





KTU STUDY MATERIALS | SYLLABUS | LIVE NOTIFICATIONS | SOLVED QUESTION PAPERS

**Website:** www.ktunotes.in

## Module 5

## **Macro Preprocessor**

Macro Instruction Definition and Expansion - One pass Macro processor Algorithm and data structures, Machine Independent Macro Processor Features, Macro processor design options

## Overview



## **Macro Instruction Definition and Expansion**

A macro instruction (abbreviated to macro) is simply a notational convenience for the programmer. A macro represents a commonly used group of statements in the source programming language.

**Expanding a macros** – Replace each macro instruction with the corresponding group of source language statements. Macro instruction allows the programmer to write a shorthand version of a program, and leave the mechanical details to be handled by macro processor.

# A macro processor

- Essentially involve the substitution of one group of characters or lines for another.
- Normally, it performs no analysis of the text it handles.
- It doesn't concern the meaning of the involved statements during macro expansion

The design of a <u>macro processor</u> generally is machine independent.

# Three examples of actual macro processors:

- A macro processor designed for use by assembler language programmers
- Used with a high-level programming language
- General-purpose macro processor, which is not tied to any particular language

#### **Basic Macro Processor Functions**

- Macro Definition
- Macro Invocation
- Macro Expansion

# Macro Definition

- Two new assembler directives are used in macro definition:
  - MACRO: identify the beginning of a macro definition
  - MEND: identify the end of a macro definition
- □ label op operands
  name MACRO parameters
  :
  body
  :
  - MEND
- Parameters: the entries in the operand field identify the parameters of the macro instruction
  - We require each parameter begins with '&'
- Body: the statements that will be generated as the expansion of the macro.
- Prototype for the macro:
  - The macro name and parameters define a pattern or prototype for the macro instructions used by the programmer

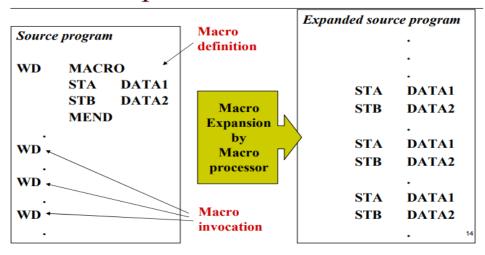
# **Macro Invocation**

- □ A *macro invocation statement* (a *macro call*) gives the **name** of the macro instruction being invoked and the **arguments** in expanding the macro.
- ☐ Macro Invocation vs. Subroutine Call.
  - Statements of the macro body are expanded each time the macro is invoked.
  - Statements of the subroutine appear only one, regardless of how many times the subroutine is called.
  - Macro invocation is more efficient than subroutine call, however, the code size is larger

# Macro Expansion

- □ Each macro invocation statement will be expanded into the statements that form the **body** of the macro.
- Arguments from the macro invocation are substituted for the parameters in the macro prototype.
  - The arguments and parameters are associated with one another according to their positions.
    - The first argument in the macro invocation corresponds to the first parameter in the macro prototype, etc.

# Macro Expansion



# Macro Expansion with Parameters Substitution

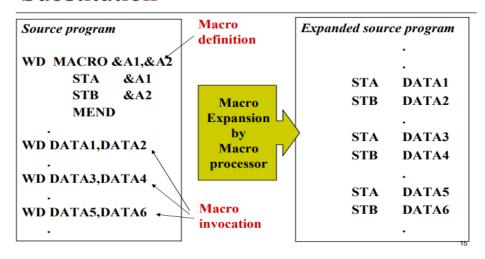
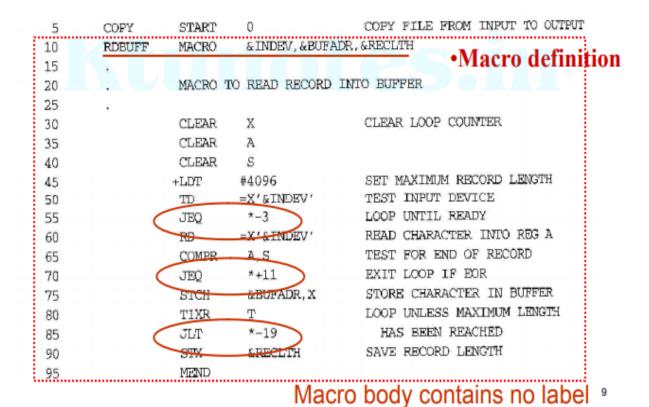


Figure next shows an example of SIC/XE program using macro instructions.



## Macro definition

100	WRBUFF	MACRÓ	&OUTDEV, &BUF	ADR,&RECLTH
105				
110		MACRO I	O WRITE RECORD	FROM BUFFER
115	,			
120		CLEAR	X	CLEAR LOOP COUNTER
125		LDT	&RECLTH	
130		LDCH	&BUFADR,X	GET CHARACTER FROM BUFFER
135		TD	=X'&OUTDEV'	TEST OUTPUT DEVICE
140		JEQ	*-3	LOOP UNTIL READY
145	`	WD	=X'&OUTDEV'	WRITE CHARACTER
150		TIXR	Т	LOOP UNTIL ALL CHARACTERS
155		JLT	*-14	HAVE BEEN WRITTEN
160		MEND		

Macro body contains no label 10

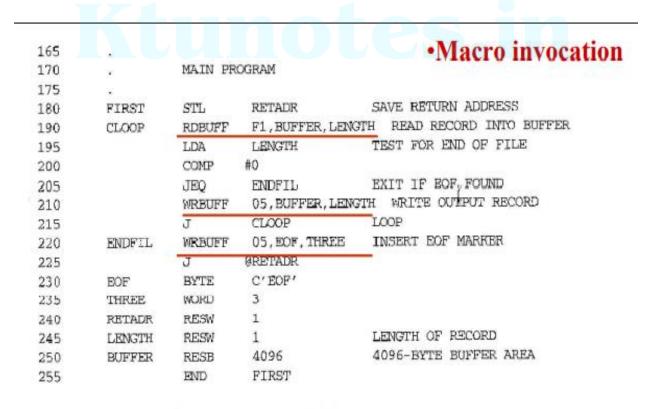


Figure 4.1 Use of macros in a SIC/XE program.

Lin	ie :	Source sta	tement	
5	COPY	START	0	
10	RDBUFF	NACRO		COPY FILE PROM INPUT TO OUTPUT
15	MAROLE	PUNCHO	&INDEV, &BUFADR	, ARGULITH
20		Manne I		and the second s
25	100000	PLACEO 1	TO READ RECORD IN	TO BUFFER
30	***	CLEAR		CHIEF TARE COLUMN
35		CLEAR	x	CLEAR LOOP COUNTER
40		100000000000000000000000000000000000000	A	
45		*LDT	<b>#</b> 4096	
50		TD	TOTAL CONTRACTOR OF THE PARTY O	SET MAXIMUM RECORD LENGTH
55		The state of the s	*X. & INDEA.	TEST INPUT DEVICE
60		JEO RD	*-3	LOOP UNTIL READY
65			-X. FINDEA.	READ CHARACTER INTO REG A
70	T- 5176	COMPR	A,S	TEST FOR END OF RECORD
75		JEQ.	*+11	EXIT LOOP IF BOR
		STCH	&BUFADR, X	STORE CHARACTER IN BUFFER
80		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
85		JLT	*-19	HAS BEEN REACHED
90		STX	&RECLTH	SAVE RECORD LEMOTH
95	-	MEND		
100	WRBUFF	MACRO	&OUTDEV, &BUFADE	R, &RECLTH
105		0.000.000		
110		MACRO T	O WRITE RECORD F	ROM BUFFER
115				
120		CLEAR	X	CLEAR LOOP COUNTER
125		LDT	ARECL/TH	
130		LDCH	&BUFADR, X	GET CHARACTER FROM BUFFER
135		TD	=X.FOOLDEA.	TEST OUTPUT DEVICE
140		JEQ	*-3	LOOP UNTIL READY
145		WD	=X. #OOLDEA.	WRITE CHARACTER
150		TIXR	T	LOOP UNTIL ALL CHARACTERS
155		JLT	*-14	HAVE BEEN WRITTEN
160		MEND		
165				
170		MAIN PR	OGRAM	
175	The same			
180	FIRST	STL	RETADR	SAVE RETURN ADDRESS
190	CLOOP	RDBUFF	F1, BUFFER, LENGT	TH READ RECORD INTO BUFFER
195		LDA	LENGTH	TEST FOR END OF FILE
200		COMP	#0	
205		JBQ	EMDFIL	EXIT IF BOF FOUND
210		WRBUFF	05, BUFFER, LENGT	TH WRITE OUTPUT RECORD
215		J	CLOOP	LOOP
220	ENDFIL	WRBUFF	05, BOF, THREE	INSERT HOF MARKER
225		J	GRETADR	
230	BOF	BYTE	C'EOF'	
235	THREE	WORD	3	
240	RETADR	RESW	1	
245	LENGTH	RESW	1	LENGTH OF RECORD
250	BUFFER	RESB	4096	4096-BYTE BUFFER AREA
255		END	FIRST	
		- 11 11 11 11 11	3.700000	

• This program defines and uses two macro instructions – RDBUFF and WRBUFF. The functions of RDBUFF macro are similar to those of the RDREC subroutine which we have studied earlier. Likewise WRBUFF macro similar to WRREC subroutine. The definitions of these macro instructions appear in the source program following the start statement. Two new assembler directives are used in macro definitions:

#### 1. MACRO

#### 2. MEND

The first MACRO statement (line 10) identifies the beginning of a macro definition. The symbol in the label field (RDBUFF) is the name of the macro, and the entries in the operand field identify the parameters of the macro instructions. In our macro language, each parameter begins with the character &, which facilitates the substitution of parameters during macro expansion. The macro name and parameters define a pattern or prototype for the macro instructions used by the programmer. Following the MACRO directive, are the statements that make up the body of the macro definition (line 15 through 90). These are the statements that will be generated as the expansion of the macro. The MEND assembler directive (line 95) marks the end of the macro definition. The definition of the WRBUFF macro(line 100 through 160) follows a similar pattern. The main program itself begins at line 180. The statement on line 190 is a macro invocation statement that gives the name of the macro instruction being invoked and the argument5s to be used in expanding the macro. A macro invocation statement is often called referred to as macro call. The macro invocation and subroutine call are different. The above figure gives the input program to a macro processor. The output generated is given in next figure.

				•Macro expansion
5	COPY	START	0	COPY FILE FROM INPUT TO OUTPUT
180	FIRST	STL	RETADR	SAVE RETURN ADDRESS
190	.CLOOP	RDBUFF	FI, BUFFER, LENGTH	READ RECORD INTO BUFFER
<b>19</b> 0a	CLOOP	CLEAR	X	CLEAR LOOP COUNTER
190b		CLEAR	A	
190c		CLEAR	S	
190d		+LDT	#4096	SET MAXIMUM RECORD LENGTH
190e		TD	=X'F1'	TEST INPUT DEVICE
190 <b>£</b>		JEQ	*-3	LOOP UNTIL READY
190g		RD	=X'F1'	READ CHARACTER INTO REG A
190h		COMPR	A,S	TEST FOR END OF RECORD
190i		JEQ	*+11	EXIT LOOP IF EOR
190j		STCH	BUFFER, X	STORE CHARACTER IN BUFFER
190k		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
1901		JLT	*-19	HAS BEEN REACHED
190m		STX	LENGTH	SAVE RECORD LENGTH

# Macro expansion

195 200	LDA COMP	LENGTH #0	TEST FOR END OF FILE
205	JEQ	ENDFIL	EXIT IF EOF FOUND
210 .	WRBUFF	05,BUFFER,LENGTH	WRITE OUTPUT RECORD
210a	CLEAR	X	CLEAR LOOP COUNTER
210b	L'DJ,	LENGTH	
210c	LDCH	BUFFER, X	GET CHARACTER FROM BUFFER
210d	מנו	=X'05'	TEST OUTPUT DEVICE
210e	JEQ	*-3	LOOP UNTIL READY
210f	MD	=X'05'	WRITE CHARACTER
210g	TIXR	T	LOOP UNTIL ALL CHARACTERS
210h	JLT	*-14	HAVE BEEN WRITTEN

The macro instruction definitions have been deleted since they are no longer needed after the macros are expanded. Each macro invocation statement has been expanded into the statements that form the body of the macro, with the arguments from the macro invocation substituted for the parameters in the macro prototype. The arguments and parameters are associated with one another according to their positions. The first argument in the macro invocation corresponds to the first parameter in the macro prototype, and so on. In expanding the macro invocation on line 190, for example, the argument F1 is substituted for the parameter &INDEV wherever it occurs in the body of the macro. Similarly, BUFFER is substituted for &BUFADR and LENGTH is substituted for &RECLTH. Lines 190a through 190m show the complete expansion of the macro invocation on line 190. The comment lines within the macro body have been deleted, but comments on individual statements have been retained. Note that the macro invocation statement itself has been included as a comment line. This serves as a documentation of the statement written by the programmer. The label on the macro invocation statement (CLOOP) has been retained as a label on the first statement generated in the macro expansion. This allows the programmer to use a macro instruction in exactly the same way as an assembler language mnemonic. The macro invocations on line 210 and 220 are expanded in the same way. Note that the two invocations of WRBUFF specify different arguments, so they produce different expansions. After macro processing, the expanded file obtained (fig 4.2) servers as the input to assembler. The macro invocation statements will be treated as comments, and the statements

generated from the macro expansions will be assembled exactly as though they had been written directly by the programmer.

ine	Sour	ce statem	ent	
5	COPY	START	0	COPY FILE FROM INPUT TO OUTPUT
80	FIRST	STL	RETADR	SAVE RETURN ADDRESS
90	.CLOOP	RDBUFF	F1. BUFFER, LENGTH	READ RECORD INTO BUPPER
90a	CLOOP	CLEAR	X	CLEAR LOOP COUNTER
90Ъ	200000	CLEAR	A	
90c		CLEAR	g	
90d		+LDT	#4096	SET MAXIMUM RECORD LENGTH
90e		TD	=X'F1'	TEST INPUT DEVICE
90f		JEQ	•-3	LOOP UNTIL READY
90g		RD	=X.E1.	READ CHARACTER INTO REG A
90h		COMPR	A,S	TEST FOR END OF RECORD
901		JEO	*+11	EXIT LOOP IF BOR
901		STCH	BUFFER, X	STORE CHARACTER IN BUFFER
90k		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
901		JLT	*-19	HAS BEEN REACHED
90m		STX	LENGTH	SAVE RECORD LENGTH
95		LDA	LENGTH	TEST FOR END OF FILE
100		COMP	+0	
05		JEQ	ENDFIL	EXIT IF BOF FOUND
110		WRECFF	05, BUFFER, LENGTH	WRITE OUTPUT RECORD
10a		CLEAR	x	CLEAR LOOP COUNTER
10b		LDT	LENGTH	
10c		LDCH	BUFFER, X	GET CHARACTER FROM BUFFER
10d		TD	=X'05'	TEST OUTPUT DEVICE
110e		JEQ	*-3	LOOP UNTIL READY
101		MD	=X'05'	WRITE CHARACTER
210g		TIXR	T	LOOP UNTIL ALL CHARACTERS
210h		JUT	*-14	HAVE BEEN WRITTEN
215		J	CLOOP	LOOP
220	.morit	WEBUFF	05, EOF, THREE	INSERT EOF MARKER
220a	ENDFIL	CLEAR	X	CLEAR LOOP COUNTER
220b		LDT	THREE	
220c		LDCH	BOF, X	GET CHARACTER FROM BUFFER
2204		TD	=X'05'	TEST OUTPUT DEVICE
120e		JEQ	*-3	LOOP UNTIL READY
20f		WD	=X'05'	WRITE CHARACTER
120g		TIXR	7	LOOP UNTIL ALL CHARACTERS
120h		JLT	*-14	HAVE BEEN WRITTEN
225		J	GRETADR	
230	EOF	BYTE	C'ROF'	
235	THREE	WORD	3	
240	RETADR	RESW	1	
245	LENGTH	RESW	1	LENGTH OF RECORD
250	BUFFER	RESB	4096	4096-BYTE BUFFER AREA
255		END	FIRST	
			rom Fig. 4.1 with mac	

Comparing the fig 2.5 and fig 4.2, there is a difference between subroutine call and macro invocation can be clearly understood. In fig 4.2, the statements from the body of the macro WRBUFF are generated twice: line 210a through 210h and lines 220a through 220h. In program 2.5, the corresponding statements appear only once in the subroutine WRREC (line 210 through 240). In general, the statements that form the expansion of a macro are generated (and assembled) each time the macro is invoked. Statements in a subroutine appear only once,

regardless of how many times the subroutine is called. **Note: the body of the macro contains no labels**.

In fig 4.1 line 140 contains the statement, JEQ \*-3 and line 155 contains JLT\*-14. The corresponding statements in the WRREC subroutine are JEQ WLOOP and JLT WLOOP where WLOOP is a label on the TD instruction that tests the output device. If such a label appeared on line 135 of the macro body, it would be generated twice —on lines 210d and 220d of fig 4.2. This would result in an error (a duplicate label definition) when the program is assembled. To avoid duplication of symbols, we have eliminated labels from the body of our macro definitions. The use of statements like JLT \*-14 is generally considered to be a poor programming practice. It is somewhat less objectionable within a macro definition; however it is still an in convenient and error prone method. We will study the solutions to avoid this problem in later class.

## **Macro Processor Algorithm and Data Structures**

- Two-pass macro processor
- One-pass macro processor

### Two-pass macro processor

- Pass1: process all macro definitions
- Pass2: expand all macro invocation statements

#### Problem

- Does not allow nested macro definitions
- Nested macro definitions The body of a macro contains definitions of other macros
- Because all macros would have to be defined during the first pass before any macro invocations were expanded.

1	MACROS	MACRO	(Defines SIC standard version macros
2	RDBUFF	MACRO	&INDEV,&BUFADR,&RECLTH
		* 11	
		5	(SIC standard version)
3		MEND	(End of RDBUFF)
4	WRBUFF	MACRO	&OUTDEV, &BUFADR, &RECLTH
			(SIC standard version)
5		MEND	(End of WRBUFF)
6		MEND	(End of MACROS)
			(a)
1	MACROX	MACRO	{Defines SIC/XE macros}
2	RDBUFF	MACRO	&INDEV, &BUFADR, &RECLTH
			{SIC/XE version}
		*	
3		MEND	(End of RDBUFF)
4	WRBUFF	MACRO	&OUTDEV, &BUFADR, &RECLTH
		*	(SIC/XE version)
		*	(SIC/AE VELSION)
5		MEND	(End of WRBUFF)
6		MEND	(End of MACROX)
			(b)
	Figure	13 Evample	of the definition of macros within a macro body.

A program that is to be run on SIC system could invoke MACROS whereas a program to be run on SIC/XE can invoke MACROX. However, defining MACROS or MACROX does not define RDBUFF and WRBUFF. These definitions are processed only when an invocation of MACROS or MACROX is expanded.

#### Solution

One-pass macro processor

#### One-pass macro processor

- One-pass processor can alternate between macro definition and macro expansion
- In One-pass macro processor, every macro must be defined before it is called
- Restriction

- The definition of a macro must appear in the source program before any statements that invoke that macro
- This restriction does not create any real inconvenience.

#### **Data Structures**

There are three data structures involved in one –pass macro processor.

- 1. DEFTAB
- 2. NAMTAB
- 3. ARGTAB

#### **DEFTAB**

- A *definition table* used to *store macro definition* including
  - macro prototype
  - macro body
- Comment lines are omitted.
- Positional notation has been used for the parameters for efficiency in substituting arguments.
  - E.g. the first parameter &INDEV has been converted to ?1 (indicating the first parameter in the prototype)

#### **NAMTAB**

- A name table used to store the macro names
- Serves as an index to DEFTAB
  - Pointers to the beginning and the end of the macro definition

#### **ARGTAB**

- An argument table is used to store the arguments used in the expansion of macro invocation
- When a macro invocation statement is recognized, the arguments are stored in ARGTAB
  according to their position in the argument list

 As the macro is expanded, arguments are substituted for the corresponding parameters in the macro body

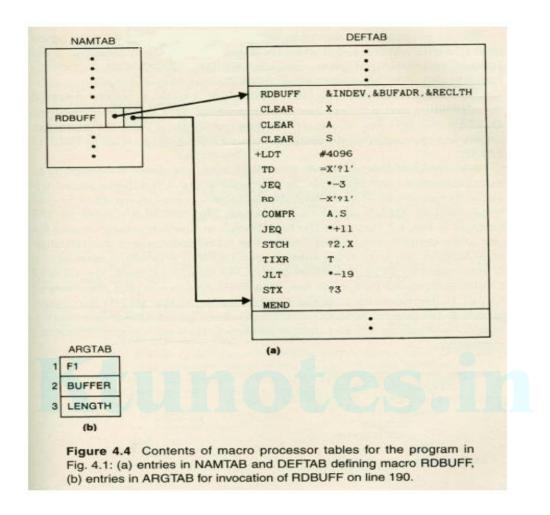


Fig. 4.4 shows positions of the contents of these tables during the processing of the program. Fig. 4.4(a) shows the definition of RDBUFF stored in DEFTAB, with an entry in NAMTAB identifying the beginning and end of the definition. Note the positional notation that has been used for parameters: The parameter &INDEV has been converted to?1(indicating the first parameter in the prototype) &BUFADR has been converted to ?2, and so on. Fig. 4.4(b) shows the ARGTAB as it would appear during expansion of the RDBUFF statement on line 190. For this invocation, the first argument is F1, the second is BUFFER etc. This scheme makes substitution of macro arguments much more efficient. When the ?n notation is recognized in a line from DEFTAB, a simple indexing operation supplies the proper argument from the ARGTAB.

## Algorithm

```
begin {macro processor}
   EXPANDING := FALSE
   while OPCODE ≠ 'END' do
       begin
          GETLINE
          PROCESSLINE
       end {while}
end {macro processor}
procedure PROCESSLINE
   begin
       search NAMTAB for OPCODE
       if found then
          EXPAND
       else if OPCODE = 'MACRO' then
       else write source line to expanded file
   end {PROCESSLINE}
```

```
procedure DEFINE
   begin
       enter macro name into NAMTAB
       enter macro prototype into DEFTAB
       LEVEL := 1
       while LEVEL > 0 do
          begin
              GETLINE
              if this is not a comment line then
                 begin
                     substitute positional notation for parameters
                     enter line into DEFTAB
                     if OPCODE = 'MACRO' then
                        LEVEL := LEVEL + 1
                     else if OPCODE = 'MEND' then
                        LEVEL := LEVEL - 1
                 end {if not comment}
          end {while}
       store in NAMTAB pointers to beginning and end of definition
   end {DEFINE}
```

```
procedure EXPAND
   begin
       EXPANDING := TRUE
       get first line of macro definition (prototype) from DEFTAB
       set up arguments from macro invocation in ARGTAB
       write macro invocation to expanded file as a comment
       while not end of macro definition do
          begin
              GETLINE
              PROCESSLINE
          end {while}
       EXPANDING := FALSE
    end {EXPAND}
procedure GETLINE
    begin
       if EXPANDING then
              get next line of macro definition from DEFTAB
              substitute arguments from ARGTAB for positional notation
           read next line from input file
    end {GETLINE}
```

## The procedure PROCESSING

## The procedure DEFINE

 Called when the beginning of a macro definition is recognized, makes the appropriate entries in DEFTAB and NAMTAB.

## The procedure EXPAND

 Called to set up the argument values in ARGTAB and expand a macro invocation statement.

## The procedure GETLINE

 Called at several points in the algorithm, gets the next line to be processed.

## EXPANDING is set to TRUE or FALSE.

The procedure DEFINE, which is called when the beginning of a macro definition is recognized, makes the appropriate entries in DEFTAB and NAMTAB. EXPAND is called to set up the argument values in ARGTAB and expand a macro invocation statement. The procedure GETLINE, which is called at several points in the algorithm, gets the next line to be processed

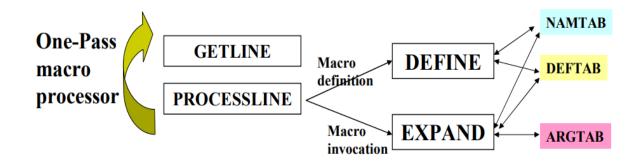
This line may come from DEFTAB (the next line of a macro being expanded), or from the input file, depending upon whether the Boolean variable EXPANDING is set to TRUE or FALSE.

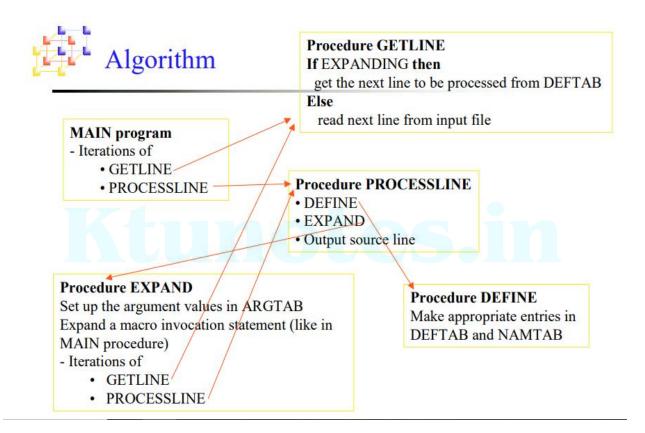
#### Handling nested macro definition

## In DEFINE procedure

- When a macro definition is being entered into DEFTAB, the normal approach is to continue until an MEND directive is reached.
- This would not work for nested macro definition because the first MEND encountered in the inner macro will terminate the whole macro definition process.
- To solve this problem, a counter LEVEL is used to keep track of the level of macro definitions.
  - Increase LEVEL by 1 each time a MACRO directive is read.
  - Decrease LEVEL by 1 each time a MEND directive is read.
  - A MEND terminates the whole macro definition process when LEVEL reaches 0.
  - This process is very much like matching left and right parentheses when scanning an arithmetic expression.

	1 2	MACROS RDBUFF	MACOR MACRO	{Defines SIC standard version macros} &INDEV,&BUFADR,&RECLTH
				{SIC standard version}
	3		MEND	{End of RDBUFF}
	4	WRBUFF	MACRO	&OUTDEV,&BUFADR,&RECLTH
	(0.5)		•	
			* 0.50	{SIC standard version}
	5		MEND	{End of WRBUFF}
			•	
	6		MEND	{End of MACROS}
1				





Most macro processors allow the definitions of commonly used macro instructions to appear in a standard library, rather than in the source program. This makes the use of such macros much more convenient. Definitions are retrieved from this library as they are needed during macro processing.

## **Machine Independent Macro-Processor Features**

- 1. Concatenation of Macro Parameters
- 2. Generation of Unique Labels

- 3. Conditional Macro Expansion
- 4. Keyword Macro Parameters

#### **Concatenation of Macro Parameters**

Most macro processors allow parameters to be concatenated with other character strings. Suppose a program contains one series of variables named by the symbols XA1, XA2, XA3 .... Another series named XB1, XB2,XB3,.. Etc. If similar processing is to be performed on each series of variables, the programmer might want to incorporate this processing into a macro instruction. The parameter to such a macro instruction could specify the series of variables to be operated on(A,B,etc,.). The macro processor would use this parameter to construct the symbols required in the macro expansion(XA1, XB1). Suppose that the parameter to such a macro instruction is named &ID.

The body of the macro definition might contain a statement like

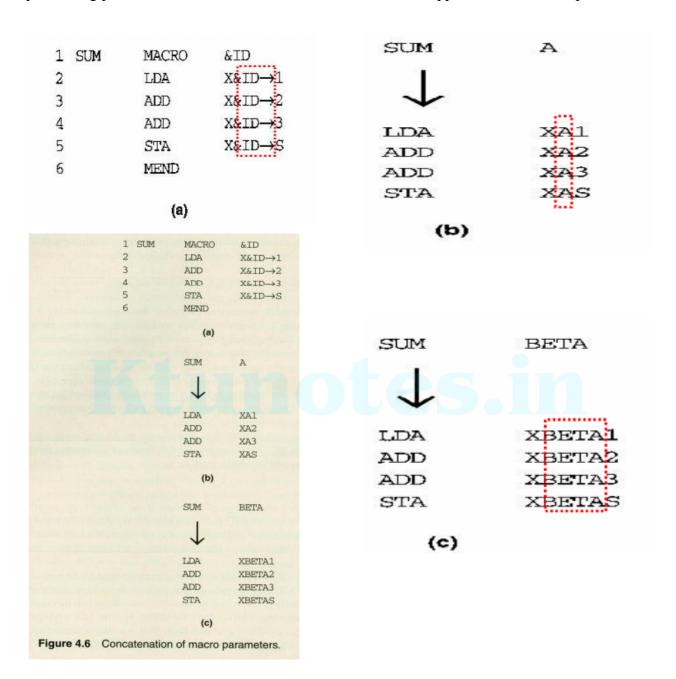
in which the parameter &ID is concatenated after the character string X and before the character string 1. There is a problem with such a statement. The beginning of the macro processor is identified by the starting symbol &. However the end of the parameter is not marked. Thus the operand in the foregoing statement could equally well represent the character string X followed by the parameter &ID1. In the particular case, the macro processor could potentially deduce the meaning that was intended. However, if the macro definition contained both &ID and &ID1 as parameters, it would be ambiguous. Most macro processor deal with this problem by providing a special **concatenation operator.** 

In SIC macro language, this operator is the character  $\rightarrow$ .

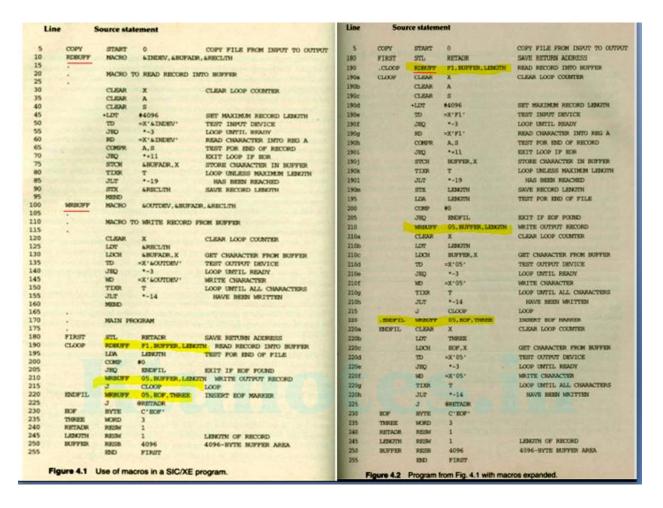
Then the previous statement would be written as:

So that end of the parameter is clearly identified.

The macro processor deletes all occurrences of the concatenation operator immediately after performing parameter substitution, so the character  $\rightarrow$  will not appear in the macro expansion.



 In fig.(a) shows a macro definition that uses the concatenation operator as previously described  In fig (b) and(c) shows macro invocation statements and the corresponding macro expansions

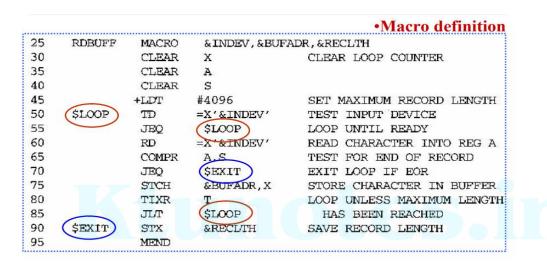


- □ Labels in the macro body may have *duplicate labels* problem
  - If the macro is invocated multiple times.
  - Use of relative addressing is very inconvenient, error-prone, and difficult to read.
    - Example
      - JEQ \*-3
      - Inconvenient, error-prone, difficult to read

#### **Generation of Unique Labels**

The macro body do not contain labels. This leads to the use of relative addressing at the source statement level. Consider for example, the definition of WRBUFF in fig 4.1. If a label were

placed on the TD instruction on line 135, this label would be defined twice-once for each invocation of WRBUFF. This duplicate definition would prevent correct assembly of the resulting expanded program. Because it was not possible to place a label on line 135 of this macro definition, the Jump instruction on line 140 and line 155 were written using the relative operands \*-3 and \*-14. This sort of relative addressing in a source statement may be acceptable for short jumps such as JEQ \*-3. However, for longer jumps spanning several instructions, such notation is very in convenient, error prone and difficult to read. Many macro processors avoid these problems by allowing the creation of special types of labels within macro instruction.



RDBUFF	F1, BU	IFFER, LENG	TH •Macro expansion
30 35 <b>4</b> 0	CLEAR CLEAR CLEAR	X A S	CLEAR LOOP COUNTER
45 50 \$AALOO 55 60 65	+LDT P TD JEQ RD COMPR	#4096 =X'F1' \$AALOOP =X'F1' A.S	SET MAXIMUM RECORD LENGTH TEST INPUT DEVICE LOOP UNTIL READY READ CHARACTER INTO REG A TEST FOR END OF RECORD
70 75 80 85 90 \$AAEXI	JEQ STCH TIXR JLT T STX	\$AAEXIT BUFFER,X T \$AALOOP LENGTH	EXIT LOOP IF EOR STORE CHARACTER IN BUFFER LOOP UNLESS MAXIMUM LENGTH HAS BEEN REACHED SAVE RECORD LENGTH
		(b)	

The figure illustrates one technique for generating unique labels within a macro expansion. A definition of the RDBUFF macro is in fig (a). Labels used within the macro body begin with the

special character \$. Fig(b) shows a macro invocation statement and the resulting macro expansion. Each symbol beginning with \$ has been modified by replacing \$ with \$AA. More generally, the character \$ will be replaced by \$xx, where xx is a two character alphanumeric counter of the number of macro instructions expanded. For the first macro expansion in a program, xx will have the value AA.

```
    $LOOP TD =X'&INDEV'
    1st call:

            $AALOOP TD =X'F1'

    2nd call:

            $ABLOOP TD =X'F1'
```

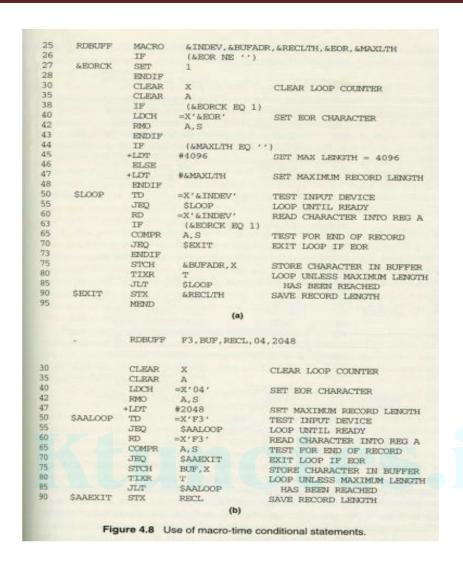
For the succeeding macro expansion, xx will be set to AB, AC etc. *If only alphabetic and numeric characters are allowed in xx, such a two-character counter provides for as many as 1296 macro expansion in a single program.* This results in the generation of unique labels for each expansion of a macro instruction.

#### **Conditional Macro Expansion**

In all our previous examples of macro instructions, each invocation of a particular macro was expanded into the same sequence of statements. These statements could be varied by the substitution of parameters, but the form of the statements, and the order in which they appeared where changed. Most macro processors can modify the sequence of statements generated for a macro expansion, depending on the arguments supplied in the macro invocation. Such a capability adds greatly to the power and flexibility of a macro language. The term conditional assembly can be used to describe this:

## □ Macro-time conditional statements

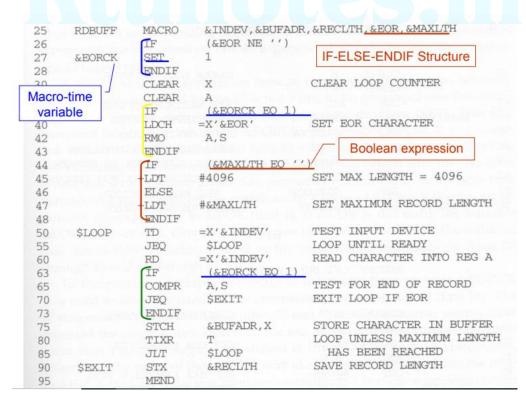
- Macro processor directives:
  - □ IF-ELSE-ENDIF
  - $\Box$  SET



# □ *Macro-time variables* (also called a *set symbol*)

- Be used to store working values during the macro expansion
- Any symbol that begins with the character & and is not a macro parameter
- Be initialized to 0
- Be changed with their values using SET
  - □ &EORCK SET 1

		RDBUFF	OE, BUFFER, L	ENGTH,,80
30		CLEAR	x	CLEAR LOOP COUNTER
35		CLEAR	A	
47		+LDT	#80	SET MAXIMUM RECORD LENGTH
50	\$ABLOOP	TD	=X'0E'	TEST INPUT DEVICE
55		JEQ	\$ABLOOP	LOOP UNTIL READY
60		RD	=X'0E'	READ CHARACTER INTO REG A
75		STCH	BUFFER, X	STORE CHARACTER IN BUFFER
80		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
87		JLT	\$ABLOOP	HAS BEEN REACHED
90	\$ABEXIT	STX	LENGTH	SAVE RECORD LENGTH
			(c)	
		RDBUFF	F1, BUFF, RLE	NG, 04
30		CLEAR	x	CLEAR LOOP COUNTER
35		CLEAR	A	
40		LDCH	=X'04'	SET EOR CHARACTER
42		RMO	A,S	
45		+LDT	#4096	SET MAX LENGTH = 4096
50	\$ACLOOP	TD	=X'F1'	TEST INPUT DEVICE
55		JEQ	\$ACLOOP	LOOP UNTIL READY
60		RD	=X'F1'	READ CHARACTER INTO REG A
65		COMPR	A,S	TEST FOR END OF RECORD
70		JEQ	\$ACEXIT	EXIT LOOP IF EOR
75		STCH	BUFF, X	STORE CHARACTER IN BUFFER
80		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
85		JLT	\$ACLOOP	HAS BEEN REACHED
nn	\$ACEXIT	STX	RLENG	SAVE RECORD LENGTH
90			(d)	
90			(a)	



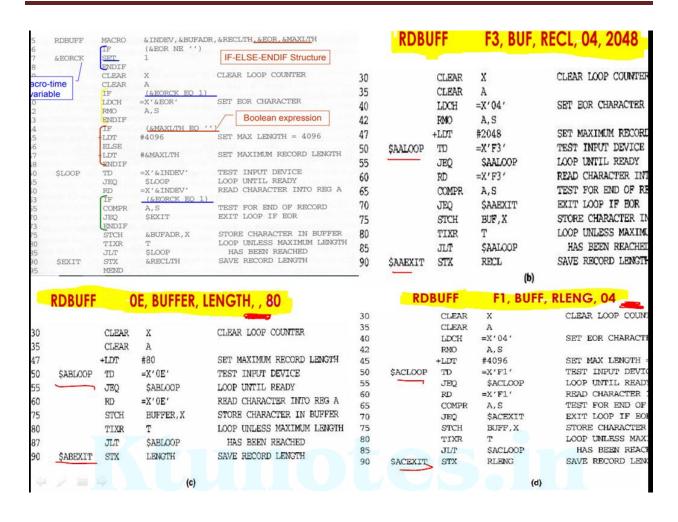


Fig 4.8(a), shows a definition of a macro RDBUFF, the logic and functions of which are similar to those previously explained. However, this definition of RDBUFF has two additional parameters:

- &EOR, which specifies a hexadecimal character code that marks the end of a record
- &MAXLTH, which specifies the maximum length record that can be read

It is possible for either or both of these parameters to be omitted in an invocation of RDBUFF. The statements from line 44 through 48, of the definition illustrate a simple macro-time conditional structure. The IF statement evaluates a Boolean expression, that is its operand. If the value of this expression is TRUE, the statements following the IF are generated until an ELSE is encountered. Otherwise, this statement are skipped, and the statement following the ELSE are generated. The ENDIF statement terminates the conditional expression that was begun by the IF

statement (As usual, the ELSE clause can be omitted entirely.). Thus if the parameter &MAXLTH is equal to the null string, (ie., if the corresponding argument was omitted in the macro invocation statement), then statement on line 45 is generated. Otherwise, the statement on line 47 is generated. Similar structure appears on lines 26 through 28. In this case however, the statement controlled by the IF is not a line to be generated into the macro expansion. Instead, it is another macro processor directive (SET). This SET statement assigns the value 1 to &EORCK.

If the value of the specified Boolean expression is FALSE, the macro processor skips ahead in DEFTAB until it finds the next ELSE or ENDIF statement. The macro processor then resumes normal macro expansion. This implementation does not allow nested IF structures. It is extremely important to understand that the testing of Boolean expressions in IF statements occur at the time macros are expanded. By the time the program is assembled, all such decisions have been made. There is only one sequence of source statement (for eg: fig 4.8(C)) and the conditional macro expansion directives have been removed. Thus macro-time IF statements correspond to options that might have been selected by the programmer in writing the source code.

They are fundamentally different from statements such as COMPR (or IF statements in a high level programming language), which test data values during program execution. The same applies to the assignment of values to macro-time variables, and to the other conditional macro-expansion directives as we have discussed earlier. The macro-time IF-ELSE-ENDIF structure provides a mechanism for either generating(once) or skipping selected statements in the macro-body. A different type of conditional macro expansion statement fig. 4.9

## □ Macro-time looping statement

- Macro processor directives:
  - □ WHILE-ENDW

## ■ Macro processor function

- %NITEMS: the number of members in an argument list
  - □ E.g. &EOR=(00,03,04)
    - => %NITEMS(&EOR) is 3
  - □ Specify member in the list: &EOR[1]

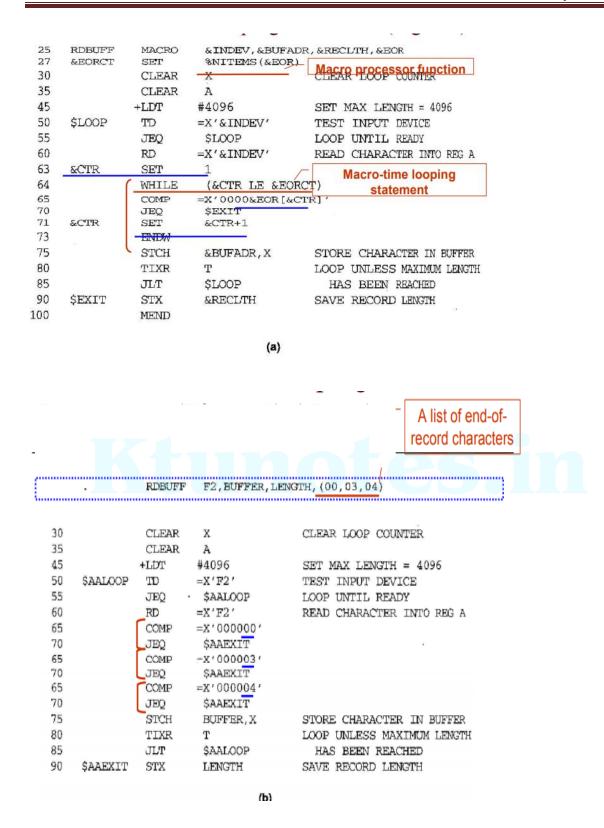


Fig.4.9(a) shows another definition of RDBUFF. The purpose and function of the macro are the same as before. With this definition, however the programmer can specify a list of end-of-record

characters. In the macro invocation statement in fig. 4.9(b) for example, there is a list (00,03,04) corresponding to the parameter &EOR. Any one of these characters is to be interpreted as marking the end of a record. To simplify the macro definition, the parameter &MAXLTH has been deleted. The maximum record length will always be 4096. The definition in Fig. 4.9(a) uses a macro-time looping statement WHILE.

The WHILE statement specifies that the following lines, until the next ENDW statement are to be generated repeatedly, as long as a particular condition is true. As before, the testing of this condition, and the looping, are done while the macro is being expanded. The conditions to be tested involve macro—time variables and arguments, not run-time data values.

## **Implementation**

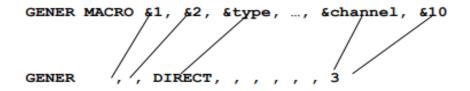
When a WHILE statement is encountered during macro expansion, the specified Boolean expression is evaluated. If the value of this expression is FALSE, the macro processor skips ahead in DEFTAB until it finds the next ENDW statement, and then resumes normal macro expansion. If the value of the Boolean expression is TRUE, the macro processor continues to process the lines from DEFTAB in the usual way until the next ENDW statement. When the ENDW is encountered, the macro processor returns to the preceding WHILE, re-evaluates the Boolean expression, and takes action based on the new value of this expression as previously described. This method of implementation does not allow for nested WHILE structures.

#### **Keyword Macro Parameters**

In macro definitions we have used positional parameters. The parameters and arguments were associated with each other according to their positions in the macro prototype and the macro invocation statement. With positional parameters, the programmer must be careful to specify the arguments in the proper order. If an argument is to be omitted, the macro invocation statement must contain a null argument(two consecutive commas) to maintain the correct argument positions. (check fig.4.8(c)). Positional parameter are quite suitable for most macro instructions.

However, if a macro has a large number of parameters, and only few of these are given values in a typical invocation, a different form of parameter specification is more useful. Such a macro may occur in a situation in which a large and complex sequence of statements- perhaps even an

entire operating system is to be generated from a macro invocation. In such a case most of the parameters may have acceptable default values. The macro invocation specifies only the changes from the default set of values. Suppose that a macro instruction GENER has 10 possible parameters, but in a particular invocation of the macro, only the third and ninth parameters are to be specified. If positional parameters were used, the macro invocation statement might look like:



Using a different form of parameter specification, called keyword parameters, each argument value is written with a keyword that names the corresponding parameter. Arguments may appear in any order. If the third parameter in the previous example, is named &TYPE and the ninth parameter is named &CHANNEL, the macro invocation statement would be:

This statement is obviously much easier to read, and much less error prone, than positional version. Fig4.10(a) shows a version of the RDBUFF macro definition using keyword parameters.

The parameter is assumed to have this default value if its name does not appear in the macro invocation statement. Thus the default value for the parameter &INDEV is F1. There is no default value for the parameter &BUFADR. Default values can simply the macro definition in many cases: For example, the macro definitions in Fig.4.10(a) and 4.8(a) both provide for setting the maximum record length to 4096 unless a different value is specified by the user. The default value is established in fig.4.10(a) takes care of this automatically.

25	RDBUFF	MACRO	&INDEV=F1,&BUR	FADR=, &RECLTH=, &EOR=04, &MAXLTH=409
26		IF	(&BOR NE '')	
27	& EORCK	SET	1	
28	100000000000000000000000000000000000000	ENDIF		
30		CLEAR	X	CLEAR LOOP COUNTER
35		CLEAR	A	The state of the s
38		IF	(&BORCK EQ 1)	
40		LDCH	=X'&BOR'	SET EOR CHARACTER
42		RMO	A.S	DEL DON CHRONCIAN
43		ENDIF	20,0	
47		+LDT	#&MAXL/TH	SET MAXIMUM RECORD LENGTH
50	SLOOP	TD	-X. FINDEA.	TEST INPUT DEVICE
55	STOOL	JEO	SLOOP	LOOP UNTIL READY
60		RD	=X, %INDEA,	
63		IF	(&EORCK BO 1)	READ CHARACTER INTO REG A
65				more non min on necess
70		COMPR	A,S	TEST FOR END OF RECORD
		JEQ	SEXIT	EXIT LOOP IF BOR
73		ENDIF	4 manual man, 14	ALL RESIDENCE AND ADDRESS OF THE PARTY OF TH
75		STCH	&BUFADR, X	STORE CHARACTER IN BUFFER
80		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
85	24.000.00	JLT	\$LOOP	HAS BEEN REACHED
90	SEXIT	STX	&RECLTH	SAVE RECORD LENGTH
95		MEND		
	17	RDBUFF	BUFADR=BUFFER,	RECLITH=LENGTH
30		CLEAR	X	CLEAR LOOP COUNTER
7070		CLEAR	A	
0.0		LDCH	=X'04'	SET EOR CHARACTER
42		RMO	A,S	
47	*****	+LDT	#4096	SET MAXIMUM RECORD LENGTH
50	\$AALOOP	TD	=X'F1'	TEST INPUT DEVICE
55		JEQ	\$AALOOP	LOOP UNTIL READY
50		RD	=X'F1'	READ CHARACTER INTO REG A
55		COMPR	A,S	TEST FOR END OF RECORD
70		JEQ	\$AAEXIT	EXIT LOOP IF EOR
75		STCH	BUFFER, X	STORE CHARACTER IN BUFFER
30		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
85	1223000000	JLT	\$AALOOP	HAS BEEN REACHED
90	SAAEXIT	STX	LENGTH	SAVE RECORD LENGTH
			(1	)

		RDBUFF	THE THE	TH, BUFADR=BUFFER, EOR=, INDEV=F3
30		CLEAR	X	CLEAR LOOP COUNTER
35		CLEAR	A	
47		+LDT	#4096	SET MAXIMUM RECORD LENGTH
50	\$ABLOOP	TD	=X'F3'	TEST INPUT DEVICE
55		JEO	\$ABLOOP	LOOP UNTIL READY
60		RD	=X'F3'	READ CHARACTER INTO REG A
75		STCH	BUFFER, X	STORE CHARACTER IN BUFFER
80		TIXR	T	LOOP UNLESS MAXIMUM LENGTH
85		JLT	\$ABLOOP	HAS BEEN REACHED
90	\$ABEXIT	STX	LENGTH	SAVE RECORD LENGTH
			(c)	
	Figure 4.1			

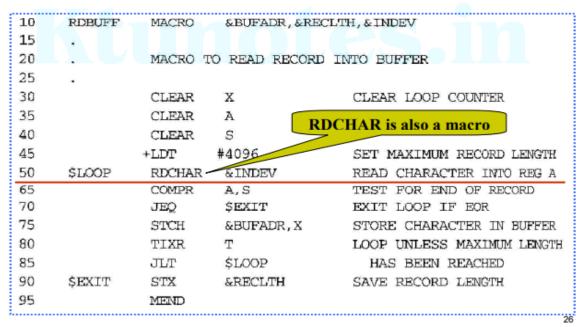
In fig4.8(a), an IF-ELSE-ENDIF structure is required to accomplish the same thing. Other part of fig 4.10 contains examples of the expansion of keyword macro invocation statements. In fig 4.10(b), all default values are accepted .In fig 4.10(c), the value of &INDEV is specified as F3, and the value of &EOR is specified as null. These values override the corresponding defaults. Note that, the arguments may appear in any order in the macro invocation statement.

## **Macro Processor Design Options**

- 1. Recursive Macro Expansion
- 2. General Purpose Macro Processors
- 3. Macro Processing within language Translators

#### **Recursive Macro Expansion**

Fig 4.11(a) example of invocation of one macro by the other- nested macro invocation.



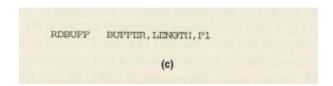
(a)

```
MACRO &IN
10
15
                MACRO TO READ CHARACTER INTO REGISTER A
20
25
                         =X'&IN'
                                          TEST INPUT DEVICE
30
                JEO.
                          *-3
                                          LOOP UNTIL READY
35
                RD
                         =X'&IN'
                                          READ CHARACTER
40
                MEND
                                   (b)
                RDBUFF
                         BUFFER, LENGTH, F1
                                   (c)
```

Figure 4.11 Example of nested macro invocation.

The definition of RDBUFF in fig 4.11(a) is same as in fig 4.1. The order of the parameters has been changed. In this case, we assume that a related macro instruction (RDCHAR) already exists. The purpose of RDCHAR Fig. 4.11(b) is to read one character from a specified device into register A, taking care of the necessary test-and-wait loop. It is convenient to use a macro like RDCHAR in the definition of RDBUFF so that the programmer who is defining RDBUFF need not worry about the details of device access and control. (RDCHAR might be written at different time, or even by a different programmer). The advantages of using RDCHAR in this way would be even greater on a more complex machine, where the code to read a single character might be longer and more complicated than our simple three-line version. The algorithm we studied does not work properly if a macro invocation statement appears within the body of a macro instruction.

Suppose if the algorithm applied to the macro invocation statement in fig.4.11(c)



The procedure EXPAND would be called when the macro was recognized. The arguments from the macro invocation would be entered into ARGTAB as follows:

Parameter	Value
1	BUFFER
2	LENGTH
3	F1
4	(unused)

The Boolean variable EXPANDING would be set to TRUE, and expansion of the macro invocation statement would be begin. The processing would proceed normally until line 50, which contains a statement invoking RDCHAR. At that point, PROCESSLINE would call EXPAND again. " This time, ARGTAB would look like:

1 F1 2 (unused)	Para	meter	Value
		1 2	
			(darabed)

The expansion of RDCHAR would also proceed normally. At the end of this expansion, however, a problem would appear. When the end of the definition of RDCHAR was recognized, EXPANDING would be set to FALSE. Thus the macro processor would "forget" that it had been in middle of expanding a macro when it encountered the RDCHAR statement. In addition, the argument from the original macro invocation (RDBUFF) would be lost because the values in ARGTAB were overwritten with the arguments from the invocation of RDCHAR.

The cause of these difficulties is the recursive call of the procedure EXPAND. When the RDBUFF macro invocation is encountered, EXPAND is called. Later it calls PROCESSLINE for line 50, which results in another call to EXPAND before a return is made from the original call. A similar problem would occur with PROCESSLINE, since this procedure too would be called recursively. For ex: there might be confusion about whether the return from PROCESSLINE should be made to the main (outermost) loop of the macro processor logic or to the loop within EXPAND.

These problems are not so difficult to solve, if the programming language (such as Pascal or C) that allows the recursive calls. Then the compiler would be sure of the previous values of any variables declared within a procedure were saved when that procedure was called recursively. It would take care of other details involving return from the procedure. If the programming language supports recursion is not available, the programmer must take care of handling such items as return addresses and values of local variables. In such case, the PROCESSLINE and EXPAND would probably not be procedures at all. Instead the same logic would be incorporated into a looping structure, with data values being saved on the stack.

#### **Solution**

- Use a Stack to save ARGTAB.
- Use a counter to identify the expansion.

## **General Purpose Macro-Processors**

# Three examples of actual macro processors:

- A macro processor designed for use by assembler language programmers
  ANSI C Macro
  Processor
- Used with a high-level programming language
- General-purpose macro processor

ELENA Macro Processor

- Not tied to any particular language
- □ Can be used with a variety of different languages.

The advantages of such a general-purpose approach to macro processing are:

- The programmer does not need to learn about a different macro facility for each compiler or assembler language, so much of the time and expense involved in training are eliminated
- The cost involved in producing a general purpose macro processor is greater than
  those for developing a language-specific processor. However this expense does not
  need to be repeated for each language. Overall saving in software development cost
  and software maintenance effort

In spite of the advantages noted, there are still relatively few general-purpose macro processors. Why?

Large number of details must be dealt with in a real programming language

- Comment identifications ( //, /\* \*/, ...)
- Grouping together terms, expressions, statements (begin\_end, { }, ...)
- Tokens (keywords, operators)
- **...**
- 1. Large number of details must be dealt with in a real programming language
  - In a special purpose macro processor, these details can be built into its logic and structure
  - In the general –purpose, the user must define the specific set of rules to be followed
- 2. In a typical programming language, there are several situations in which normal macro parameter substitution should not occur

Eg:. comments should usually be ignored by a macro processor

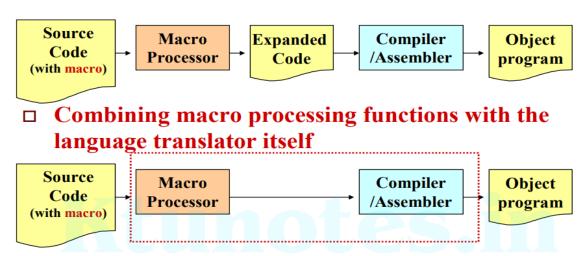
- 3. Another difference between programming languages is related to their facilities for grouping together terms, expressions, or statements
  - E.g. Some languages use keywords such as begin and end for grouping statements.
     Others use special characters such as { and }.
- 4. A more general problem involves the tokens of the programming language
- E.g. identifiers, constants, operators, and keywords
- E.g. blanks
- 5. Another potential problem with general purpose macro processors involves the syntax used for macro definitions and invocation statements. With most special-purpose macro processors, macro invocations are very similar in form to statements in the source programming language.

## **Macro Processing within Language Translators**

- Preprocessors
  - They process macro definition and expand macro invocations, producing an expanded version of the source program
  - This expanded program is then used as input to an assembler or compiler
- Combining the macro processing function with the language translator itself
- Achieved using Line –by –line macro processor
  - The macro processor reads the source program statements
  - Process the statement
  - The output lines are passed to the language translator as they are generated, instead of being written to an expanded source file
  - Thus macro processor operates as a sort of input routine for the assembler or compiler
- Advantages of line-by line macro processor:
  - It avoids making an extra pass over the source program.
  - **Data structures** required by the macro processor and the language translator can be combined
    - □ E.g., OPTAB and NAMTAB)
  - Utility subroutines can be used by both macro processor and the language translator.
    - □ Scanning input lines
    - Searching tables
    - Data format conversion
  - It is easier to give diagnostic messages related to the source statements.
    - i.e., the source statement error can be quickly identified without need to backtrack the source
- Line –by-line macro processor may use some of the same utility routines as the language translator, the functions of macro processing and program translation are relatively independent

- The main form of communication between the two functions is the passing of source statements from one to the other
- It is possible to have even closer cooperation between the macro processor and the assembler or compiler
- Such a scheme can be thought of as a language translator with an integrated macro processor

## □ Preprocessors



## **Integrated Macro Processor**

- Integrate a macro processor with a language translator (e.g., compiler)
- Advantages:
- An integrated macro processor can potentially make use of any information about the source program that is extracted by the language translator.
- An integrated macro processor can support macro instructions that depend upon the context in which they occur.
- Since the Macro Processor may recognize the meaning of source language

## **Drawbacks of Line-by-line or Integrated Macro Processor**

- They must be specially designed and written to work with a particular implementation of an assembler or compiler.
- The costs of macro processor development are added to the costs of the language translator, which results in a more expensive software.
- The assembler or compiler will be considerably larger and more complex than it would be if a macro preprocessor were used
- The size may be a problem if the translator is to run on a computer with limited memory
- In any case, the additional complexity will add to the overhead of language translation
- Decision about what type of macro processor to use should be based on considerations such as the frequency and complexity of macro processing that is anticipated, and other characteristics of computing environment

