SQL - Basics

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SQL Overview

- Structured Query Language
 - standard query language for relational system
 - developed in IBM Almaden (system R)
- Some features
 - Declarative: specify the properties that should hold in the result, not how to obtain the result
 - Complex queries have procedural elements
 - Set/Bag semantics
- International Standards
 - SQL1 (1986)
 - SQL2 (SQL-92)
 - SQL3 (SQL-99)
 - Later extensions in 2003, 2006, 2008, 2011, 2016



SQL Overview (cont)

- Data definition component
 - CREATE TABLE table-name (col-defs, constraints)
 - DROP TABLE table-name
 - ALTER TABLE table-name action
 - ✓ modifies the definition of a table where action is
 - ✓ ADD (col-def)
 - ✓ MODIFY (col-def)
 - ✓ ADD costraint
 - ✓ etc.
- Data update component
 - INSERT INTO table-name ...
 - DELETE FROM table-name ...
 - UPDATE table-name ...
- Query Language



Example Tables

```
branch (bname, address, city, assets)
customer(cname, street, city)
deposit(accno, cname, bname, balance)
loan(accno, cname, bname, amount)
```

```
CREATE TABLE branch (
bname CHAR(20),
address CHAR(30), (in Oracle, a semicolon city CHAR(15), always marks the end of a SQL statement.)
);
```



Simple Queries

branch (bname, address, city, assets)

 Find the names of all branches with assets greater than \$2,500,000.

SELECT bname

FROM branch

WHERE assets > 2500000

 Find the names of all branches in Edmonton with assets greater than \$2,500,000.

SELECT bname

FROM branch

WHERE assets>2500000 AND city='Edmonton'

 Simple predicates can be combined using AND, OR, NOT.



Querying two Relations

```
customer(cname, street, city)
deposit(accno, cname, bname, balance)
```

 List the name and the city of every customer who has an account with balance over \$2,000.

```
SELECT customer.cname, city
FROM customer, deposit
WHERE balance > 2000
AND customer.cname = deposit.cname
```



Queries With Tuple Variables

```
loan(accno, cname, bname, amount)
deposit(accno, cname, bname, balance)
```

Find customers who have both loans and deposits.

```
SELECT loan.cname
FROM loan, deposit
WHERE loan.cname = deposit.cname
```

Equivalently using tuple variables:

```
SELECT l.cname
FROM loan l, deposit d
WHERE l.cname = d.cname
```

 Range variables are really needed if the same relation appears twice in the FROM clause.



A Simple Evaluation Alg.

```
SELECT ...
FROM R1 r1, R2 r2, ...
WHERE C
```

- Tuple variables r1, r2, ... respectively range over rows of R1, R2, ...
- Evaluation strategy:
 - FROM clause produces Cartesian product of listed tables
 - WHERE clause produces table containing only rows satisfying condition
 - SELECT clause retains listed columns

```
for every tuple r1 in R1, r2 in R2, ...
let r := r1, r2, ...
if C(r) then
output the desired columns
```



From SQL to Relational Algebra

```
FROM loan l, deposit d
WHERE l.cname = d.cname
```

Equivalent to:

```
temp = loan \times deposit [a1,cn1,bn1,b1,a2,cn2,bn2,a2]
\pi_{cn1} \sigma_{cn1=cn2} (temp)
```

This is a simple evaluation algorithm for SELECT.



Queries With Set/Bag Results

- Find all cities of customers.
 - **SELECT** city **FROM** customer
- Result?
- To get rid of duplicates, we need
 - SELECT DISTINCT city
 FROM customer



Duplicates

- Duplicate rows not allowed in a relation
- However, duplicate elimination from query result is costly and not automatically done; it must be explicitly requested:

```
SELECT DISTINCT .....
FROM .....
```



Working with Strings

- Equality and comparison operators apply to strings (based on lexical ordering)
 - E.g. WHERE cname < 'P'
- Concatenate operator applies to strings
 - E.g. WHERE *bname* || '--' || *address* =
- Expressions can also be used in SELECT clause
 - E.g. SELECT bname || '--' || address AS NameAdd FROM branch



Partial Matching

customer(cname, street, city)

 Find every customer whose address starts with "Computing Science".

```
SELECT *
FROM customer
WHERE address LIKE 'Computing Science%'
```

- Expression: col-name [NOT] LIKE pattern
 - Pattern may include wildcard characters '%'
 matching any string and '_' (underscore) matching
 any single character.



Naming the Results

deposit(accno, cname, bname, balance)

 For every deposit holder who has over \$1000, find the customer name and the balance over \$1000.

```
SELECT cname, (balance - 1000) as bal
FROM deposit
WHERE balance > 1000
```



Ordering the Results

branch (bname, address, city, assets)

 Find the names and assets of all branches with assets greater than \$2,500,000 and order the result in ascending order of asset values.

SELECT bname, assets

FROM branch

WHERE assets > 2500000

ORDER BY assets

 Default is ascending order; a descending order can be specified by the DESC keyword.



Queries Involving Set Operators

- set union : Q1 UNION Q2
 - the set of tuples in Q1 or Q2
- set difference : Q1 EXCEPT Q2
 - the set of tuples in Q1 but not in Q2
- set intersection : Q1 INTERSECT Q2
 - the set of tuples in both Q1 and Q2
- Q1 and Q2 must be union-compatibles
 - same number/types of attributes.



Queries With Set Operators

List deposit-holders who have no loans.

```
(SELECT cname deposit)

EXCEPT

(SELECT cname from loan)
```

 List cities where there is either a customer or a branch.

```
(SELECT city
FROM customer)
UNION

(SELECT city
FROM branch)
```



Queries With Set Operators

 Find all cities that have both customers and branches in.

```
(SELECT city
FROM customer)

INTERSECT

(SELECT city
FROM branch)
```

Find every city that has a branch but no customer.



Queries Over Multiple Relations

```
branch (bname, address, city, assets)
customer(cname, street, city)
deposit(accno, cname, bname, balance)
```

```
FROM branch.bname, assets
FROM branch, customer, deposit
WHERE customer.city = 'Jasper'
AND customer.cname = deposit.cname
AND deposit.bname = branch.bname
```

What does the query do?

Find the name and asset of every branch that has a deposit account holders who lives in Jasper.



Queries With Nested Structures

Queries within the WHERE clause of an outer query

```
SELECT
FROM
WHERE OPERATOR (SELECT ... FROM ... WHERE)
```

- There can be multiple levels of nesting
- Operators: IN, (NOT) EXISTS, < ALL, ...
- Avoid nesting as much as possible.



Nested Structures Using "IN"

```
FROM
WHERE expr (expr list) IN
(set of values)
```

• E.g.

```
... WHERE province IN ('AB', 'BC')
... WHERE province NOT IN ('AB', 'BC')
```



Example

deposit(accno, cname, bname, balance)

What does the query do?

Find every customer who has a deposit in some branch at which John Doe has a deposit.



The Same Example Without Nesting

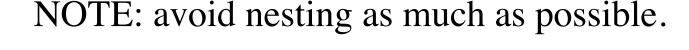
deposit(accno, cname, bname, balance)

SELECT d1.cname

FROM deposit d1, deposit d2

WHERE d2.cname = 'John Doe'

AND d1.bname = d2.bname





Nested Structures Using "< ALL",...

```
SELECT
FROM
WHERE expr < ALL (set of values)</pre>
```

- Other forms: "<= ALL", "= ALL", ">= ALL", ">= ALL",
- Op ALL (set of values) evaluates to true iff the comparison evaluates to true for every value in the set.
- Op SOME (set of values) evaluates to true iff the the comparison evaluates to true for at least one value in the set.



Nested Structures Using ">ALL", ...

branch (bname, address, city, assets)

 Find branches that have assets greater than the assets of all branches in Calgary.

```
FROM branch
WHERE assets > ALL

(SELECT assets
FROM branch
WHERE city = 'Calgary')
```



Nested Structures Using "EXISTS"

```
SELECT
FROM
WHERE (NOT) EXISTS (SELECT ... )
```

- EXISTS (SELECT ...)
 evaluates to true iff the result of the subquery
 contains at least one row.
- The expression is evaluated for every iteration of the outer query.



"EXISTS" Construct Example

```
branch (bname, address, city, assets)
customer(cname, street, city)
```

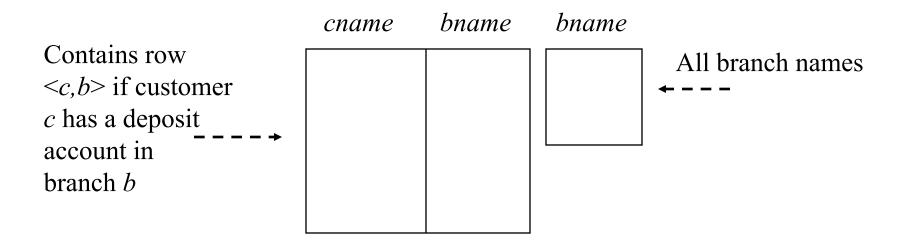
 Find the names of customers who live in a city with no bank branches.

- Find the names of customers who live in a city which has a bank branch.
 - Change NOT EXISTS to EXISTS. (can also write it using join)



Division

- Query type: Find the subset of items in one set that are related to all items in another set
- Example: Find customers who have deposit accounts in all branches.



 $\pi_{\text{cname, bname}}$ (deposit) / π_{bname} (branch)



Division

- Strategy for implementing division in SQL:
 - Find set, A, of all branches in which a particular customer, c, has a deposit account.
 - Find set, B, of all branches.
 - Output c if A

 B, or, equivalently, if B

 A is empty



Division – SQL Solution

branch (bname, address, city, assets)
deposit(accno, cname, bname, balance)

```
SELECT c.cname
FROM customer c
WHERE NOT EXISTS
  (SELECT b.bname
                         -- set B of all branches
   FROM branch b
      EXCEPT
   SELECT d.bname
                          -- set A of branches in which
                           -- customer c has a deposit account
   FROM deposit d
   WHERE d.cname=c.cname
                              -- global variable
```



Division – 2nd SQL Solution

- Find customers who have deposit accounts in all branches.
 - Same as "find all customers such that there is no branch where they do not have deposits in."

```
SELECT c.cname
          FROM
                  customer c
          WHERE NOT EXISTS
               (SELECT
                           bname
                                                       branches where
                FROM
                           branch b
                WHERE NOT EXISTS
                                                       c has no deposits in
                   	ilde{	iny SELECT} 	ilde{	iny *}
Deposits of c in
                   FROM
                          deposit
branch b
                   WHERE b.bname=deposit.bname
                   AND
                           c.cname = deposit.cname))
```



Division – Another Example

- Find professors who have taught courses in all departments.
- Strategy for implementing division in SQL:
 - Find set, A, of all departments in which a particular professor, p, has taught a course
 - Find set, B, of all departments
 - Output p if A

 B, or, equivalently, if B

 empty



Division – SQL Solution

Professor (Id, Name, DeptId)
Department (DeptId, Name)
Course (DeptId, CrsCode, CrsName, Descr)
Teaching (ProfId, CrsCode, Semester)

```
SELECT P.Id
FROM Professor P
WHERE NOT EXISTS
(SELECT D.DeptId -- set B of all dept Ids
FROM Department D
EXCEPT
SELECT C.DeptId -- set A of dept Ids of depts in
```

FROM Teaching T, Course C
WHERE T.ProfId=P.Id -- global variable
AND T.CrsCode=C.CrsCode)

-- which P has taught a course



SQL and SQLite

- SQLite: pretty modern but light language
 - Released in 2000 (Oracle was 1977, DB2 was 1983)
 - It is light (intended for small devices)
 - Tightly integrated with applications (software library rather than a stand-alone system/process to communicate)
 - Locks the whole file/database during writes (not a good choice for write-intensive/concurrent transactions)
 - Supports both in-memory and on-disk databases
- Differences with SQL-92
 - Does not support op ALL in nested queries
 - Does not support RIGHT OUTER JOIN and FULL OUTER JOIN
 - Views are read-only
 - ALTER TABLE command is very limited
 - Foreign keys constraints are not enforced by default
 - More left to be explored!



Basic Data Types in SQLite

■ **INTEGER** 1, 2, 3, 4, 6, 8 bytes integer

(depending on the magnitude)

REAL 8 bytes floating point number

TEXT stored using database encoding (e.g. UTF-8)

BLOB stored as is

NULL null value

NUMERIC can store all other types; stored value is converted to INTEGER or REAL if the conversion is lossless and reversible



Other Types in SQLite

INT,SMALLINT INTEGER

CHAR(n), VARCHAR(n) TEXT

DOUBLE, FLOAT REAL

DECIMAL
 NUMERIC

DATE, DATETIME NUMERIC

