CMPUT 325: NON-PROCEDURAL PROGRAMMING LANGUAGES

A List Interpreter in LISP

Md Solimul Chowdhury

mdsolimu@ualberta.ca

INTRODUCTION

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- In assignment 2, we are implementing an interpreter for a language named FL.
 - Primitive functions.
 - User-defined functions.
- We will have a walkthrough
 - With an implementation of a similar interpreter
 - A restricted version of <u>list interpreter in LISP</u>
- Should help you in implementing your current assignment.
- Implementation of the List Interpreter has been posted
 - o in the Instructor, TA's and Labs page in eclass)

https://eclass.srv.ualberta.ca/pluginfile.php/6714341/mod page/content/89/i
nterpreter.lisp

A LIST INTERPRETER IN LISP

• The List interpreter implements 16 functions

```
OPERATIONS
                     - a variable
 (xquote s)
                     - a constant
  (+ e1 e2)
  (- e1 e2)
  (* e1 e2)
  (/ e1 e2)
 (rem e1 e2)
 (atom e1)
 (car e1)
 (cdr e1)
 (cons e1 e2)
 (eq e1 e2)
 (leq e1 e2)
OTHERS
 (if e1 e2 e3)
 (e e1 e2 ... en)
 (let (x1 x2 ... xk) (e1 e2 ... ek) e)
 (letrec (x1 x2 ... xk) (e1 e2 ... ek) e)
 (lambda (x1 x2 ... xn) e )
```

A LIST INTERPRETER IN LISP

- You may categorize these functions into 2 types:
 - OPERATIONS
 - Takes at <u>most two parameters</u>
 - OTHERS
 - Takes more than two parameters

- Categorization helps to reduce the
 - Complexity of Implementations.

```
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  (atom e1)
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  (cons e1 e2)
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  (if e1 e2 e3)
  (e e1 e2 ... en)
 (let (x1 x2 ... xk) (e1 e2 ... ek) e)
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  (lambda (x1 x2 ... xn) e )
```

A LIST INTERPRETER IN LISP: SOME CAVEATS

- Calling function to the interpreter is: startEval
- Accepts an expression **e** as argument

```
(startEval 'e)
```

- quote is needed here
- has <u>no notion of a built-in list</u> data type.

```
(startEval '(car '(1 2) ) ) → NIL → wrong results.
```

Have to use cons.

```
(starteval '(car (cons 1 (cons 2 nil)))) → 1
```

SOME EXAMPLES

```
(starteval '(xquote (1 2)))
(starteval '(atom x))
(starteval '(car (cons 1 (cons 2 nil))))
(starteval '(cdr (cons 1 (cons 2 nil))))
(starteval '((lambda (x) (+ x 1)) 5))
(starteval '(+ 5 2))
```

IMPLEMENTATION OF THE LIST INTERPRETER

• **starteval** is the function that implements the interpreter.

```
(defun startEval (e)
      (xeval e nil nil)
)
```

- parameters: **e**, expression of the form shown in the previous slides
- Passes e to a helper function xeval that evaluates e
 - recursively, if needed
 - xeval is the working horse for the list interpreter
 - Comprised of a series of nested if-else-if ... expressions to handle various cases.

IMPLEMENTATION OF THE LIST INTERPRETER: TRIVIAL CASES

```
(if (eq e t)
   ; t is bound to itself
(if (null e)
  ; nil is bound to itself
 nil
(if (numberp e)
  ; numbers are bound to themselves
```

IMPLEMENTATION OF THE LIST INTERPRETER

```
(if (atom e)
  ; a variable - return result of searching context
    for variable represented by 'e'.
(xassoc e n v) -> e.g., finding location/value
       (e1 (car (cdr e)))
       (e2 (car (cdr (cdr e))))
       (e3 (car (cdr (cdr e)))))
```

IMPLEMENTATION OF THE LIST INTERPRETER: OPERATIONS

```
(if (eq func-name 'xquote)
    ; return the first argument unevaluated
    e1
 (if (member func-name '(+ - * / rem atom car cdr cons eq leq))
    ; Functions with 1 or 2 expressions as arguments,
where each argument needs to be evaluated.
    (let ( (ev-e1 (xeval e1 n v)) Recursive calls
                                            For Grounding
          (ev-e2 (xeval e2 n v)≯
                                        e, and ez
       (if (eq func-name '+)
          ; return the sum of the two evaluated arguments.
    obviously, the evaluated arguments should be integers.
          (+ ev-e1 ev-e2)
```

IMPLEMENTATION OF THE LIST INTERPRETER: OPERATIONS

Similarly, cdr

```
(if (eq func-name 'cdr)
           ; return the list represented by ev-e1 with its first
           ; element removed. If ev-e1 is an atom,
           ; return nil (not the constant nil, just nil -
           ; this will allow callers to test for execution
           ; errors)
                                ev-el: (xevol \ e_i \ n \ v)
e_i: (cdr (car \ e))
          (if (atom ev-e1)
              nil
              (cdr ev-e1)
```

IMPLEMENTATION OF THE LIST INTERPRETER: OTHERS

```
Functions with other argument formats
     (let ()
(if (eq func-name 'if)
              ; If first argument evaluate to true,
                   return evaluated second argument,
                   otherwise return evaluted third argument.
             (if (xeval e1 n v) \rightarrow if e_l = T, (x \neq val e_l \land v) = T

(xeval e2 n v) e_l = T, (x \neq val e_l \land v) = T

(xeval e3 n v)
```

IMPLEMENTATION OF THE LIST INTERPRETER: OTHERS

```
(if (eq func-name 'let)
   ; Simply add the names and values defined in the let
     to the current name and value lists, then evaluate
     the body of the let.
                                (let (n1, ..., nx) (e1,..., eu) e
  (let ((new-names (cadr e))
        (new-values (caddr e))←
        (body (cadddr e)))
        (let (
             (ev-parms (evlis new-values n v)) fretions
              (let
              ((all-values (cons ev-parms v)))
              (xeval body all-names all-values) -> call
```