CMPUT 379 – Operating System Concepts Practice Final Exam - Fall 2019 Department of Computing Science, University of Alberta Instructor: Omid Ardakanian

Full Name:	
Student ID Number:	

Instructions

- Print your full name and ID clearly above.
- You have 120 minutes to complete the exam.
- There should be 6 questions and 8 pages in this exam booklet. You are responsible for checking that your exam booklet is complete.
- It is a closed-book exam. You are not allowed to bring printed or handwritten notes. Electronic devices such as laptops, calculators, phones, etc. are strictly prohibited.
- Place answers in the spaces provided on the question pages. Keep your answers brief. Think about each question a bit before answering. If required, you may use the backside of a sheet to provide an answer but clearly indicate when you have done so.
- This exam is worth 100 points and counts 35% toward your final grade in this course. The weight of each question is indicated in square brackets by the question number.

Question	1	2	3	4	5	6	Total
Out of	15	10	30	15	10	20	100

Question 1 [15 points]: Operating System Concepts

Choose either True or False for the questions below. You do not need to provide justifications.

	\mathbf{r}
(1)	An operating system provides a specific set of services to user programs through system calls regardless of the programming language they are written in.
	True False
(2)	Named pipe cannot be used by two processes that do not have a parent-child relationship.
	True False
(3)	Round Robin scheduling may lead to starvation.
	True False
(4)	A thread pool can prevent an overly popular web site from crashing the hosting web server.
	True False
(5)	Deadlock will eventually happen when a system is in an unsafe state.
	True False
(6)	Increasing the page size will increase external fragmentation.
	True False
(7)	The number of page faults will certainly decrease as we increase the number of page frames.
	True False
(8)	A Translation Lookaside Buffer (TLB) caches frequently accessed page table entries.
	True False
(9)	Thrashing occurs when the sum of the working sets of all processes exceeds the available memory.
	True False
(10)	A disk can have multiple partitions, each having a different file system.
	True False
(11)	A file descriptor is an index into the system-wide open file table.
	True False
(12)	Indexed allocation has a higher block-pointer overhead than linked allocation for small files.
	True False
(13)	Random-access I/O is faster on NVM drives than HDDs.
	True False
(14)	The SSTF algorithm for disk scheduling yields the minimum total length of disk seeks needed to
	service a given set of I/O requests.
	True False
(15)	The storage overhead of RAID 1 is higher than RAID 5.
	True False

Question 2 [10 points]: Multiple Choice QuestionsYou must select the best answer from the choices given.

Part 2.1 [2 points]: Which scheduling algorithm makes the most sense for NVM devices?
○ SCAN.
○ Shortest Seek Time First (SSTF).
○ C-SCAN.
○ First Come, First Served (FCFS).
Part 2.2 [2 points]: Which memory allocation strategy picks the largest hole?
○ First fit.
○ Best fit.
○ Worst fit.
○ None of the above.
Part 2.3 [2 points]: Which of the following page replacement algorithms approximates the Least Recently Used (LRU) page replacement algorithm more closely?
○ First-In First-Out.
○ Most Frequently Used.
○ Second Chance.
○ Additional-Reference-Bits.
Part 2.4 [2 points]: Which of the following scheduling algorithms is non-preemptive?
○ First-Come First-Served.
 Shortest Remaining Time First.
○ Round Robin.
 Multilevel Feedback Queue.
Part 2.5 [2 points]: How does the effective memory access time on a TLB miss of a system with a 4-level page table compare with that of a system with a 2-level page table?
\bigcirc 5/3x slower.
○ 2x slower.
\bigcirc 5/3x faster.
○ 2x faster.

Question 3 [30 points]: Short-Answer Questions

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		dding more m	emory to a cor	nputer make its	processes run fa	ster?
3.2 [5 points]]: Why could ac					
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Part 3.3 [5 points	s]: What are the differen	ce between segment	tation and paging?	Give two differences.
Each inode contai	ins 16 direct blocks and	em which has 32-bit 2 single indirect bl	file pointers and bocks. What is the	plocks of size 512 bytes. maximum possible file
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Question 4 [15 points]: Process Synchronization

Part 4.1 [10 points]: Build a mutex lock using a POSIX semaphore by implementing the three functions defined blow.

Hint: Recall that a POSIX semaphore must be initialized using sem_init before it can be operated on using sem_post and sem_wait.

```
struct lock_t {
  sem_t asem;
};

void lock_init(lock_t *lock) {

void acquire(lock_t *lock) {

}

void release(lock_t *lock) {
```

Part 4.2 [5 points]: Briefly explain the difference in behaviour of sem_post and pthread_cond_signal when no threads are waiting in the corresponding semaphore or condition variable.

Question 5 [10 points]: Page Replacement Policies

Consider the following page reference string in a demand-paged virtual memory system with 3 page frames that are initially empty:

12, 8, 3, 7, 8, 3, 9, 10, 3, 5, 7, 10, 3, 7, 5, 7, 12, 10, 7, 5, 12

Part 5.1 [5 points]: How many times will page fault occur under the LRU policy? Use the following table to keep track of the pages that are loaded in the page frames in each step. If there is a tie, break it by replacing the oldest page.

Request	12	8	3	7	8	3	9	10	3	5	7	10	3	7	5	7	12	10	7	5	12
Frame 1																					
Frame 2																					
Frame 3																					

Part 5.2 [5 points]: How many times will page fault occur under the optimal policy? Use the following table to keep track of the pages that are loaded in the page frames in each step. If there is a tie, break it by replacing the oldest page.

Request	12	8	3	7	8	3	9	10	3	5	7	10	3	7	5	7	12	10	7	5	12
Frame 1																					
Frame 2																					
Frame 3																					

Question 6 [20 points]: File System and Disk

Part 6.1 [10 points]: Consider a file that is allocated 100 data blocks on disk which are organized linked list, i.e., the pointer to the next file block is stored at the end of each block. Suppose a data be is appended to the end of the file which causes the file size to grow to 101 blocks. How many disk is	lock
and write operations are required to update the file and the file control block?	
You should assume that each operation reads or writes a single disk block, both the file and the file	con-
trol block are initially on disk, and the file buffer cache is empty.	
Part 6.2 [10 points]: Answer the same question this time for the case that an index block contain block pointers rather than storing them at the end of data blocks (i.e., indexed allocation is used ins of linked allocation).	