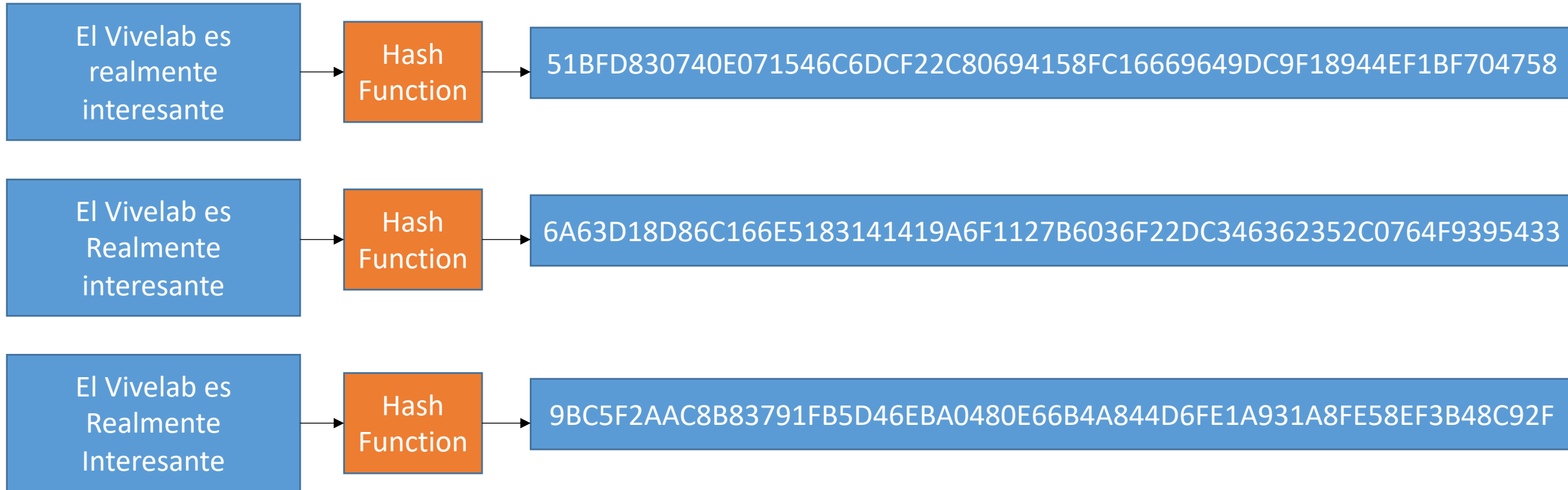


Introducción a la Criptografía y a la Seguridad de la Información

Part 7a
Cryptographic Hash
Functions

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What is a hash function?



A formal definition of hash function

- Deterministic algorithm
- No encryption
- Not signature

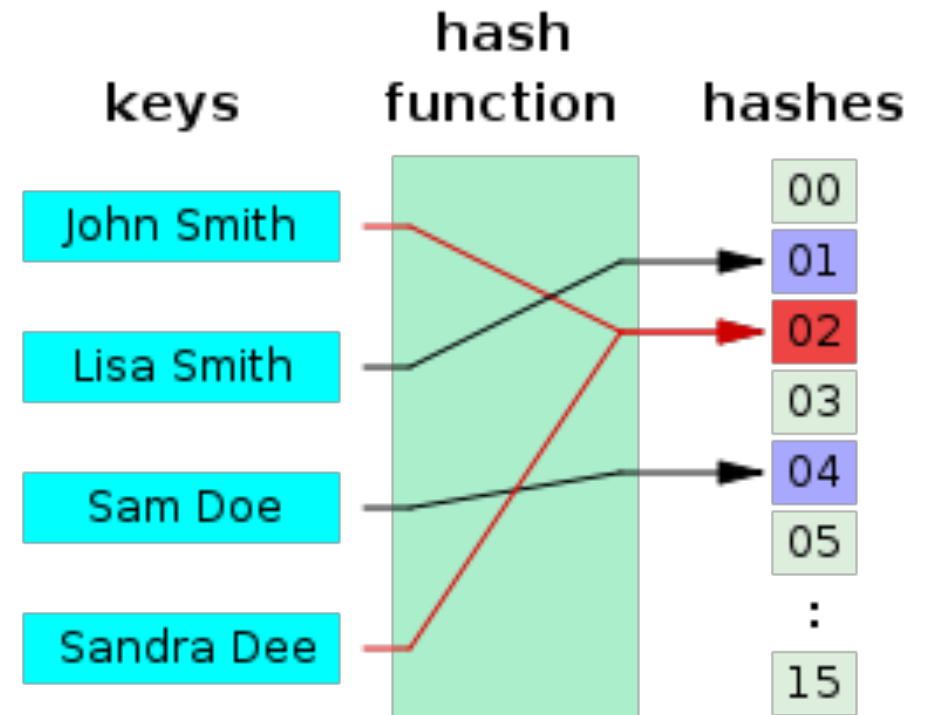
$$H : \left\{ \begin{array}{ll} \text{any bit string} & \rightarrow \text{m bits} \\ x & \mapsto y = H(x) \end{array} \right.$$

Properties of a (cryptographic) hash function

- Deterministic
- Uniformity
- Defined range
- Non-invertibility
- **Pre-image resistance:** Given a hash value h it should be difficult to find any message m such that $h = \text{hash}(m)$
- **Second pre-image resistance:** Given an input $m1$, it should be difficult to find a different input $m2$ such that $\text{hash}(m1) = \text{hash}(m2)$
- **Collision resistance:** It should be difficult to find two different messages $m1$ and $m2$ such that $\text{hash}(m1) = \text{hash}(m2)$

Collisions

- A hash function that maps names to integers form 0 to 15.
- There is a collision between “John Smith” and “Sandra Dee”



MD5 hash function

$$F(B, C, D) = (B \wedge C) \vee (\neg B \wedge D)$$

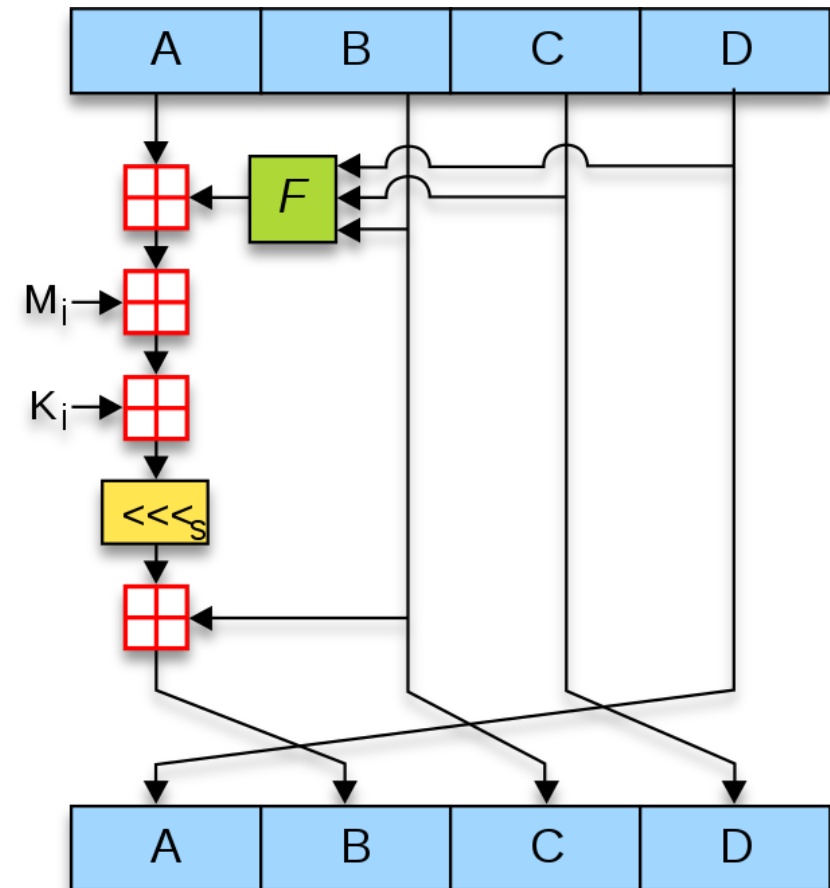
$$G(B, C, D) = (B \wedge D) \vee (C \wedge \neg D)$$

$$H(B, C, D) = B \oplus C \oplus D$$

$$I(B, C, D) = C \oplus (B \vee \neg D)$$

$\oplus, \wedge, \vee, \neg$ denote the **XOR**, **AND**, **OR** and **NOT** operations respectively.

Figure 1. One MD5 operation. MD5 consists of 64 of these operations, grouped in four rounds of 16 operations. F is a nonlinear function; one function is used in each round. M_i denotes a 32-bit block of the message input, and K_i denotes a 32-bit constant, different for each operation. \lll_s denotes a left bit rotation by s places; s varies for each operation. \boxplus denotes addition modulo 2^{32} .



SHA-1 hash function

One iteration within the SHA-1 compression function:

A, B, C, D and E are 32-bit **words** of the state;

F is a nonlinear function that varies;

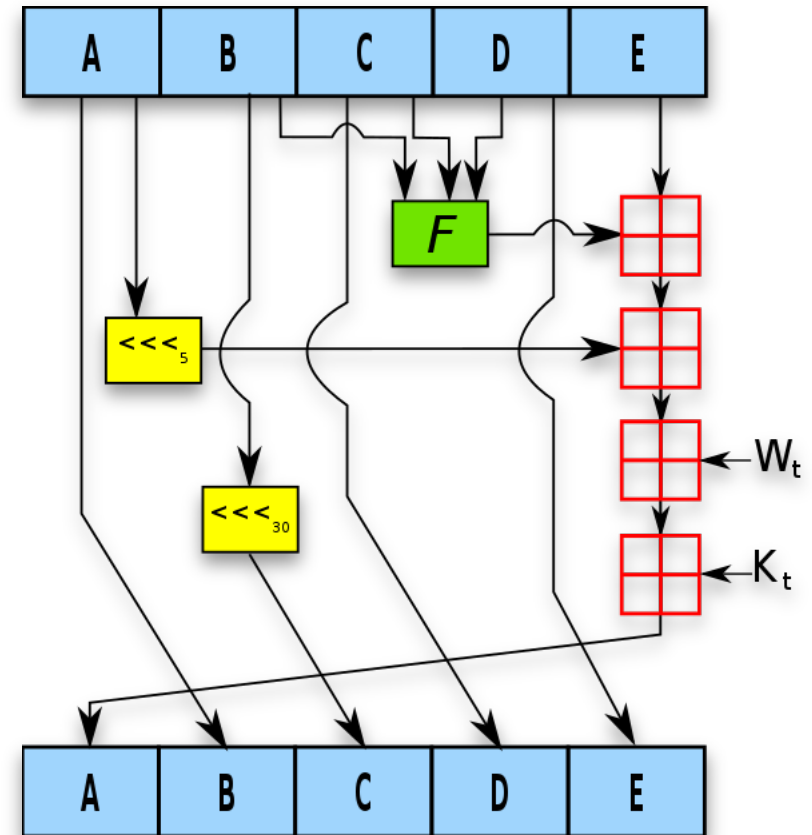
\lll_n denotes a left bit rotation by n places;

n varies for each operation;

W_t is the expanded message word of round t ;

K_t is the round constant of round t ;

\boxplus denotes addition modulo 2^{32} .



Comparison

Algorithm and variant		Output size (bits)	Internal state size (bits)	Block size (bits)	Rounds	Operations	Security (in bits) against collision attacks	Capacity against length extension attacks	Performance on Skylake (median cpb) ^[52]		First Published
									long messages	8 bytes	
MD5 (as reference)		128	128 (4 × 32)	512	64	And, Xor, Rot, Add (mod 2 ³²), Or	≤18 (collisions found) ^[53]	0	4.99	55.00	1992
SHA-0		160	160 (5 × 32)	512	80	And, Xor, Rot, Add (mod 2 ³²), Or	<34 (collisions found)	0	≈ SHA-1	≈ SHA-1	1993
SHA-1							<63 (collisions found) ^[54]		3.47	52.00	1995
SHA-2	SHA-224	224	256 (8 × 32)	512	64	And, Xor, Rot, Add (mod 2 ³²), Or, Shr	112	32	7.62	84.50	2004
	SHA-256	256					128	0	7.63	85.25	2001
	SHA-384	384	512 (8 × 64)	1024	80	And, Xor, Rot, Add (mod 2 ⁶⁴), Or, Shr	192	128 (≤ 384)	5.12	135.75	2001
	SHA-512	512					256	0	5.06	135.50	
	SHA-512/224	224					112	288	≈ SHA-384	≈ SHA-384	2012
SHA-512/256	256	128	256								
SHA-3	SHA3-224	224	1600 (5 × 5 × 64)	1152	24 ^[55]	And, Xor, Rot, Not	112	448	8.12	154.25	2015
	SHA3-256	256		1088			128	512	8.59	155.50	
	SHA3-384	384		832			192	768	11.06	164.00	
	SHA3-512	512		576			256	1024	15.88	164.00	
	SHAKE128	d (arbitrary)	1344	min(d/2, 128)		256	7.08	155.25			
	SHAKE256	d (arbitrary)	1088	min(d/2, 256)		512	8.59	155.50			

Hash online

- Test some hash functions in:

<http://www.convertstring.com/Hash>

Use of hash functions in Blockchain

- <https://colab.research.google.com/drive/171hbYK3ERQzXW1yNLQKMSLoNwL1fzgdb#scrollTo=wu-qTT0QGz2>

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