

# COS80013

# Internet Security

Week 5

March 31 2025

Yasas Supeksala



# Acknowledgement of Country

We respectfully acknowledge the Wurundjeri People of the Kulin Nation, who are the Traditional Owners of the land on which Swinburne's Australian campuses are located in Melbourne's east and outer-east, and pay our respect to their Elders past, present and emerging.

We are honoured to recognise our connection to Wurundjeri Country, history, culture, and spirituality through these locations, and strive to ensure that we operate in a manner that respects and honours the Elders and Ancestors of these lands.

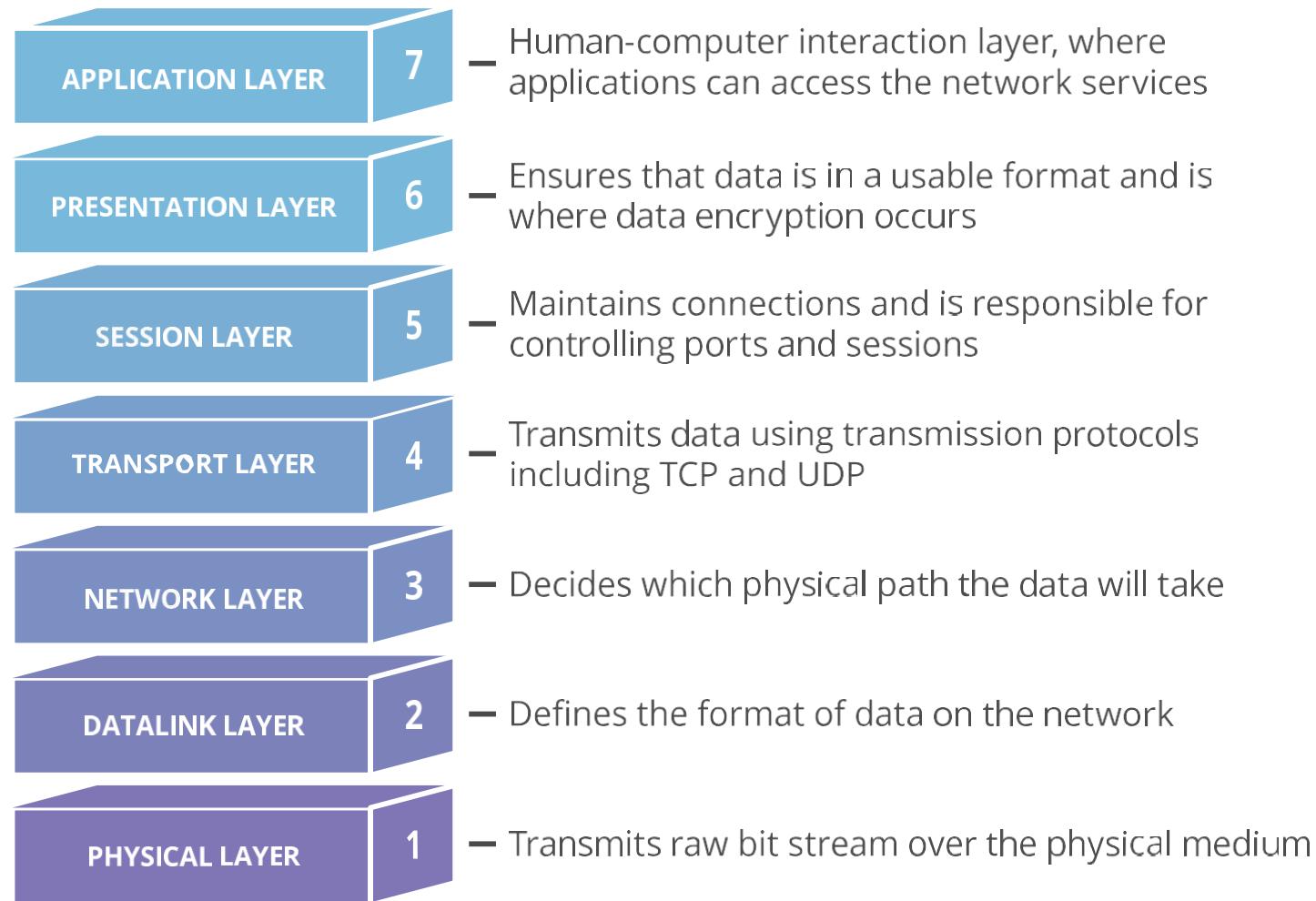
We also respectfully acknowledge Swinburne's Aboriginal and Torres Strait Islander staff, students, alumni, partners and visitors.

We also acknowledge and respect the Traditional Owners of lands across Australia, their Elders, Ancestors, cultures, and heritage, and recognise the continuing sovereignties of all Aboriginal and Torres Strait Islander Nations.



# Network Basics

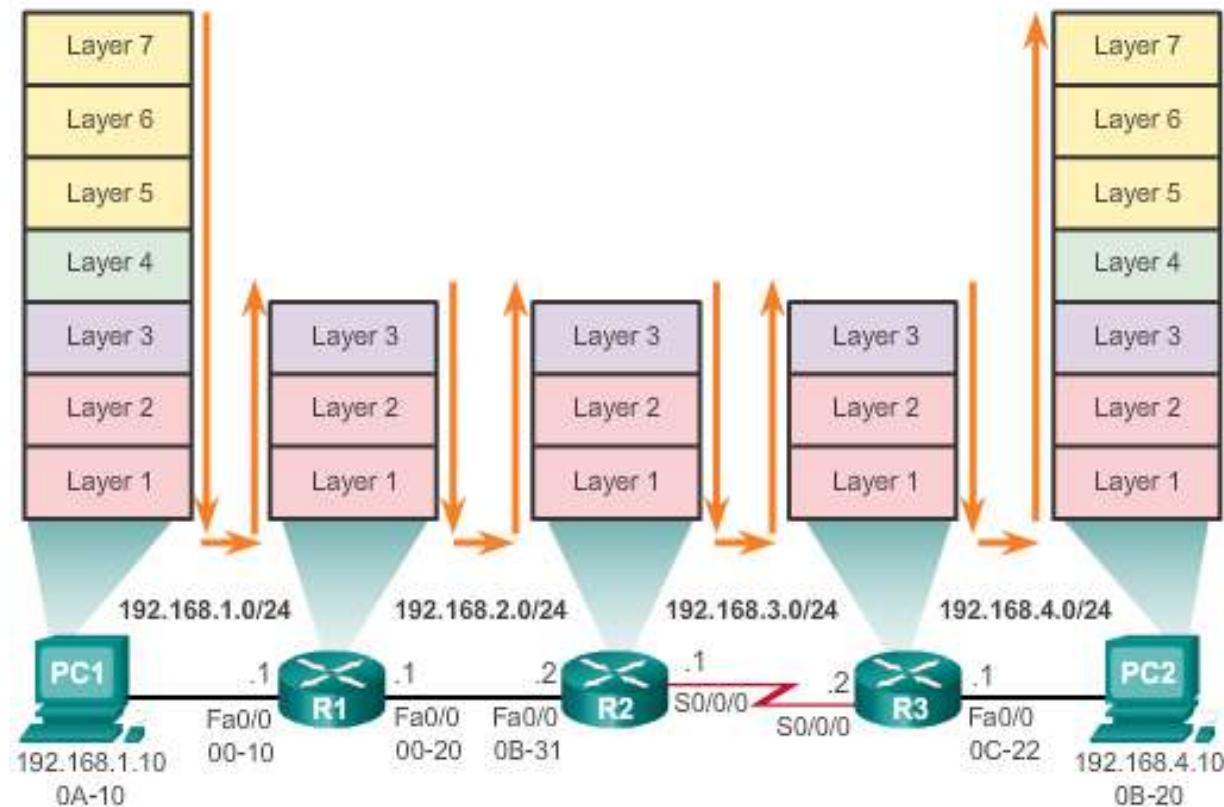
# OSI Model



(img source: <https://www.cloudflare.com/learning/ddos/glossary/open-systems-interconnection-model-osi/>)

# Review

## Encapsulating and De-Encapsulating Packets



# Network Tools

Ping

Traceroute/Tracert

ipconfig/ifconfig

Netstat

Packet Sniffing

# Network Tools

## Ping

- A network diagnostic utility for testing network connections.
- Sends an ICMP "echo request" packet to an IP address and waits for an ICMP "echo response" packet and reports the time delay in milliseconds.
- If a domain name is used, ping will initiate and use the results of a DNS query.
- Ping (and other ICMP requests) may be blocked at firewalls to prevent network reconnaissance.

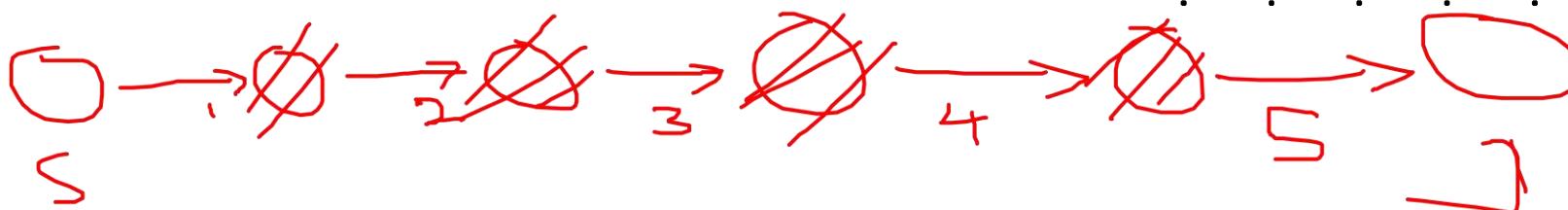
```
File Edit View Search Terminal Help
$ ping -i 2 -c 10 www.geeksforgeeks.org
PING d13vvqr7dxay1j.cloudfront.net (52.222.128.37) 56(84) bytes of data.
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=1 ttl=244 time=320 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=2 ttl=244 time=100 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=3 ttl=244 time=2133 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=4 ttl=244 time=844 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=5 ttl=244 time=926 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=6 ttl=244 time=1704 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=7 ttl=244 time=1496 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=8 ttl=244 time=1496 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=9 ttl=244 time=1496 ms
64 bytes from server-52-222-128-37.bom51.r.cloudfront.net (52.222.128.37): icmp_seq=10 ttl=244 time=1496 ms
--- d13vvqr7dxay1j.cloudfront.net ping statistics ---
10 packets transmitted, 7 received, 30% packet loss, time 20434ms
rtt min/avg/max/mdev = 100.846/1075.329/2133.966/685.503 ms, pipe 2
$ 
```

# Traceroute/Tracert

- A network tool for tracking the path taken by IP packets.
- Sends a sequence of UDP or TCP packets with incrementally increasing TTL (time to live) values.
- Collects the resulting ICMP "time exceeded" packets.
- Displays the source IP addresses and host names of the ICMP packets in sequence to show the path taken by the original packets.

```
C:\Windows\system32\cmd.exe
C:\>tracert amazon.com
Tracing route to amazon.com [54.239.25.200]
over a maximum of 30 hops:
  1  6 ms   4 ms   2 ms  10.192.4.1
  2  <1 ms   <1 ms   <1 ms  67-198-47-97.static.grandennetworks.net [67.198.4.2.97]
  3  1 ms   1 ms   1 ms  ae1-708.austtxsuk001.aggr02.austtx.grandecom.net
  4  2 ms   2 ms   2 ms  24-155-121-210.static.grandennetworks.net [24.155.1.210]
  5  22 ms  22 ms  22 ms  24-155-121-2.static.grandennetworks.net [24.155.1.2]
  6  15 ms  15 ms  15 ms  ae0-0.core01.gf01.dlistx.grandecom.net [24.155.1.821]
  7  15 ms  15 ms  15 ms  52.95.217.96
  8  73 ms  84 ms  61 ms  176.32.125.188
  9  56 ms  55 ms  55 ms  176.32.125.241
  10  56 ms  57 ms  55 ms  176.32.125.248
  11  48 ms  48 ms  48 ms  54.248.229.172
  12  55 ms  55 ms  57 ms  54.248.228.175
  13  63 ms  58 ms  56 ms  54.239.109.132
  14  58 ms  58 ms  49 ms  54.239.111.117
  15  55 ms  55 ms  55 ms  205.251.244.238
  16  *      *      *      Request timed out.
  17  *      *      *      Request timed out.
  18  *      *      *      Request timed out.
  19  *      *      *      Request timed out.
  20  *      *      *      Request timed out.
  21  *      *      *      Request timed out.
  22  *      *      *      Request timed out.
  23  *      *      *      Request timed out.
  24  *      *      *      Request timed out.
  25  *      *      *      Request timed out.
  26  *      *      *      Request timed out.
  27  *      *      *      Request timed out.
  28  *      *      *      Request timed out.
  29  *      *      *      Request timed out.
  30  *      *      *      Request timed out.

Trace complete.
```



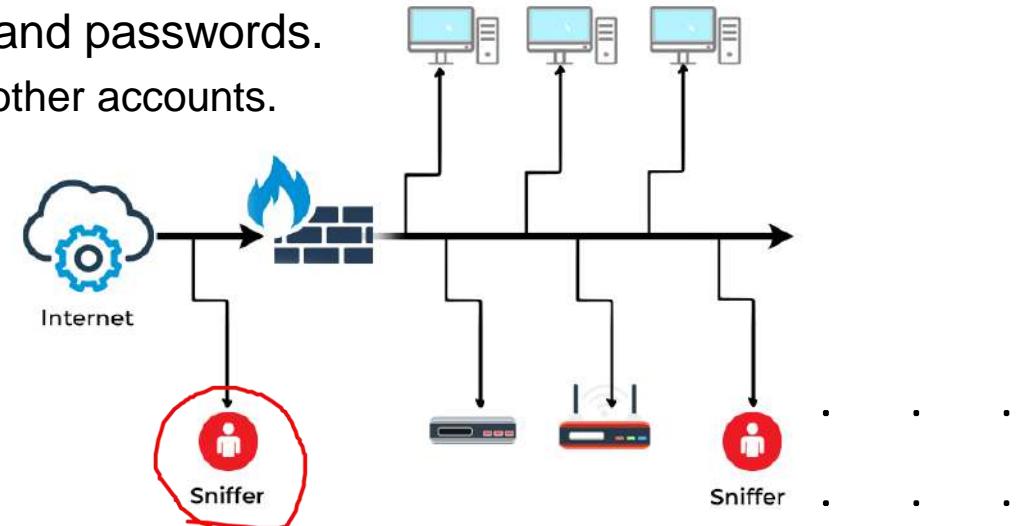
# Netstat

- Displays all current network connections including TCP and UDP and other protocols.
- Can indicate the presence of trojans and spyware "phoning home".

```
admin@tecmint ~ $ sudo netstat -ltup
Active Internet connections (only servers)
Proto Recv-Q Send-Q Local Address           Foreign Address         State      PID/Program name
tcp        0      0 *:http                  *:*                   LISTEN     1423/nginx -g daemo
tcp        0      0 tecmint:domain          *:*                   LISTEN     2992/dnsmasq
tcp        0      0 *:ssh                   *:*                   LISTEN     1409/sshd
tcp        0      0 localhost:ipp           *:*                   LISTEN     2738/cupsd
tcp        0      0 *:https                *:*                   LISTEN     1423/nginx -g daemo
tcp6       0      0 [::]:http              [::]:*                LISTEN     1423/nginx -g daemo
tcp6       0      0 [::]:ssh              [::]:*                LISTEN     1409/sshd
tcp6       0      0 ip6-localhost:ipp      [::]:*                LISTEN     2738/cupsd
tcp6       0      0 [::]:https             [::]:*                LISTEN     1423/nginx -g daemo
udp        0      0 *:ipp                 *:*                   2740/cups-browsed
udp        0      0 *:mdns                *:*                   1022/avahi-daemon:
udp        0      0 *:36390               *:*                   2992/dnsmasq
udp        0      0 *:59072               *:*                   1022/avahi-daemon:
udp        0      0 tecmint:domain          *:*                   2992/dnsmasq
udp        0      0 *:bootpc              *:*                   2982/dhclient
udp        0      0 tecmint:ntp              *:*                   1465/ntpd
udp        0      0 localhost:ntp            *:*                   1465/ntpd
udp        0      0 *:ntp                 *:*                   1465/ntpd
udp6       0      0 [::]:43740              [::]:*                1022/avahi-daemon:
udp6       0      0 [::]:mdns              [::]:*                1022/avahi-daemon:
udp6       0      0 fe80:::dd8c:3d40:817:ntp [::]:*                1465/ntpd
udp6       0      0 ip6-localhost:ntp      [::]:*                1465/ntpd
udp6       0      0 [::]:ntp               [::]:*                1465/ntpd
admin@tecmint ~ $
```

# Packet Sniffing

- Packet sniffers record IP packets on the network. They were originally designed to help diagnose problems in networks.
  - Good for picking up MAC addresses, IP addresses
- Many internet-based services expect to receive user names and passwords in plain text.
  - Telnet, FTP, SNMP
- Computer users get lazy and re-use the same user names and passwords.
  - If you can get their FTP password, you can probably use it on other accounts.
- Popular Sniffers:
  - TCPDump, Snort, Wireshark. Windows and Linux versions.
  - reviews: <http://sectools.org/sniffers.html>



# Packet Sniffers

## Wireshark

- Text (tethereal) and GUI (wireshark) versions available for Windows and Linux
- Lists packet contents on the screen and logs them to a file for analysis later.
- Summarises types of packets intercepted.



## Snort

- Text-based IDS with packet sniffing and logging abilities.
- Separates packets into IP address and port number
- Easy to search packets using *grep* (linux)

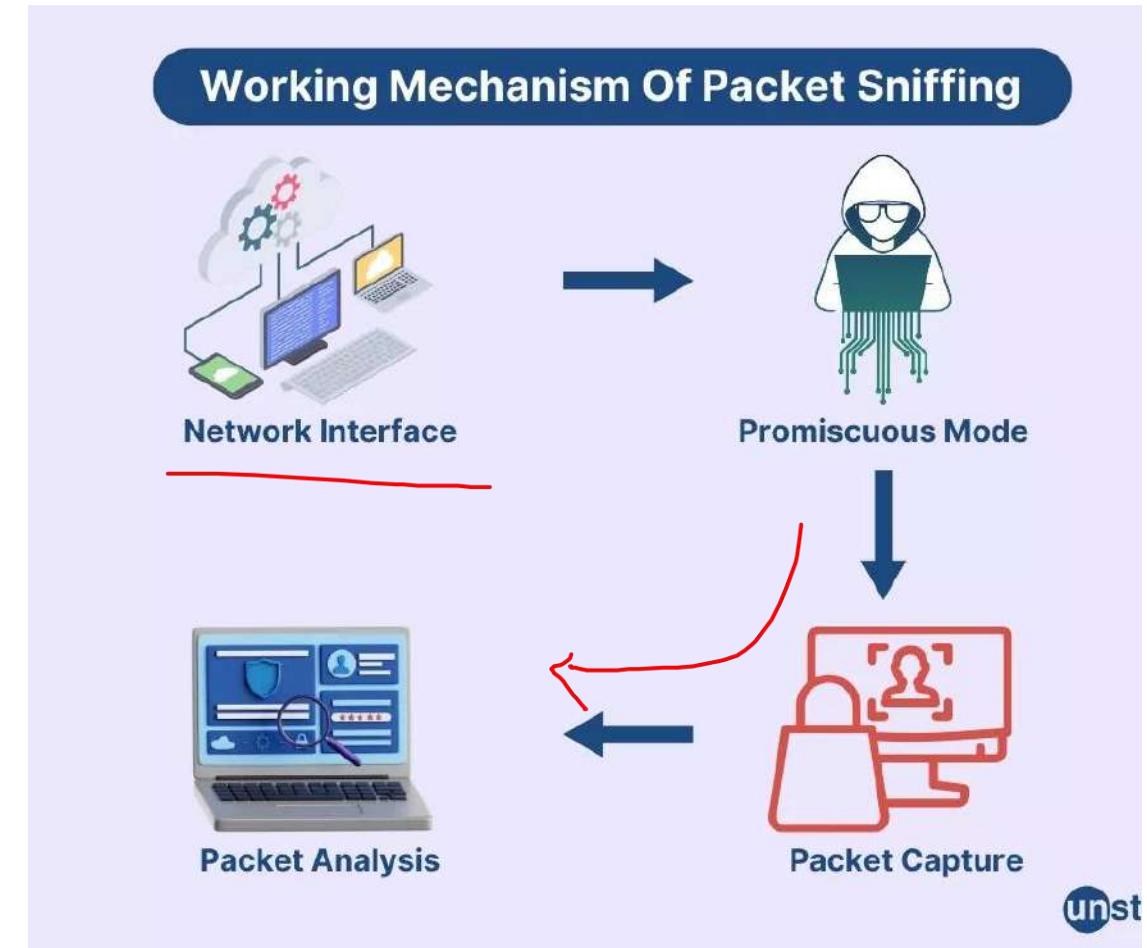


# Network Attacks

# Common Attacks

Sniffing a NIC/hub/wifi broadcast.

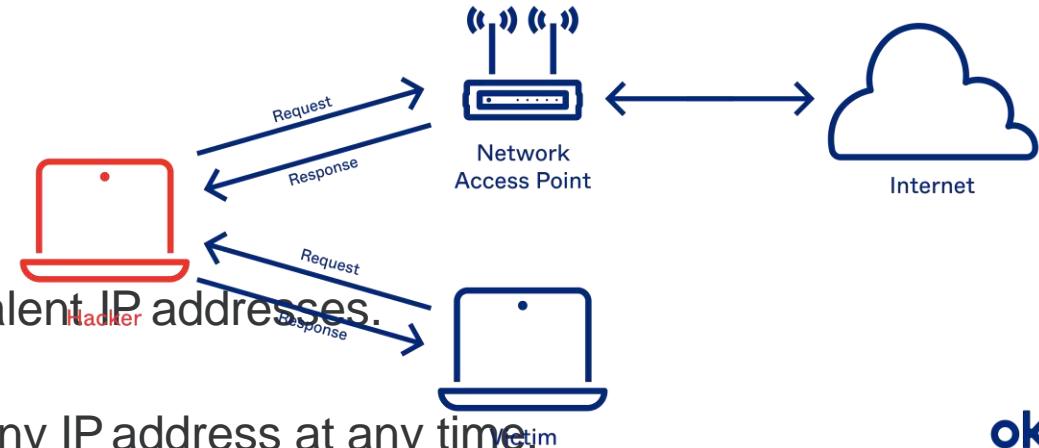
- Sniffing means capturing IP packets and displaying/analysing them.
- Common sniffing programs include Snort and Wireshark.
- Local NIC is set to Promiscuous Mode.
- Sniffing program records packets presented to NIC, including all packets passing through the hub.
- Switches are also susceptible to ARP cache poisoning and MAC flooding.  
Switches act like hubs when overloaded.



## Common Attacks

ARP cache poisoning.

- ARP maintains a table of MAC addresses and their equivalent IP addresses.
- ARP sends out queries: "Who has IP 239.254.2.15"
- Replies are added to the ARP table, and can come from any IP address at any time.



okta

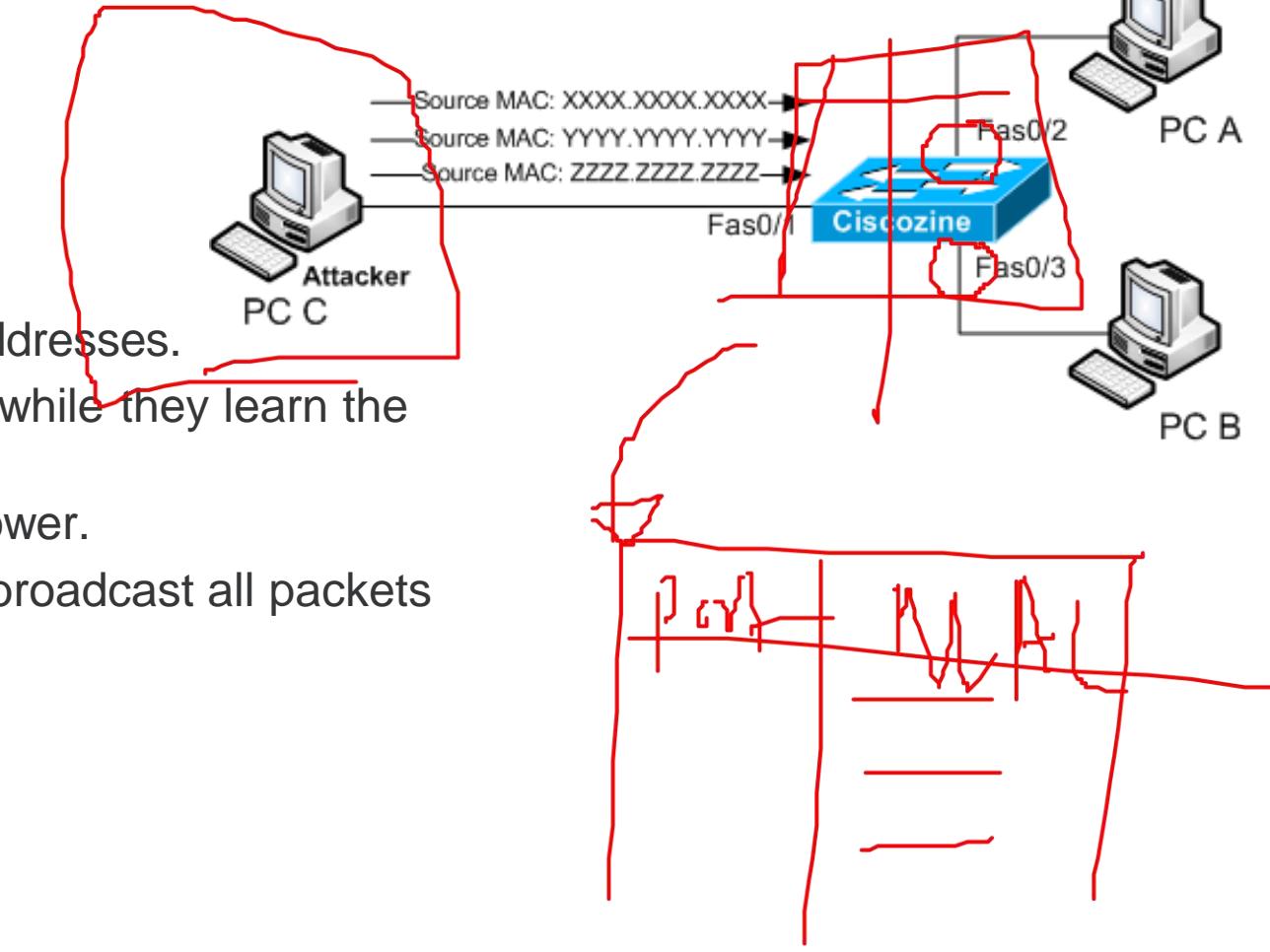
MITM

- Attacker attaches his PC to the network.
  - sniffs to find target IP addresses.
  - Sends a home-made ARP reply nominating his MAC as corresponding to the requested IP address.
  - ARP adds the attackers MAC address to the ARP table.
  - Future traffic to the target goes to the attacker.
  - Attacker reads packets, doctors them and sends them on to the target who is none the wiser.
- .....

# Common Attacks

MAC flooding.

- Switches use ARP to map MAC addresses to IP addresses.
- When booted up, switches operate in hub mode while they learn the MAC addresses attached to each port.
- Switches have limited memory and computing power.
- When overloaded, they revert to Hub-mode and broadcast all packets to all interfaces.



How it works

- Attacker attaches his PC to the network.
- Sends out multiple invalid ARP responses.
- Switch overloads and reverts to hub-mode (all lights start flashing).
- Attacker collects all packets from the network.

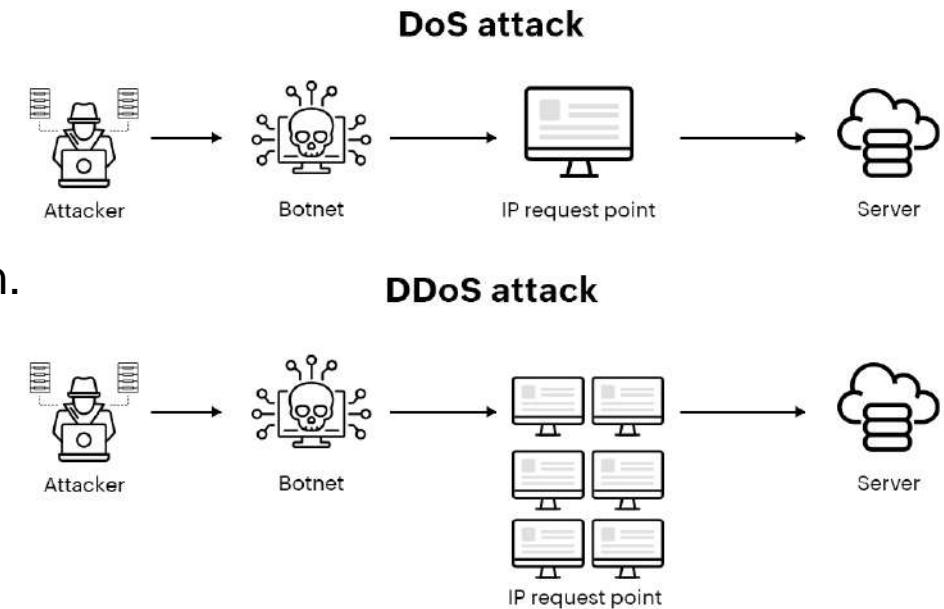
# DoS (Denial of Service)

DOS attacks are aimed at servers on the internet

- web servers (HTTP), name servers (DNS), FTP servers
- Can be launched at specific machines if the IP address is known.

The goal is to Deny Service to legitimate customers/users.

- Customers go elsewhere
- Organisations lose money/trade/reputation



# DoS (Denial of Service)

## Ping of death

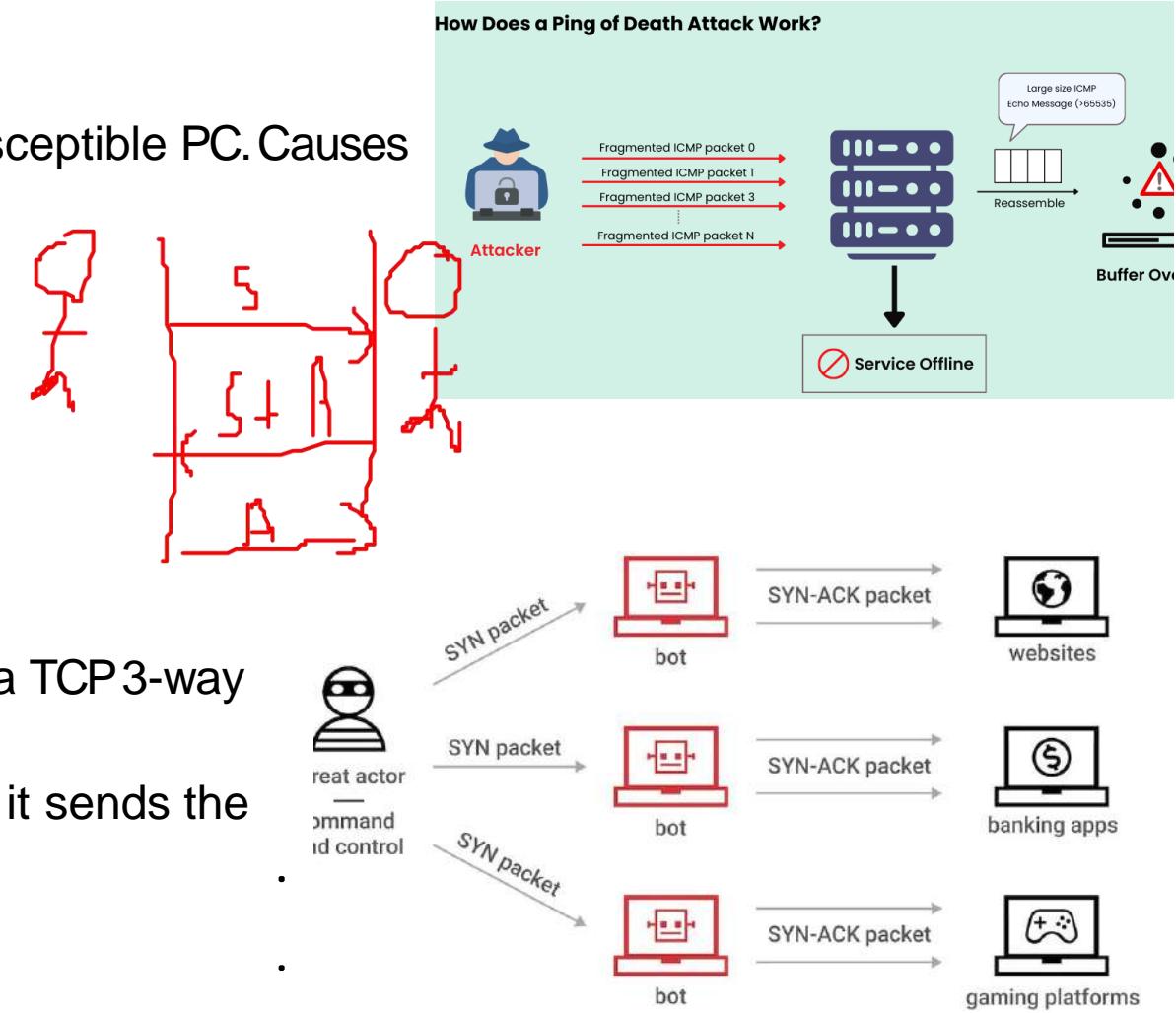
- an over-sized ping packet (>65535 bytes) is sent to a susceptible PC. Causes a re-boot on susceptible machines.

## ICMP

- attacker sends a stream of ICMP echo request packets.

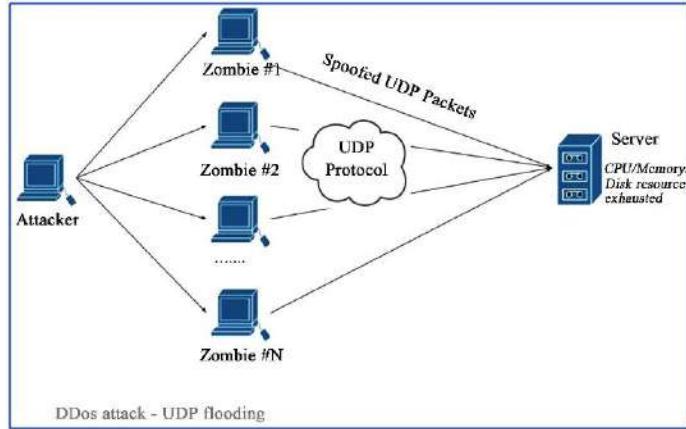
## SYN Flooding

- a resource depletion attack
- attacker sends a stream of SYN packets (the first part of a TCP 3-way handshake).
- Each SYN packet consumes memory on the server while it sends the SYN\_ACK packet and waits for the returning ACK packet.
- The attacker never sends the ACK packet.



# DoS (Denial of Service)

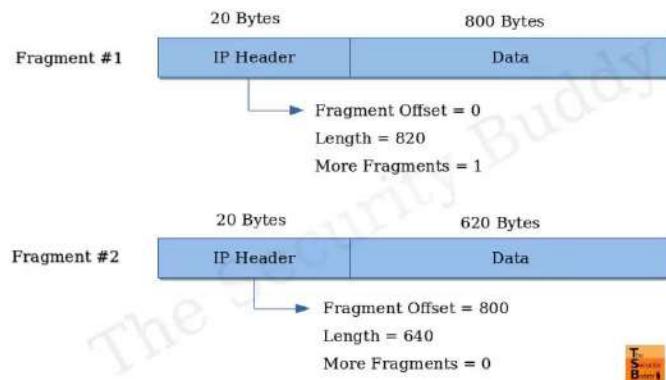
UDP  
flood



UDP packets are sent to random ports of the victim as fast as possible.

The target machine will respond to each packet with a ICMP destination unreachable packet.

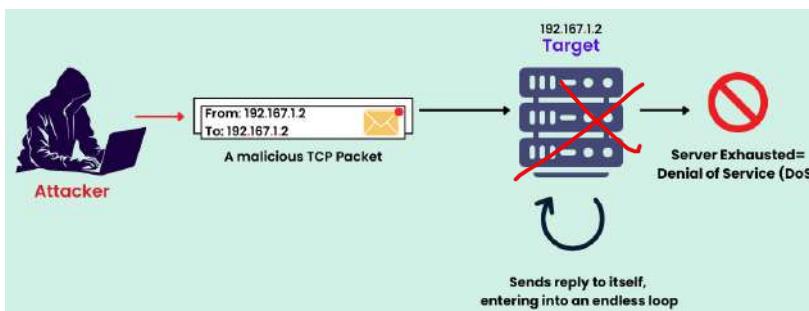
Teardrop



A fragmented packet is sent to the attacker, but the parts have been corrupted and can not be put back together.

Server consumes resources requesting retransmit.

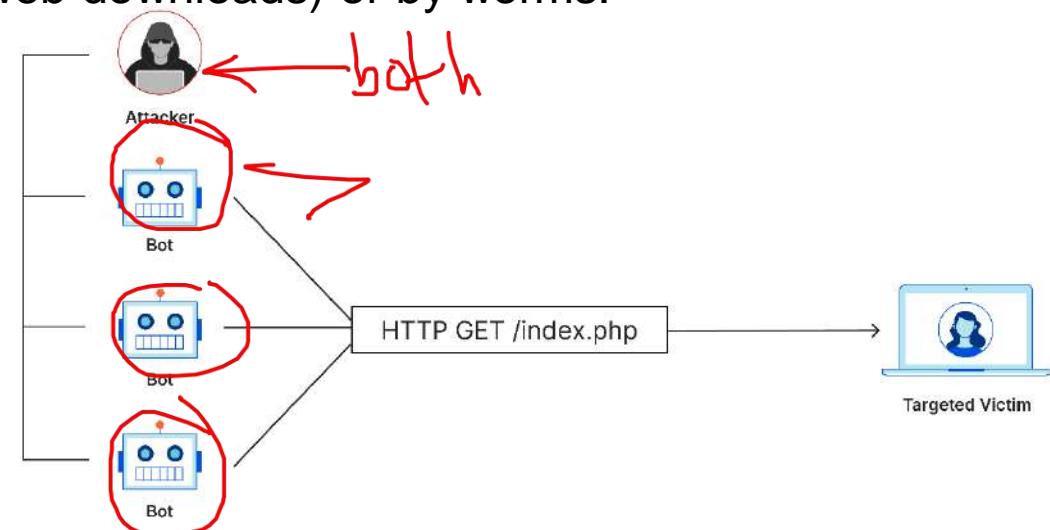
Land  
attack  
→



A packet with a spoofed source IP address is sent to the victim. The return address is the same as the destination address, so the victim's machine ties it self in knots by answering back to itself.

# DDoS

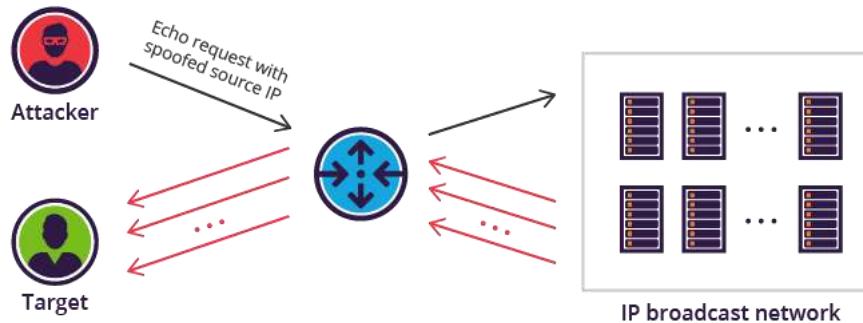
- A Distributed DOS involves simultaneous attacks on a single site by 'zombie' machines infected with a 'bot'. Each bot is controlled by a 'bot-herder' originally through IRC channels and now through http. A bot-herder's 'bot-net' may range in size from a few hundred to millions of infected PCs.
- Each bot is actually a small server program which has installed itself and is capable of launching DOS attacks, sending spam, infecting other machines, or all of the above.
- Spread by trojans (e-mail, web downloads) or by worms.



# Smurf Attack (DDoS)

Amplification attack

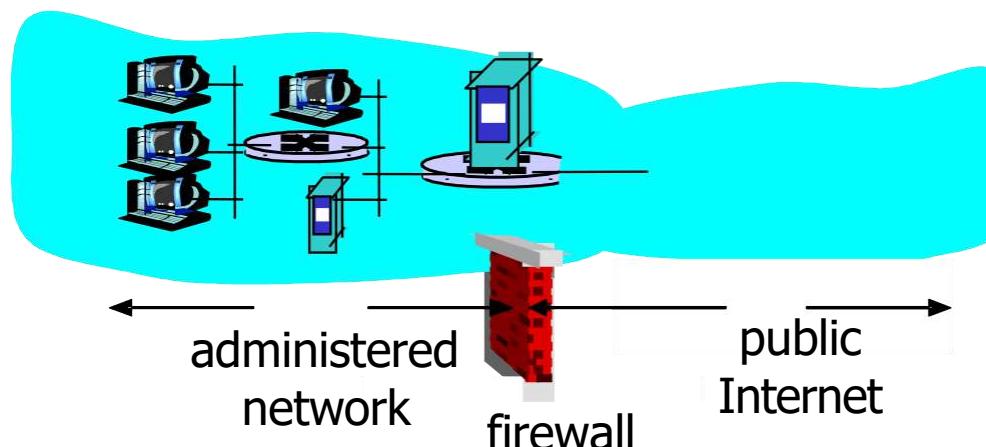
- An attacker sends a ping (ICMP echo request) to the broadcast address of the victim's network.
- This sends the ping to all IPs on the subnet. The original ping has a spoofed return IP address which is the address of the intended victim.
- Each IP reached by the original ping replies to the victim.



# Firewalls

# Firewalls: Why

- prevent denial of service attacks:
  - SYN flooding: attacker establishes many bogus TCP connections, no resources left for “real” connections
- prevent illegal modification/access of internal data.
  - e.g., attacker replaces CIA’s homepage with something else
- allow only authorized access to inside network (set of authenticated users/hosts)
- three types of firewalls:
  - stateless packet filters
  - stateful packet filters
  - application gateways

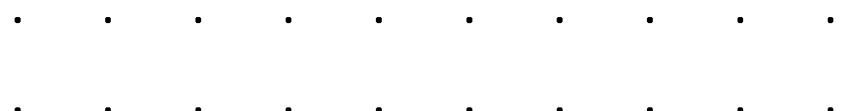


# Firewalls and port blocking

- A firewall filters incoming traffic according to a set of rules depending on things like:
  - the destination IP address or host +domain name
  - the source IP address or domain name
  - the protocol being used (bound to specific ports)
  - the port number of the destination
  - the process (program name) listening at the destination
  - the contents of a packet (high end firewalls and IDS)
- The primary defence offered by a firewall is to block particular destination ports or source IP addresses.
  - Some firewalls use NAT traversal to 'hide' the inside of the network.
  - • • • •

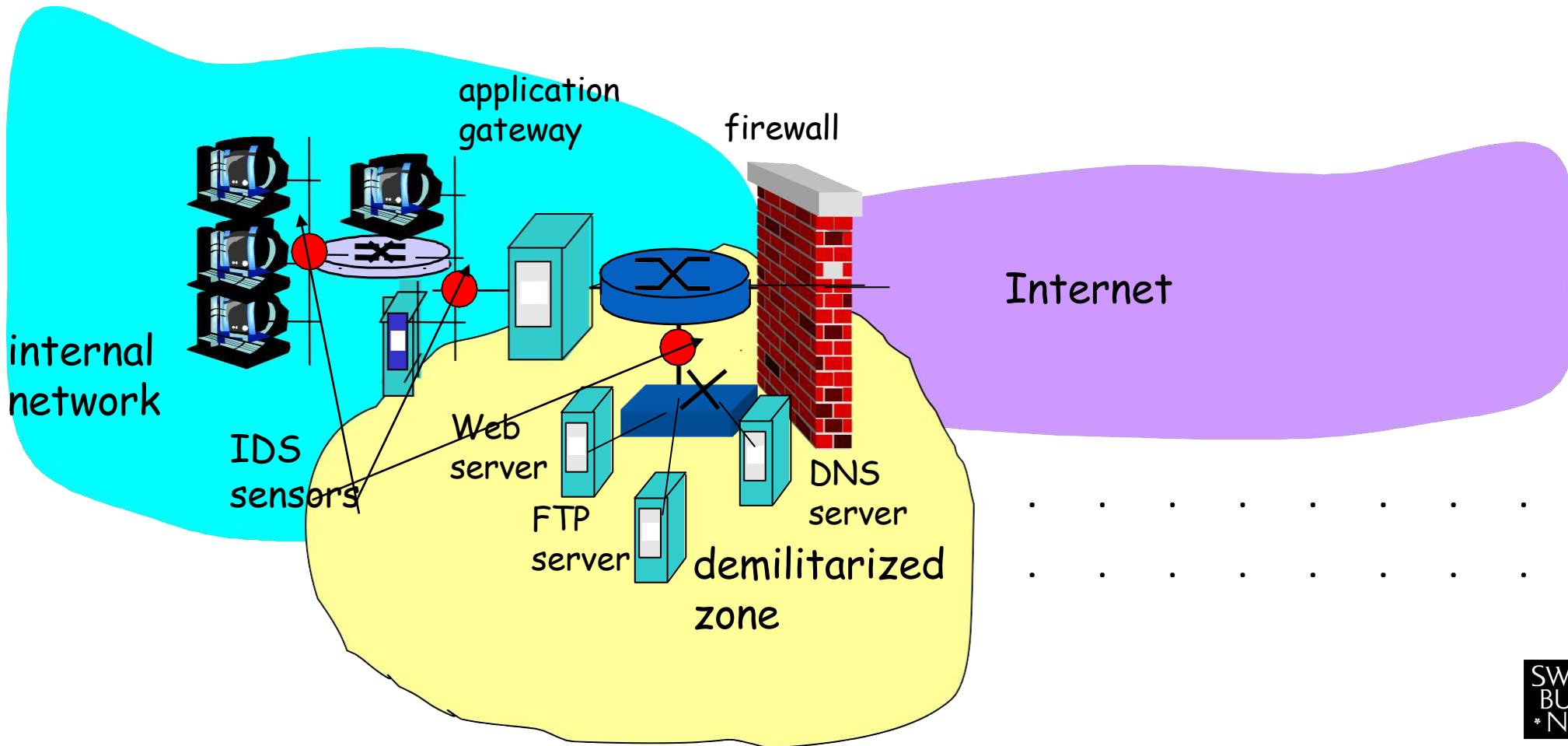
# Proxy servers and IP addresses

- Proxy servers perform three functions
  - Filter packets based on content, source IP address or domain name.
  - Cache downloads to speed up repeated downloading of web pages, media files and archived material.
  - Perform NAT traversal to facilitate the sharing of one external IP address by an internal network of hosts.



# Intrusion Detection Systems (IDS)

- a system that monitors network traffic for suspicious activity



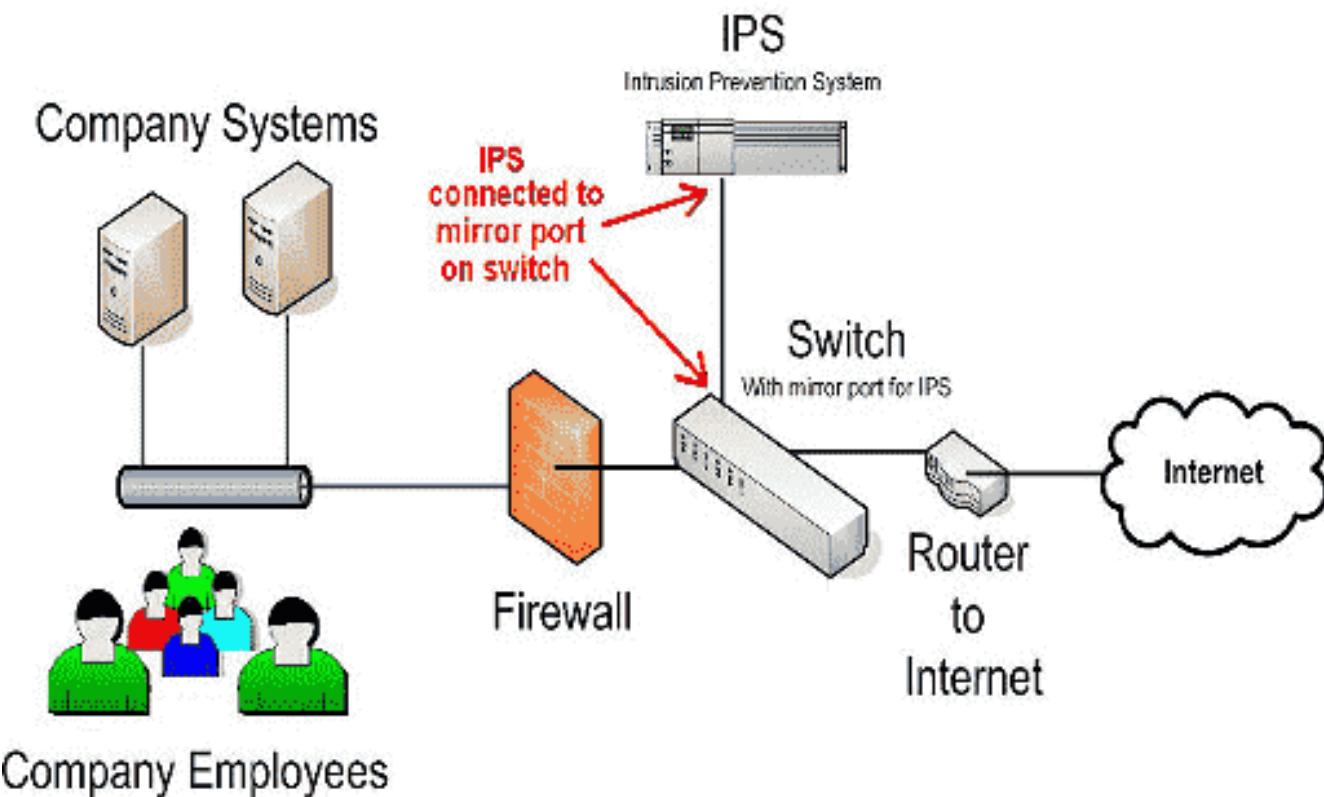
# Intrusion Prevention System (IPS)

- To actively prevent security threats from entering a network or system

There are several types of IPS

- Network-based- Protect your computer network
  - Wireless- Protect wireless networks only
  - Network behavior- Examine network traffic
  - Host-based- Come as installed software to protect a single computer
- • • • • • • • • • • •

# IPS



# IDS VS IPS

- A security system
- To actively prevent security threats from entering a network or system
- IDS - monitors network traffic → generates alerts
- IPS actively block → prevent network traffic.
- 3 Different
  - Active vs Passive
  - Detection vs. Prevention
  - Flexibility vs. Automation



# Thank you

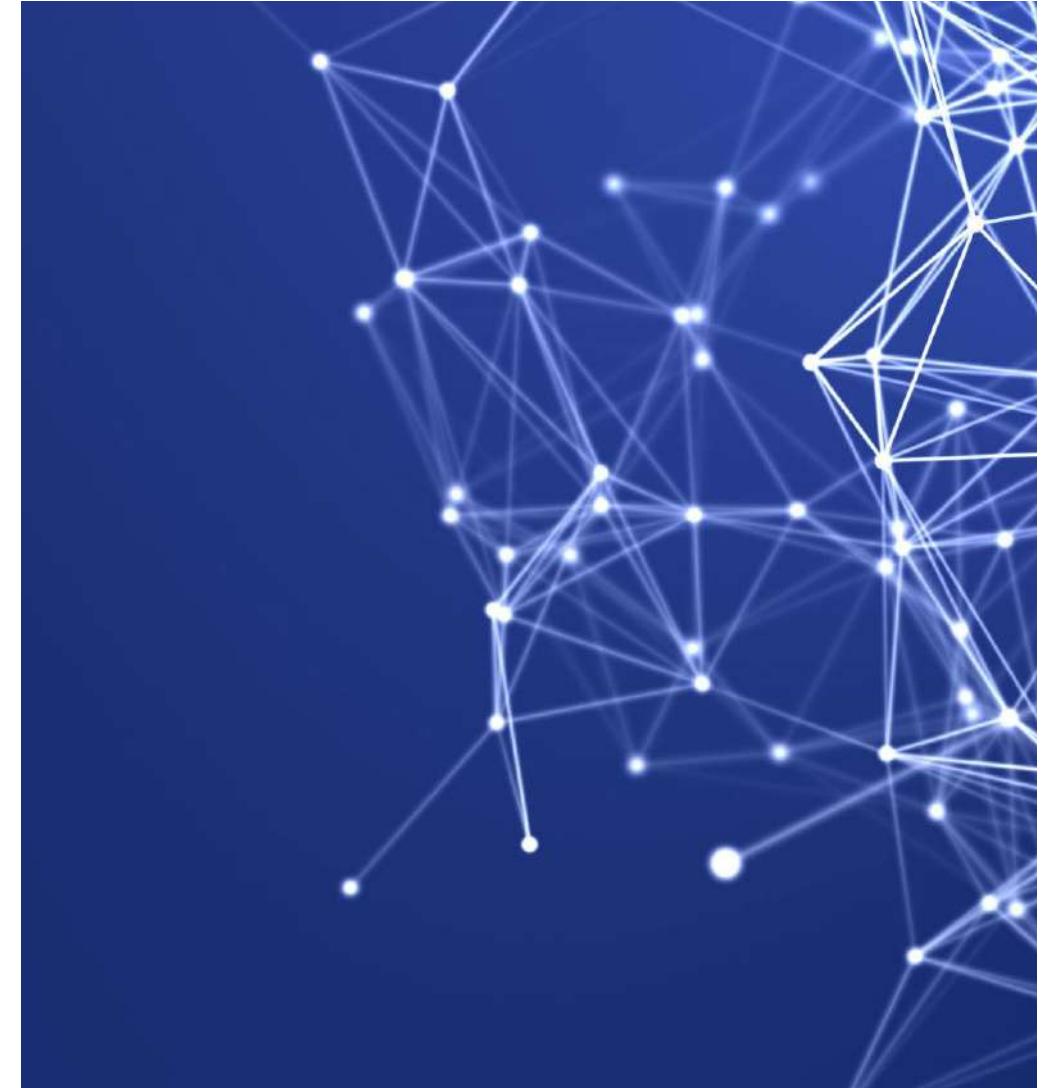
# COS80013

# Internet Security

Week 5

**Presented by Dr Rory Coulter**

31 March 2025



# Week 5 Class

# Quiz

## Week 8 in class quiz

Weeks 1 to 6 inclusive

- Weeks 1 to 6 lecture, class, supplementary and lab content
- 30 questions, concepts than specifics
  - E.g., know MITRE ATT&CK tactics rather than command line options
- 1 hour
- Turn up to your class, log in to Canvas, take quiz, lab machine only
- Closed book, MCQ
- A practice quiz will be made available

- Answers will be provided 2 weeks after
- Any funny questions, we will review, don't need to email us
- Medical certificate required if you can't attend
- Attendance will be taken, StudentID cards please
- Attend your assigned class
- We will audit when a quiz was completed and when your class is

# Lab

## Network reconnaissance and a denial of service attack

### Lab Walkthrough

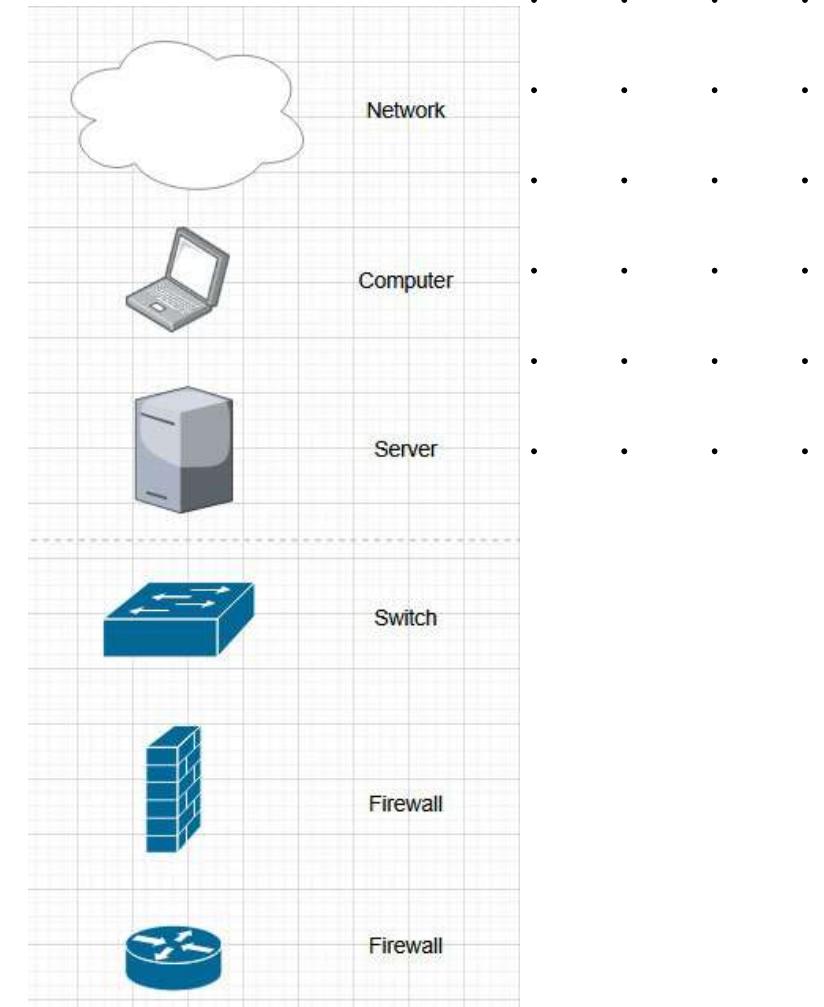
- Commands
  - Telnet
  - Ping
  - traceroute
- Purpose:
  - What hosts exist
  - What service is it running
  - What is the path does traffic traverse
- DoS:
  - Compile and run attack

# Class Activity

## Network challenges, arrangements and security considerations

Note: this is not a complete, accurate, correct nor architecturally correct – naming conventions will also incorrect

- Define a range of TTPs
  - Walk through different network types from flat to segmentation and discuss security concepts
  - Consider the network, administration and security considerations of the various tools, methodologies used to help secure networks and the assets within them
- First introduce assets and tools
  - Close with a case study



# Initial Relevant Tactics

## Tactics when considering “network” security

Four key tactics emerge when we first consider networking\*

- Initial Access [TA0001]
  - Command and Control (C2) [TA0011]
  - “*consists of techniques that adversaries may use to communicate with systems under their control within a victim network*”
- Lateral Movement [TA0008]
  - Exfiltration [TA0010]
  - “*consists of techniques that adversaries may use to steal data from your network*”

\* Though there are several more we actually should consider

# Network 1

## Characteristic: All hosts are in the same network

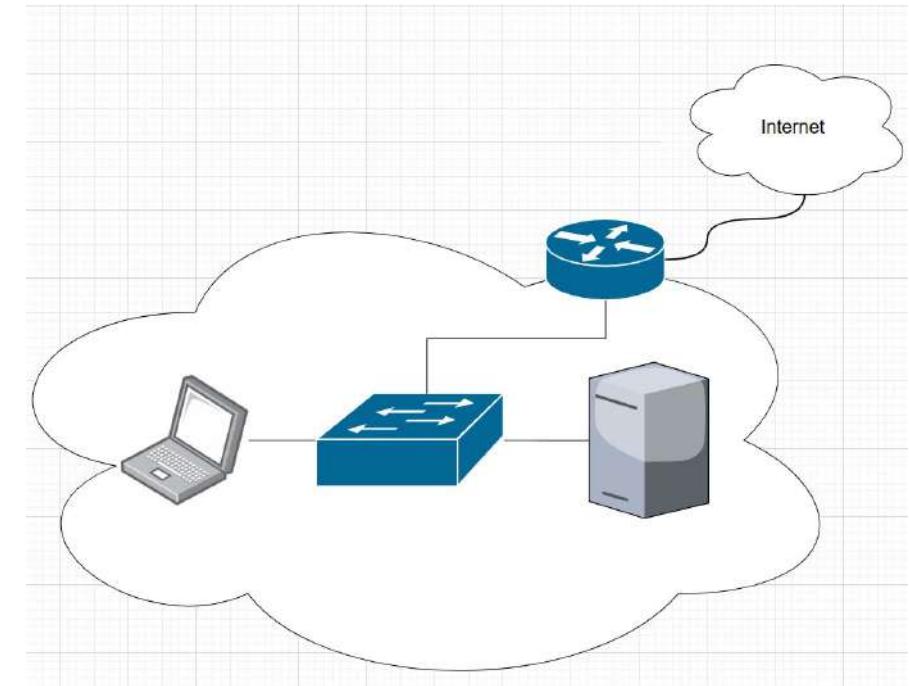
Flat network

Administration:

- Hosts: Easy, all hosts can talk to each other, management is straightforward, /24, /16, /8
- Access: Direct to host or via exposed port (e.g., 80, 443)
- Intra: Direct link to all assets
- Outbound: Connection initiated internally so let it through
- Note: Things are likely going to work smoothly

Security:

- Initial Access: Email, web view, exposed ports
- Lateral Movement: Easy (malware, actors)
- C2: Easy, allow outbound
- Exfiltration: Easy, allow outbound
- Note: Assume router is blocking popular network. But, first big concern, hosts talking directly to each other (between type)



# Network 2

## Characteristic: Separate hosts by type

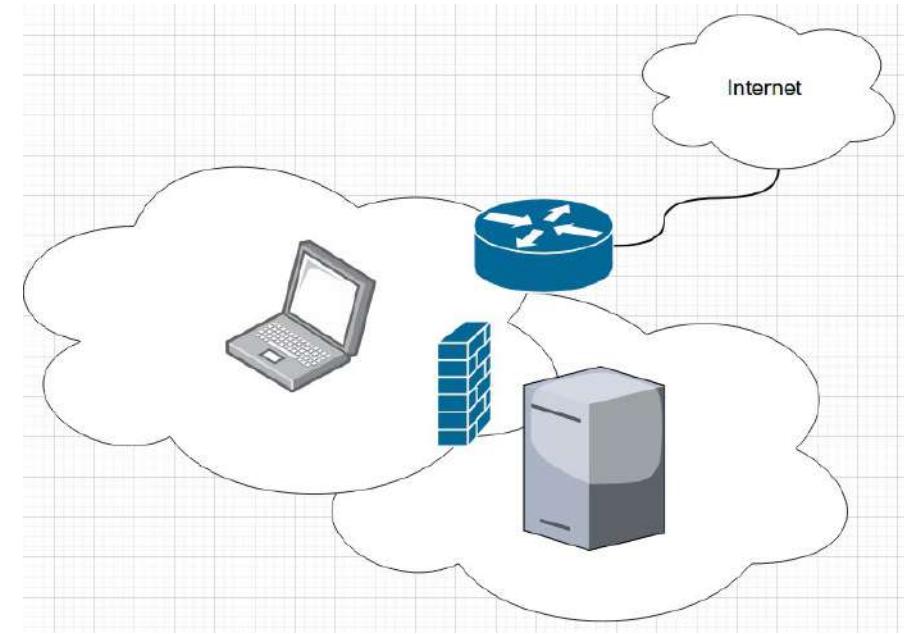
Separated network

Administration:

- Hosts: Still easy, can have a large amount, but routing or VLAN is now needed
- Access: Still direct to host or via exposed port (e.g., 80, 443)
- Intra: Direct link to all assets removed, just expose ports required (FTP, SMB, AD, etc.)
- Outbound: Connection initiated internally so let it through
- Note: L2 to L3

Security:

- Initial Access: Email, web view, exposed ports
- Lateral Movement: Still easy
- C2: Still easy, allow outbound
- Exfiltration: Still easy, allow outbound
- Note: But second but concern, all servers are together



Note: Different rules into each network – consider it a firewall per network

# Network 3

## Characteristic: Separate hosts by type, have a DMZ

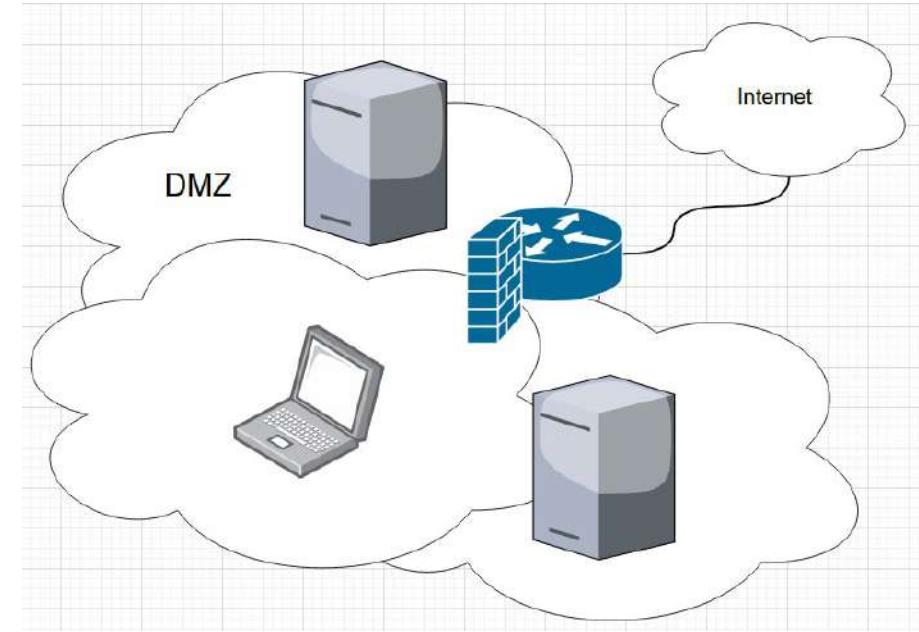
Separated network

Administration:

- Hosts: Internet-facing in the DMZ
- Access: Still direct to host or via exposed port (e.g., 80, 443)
- Intra: Expose ports to servers (FTP, SMB, AD, etc.)
- Outbound: Connection initiated internally so let it through
- Note: Should we allow connections from other LANs to DMZ?

Security:

- Initial Access: Email, web view, exposed ports
- Lateral Movement: Not as easy, but still easy
- C2: Still easy, allow outbound
- Exfiltration: Still easy, allow outbound
- Note: How can we isolate lateral movement even more?



# Network 4

## Characteristic: Remove the ability to jump directly between zones

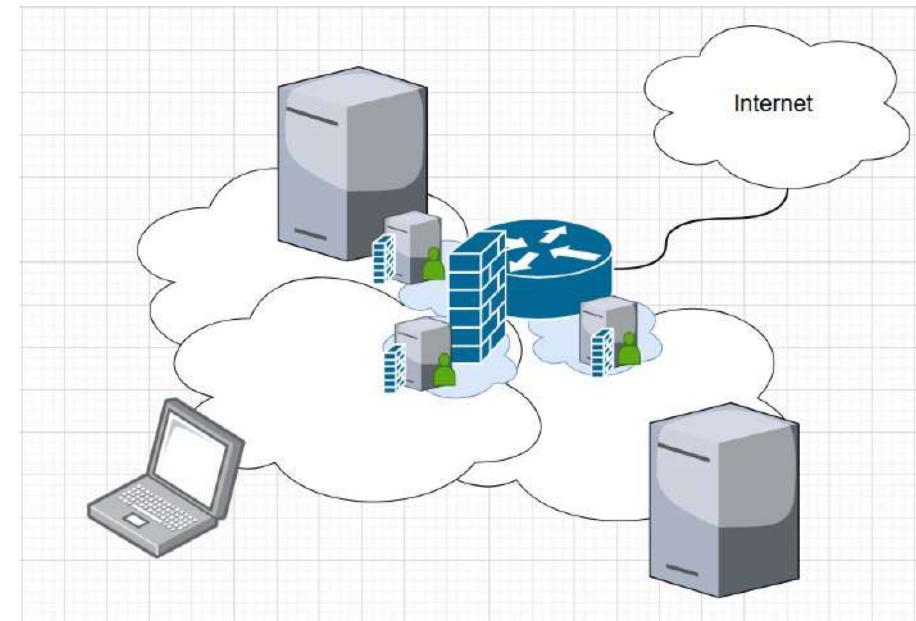
Separated network

Administration:

- Hosts: New jump host, much more admin
- Access: Via dedicated host
- Intra: Expose ports to servers (FTP, SMB, AD, etc.), remote access via jump host
- Outbound: Connection initiated internally so let it through
- Note: Should assets get equal Internet access?

Security:

- Initial Access: Email, web view, exposed ports
- Lateral Movement: Getting harder
- C2: Still easy, allow outbound
- Exfiltration: Still easy, allow outbound
- Note: Traffic is still getting out freely, possibly something we can do about it?



# Network 5

## Characteristic: Proxy traffic, Email gateway

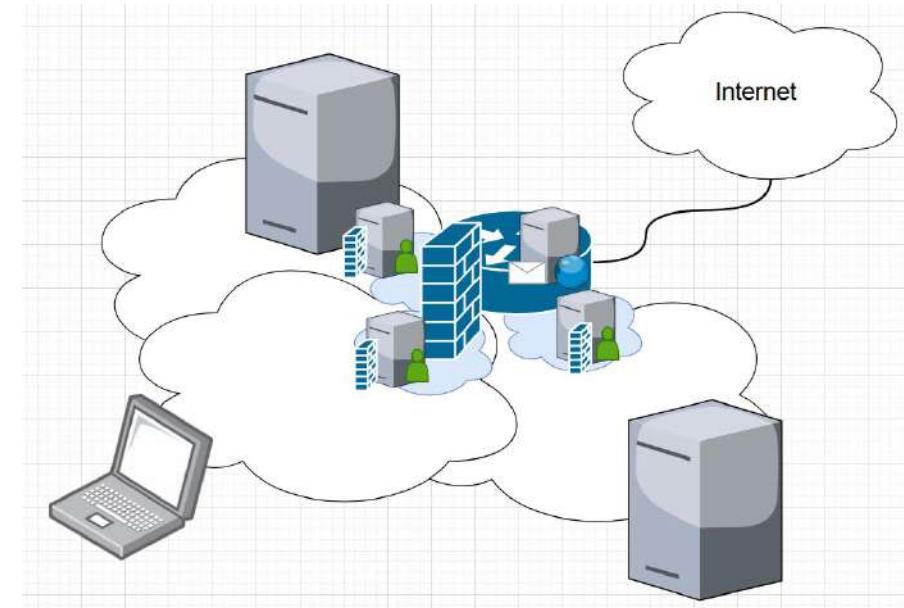
Separated network

Administration:

- Hosts: Proxy and Email gateway need rules
- Access: Via dedicated host
- Intra: Expose ports to servers (FTP, SMB, AD, etc.), remote access via jump host
- Outbound: Pass through proxy, apply better rules
- Note: What about the secret network or operational technology?

Security:

- Initial Access: Reducing vectors of email, web view, exposed ports
- Lateral Movement: Same as before
- C2: Restrictive
- Exfiltration: A little harder
- Note: How can we restrict communication further?



# Network 6

## Characteristic: Air gapped

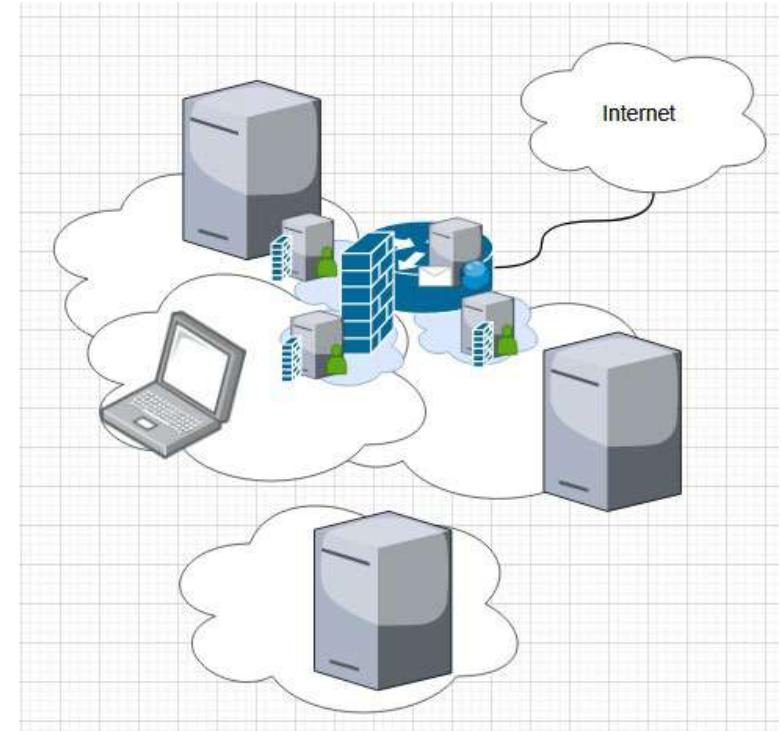
Separated network

Administration:

- Hosts: Entirely new setup now
- Access: None
- Intra: None
- Outbound: Layers of security
- Note: Administration is now complex

Security:

- Initial Access: Physical access
- Lateral Movement: Very complex if physically separated
- C2: Restrictive
- Exfiltration: Restrictive
- Note: How can we restrict communication further?



# Not Just Traffic

## Threat vectors of network-based assets

Devices themselves are also susceptible

Initial Access	Execution	Persistence	Privilege Escalation	Defense Evasion	Credential Access	Discovery	Lateral Movement	Collection	Command and Control	Exfiltration	Impact
2 techniques	2 techniques	10 techniques	4 techniques	11 techniques	6 techniques	11 techniques	1 techniques	4 techniques	7 techniques	2 techniques	4 techniques
Exploit Public-Facing Application Valid Accounts (2)	Command and Scripting Interpreter (3) Software Deployment Tools	Account Manipulation (1) Boot or Logon Autostart Execution Boot or Logon Initialization Scripts (1) Create Account (1) Modify Authentication Process (1) Power Settings Pre-OS Boot (3) Server Software Component (1) Traffic Signalling (1) Valid Accounts (2)	Account Manipulation (1) Boot or Logon Autostart Execution Boot or Logon Initialization Scripts (1) Valid Accounts (2)	Direct Volume Access Impair Defenses (3) Indicator Removal (3) Modify Authentication Process (1) Modify System Image (2) Network Boundary Bridging (1) Obfuscated Files or Information (1) Pre-OS Boot (3) Traffic Signalling (1) Valid Accounts (2) Weaken Encryption (2)	Adversary-in-the-Middle (2) Brute Force (4) Input Capture (1)	File and Directory Discovery Network Service Discovery Network Sniffing Password Policy Discovery Process Discovery Remote System Discovery System Information Discovery System Network Configuration Discovery System Network Connections Discovery System Owner/User Discovery System Time Discovery	Software Deployment Tools	Adversary-in-the-Middle (1) Data from Configuration Repository (2) Data from Local System Input Capture (1)	Application Layer Protocol (5) Encrypted Channel (2) Hide Infrastructure Ingress Tool Transfer Non-Application Layer Protocol Proxy (3) Traffic Signalling (1)	Automated Exfiltration (1) Exfiltration Over Alternative Protocol (1)	Disk Wipe (2) Firmware Corruption Inhibit System Recovery System Shutdown/Reboot

SOURCE: <https://attack.mitre.org/matrices/enterprise/network/>

# Air-Gap Case Study

## Stuxnet

### 1) Understand the story

- Background (2010):
  - Iran's Natanz uranium enrichment plant faced sudden and unexplained centrifuge failures
  - A sophisticated digital weapon designed to physically destroy equipment, specifically targeting Siemens industrial control systems
  - Main Goal: To sabotage the centrifuges at Natanz, a facility key to Iran's nuclear program
  - Siemens PLCs (Programmable Logic
- Controllers) used for controlling centrifuge speed.
- The system was air-gapped, meaning it was not directly connected to the Internet
  - Spread through USB flash drives using Autorun and local network exploits
  - Targeted systems of companies connected to the nuclear program, enabling the weapon to be introduced into Natanz indirectly

# Stuxnet

## Challenges

Continued

- Stuxnet overcame this security measure by using physical media (e.g., USB flash drives) to bridge the gap between isolated systems and the outside world
- The attackers first infected computers outside the air-gapped system, particularly targeting companies connected to the Iranian nuclear program
- These companies were involved in industrial control systems and had a connection to Natanz. Some were identified as contractors and suppliers of Siemens equipment used at the plant
- Upon inserting the USB drives into computers inside Natanz, Stuxnet spread across the internal network
- The infected computers at Natanz then communicated with Siemens PLCs, which controlled the centrifuges
- This allowed the attackers to manipulate the centrifuges' speed and pressure remotely, causing them to break down without detection



# Thank you

# **COS80013 Internet Security**

**Lecture Week 5A**



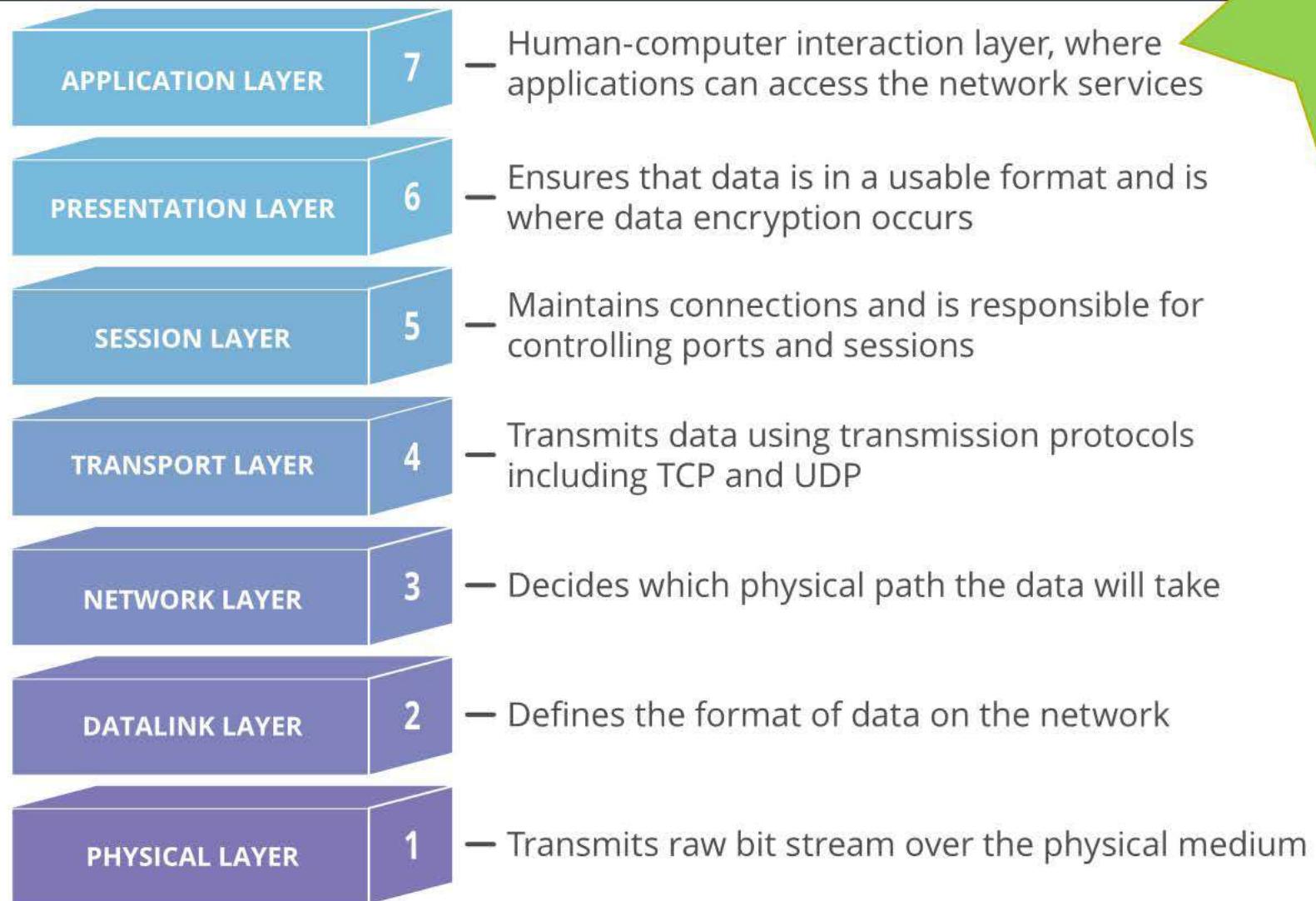
# Networking



- Network attacks
- Defence Tools

# OSI

Revision



(img source: <https://www.cloudflare.com/learning/ddos/glossary/open-systems-interconnection-model-osi/>)

# Network Attacks

- All layers of the OSI model are susceptible to attacks.
- Application layer:
  - viruses, worms, trojans
- Session layer:
  - null session attacks on Windows PCs
- Presentation layer
  - DNS DOS, DNS cache poisoning, Zone transfer
- Network layer
  - ARP cache poisoning, ARP spoofing
- Physical / Data link layers
  - Sniffing, playback attacks
- These and other layers have and will be attacked.

# MAC addresses are not fixed

**Android phones and tablets change MAC – on re-boot.**

<https://forums.androidcentral.com/verizon-galaxy-nexus/147857-wifi-mac-address-changes-every-reboot.html>

“There was kernel code identified that will randomize a MAC address”

**Apple things change MAC address before they do wi-fi-association.**

These “features” are to prevent tracking by

WiFi.

<https://www.theverge.com/2014/6/9/5792970/ios-8-strikes-an-unexpected-blow-against-location-tracking>

Also breaks MAC filtering, some other security protocols

# Packet Sniffing

- Packet sniffers record IP packets on the network. They were originally designed to help diagnose problems in networks.
  - Good for picking up MAC addresses, IP addresses
- Many internet-based services expect to receive user names and passwords in plain text.
  - Telnet, FTP, SNMP
- Computer users get lazy and re-use the same user names and passwords.
  - If you can get their FTP password, you can probably use it on other accounts.
- Popular Sniffers:
  - TCPDump, Snort, Wireshark. Windows and Linux versions.
  - reviews: <http://sectools.org/sniffers.html>

# Packet Sniffers

## Wireshark

- Text (tethereal) and GUI (wireshark) versions available for Windows and Linux
- Lists packet contents on the screen and logs them to a file for analysis later.
- Summarises types of packets intercepted.

## Snort

- Text-based IDS with packet sniffing and logging abilities.
- Separates packets into IP address and port number
- Easy to search packets using *grep* (linux)

# Common Attacks

## ARP cache poisoning.

- ARP maintains a table of MAC addresses and their equivalent IP addresses.
- ARP sends out queries: "Who has IP 239.254.2.15"
- Replies are added to the ARP table, and can come from any IP address at any time.

## How it works (MITM)

- Attacker attaches his PC to the network.
- sniffs to find target IP addresses.
- Sends a home-made ARP reply nominating his MAC as corresponding to the requested IP address.
- ARP adds the attackers MAC address to the ARP table.
- Future traffic to the target goes to the attacker.
- Attacker reads packets, doctors them and sends them on to the target who is none the wiser.

# Common Attacks

## MAC flooding.

- Switches use ARP to map MAC addresses to IP addresses.
- When booted up, switches operate in hub mode while they learn the MAC addresses attached to each port.
- Switches have limited memory and computing power.
- When overloaded, they revert to Hub-mode and broadcast all packets to all interfaces.

## How it works

- Attacker attaches his PC to the network.
- Sends out multiple invalid ARP responses.
- Switch overloads and reverts to hub-mode (all lights start flashing).
- Attacker collects all packets from the network.

# DNS Attacks

## DNS cache poisoning

- Attacker sends many DNS queries and replies at the same time. The replies spoof the IP of the authoritative name server.
- Need to guess the DNS query ID (0-65535)
  - may be predictable
  - birthday paradox increases chances of a correct guess to ~50%

## Prevention:

- Use random identifiers for queries
- Always check identifiers
- Port randomization for DNS requests
- Deploy DNSSEC
  - Challenging because it is still being deployed and requires reciprocity

# **DoS (Denial of Service)**

**DOS attacks are aimed at servers on the internet**

- web servers (HTTP), name servers (DNS), FTP servers
- Can be launched at specific machines if the IP address is known.

**The goal is to Deny Service to legitimate customers/users.**

- Customers go elsewhere
- Organisations lose money/trade/reputation

# ICMP Flooding (DOS)

## ICMP (ping)

- attacker sends a stream of ICMP echo request packets.
- Solution:
- Change your IP address or block ICMP traffic at the ISP or router.
- Buy a router/appliance that drops ping packets really fast.

# SYN Flooding (DOS)

## SYN Flooding

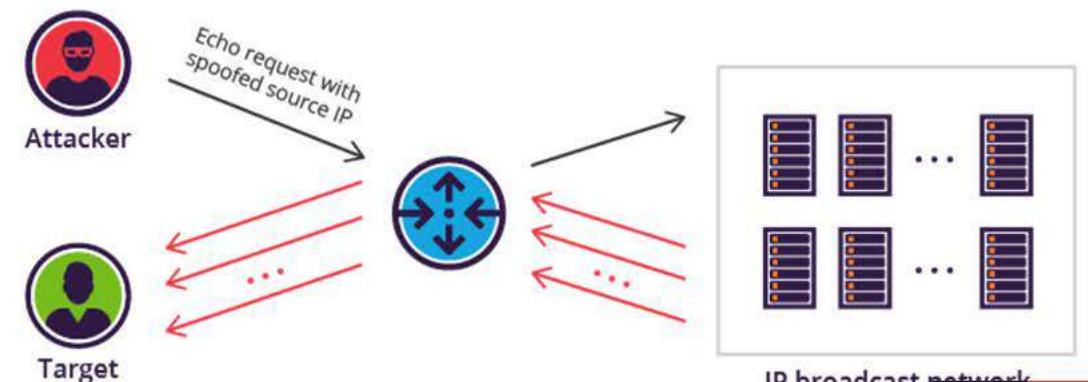
- a resource depletion attack
- attacker sends a stream of SYN packets (the first part of a TCP 3-way handshake).
- Each SYN packet consumes memory on the server while it sends the SYN\_ACK packet and waits for the returning ACK packet.
- The attacker never sends the ACK packet.
- Syn attack solutions:
  - SYN cookies – Instead of storing the SYN pack info in memory, it is returned as a SYN cookie to the attacker's machine. When the attacker's machine returns an ACK packet (which it won't) the memory for the connection is allocated.
  - RST cookies – After receiving at attacker's SYN packet, the target returns a SYN\_ACK packet containing the wrong info. The attacker should then send a RST packet (which he won't), after which connections from that client are now accepted.
  - Modern switches detect and ignore SYN floods

# Smurf Attack (DOS)

- Amplification attack
  - An attacker sends a ping (ICMP echo request) to the *broadcast* address of the victim's network.
  - This sends the ping to all IPs on the subnet. The original ping has a spoofed return IP address which is the address of the intended victim.
  - Each IP reached by the original ping replies to the victim.

Solution:

- Disable/filter broadcast address
- Modern switches detect and prevent Smurf attacks



# UDP Flood (DOS)

UDP flood

- UDP packets are sent to random ports of the victim as fast as possible.
- The target machine will respond to each packet with a ICMP destination unreachable packet.

Solution:

- turn off ICMP destination unreachable.

# Ping of Death (DOS)

## Ping of death

- an over-sized ping packet (>65535 bytes) is sent to a susceptible PC.  
Causes a re-boot on susceptible machines.

## Solution:

- Use a newer OS / product or patch the old one.

Router vulnerability: <http://www.securityfocus.com/bid/1240>  
ping -t -l 65500 example.victim.com

# Teardrop (DoS)

## Teardrop

- A fragmented packet is sent to the attacker, but the parts have been corrupted and can not be put back together.
- Server consumes resources requesting retransmit.
- May crash the server.

## Solution:

- Use a newer OS / product or patch the old one.  
TCP/IP vulnerability: <http://www.securityfocus.com/bid/124/exploit>

# LANd attack (DOS)

## LANd attack

- Amplification attack
- A packet with a spoofed source IP address is sent to the victim. The return address is the same as the destination addresses, so the victim's machine ties it self in knots by answering back to itself.

## Solution:

- Update/patch your OS.
- Smarter switches and hubs detect this and block it.

[http://www.securityfocus.com/bid/13658  
exploit](http://www.securityfocus.com/bid/13658/exploit)

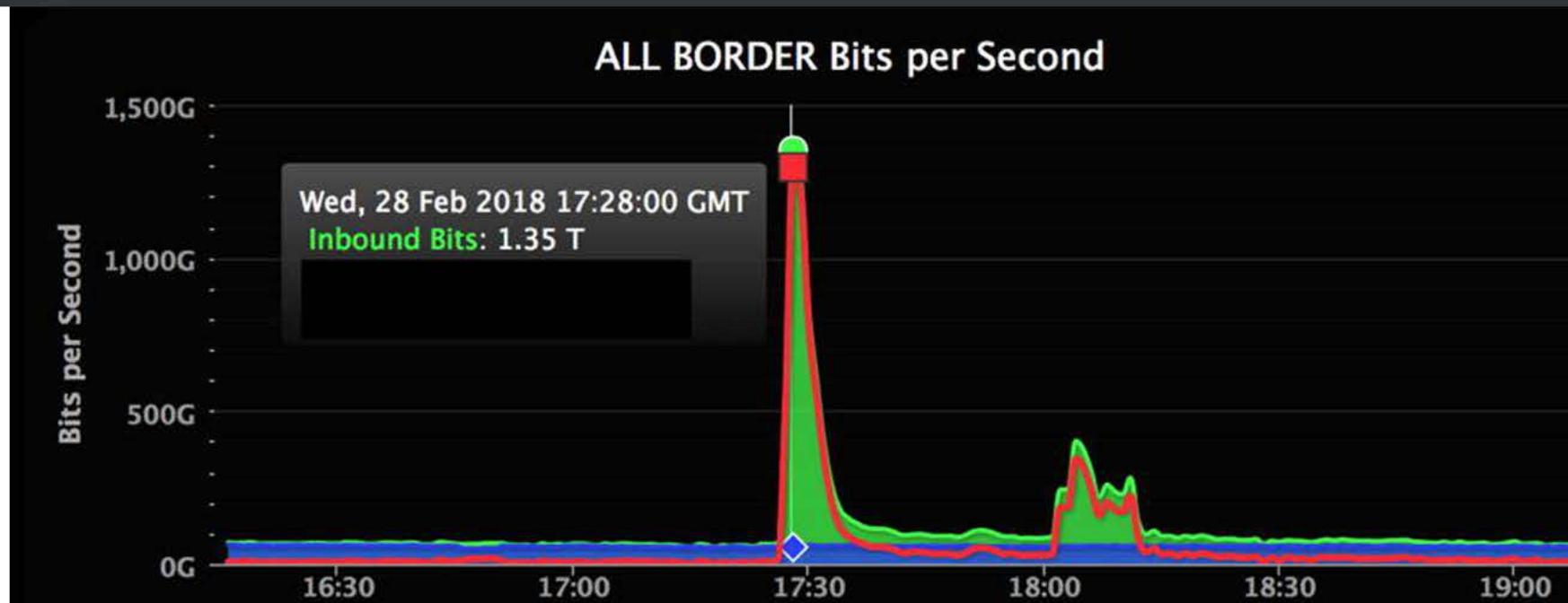
# DDoS

- A *Distributed DOS* involves simultaneous attacks on a single site by 'zombie' machines infected with a 'bot'. Each bot is controlled by a 'bot-herder' originally through IRC channels and now through http. A bot-herder's 'bot-net' may range in size from a few hundred to millions of infected PCs.

<http://blog.washingtonpost.com/securityfix/2006/03/post.html>

- Each bot is actually a small server program which has installed itself and is capable of launching DOS attacks, sending spam, infecting other machines, or all of the above.
- Spread by trojans (e-mail, web downloads) or by worms.

# DDoS attack example



June 25, 2020

4

## Re-Hash: The Largest DDoS Attacks in History

Amazon reported sustaining a 2.3 Tbps DDoS attack in 2020 – here's what to know about the largest DDoS attacks on record & how they're measured

# Stopping DDoS

DDOS targets are chosen for a reason

- Find out who your enemies are.
- Are you a victim of extortion?
- Have you offended a script kiddie?
- Is there a political motive?

Options:

- Get a new IP address.
- Mitigation strategies: rate limiting, blackholing (not very effective)
- READ the NOTES
- <http://www.symantec.com/connect/articles/closing-floodgates-ddos-mitigation-techniques>

# Stopping DDoS

DDOS mitigation

Use cloud services to absorb/filter your traffic

- Very cost-effective
  - scale up during an attack
  - only pay for the bandwidth you need.

<https://www.cloudflare.com/features-security>

<http://www.spamhaus.org/news/article/695/answers-about-recent-ddos-attack-on-spamhaus>

<http://www.prolexic.com/services-dos-and-ddos-mitigation.html>

<http://ddos-protection-services-review.toptenreviews.com/>

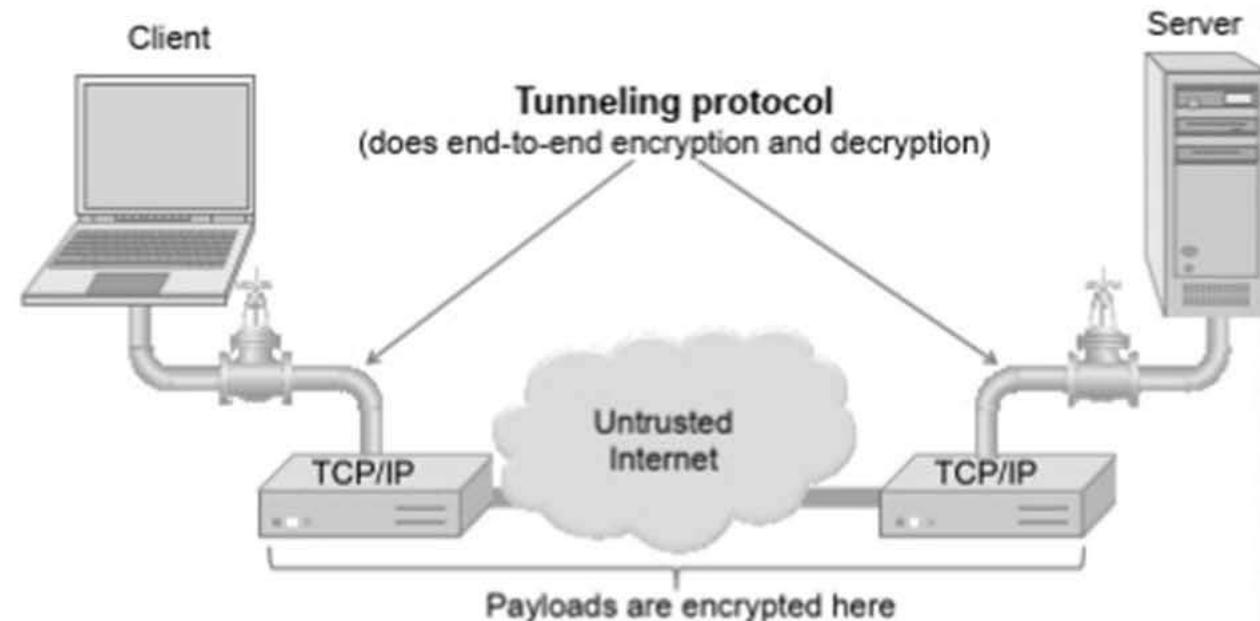
# Defence Strategies

- Identify attack early
- overprovision bandwidth
- defend at the network perimeter
- talk to your ISP
  - filtering, rate-limiting
  - port knocking
- call a DDoS mitigation specialist
  - some services do not need to be protected
  - some services should not be on the Internet (use a private corporate network)

# Tunneling

- **SSH, VPN, SSL, TLS, IPSec**

As the Internet is not secure by default, if someone is eavesdropping on a TCP connection, they can often see the complete contents of the payloads in this session. One way to prevent such eavesdropping without changing the software performing the communication is to use a tunneling protocol.



# SSH, SSL, HTTPS?

- All form encrypted tunnels.
- SSH - secure sockets handler
  - specifies the protocols that occur in the encrypted tunnel.
  - includes multiplexing and **user authentication**
  - an application which replaces Telnet (sort-of application layer).
- SSL - Secure Sockets Layer
  - uses certificate (X.509)-based key exchange
  - encrypts **data only** (sort-of transport layer)
- HTTPS
  - HTTP tunneled through SSL (now referred to as TLS)
  - <http://security.stackexchange.com/questions/1599/what-is-the-difference-between-ssl-vs-ssh-which-is-more-secure>

# SSH

- Secure Shell (SSH), secure copy protocol (SCP) and secure file transfer protocol (SFTP)
- Application layer protocol that encrypts application layer payload and then sends it by TCP.
- Uses a range of crypto (keys, certificates)
- Nullifies sniffing attacks. **BUT requires chain of trust**
- Can use server-side certificates only (like https) or server and client certificates.
- Built-in username/password authentication

# SSL..

1. TCP connection established
2. Client requests connection to server
3. Server sends **certificate (X509)** and **public key**
4. Client compares key with those in its cache, CA.  
**User asked to accept/cache key**
5. Client creates session key and sends it to server  
encrypted with server's public key
  - Server may request client public key (**optional**)
  - Server may check client's public key (cache/CA)
6. All subsequent traffic encrypted by session key

# SSH vs SSL

- SSH: Encrypt-then-MAC
  - Send ciphertext and MAC
  - Proves integrity
  - Encryption is a bit weaker  
(easier to brute-force key)
- SSL: MAC-then-Encrypt
  - MAC added to plain text and then both are encrypted
  - Vulnerable to some known plain text attacks (BEAST attack).

Message  
Authentication  
Code  
Think:  
Checksum

# SSL Vulnerability

- Heartbleed (2014)
  - Affects some OpenSSL libraries
  - Problem with the "Heartbeat" keep-alive packet
  - Allows user to scrape 64kB of RAM **before authentication.**
  - Get other people's logins
- Fix:
  - Patch OpenSSL or use a different SSL library.
- See Podcast, Slides (Blackboard)

# IP Security (IPsec)

- Suite of protocols from Internet Engineering Task Force (IETF) providing encryption and authentication at the IP layer
  - Arose from needs identified in RFC 1636
  - Specifications in:
    - RFC 2401: Security architecture
    - RFC 2402: Authentication
    - RFC 2406: Encryption
    - RFC 2408: Key management
- Objective is to encrypt and/or authenticate all traffic at the IP level.

# IP Security Issues

- Eavesdropping
- Modification of packets in transit
- Identity spoofing (forged source IP addresses)
- Denial of service
- Many solutions are application-specific
  - TLS for Web, S/MIME for email, SSH for remote login
- **IPSec aims to provide a framework of open standards for secure communications over IP**
  - Protect every protocol running on top of IPv4 and IPv6

# IPsec Benefits

- Provides a level of security for all applications.
  - Allows deployment of new/emerging applications that may not have their own security.
- Transparent to transport layer.
- Transparent to end-users.
  - No need for training, key issue, key revocation, etc.
- Can be provided to individual users where needed (e.g. off-site workers).
- Extensible to new, stronger, cryptographic methods as these become available.

# IPsec Drawbacks

- Processing performance overhead
  - Protection is applied to all traffic, though only a small portion may be security-sensitive
- Blocks access to non-IPsec hosts
- Hosts must have security association
  - Not great for short-lived connections
- Not practical for broadcast

# Uses of IPsec

- Virtual Private Network (VPN) establishment
  - For connecting remote offices and users using public Internet
- Low-cost remote access
  - e.g. teleworker gains secure access to company network via local call to ISP
- Extranet connectivity
  - Secure communication with partners, suppliers, etc.

# VPN

- Virtual Private Network
- Tunnels all traffic (all packets) over a public network.
- Works with non-routable packets
- Encrypts IP packets and sends them over public TCP/IP as encrypted payloads.
- Many implementations, some vulnerable
  - PPTP (point to point tunneling protocol)
  - L2TP (layer 2 tunneling protocol)

# VPN

- Tunneling provides separation between the outside of the tunnel (internet) and the inside (trusted private network).
- Provides no protection within the private network.
- If an **endpoint gets infected**, malware can pass through the tunnel to the rest of the private network without any barriers.
- Equivalent to LAN access to the trusted network.

# DNS Attacks

# DNS and Hosts

- Before the internet got big, people typed in the IP address of each site they wanted to visit.
- This became tedious, so a scheme was devised whereby the host and domain name of frequently visited web sites was added to a file (HOSTS) which is used to look up the IP address which corresponds to a **host.domain** name. Users could download and install updated versions of *HOSTS* files from popular web sites.
- DNS automated the domain name lookup process using dedicated DNS servers which could communicate with each other and any interested PC through UDP port 53.

# DNS and Hosts...

- The contents of a HOSTS file take precedence over DNS queries. HOSTS should be empty or just contain the loopback address (127.0.0.1) on modern PCs.
- A common strategy used by spy/adware, browser hijackers and **pharmers** is to add incorrect entries to the HOSTS file so that correctly entered URIs direct you to the wrong web site.
- A common strategy used for web filtering is to add the domain names of forbidden sites to the hosts file with the loopback address.
- On Linux: see **/etc/hosts**
- On Win: see **\Windows\System32\drivers\etc\hosts**

# DNS Attacks

- DNS cache poisoning
  - Attacker sends many DNS queries and replies at the same time. The replies spoof the IP of the authoritative name server.
  - Need to guess the DNS query ID (0-65535)
    - may be predictable
    - birthday paradox increases chances of a correct guess to ~50%
- Kaminsky attack
  - Query non-existent hosts on the target domain.
    - Easier to guess the DNS query ID
    - DNS client caches ADDITIONAL record (including IP of target host)

# DNS Cache Poisoning Prevention

- Use random identifiers for queries
- Always check identifiers
- Port randomization for DNS requests
- Deploy DNSSEC
  - Challenging because it is still being deployed and requires reciprocity

# Kaminsky Defences

- 2008 – patches which increases randomness of DNS IDs
- Local DNS servers only accept queries from inside the network
- Source-port randomisation
- UDP packet ID combined with DNS ID to enlarge ID range to 32bits (3 billion)
- DNSSec

# DNSSec

- DNSSec adds two important features to the DNS protocol
  - Data origin authentication (verify where data came from) & Data integrity protection (verify data didn't been modified during transit)
- DNS queries, replies are digitally signed.
  - prevents replies from being spoofed (fake source address, ID)
  - requires public key crypto, chain of trust
    - But certificates, private keys get stolen
    - DNSSec adoption very small (1%)
    - <https://www.eweek.com/security/dnssec-adoption-needs-to-grow-to-secure-core-internet-protocols/>

# NSLookup (Win/Linux/Mac)

- NSLookup queries the domain name system to find the IP address of a domain name.
- NSLookup can
  - Resolve DNS Issues
  - Search for optimal mail servers

# NSLookup

- NSLookup works for different levels of the DNS system.
- For the details of the `.edu.au` name servers, try

```
c:\temp>nslookup -type=mx edu.au
```

```
Server: venus.it.swin.edu.au
```

```
Address: 136.186.5.30
```

`edu.au`

primary name server = ns1.ausregistry.net

responsible mail addr = dns.ausregistry.net.au

serial = 2003071563

refresh = 14400 (4 hours)

retry = 3600 (1 hour)

expire = 3600000 (41 days 16 hours)

default TTL = 86400 (1 day)

# Reverse DNS Lookup

- **Use this to verify that an IP points to a domain name.**
  - "The reverse DNS entry for an IP is found by **reversing the IP**, adding it to "**in-addr.arpa**", and looking up the PTR record. So, the reverse DNS entry for 66.249.72.14 is found by looking up the PTR record for **14.72.249.66.in-addr.arpa**."  
[DNSStuff.com](http://DNSStuff.com)

```
c:\temp>nslookup -type=ptr 14.72.249.66.in-addr.arpa
```

Server: venus.it.swin.edu.au

Address: 136.186.5.30

Non-authoritative answer:

```
14.72.249.66.in-addr.arpa      name = crawl-66-249-72-14.googlebot.com
```

```
72.249.66.in-addr.arpa  nameserver = ns4.google.com
```

```
72.249.66.in-addr.arpa  nameserver = ns1.google.com
```

```
72.249.66.in-addr.arpa  nameserver = ns2.google.com
```

```
72.249.66.in-addr.arpa  nameserver = ns3.google.com
```

```
ns1.google.com  internet address = 216.239.32.10
```

```
ns2.google.com  internet address = 216.239.34.10
```

```
ns3.google.com  internet address = 216.239.36.10
```

```
ns4.google.com  internet address = 216.239.38.10
```

# Whois

- Gets name server info + details of system administrators.
- Difficult to use (have to get the right whois server).

```
[jhamlynharris@mercury jhamlynharris]$  
whois -h whois.melbourneit.com  
telstra.COM  
[whois.melbourneit.com]
```

Domain Name..... telstra.com  
Creation Date..... 1995-09-14  
Registration Date.... 2001-08-28  
Expiry Date..... 2009-09-13  
Organisation Name.... Telstra Corporation  
Organisation Address. 18/300  
Organisation Address.  
Organisation Address. MELBOURNE  
Organisation Address. 3000  
Organisation Address. VIC  
Organisation Address. AUSTRALIA

Admin Name..... Domains Administrator  
Admin Address..... 18/300  
Admin Address.....  
Admin Address..... MELBOURNE

Admin Address..... 3000  
Admin Address..... VIC  
Admin Address..... AUSTRALIA  
Admin Email..... ~~cdp@tppinternet.com~~  
Admin Phone..... +61.883084046  
Admin Fax.....

Tech Name..... Domains Administrator  
Tech Address..... 18/300  
Tech Address.....  
Tech Address..... MELBOURNE  
Tech Address..... 3000  
Tech Address..... VIC  
Tech Address..... AUSTRALIA  
Tech Email.....  
corpdomains@team.telstra.com  
Tech Phone..... +61.883084046  
Tech Fax.....  
Name Server..... dns0.telstra.net  
Name Server..... dns1.telstra.net  
Name Server..... sec1.apnic.net  
Name Server..... sec3.apnic.net

```
[jhamlynharris@mercury jhamlynharris]$
```

- On the web at <http://www.whois.net/>
- and on Linux computers.
  - try *whois -h whois.aunic.net*  
*swin.edu.au*