

# CS143: Database Systems

## Homework #6

- Suppose that we decompose the schema  $R(A, B, C, D, E, F)$  into  $(A, B, C, F)$  and  $(A, D, E)$ . When the following set of functional dependencies hold, is the decomposition lossless?  
 $A \rightarrow BC, CD \rightarrow E, B \rightarrow D, E \rightarrow A$   
 Explain your answer.
- List non-trivial functional dependencies satisfied by the following relation. You do not need to find all functional dependencies. It is enough to identify a set of functional dependencies that imply all functional dependencies that is satisfied by the relation.

A	B	C
$a_1$	$b_1$	$c_2$
$a_1$	$b_1$	$c_2$
$a_2$	$b_1$	$c_1$
$a_2$	$b_1$	$c_3$

- Assume *Student* and *Class* entity sets that we have used in the class. The *Student* and *Class* sets are connected by Take relationship set. We now convert the Take relationship set into a table **Take(sid, dept, cnum)** using our standard translation algorithm, where sid is the key for a student and (dept, cnum) is the key for a class.  
 Explain how functional dependencies can be used to indicate the following:
  - A one-to-one relationship exists between entity sets *Student* and *Class*.
  - A many-to-one relationship exists between entity sets *Student* and *Class*.
- Assume the following set of functional dependencies hold for the relation  $R(A, B, C, D, E)$ :  
 $A \rightarrow BC, CD \rightarrow E, B \rightarrow D, E \rightarrow A$ 
  - Is  $E$  a key for  $R$ ? Explain your answer.
  - Is  $BC$  a key for  $R$ ? Explain your answer.
- Assume the following set of functional dependencies hold for the relation  $R(A, B, C, D, E, F)$ :  
 $A \rightarrow BC, C \rightarrow E, B \rightarrow D$   
 Is it in **BCNF**? Explain your answer. If it is not, normalize it into a set of relations in **BCNF**.
- Suppose we have a relation  $R(A, B, C, D)$  with a MVD  $A \twoheadrightarrow BC$ . If we know that the tuples  $(a, b_1, c_1, d_1), (a, b_2, c_2, d_2)$  and  $(a, b_3, c_3, d_3)$  are in the current instance of  $R$ , what other tuples do we know must also be in  $R$ ?
- For relation  $R(A, B, C, D, E, F)$ , suppose a FD  $AB \rightarrow E$  and two MVDs  $AB \twoheadrightarrow C$  and  $A \twoheadrightarrow B$  hold. Is it in 4NF? Explain your answer. If not, normalize it into 4NF.