

Lecture 2 DAGs

Def: $G = ([m], E)$ a DAG

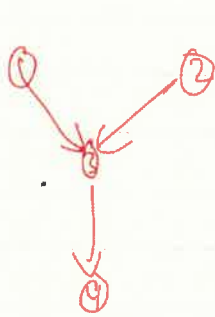
A path $\langle i_0, i_1, \dots, i_\ell \rangle$ in G is d-connecting given $C \subseteq [m]$, if

① $\forall i_h$ s.t. $i_{h-1} \rightarrow i_h \leftarrow i_{h+1}$ (collider node)

we have $i_h \in C$ or $i_h \in An(C)$
i.e. $i_h \rightarrow \dots \rightarrow x \in C$

② if $\langle i_{h-1}, i_h, i_{h+1} \rangle$ is not a collider then $i_h \notin C$.

$A, B \subseteq [m]$ are d-separated given C if there is no d-connecting path from $a \in A$ to $b \in B$.



① & ② d-connected given ③
① & 2 d-connected given 4

1 & 2 d-separated given \emptyset
2 & 4 d-connected given \emptyset
2 & 4 d-separated given 3

Def: $G = (\mathcal{I}, E)$ a DAG

A distribution P over X_1, \dots, X_m is said to satisfy

(DL) Directed Local Markov Property ^{w.r.t G} if $\forall i \in \mathcal{I}$

$$X_i \perp\!\!\!\perp \mathcal{I}_{nd_G(i) \setminus pa_G(i)} \mid pa_G(i)$$

(DG) Directed Global Markov Property if $\mathcal{I}_A \perp\!\!\!\perp \mathcal{I}_B \mid \mathcal{I}_C$

holds in P whenever A, B d-separated given C .

A distribution P over X_1, \dots, X_m with density f factorizes according to G if

$$(DF) \quad f(x) = \prod_{i \in \mathcal{I}} f(x_i \mid x_{pa_G(i)})$$

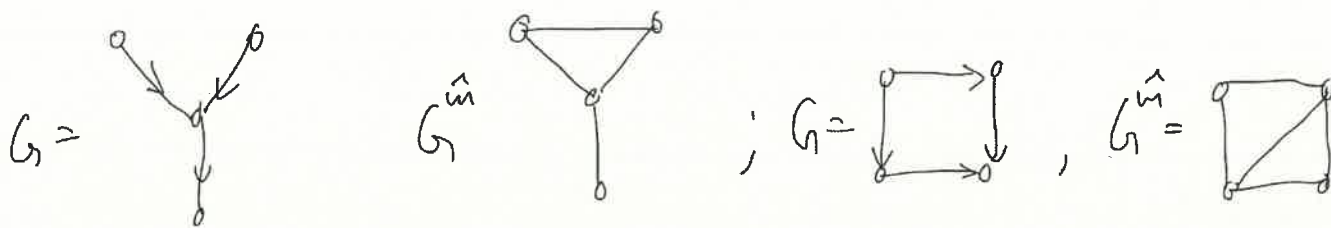
We'll prove that these are equivalent.

→ That is, the graphs are good models of the ~~the~~ distributions

Def: $G = (\mathcal{I}, E)$ a DAG. The moral graph of G , denoted G^m is the undirected graph $G^m = (\mathcal{I}, E')$

$$E' = \{ij \mid i \rightarrow j \in E \text{ or } ch_G(i) \cap ch_G(j) \neq \emptyset\}$$

We will use the results from the first lecture about undirected graphs.



Prop 1: If P factorize according to a DAG G (DF)
Then it obeys (G) wrt G^m

Proof: P has density f with

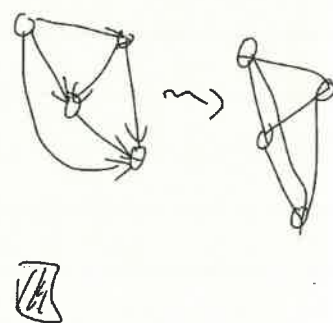
$$P(x) = \prod_{i \in [n]} f(x_i | x_{pa_G(i)})$$

the set of parents of any $i \in [n]$ will form a clique in G^m we may take functions.

$$\psi_{pa_G(i) \cup \{i\}}(x) = f(x_i | x_{pa_G(i)})$$

G DAG so all maximal cliques are of the form $\{i\} \cup pa_G(i)$ for some i .

Hence P factorizes wrt G^m
Then we use $(P) \Rightarrow (G)$



$$(DP) \text{ for } G \Rightarrow (P) \text{ for } G^m \Rightarrow (G) \text{ for } G^m$$

Prop 2: If P factorizes according to a DAG $G = ([n], E)$ and $A \subseteq [n]$ which is ancestrally closed ($An(A) = A$) then the marginal distribution P_A factorizes over $G|_A$.

Proof: Exercise just a restriction

(DM) Marginalization Markov Property

$X_A \perp\!\!\!\perp X_B \mid X_C$ holds if A, B separated by C in $(G|_{An(A \cup B \cup C)})$

(DE) \Rightarrow (DM)

Prop 3: P factorizes according to $G = ([n], E)$
Then (DM) holds.

Proof: $P_{An(A \cup B \cup C)}$ factorizes according to $G_{An(A \cup B \cup C)}$ ~~just a restriction~~

- by Prop 2.
- Prop 1 $\Rightarrow P_{An(A \cup B \cup C)}$ obeys (G) w.r.t $(G_{An(A \cup B \cup C)})$
- Hence $A \perp\!\!\!\perp B \mid C$ in $P_{An(A \cup B \cup C)}$ hence in P . \square

from Lecture 1



Prop 4 (DM) is equiv. to (DG)

Proof: Must prove

Exercise

(1) A, B d-separated given C in G



(2) A, B separated in $(G / An(A \cup B \cup C))^m$ given C .

Exercise: \uparrow suppose (1) false $\Rightarrow \exists$ d-connecting path $a=i_0, \dots, i_m=b$ given C
 $\Rightarrow \{i_1, \dots, i_{m-1}\} \subseteq An(A \cup B \cup C) \Rightarrow \exists$ path circumventing C .

Theorem $G = (\mathcal{U}, E)$ a DAG, \mathcal{P} disk. X_1, \dots, X_m
 $(DF) \Leftrightarrow (DG) \Leftrightarrow (DL)$

Proof:

$(DF) \Rightarrow (DG)$ Prop 3 & Prop 4

$(DL) \Rightarrow (DF)$ By induction over m .

Base case: $m=1, 2$

I.H. Assume holds for G with less than m vertices.

$G = (\mathcal{U}, E)$ DAG, let i be a sink in G .

$$\begin{aligned} \frac{f(x)}{f(x_{\mathcal{U} \setminus i})} &= f(x_i | x_{\mathcal{U} \setminus i}) = \\ &= f(x_i | x_{pa(i)}, x_{\mathcal{U} \setminus pa(i) \cup i}) \\ (D2) &= f(x_i | x_{pa(i)}) \end{aligned}$$

$x_{\mathcal{U} \setminus i} \setminus pa(i)$

$$(DG) \Rightarrow (DL)$$

See Laevitz

Let $i \in [m]$ need to show $\sum_i \nmid \sum_{nd(i) \setminus pa(i)} \mid \sum_{pa(i)}$

$(DG) \Rightarrow$ suffices to show i & $nd(i) \setminus pa(i)$ disjunct given $pa(i)$.

For contradiction.

Suppose \exists d-connecting path $\pi = \langle i_0 = i, \dots, i_m = j \rangle$ given $pa(i)$, $j \in nd(i)$

Note $i_1 \leftarrow i_0$ not possible since then $i \in pa(i)$

If π contains no collider $i \rightarrow \dots \rightarrow j$ contradicting $j \in nd(i)$

If π contains collider $i_{u-1} \rightarrow i_u \leftarrow i_{u+1} \Rightarrow i_u \in An(pa(i))$



↑ cycle!

