

12

Merge Sort

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CSA221 : Data Structures and Algorithms

MergeSort

Solving a Jigsaw Puzzle :

DIVIDE

Identify key areas: like the four corners, the edges & the middle section.

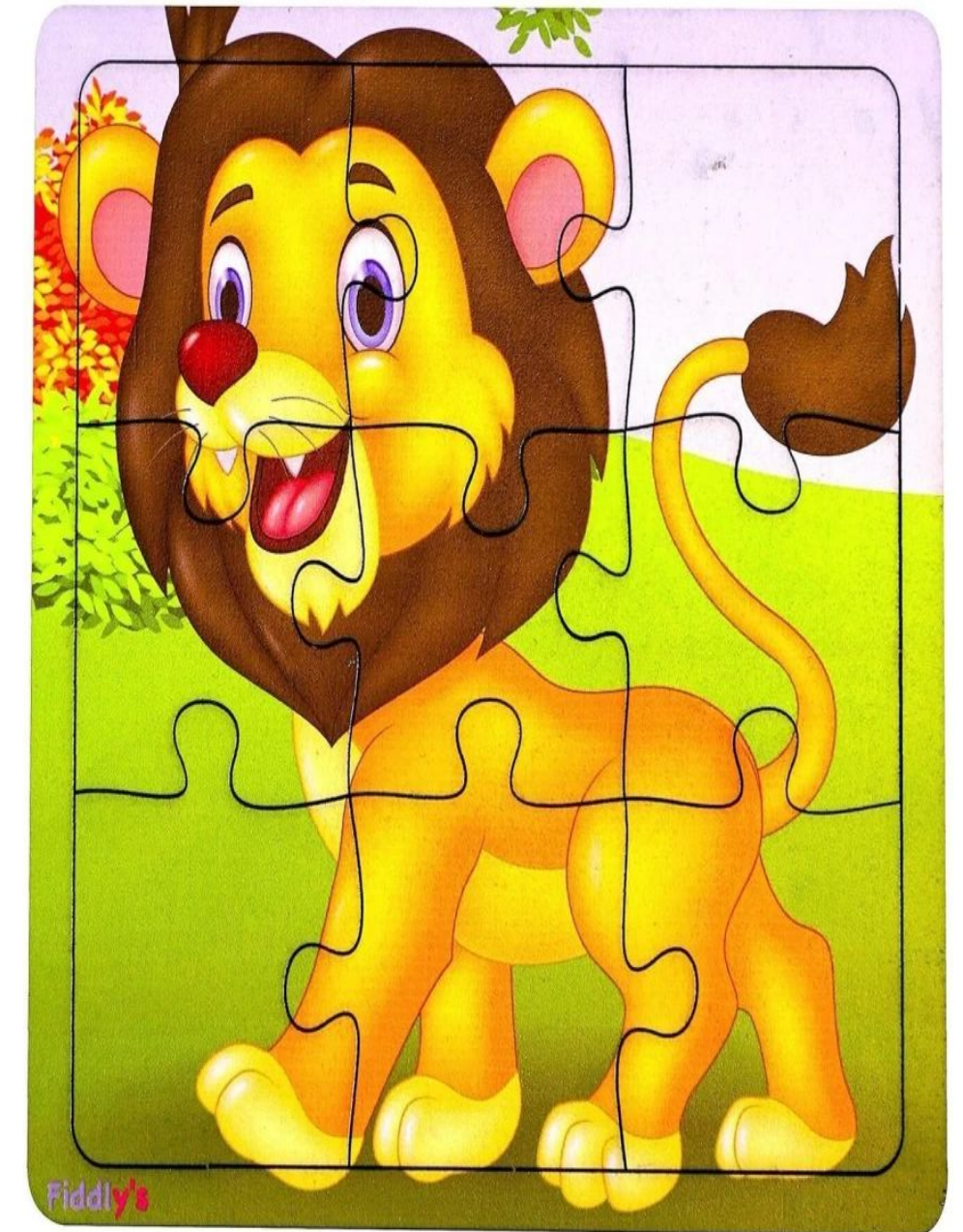
Subdivide the sections into smaller sections further if needed

Conquer

- **Focus on one section at a time**
- **Use trial and error** to complete the individual section
- **Utilize reference image** : If available

Combine

- **Connect sections**
- **Fill in gaps**



Merge Sort :

➡ It is reliable sorting algorithm that can handle **large datasets with minimal complexity.**

➡ It uses the **divide ,conquer and combine** approach to break the problem into smaller subproblems, sort them individually, and then merge them in a sorted manner.

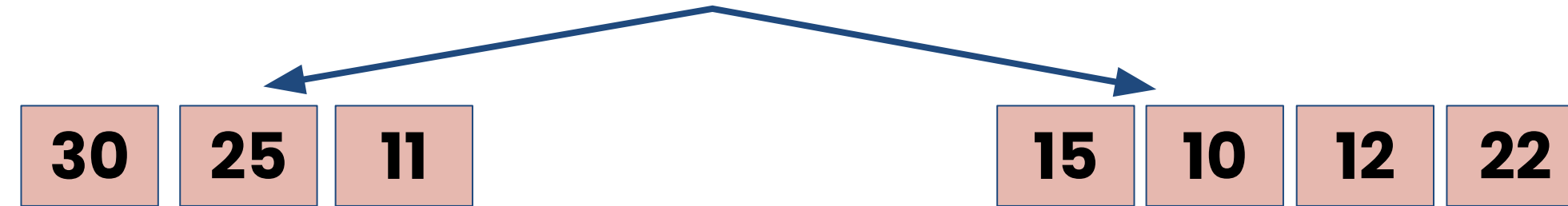
Example

| Index | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|-----------|----|----|----|----|----|----|----|
| Data Item | 30 | 25 | 11 | 15 | 10 | 12 | 22 |

Example

| Index | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
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| Data Item | 30 | 25 | 11 | 15 | 10 | 12 | 22 |

Divide Phase

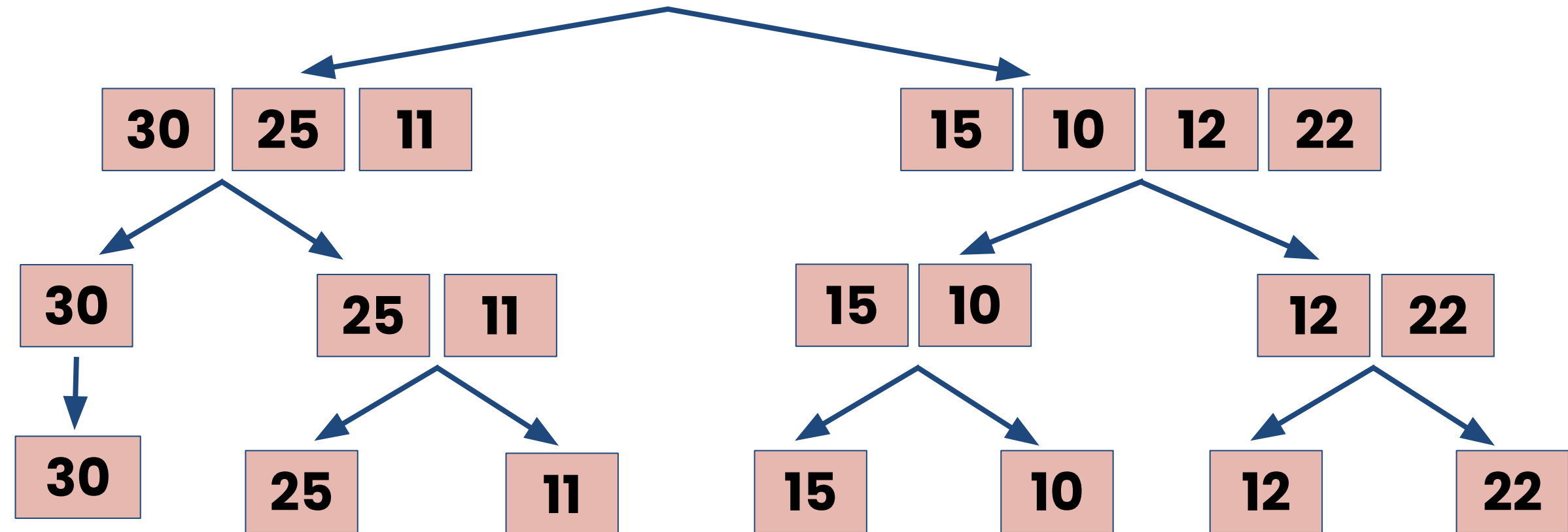


Example

| Index | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|-----------|----|----|----|----|----|----|----|
| Data Item | 30 | 25 | 11 | 15 | 10 | 12 | 22 |

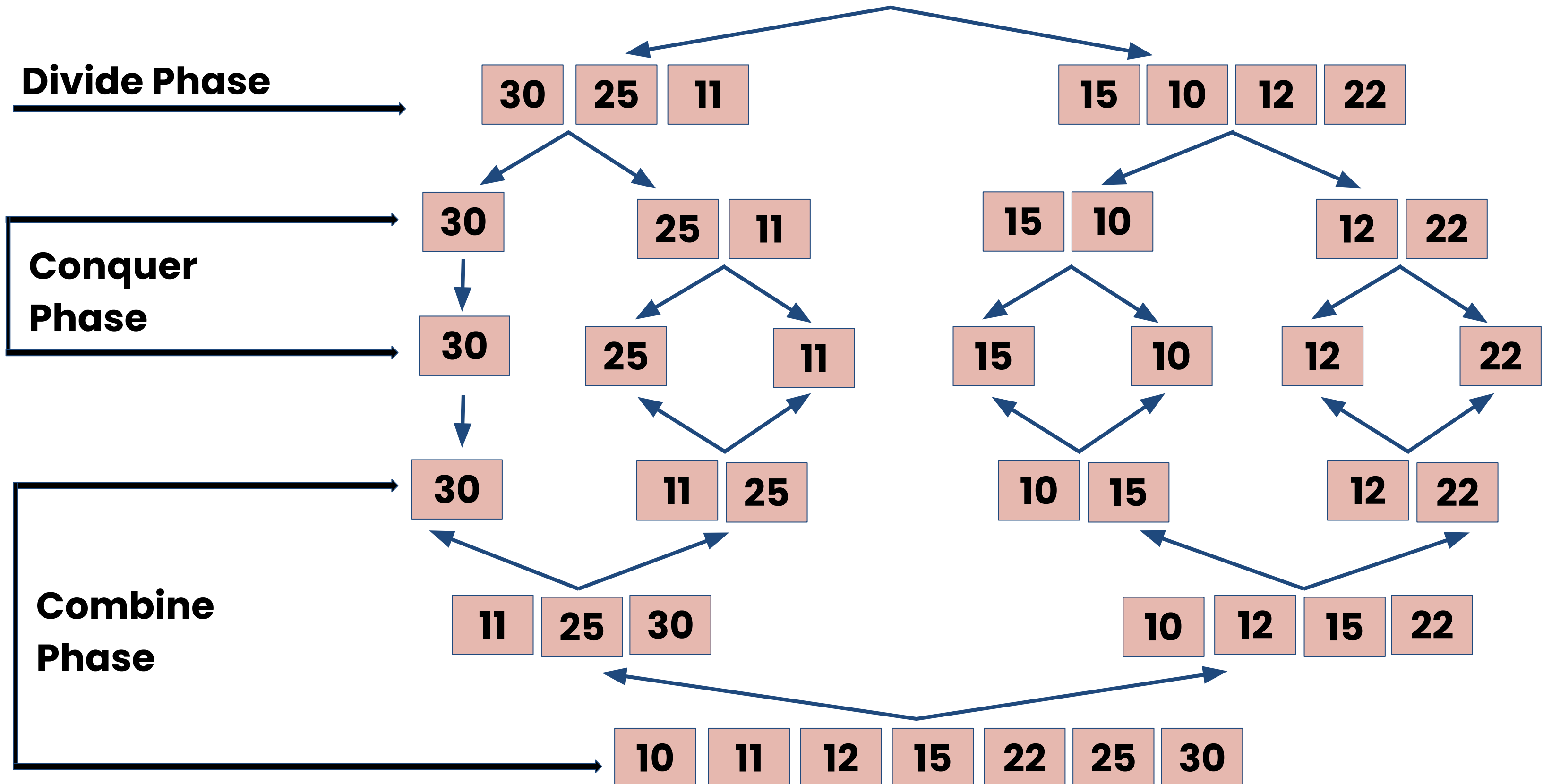
Divide Phase

**Conquer
Phase**



Example

| Index | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|-----------|----|----|----|----|----|----|----|
| Data Item | 30 | 25 | 11 | 15 | 10 | 12 | 22 |



Advantages :

1. **Stability**: Maintains the relative order of equal elements.
2. **Efficiency**: Performs well with **$O(n \log n)$** time complexity for all cases.
3. **Scalability**: Handles large datasets effectively, especially when implemented for external sorting using disk storage.
4. **Predictable Performance**: Unlike quicksort, it doesn't degrade to $O(n^2)$ in the worst case.

Merge Sort is particularly suited for applications where consistent performance and large-scale data handling are essential.

Divide-Conquer and Combine approach

Divide the subarray $A[p:r]$ to be sorted into two adjacent subarrays, each of half the size. To do so, compute the **midpoint** q of $A[p:r]$ (taking the average of p and r), and divide $A[p:r]$ into subarrays $A[p:q]$ and $A[q+1:r]$.

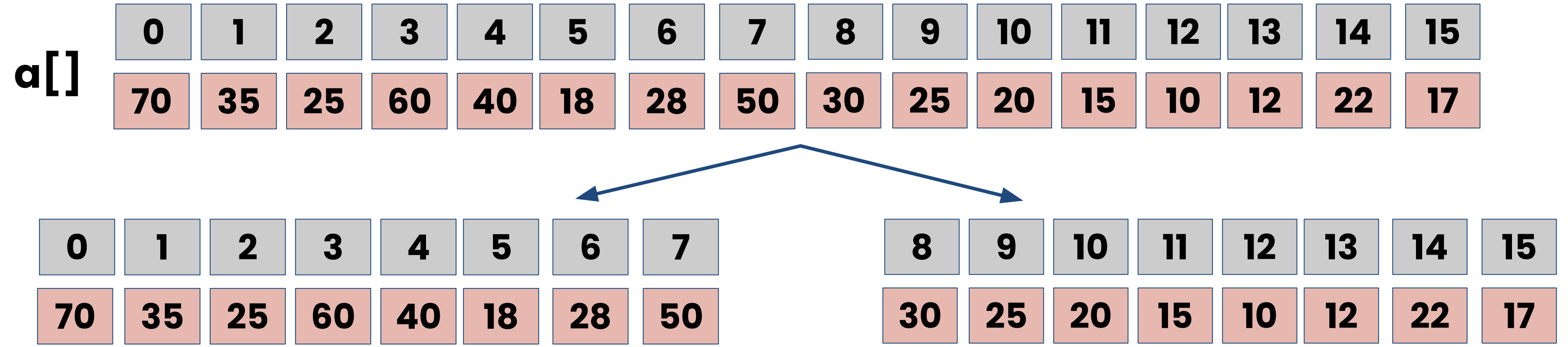
Conquer by sorting each of the two subarrays $A[p:q]$ & $A[q+1:r]$ **recursively** using merge sort.

Combine by **merging** the two sorted subarrays $A[p:q]$ & $A[q+1:r]$ back into $A[p:r]$, producing the sorted answer.

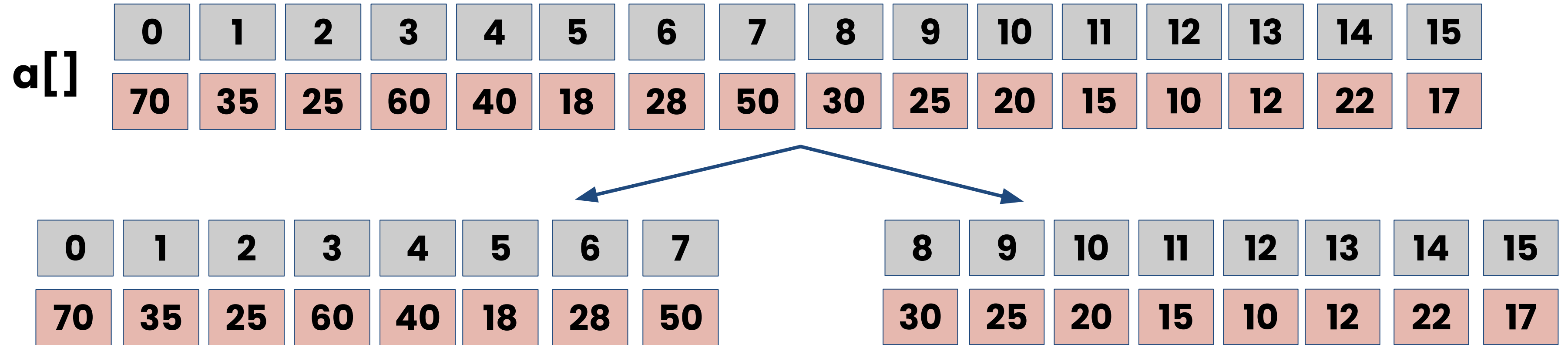
Suppose we have list of 16 elements

| | | | | | | | | | | | | | | | | |
|-----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| a[] | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| | 70 | 35 | 25 | 60 | 40 | 18 | 28 | 50 | 30 | 25 | 20 | 15 | 10 | 12 | 22 | 17 |

Suppose we have list of 16 elements

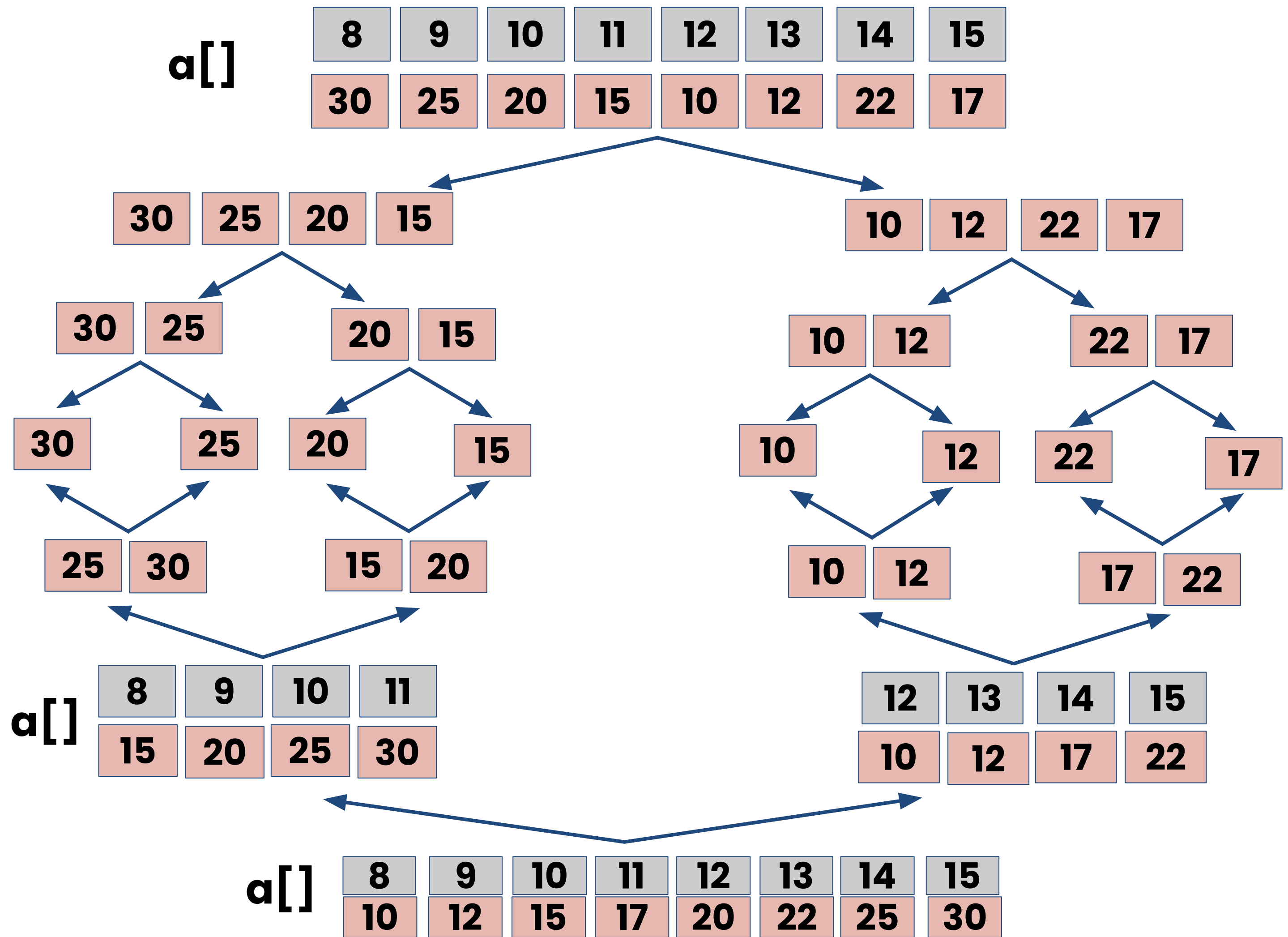


Suppose we have list of 16 elements



**Consider the indexes from 8 to 15 to
understand the merging process**

| | | | | | | | |
|----|----|----|----|----|----|----|----|
| 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| 30 | 25 | 20 | 15 | 10 | 12 | 22 | 17 |



Merging Process:

List L[]

| | | | |
|----|----|----|----|
| 0 | 1 | 2 | 3 |
| 15 | 20 | 25 | 30 |

| | | | | |
|----|----|----|----|-----------------|
| 0 | 1 | 2 | 3 | List R[] |
| 10 | 12 | 17 | 22 | |

Merging Process:

List L[]

| | | | |
|----|----|----|----|
| 0 | 1 | 2 | 3 |
| 15 | 20 | 25 | 30 |

↑

i=0

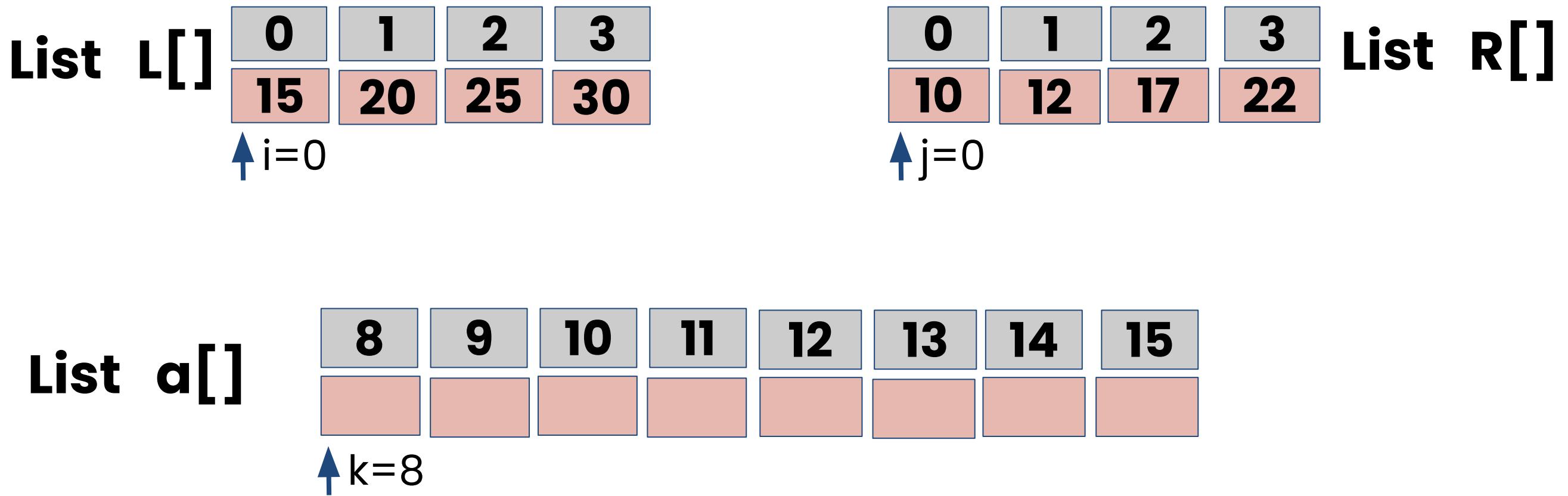
| | | | |
|----|----|----|----|
| 0 | 1 | 2 | 3 |
| 10 | 12 | 17 | 22 |

↑

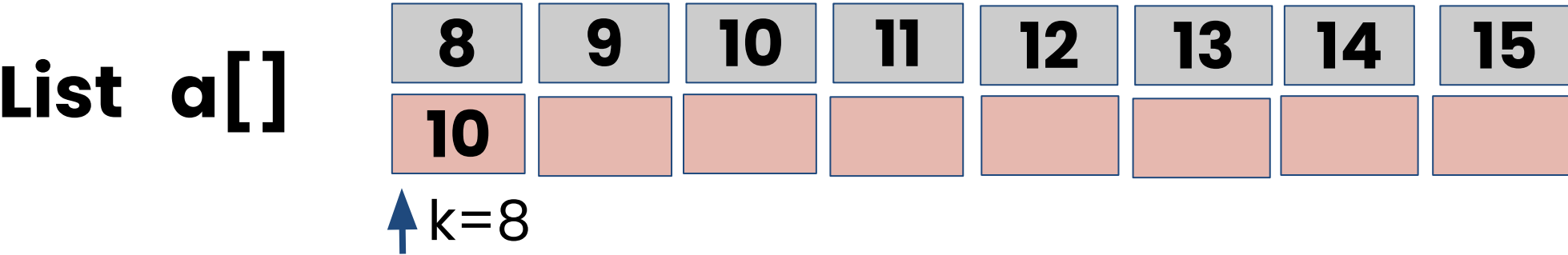
j=0

List R[]

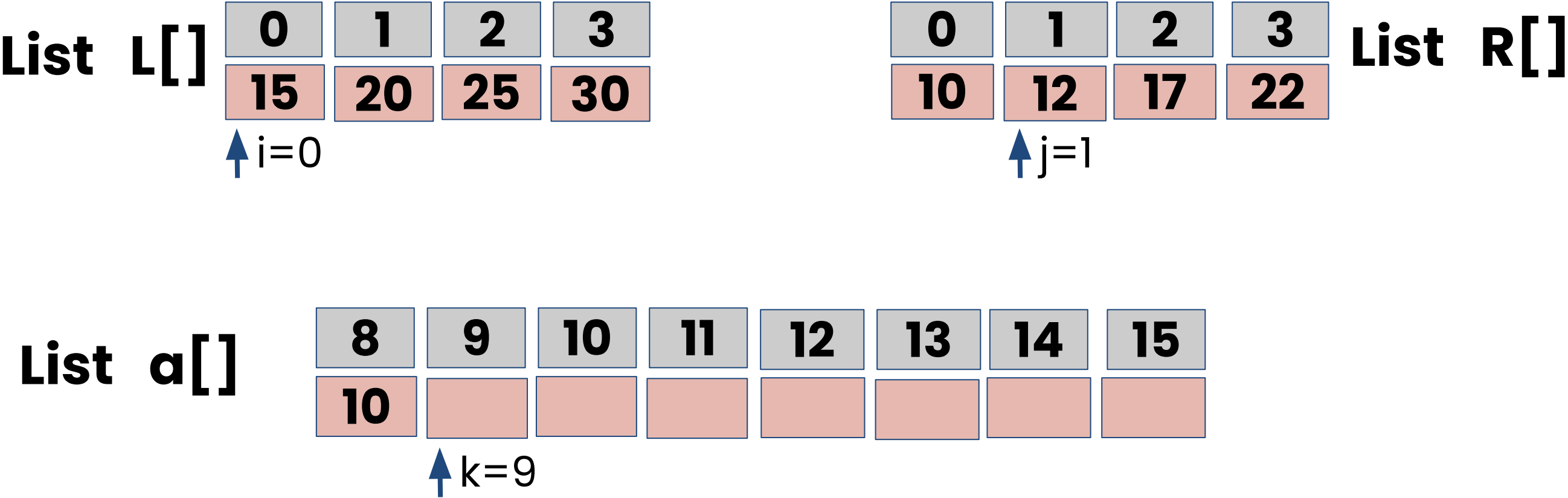
Merging Process:



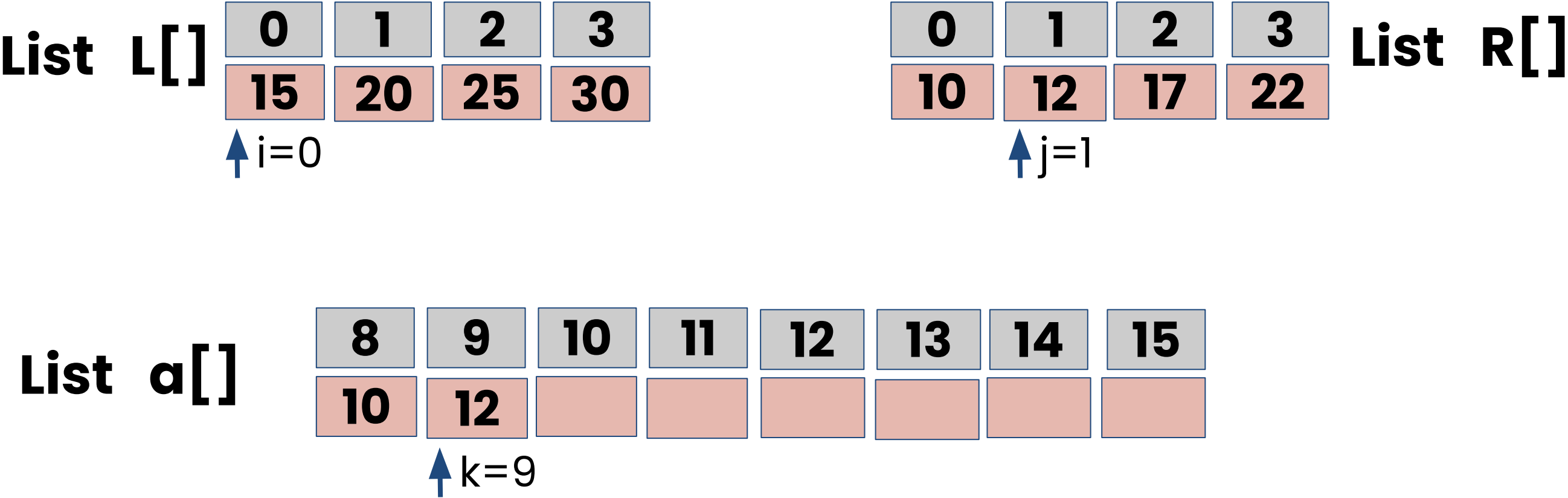
Merging Process:



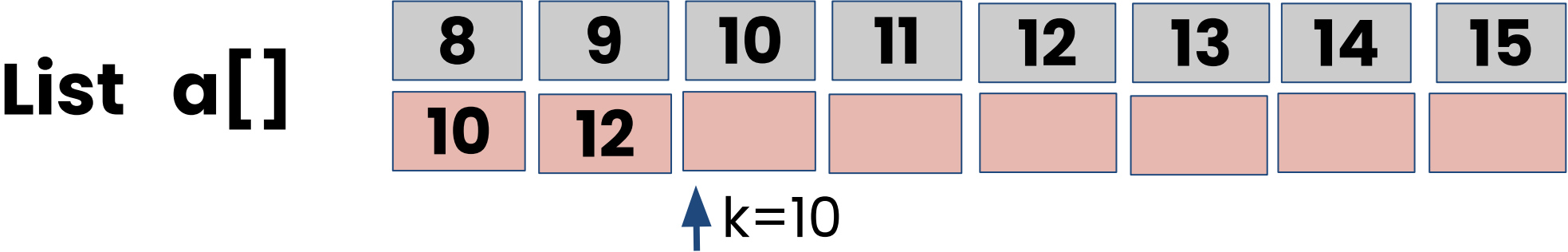
Merging Process:



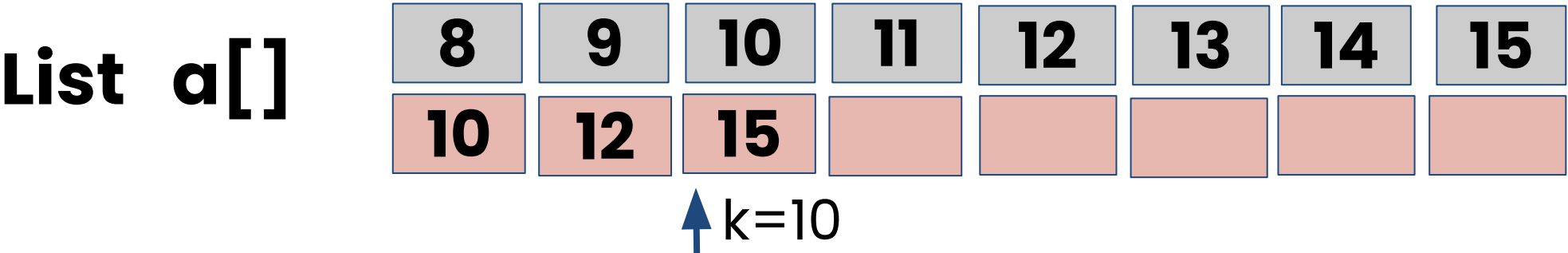
Merging Process:



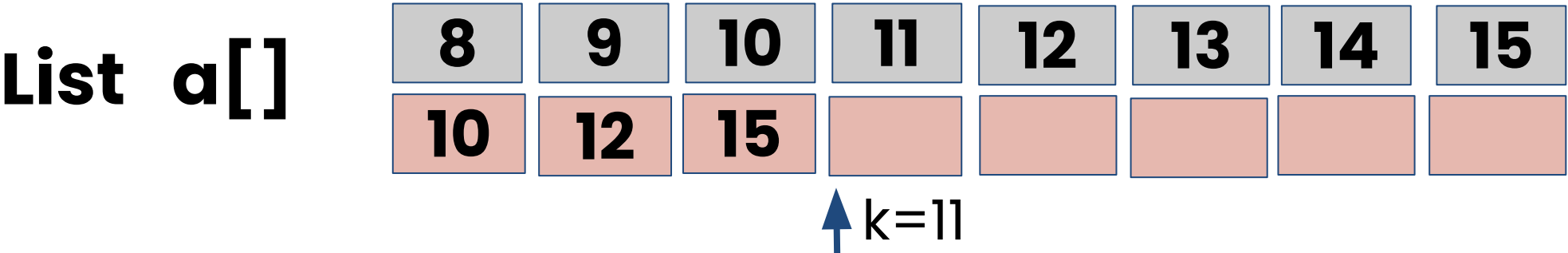
Merging Process:



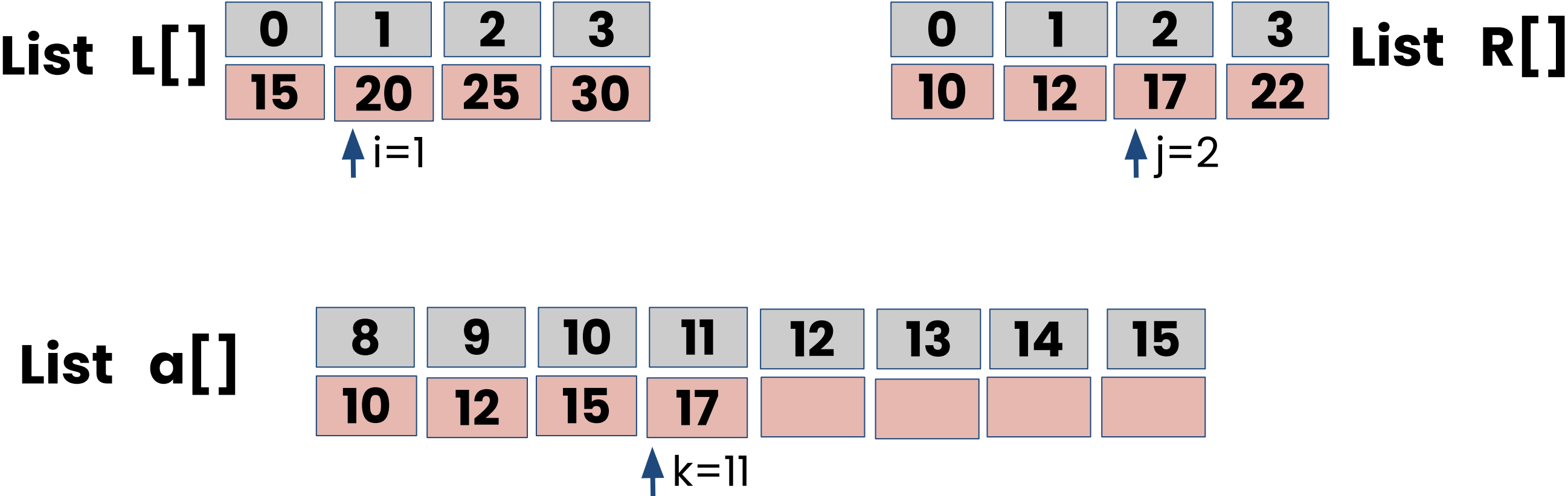
Merging Process:



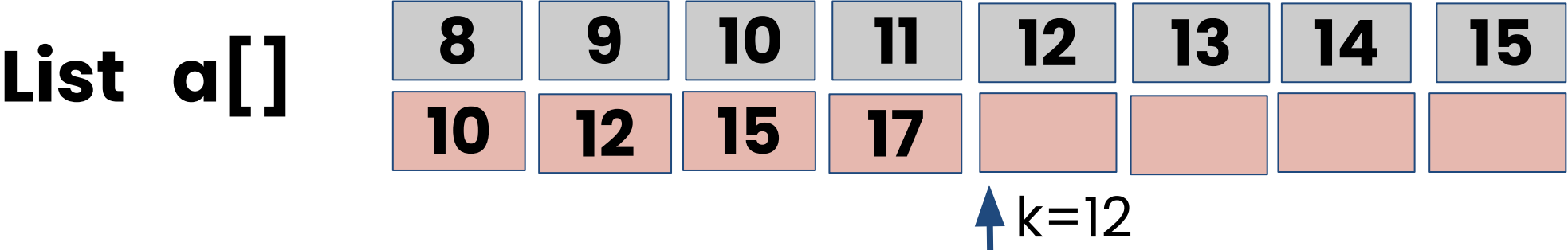
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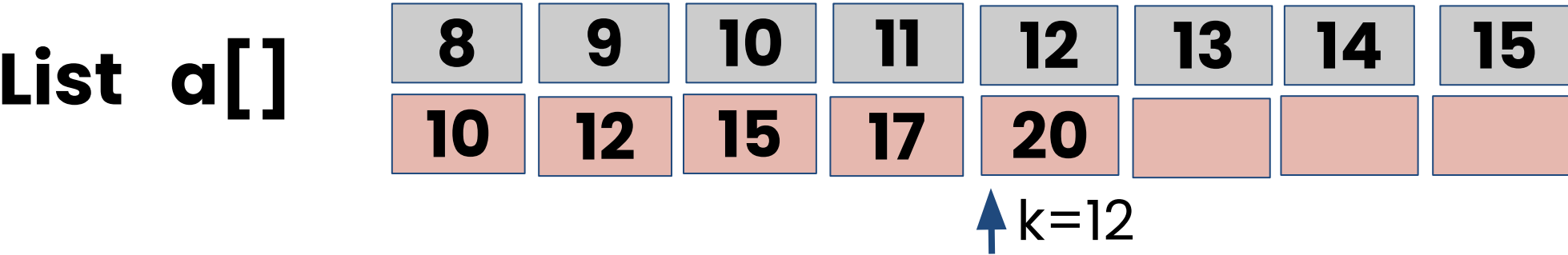
Merging Process:



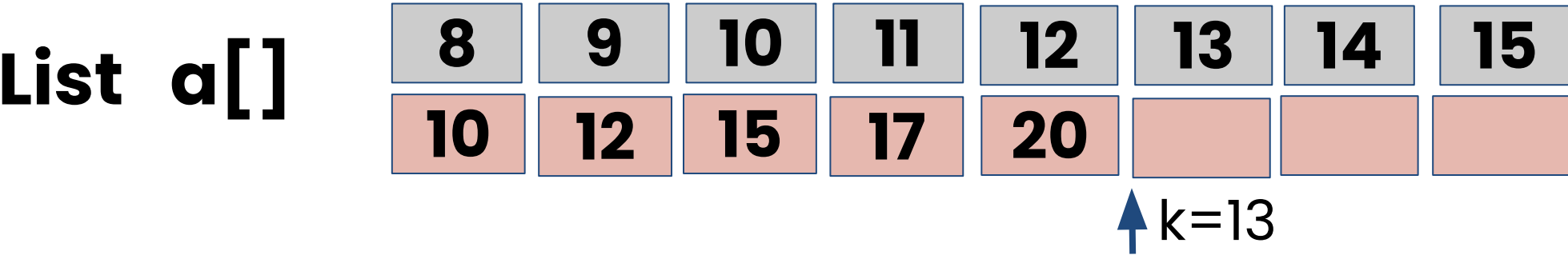
Merging Process:



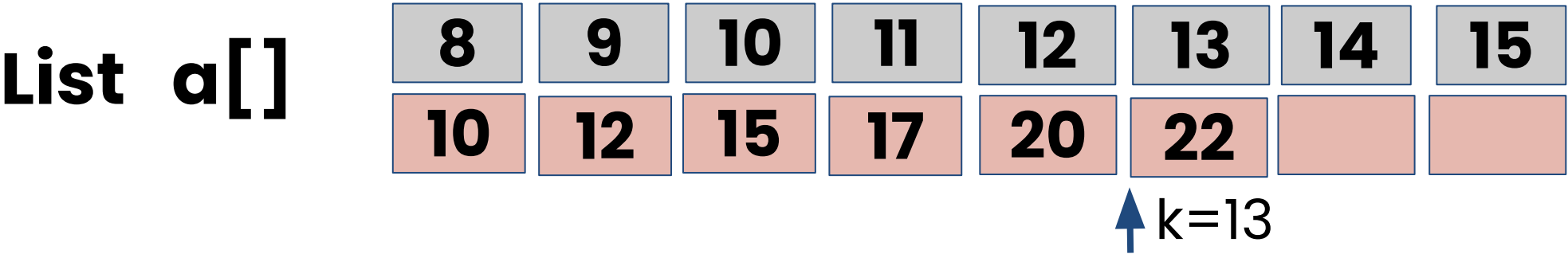
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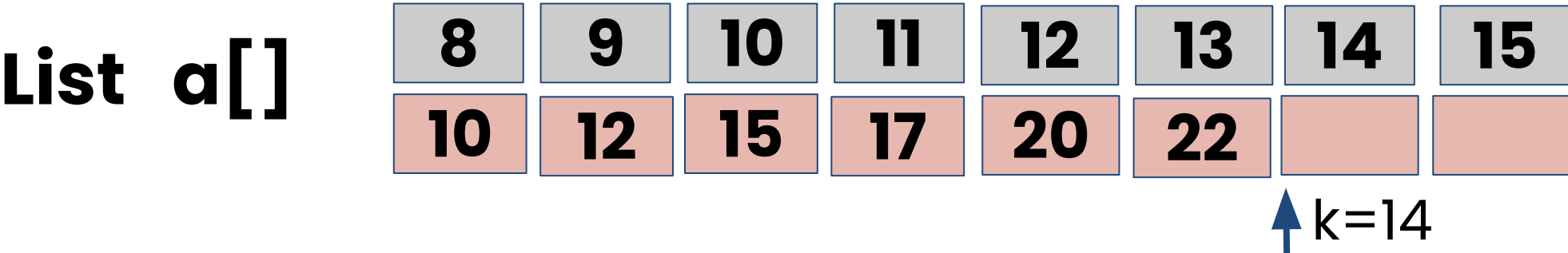
Merging Process:



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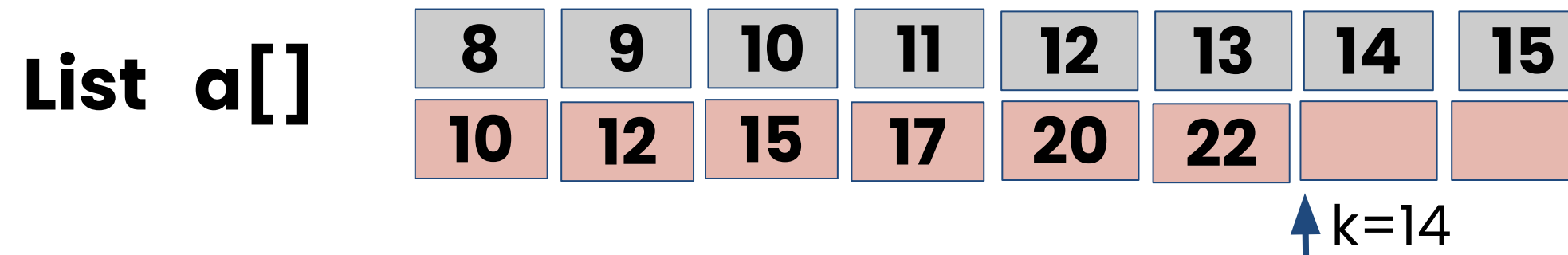
Merging Process:



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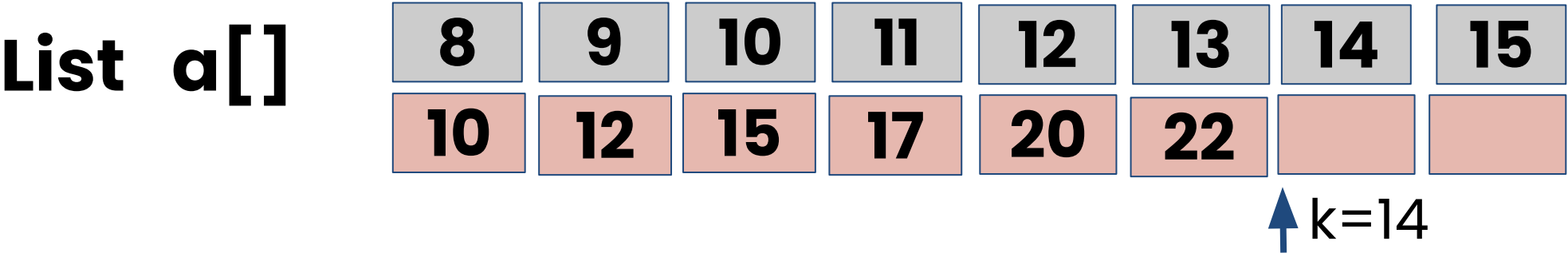
as $j=4$, means List R[] has been completely scanned,
now **copy** the remaining elements of List L[] to List a[]



Merging Process:



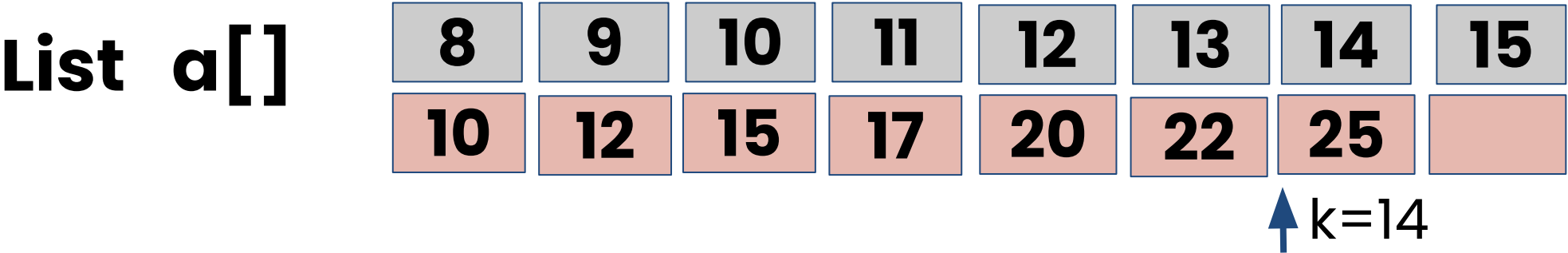
copy L[2] to a[14]



Merging Process:



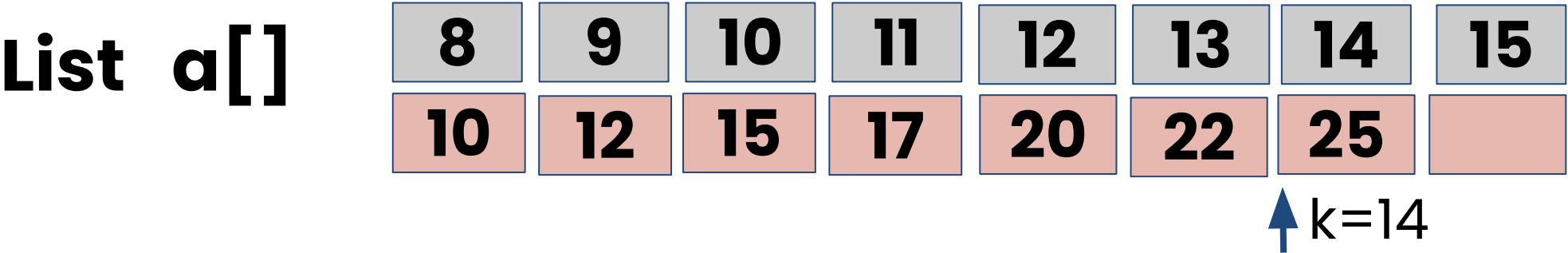
copy L[2] to a[14]



Merging Process:



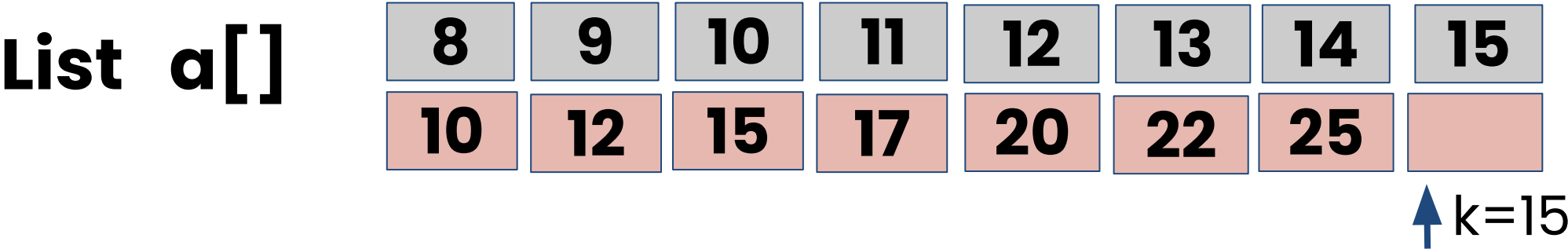
Increment pointer i & k



Merging Process:



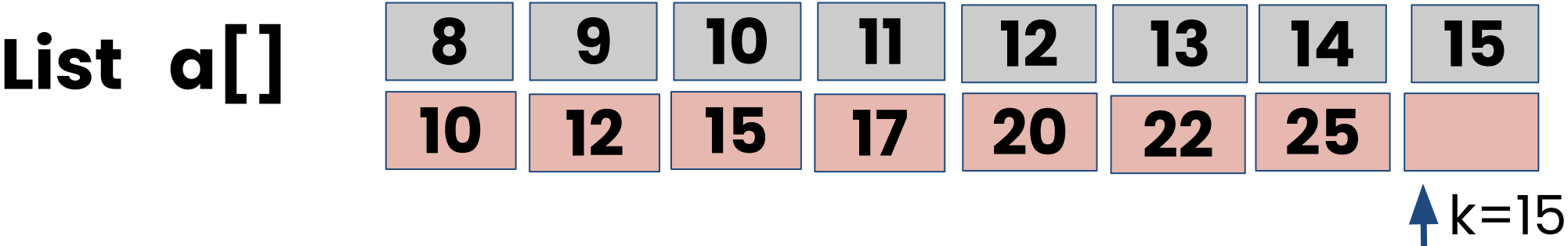
Increment pointer i & k



Merging Process:



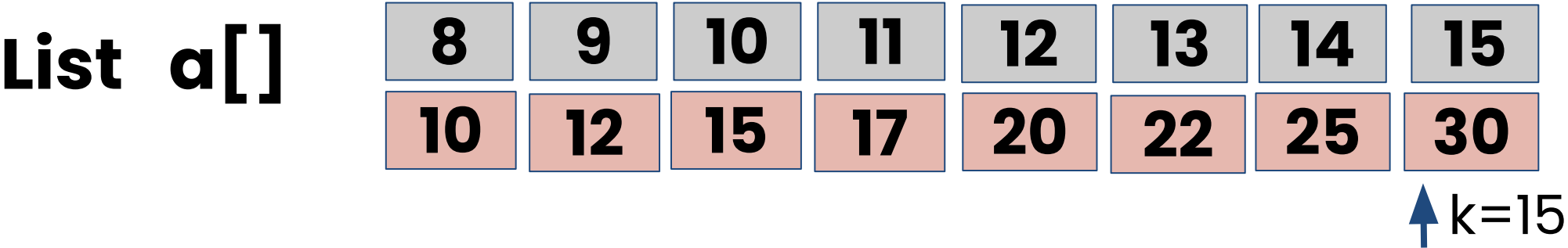
copy L[3] to a[15]



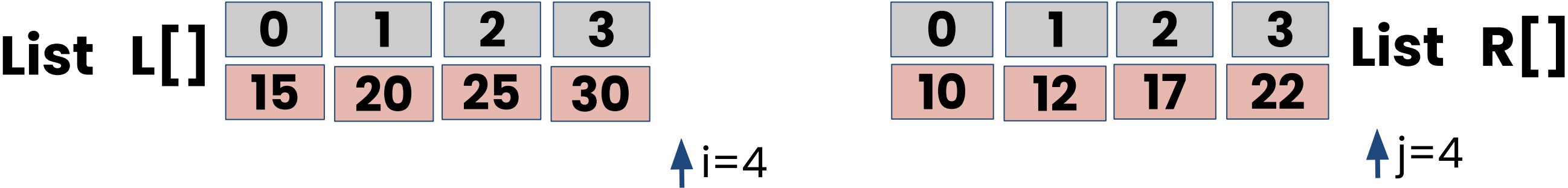
Merging Process:



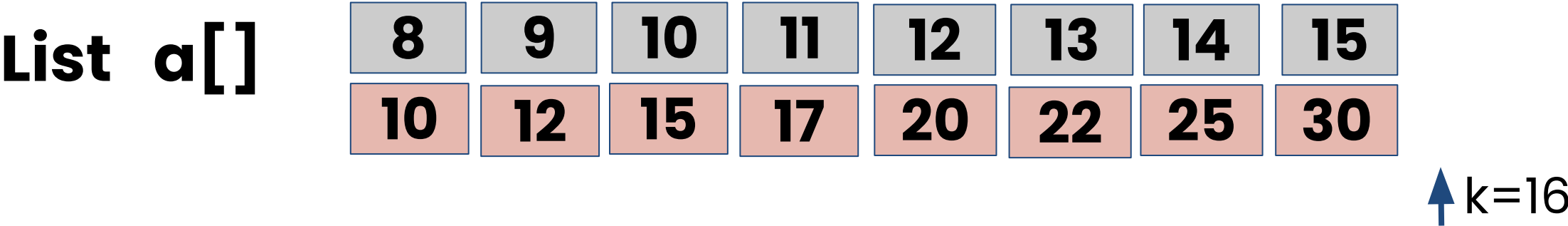
copy L[3] to a[15]



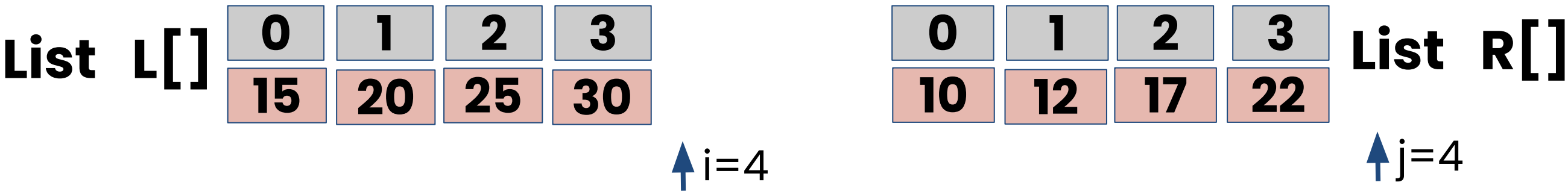
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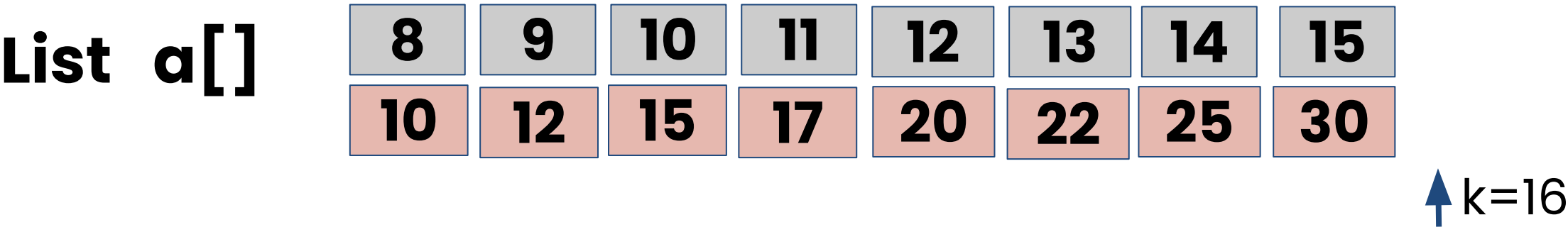
Increment pointer i & k



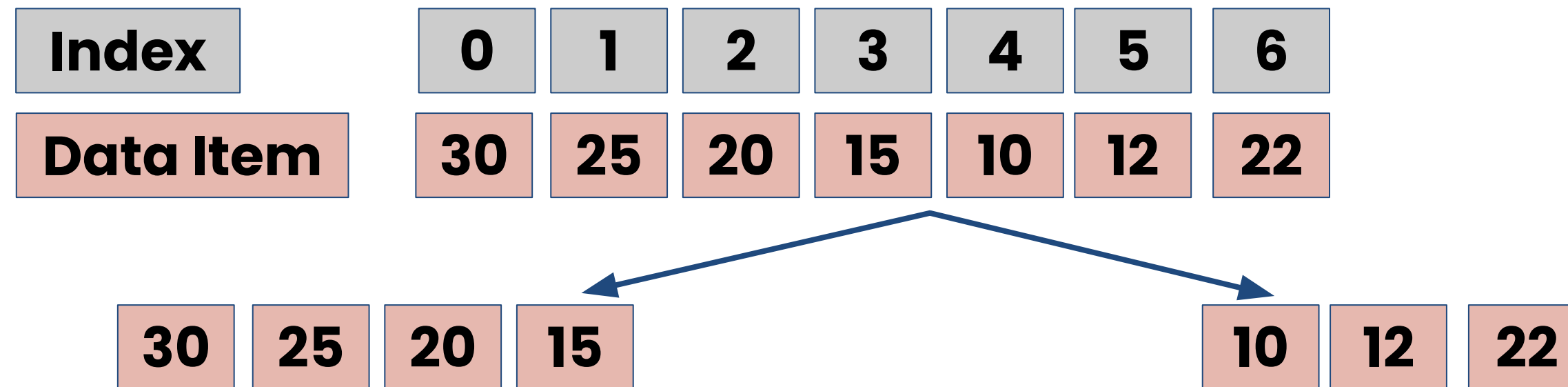
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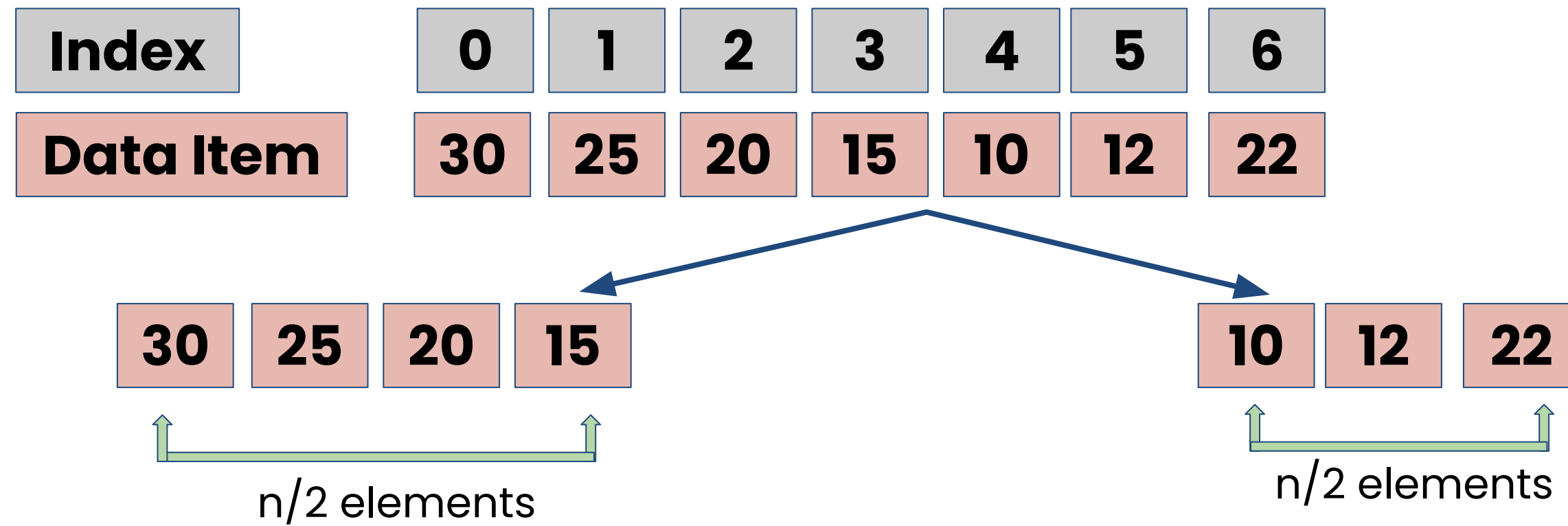
as **i==4** , means List L[] has also been completely scanned .



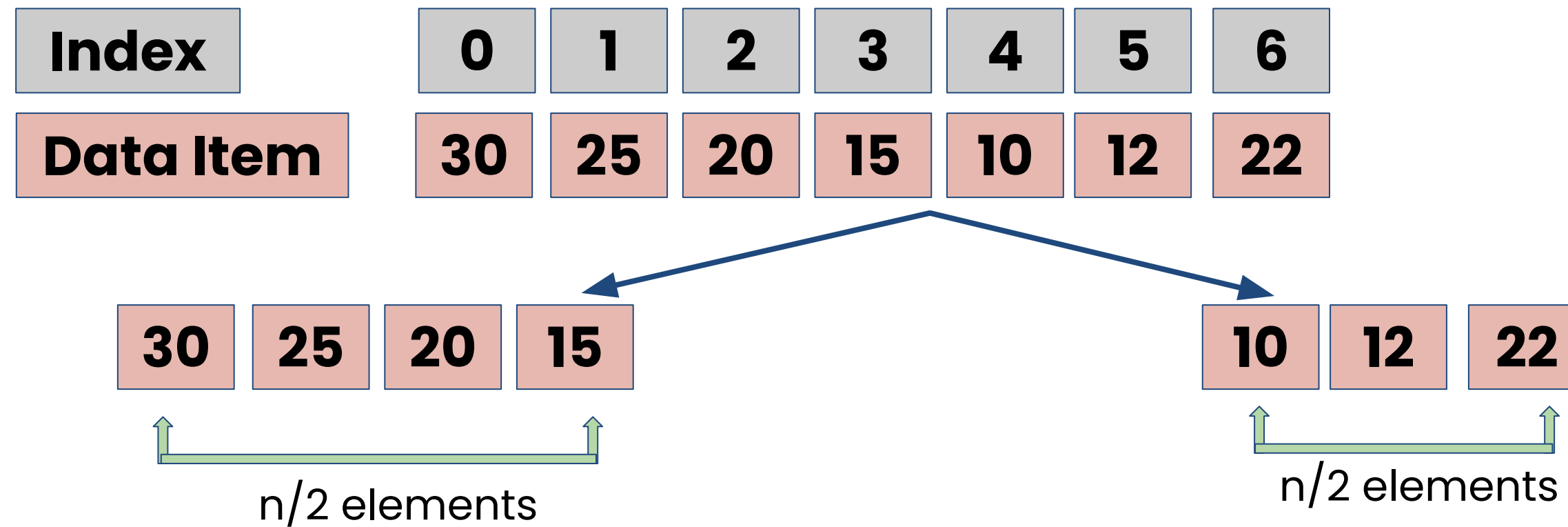
Performance Analysis



Performance Analysis



Performance Analysis



Recurrence Relation

$$T(n) = T(n/2) + T(n/2) + n$$

$$T(n) = 2T(n/2) + n$$

$$T(n) = \mathbf{O(n \log n)}$$

END