

Lambda Calculus: Notational Simplification and Examples

We will write λ for the word *lambda* in an expression, and use a vertical bar to separate the parameter of a lambda function from its body. For example,

$$((\lambda x \mid (x \ y)) \ a)$$

It is of the form $(Exp1 \ Exp2)$ where $Exp1 = (\lambda x \mid (x \ y))$ and $Exp2 = a$.

Since we will use single letters to denote identifiers, we need not leave space between two identifiers. For example, the above expression could be written as

$$((\lambda x \mid (xy)) \ a)$$

We may omit some parentheses, for the above example for instance, we may write

$$(\lambda x \mid xy) \ a$$

We omit parentheses by left association for expressions of more than two components, i.e.

$$N \ M \ Q = ((N \ M) \ Q)$$

Note that under this convention, $N \ M \ Q \neq N \ (M \ Q)$.

In what follows, a reduction step indicated by \rightarrow means beta-reduction if not said otherwise.

We extend the syntax of lambda calculus by allowing multiple parameters such as

$$(\lambda xy \mid \dots)$$

In this case, when it comes to substitution, all occurrences of x and y are simultaneously replaced by their arguments. That is, they do not interfere with each other even in the case of name conflicts. For example,

$$\begin{aligned} & (\lambda xy \mid xy) \ (\lambda y \mid y) \ (\lambda x \mid x) \\ & \rightarrow (\lambda y \mid y) \ (\lambda x \mid x) \\ & \rightarrow (\lambda x \mid x) \end{aligned}$$

In the first reduction step above, the occurrence of x in xy is replaced by $(\lambda y \mid y)$ and the occurrence y in xy by $(\lambda x \mid x)$. The two replacements happen simultaneously. Better yet, simultaneous replacement doesn't cause name conflict

As we discussed earlier, an n-ary function can be viewed equivalently as a unary function, each parameter being dealt with at a time. So, a function with multiple parameters such as $(\lambda xy|...)$ can also be viewed equivalently as $(\lambda x|(\lambda y|...))$. Thus,

$$\begin{aligned} & (\lambda xy|xy) (\lambda y|y) (\lambda x|x) = (\lambda x|(\lambda y|xy)) (\lambda y|y) (\lambda x|x) \\ & \rightarrow (\lambda y|(\lambda y|y) y) (\lambda x|x) \\ & \rightarrow (\lambda y|y) (\lambda x|x) \\ & \rightarrow (\lambda x|x) \end{aligned}$$

Reducing one parameter at a time is convenient when not enough actual parameters are available.

$$\begin{aligned} & (\lambda xz|xz)(\lambda y|x) \\ & \rightarrow (\lambda z|(\lambda y|x) z) \\ & \rightarrow (\lambda z|x) \end{aligned}$$

Note that $(\lambda y|x)z$ reduces to x .

Additional Exercises

Reduce the following λ -expressions to their normal forms. Show all the reduction steps.

(a)

$$\begin{aligned} & (\lambda x | (\lambda xy | xy)xy)s \\ & \rightarrow (\lambda xy | xy)sy \\ & \rightarrow sy \end{aligned}$$

(b)

$$\begin{aligned} & (\lambda xy | xxy)(\lambda x | xy) \\ & \rightarrow (\lambda w | (\lambda x | xy)(\lambda x | xy)w) \\ & \rightarrow (\lambda w | (\lambda x | xy)yw) \\ & \rightarrow (\lambda w | yyw) \end{aligned}$$

(c)

$$\begin{aligned} & (\lambda xy | x(xy)) (\lambda z | sz) \\ & \rightarrow (\lambda y | (\lambda z | sz)((\lambda z | sz)y)) \\ & \rightarrow (\lambda y | (\lambda z | sz)(sy)) \\ & \rightarrow (\lambda y | s(sy)) \end{aligned}$$

Note that $(\lambda y | s(sy))$ doesn't equal $(\lambda y | ssy)$. By the simplified notation, $ssy = ((ss)y)$

(c) This example is more involved notationally.

$$\begin{aligned} & ((\lambda xy \mid x (\lambda y \mid (+ y 2))) (\lambda fx \mid f (f x))) 5 \\ & \rightarrow (\lambda y \mid (\lambda fx \mid f (f x)) (\lambda y \mid (+ y 2))) 5 \\ & \rightarrow (\lambda fx \mid f (f x)) (\lambda y \mid (+ y 2)) \\ & \rightarrow (\lambda x \mid (\lambda y \mid (+ y 2))((\lambda y \mid (+ y 2))x)) \\ & \rightarrow (\lambda x \mid (\lambda y \mid (+ y 2))(+ x 2)) \\ & \rightarrow (\lambda x \mid (+ (+ x 2) 2)) \end{aligned}$$