2c.)

$$\begin{array}{r}
 1011010111 \\
 10011 \overline{) 10101000000000} \\
 \underline{10011} \\
 11000 \\
 \underline{10011} \\
 10110 \\
 \underline{10011} \\
 10100 \\
 \underline{10011} \\
 11100 \\
 \underline{10011} \\
 11110 \\
 \underline{10011} \\
 11010 \\
 \underline{10011} \\
 10011 \quad R
 \end{array}$$