





### Quick Recap:

- Structure of Linked List
- Insert Node at Beginning
- Insert Node at End
- Delete First Node
- Delete Last Node

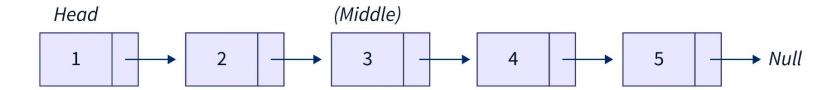




### **Find Middle in Linked List**

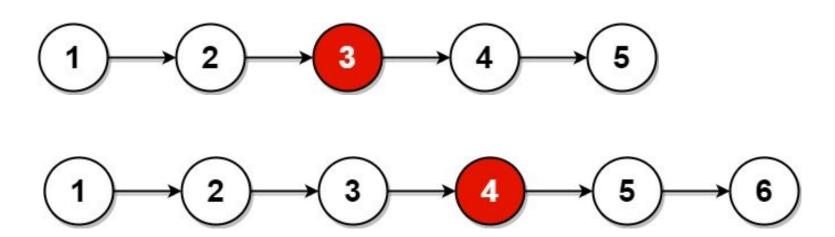


### **Middle of Linked List:**





#### **Middle of Linked List:**



Return second middle, if there are two middle nodes.



#### **Brute Force Method:**





#### **Brute Force Method:**



- 1. **Traverse the List:** Iterate through the linked list to count the total number of nodes (n).
- Find Middle Index: Compute mid = n / 2 (integer division).
- 3. **Traverse Again:** Start from the head and move mid steps to reach the middle node.
- 4. **Return Middle Node:** The node at the mid position is the middle node.



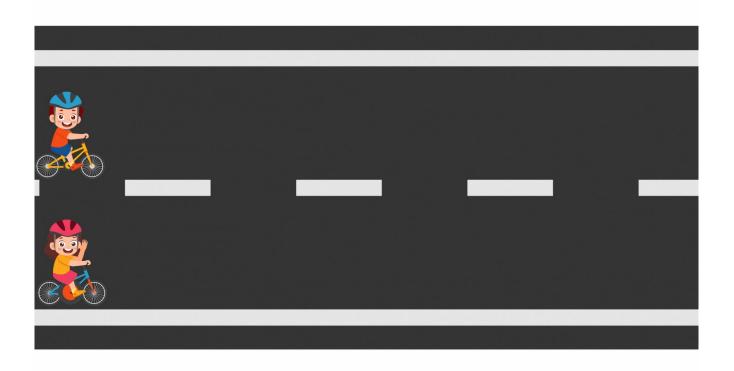
#### **Middle of Linked List:**

- Can we find the middle in a single pass?
- Is it possible to find the middle node without computing the length?



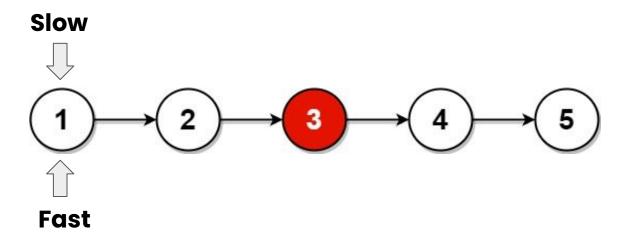


### **Two Pointers Intuition:**



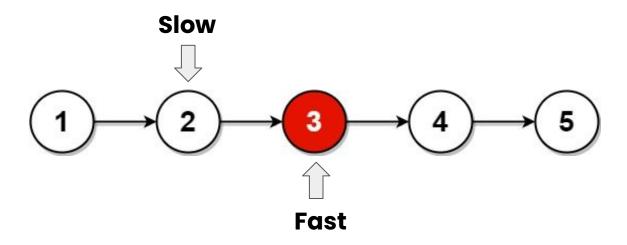


#### For Odd number of Nodes:



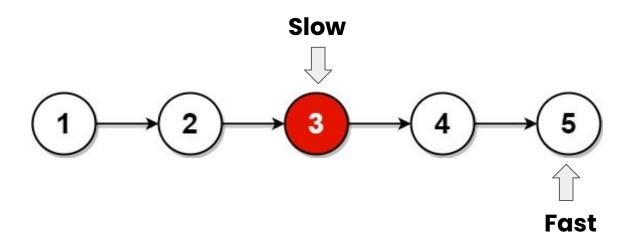


# Two Pointers Method (Slow and Fast ):





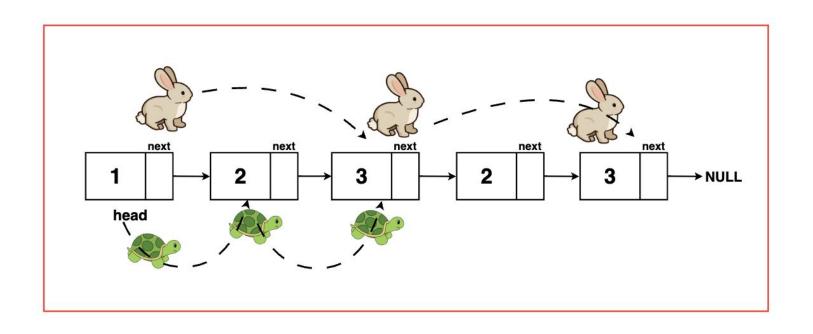
# Two Pointers Method (Slow and Fast):



Stop when Fast Pointer reaches last node.

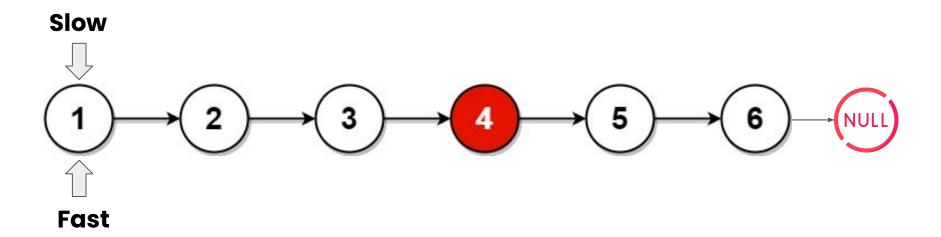


### Called as Hare and Tortoise Method:



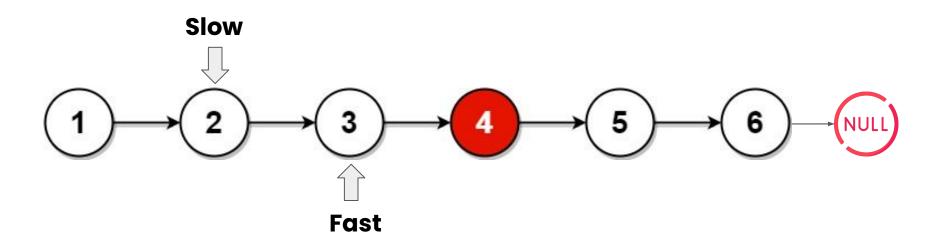


### For Even number of Nodes:



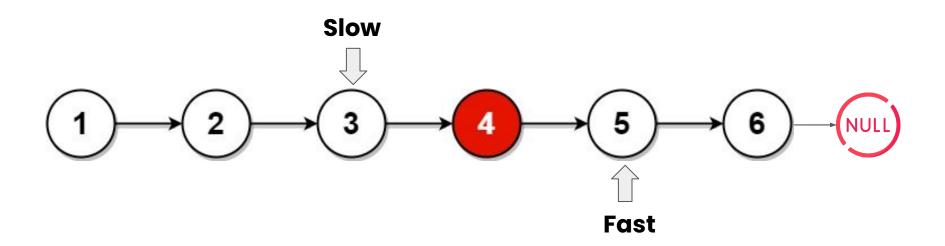


# Two Pointers Method (Slow and Fast ):



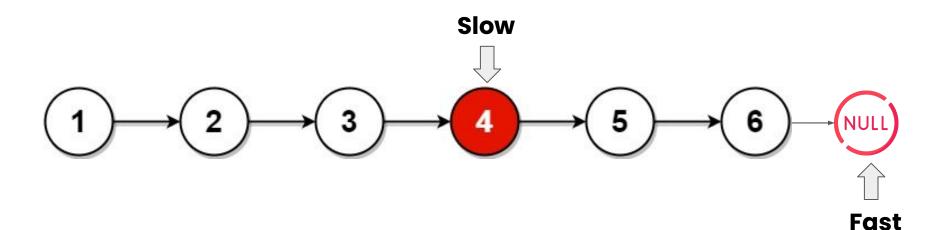


# Two Pointers Method (Slow and Fast ):





# Two Pointers Method (Slow and Fast):



Stop when Fast Pointer reaches None.



## Two Pointers Method (Slow and Fast):



- 1. **Use Two Pointers:** Initialize slow and fast pointers at the head.
- 2. Move Pointers:
  - a. slow moves **one step** at a time.
  - b. fast moves **two steps** at a time.
- 3. Stop When Fast Reaches End:
  - a. If fast or fast.next becomes NULL, slow is at the middle.
- 4. **Return Middle Node:** The slow pointer now points to the middle node.



### Code Implementation:

```
# Optimal Approach - Two Pointer Method
def find_middle(head):
   slow = fast = head
   while fast and fast.next:
       slow = slow.next # Moves one step
       fast = fast.next.next # Moves two steps
   return slow # Middle node
```

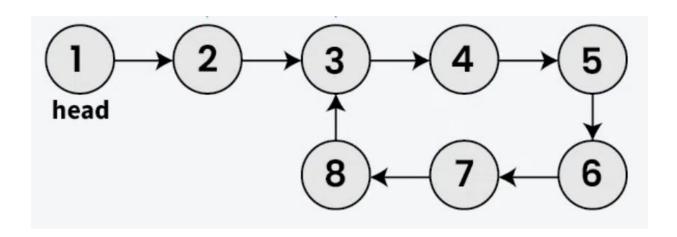


## Floyd's Cycle Detection



#### **Problem Statement:**

Given a singly linked list, check if the linked list has a loop (cycle) or not. A loop means that the last node of the linked list is connected back to a node in the same list.

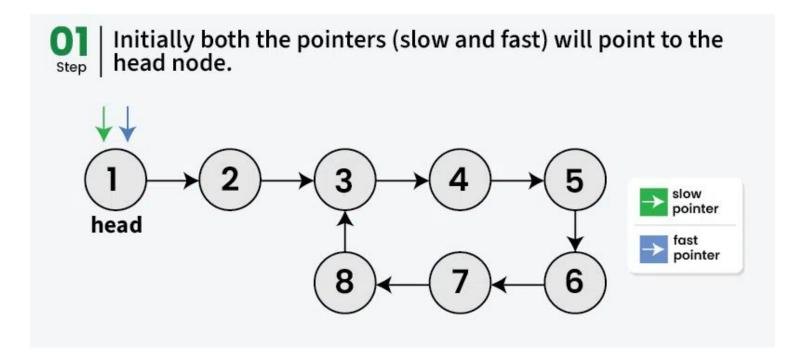




## How will you solve the problem?









slow will move one step ahead and fast move two steps ahead until slow is not equal to fast. slow pointer head fast pointer



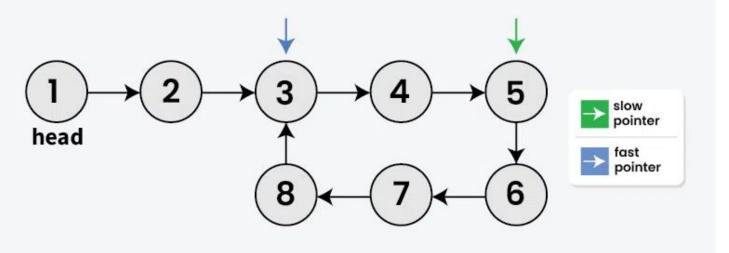
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ostep | slow will move one step ahead and fast move two steps ahead.

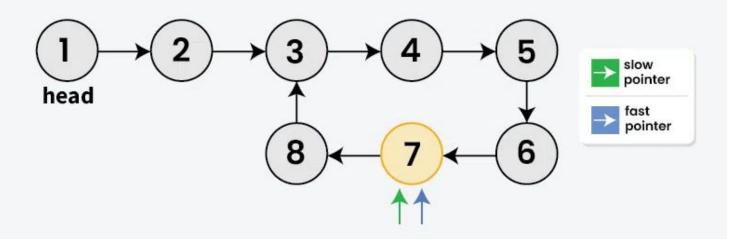




slow will move one step ahead and fast move two steps ahead. Step slow pointer head fast pointer



07 Step As slow and fast pointer points to the same node, there is a cycle in the Linked List.





## Floyd's Cycle Detection Algorithm:

- 1. Use Two Pointers: Initialize slow and fast pointers at the head.
- 2. Move Pointers:
  - o **slow** moves **one step** at a time.
  - fast moves two steps at a time.
- 3. Check for Cycle:
  - o If slow == fast, a cycle exists.
- 4. Stop Condition:
  - If fast or fast.next becomes NULL, no cycle exists.



### Code Implementation:

```
# Floyd's Cycle Detection Algorithm
def has_cycle(head):
   slow = fast = head
   while fast and fast.next:
       slow = slow.next # Moves one step
       fast = fast.next.next # Moves two steps
       if slow == fast: # If they meet, cycle exists
           return True
   return False # No cycle found
```



### **END**