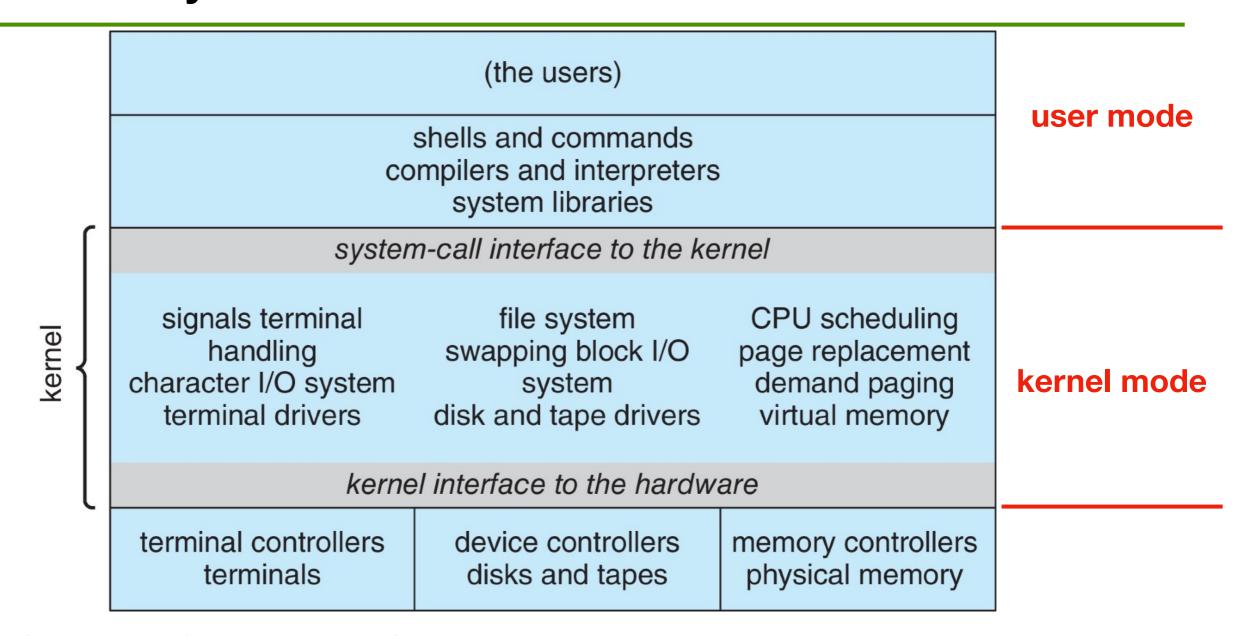
Operating System Concepts

Lecture 4a: Operating System Structure

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MWF 12:00-12:50 VVC 2 215

UNIX system has a monolithic structure

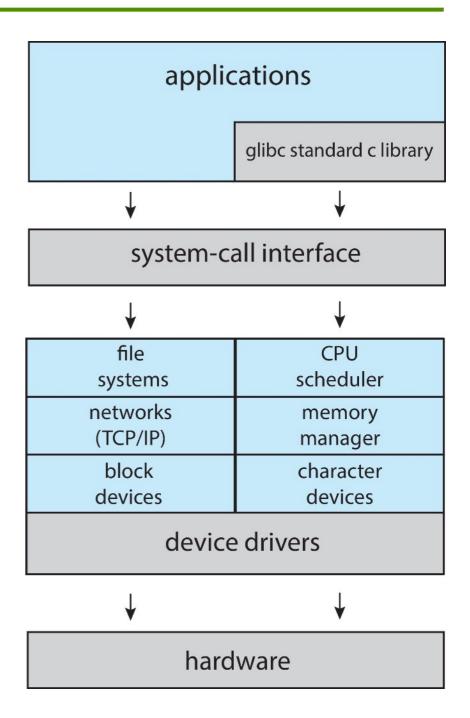


<u>Definition</u>: all functionality of the kernel placed into a single static binary file that runs in a single address space

- faster communication with the kernel
- little overhead in the system call interface

Linux system structure is also monolithic

- core (~ 30% of Linux source code) + dynamically loaded kernel modules
 - examples are device drivers, file systems, network protocols
- modules can be loaded at boot time or during run time
 - adding new modules does not require recompiling the kernel
 - each module talks to other modules through known interfaces



Linux system structure is also monolithic

Applications use the GNU version of the standard C library (glibc)

which provides the functionality required by POSIX

- glibc provides a wrapper around system calls
- glibc is the system-call interface to the kernel
- Linux kernel source: https://www.kernel.org/
 - Directory structure:

include: public headers

kernel: core kernel components (e.g., scheduler)

arch: hardware-dependent code

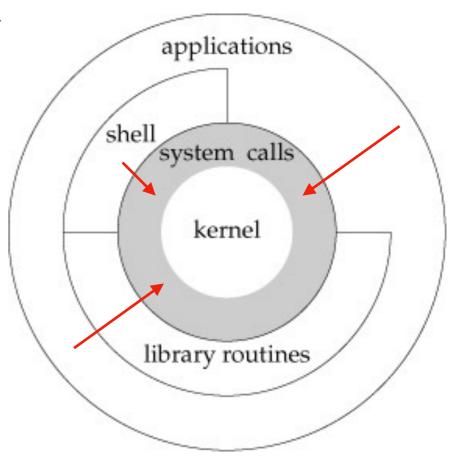
fs: file systems

mm: memory management

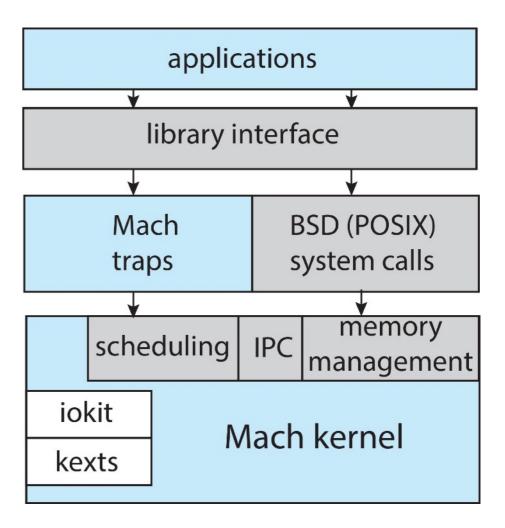
ipc: interprocess communication

drivers: device drivers

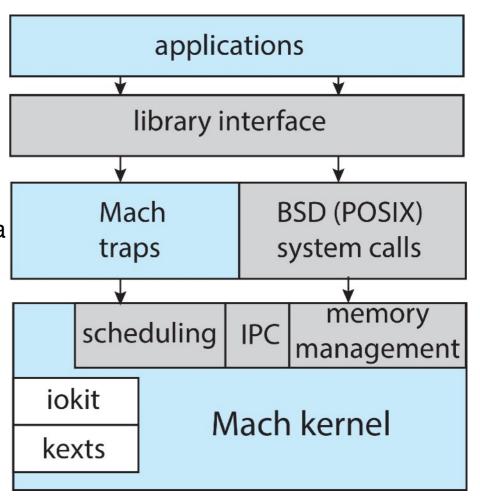
usr: user-space code
lib: common libraries



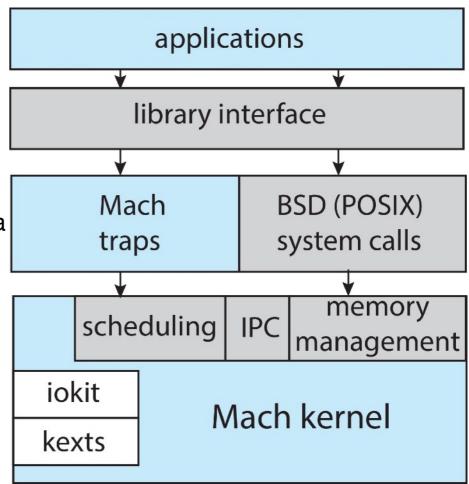
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- Layered structure
 - advantages: modularity, simplicity, portability, ease of design/ debugging
 - disadvantage: communication overhead between layers, extra copying, book-keeping

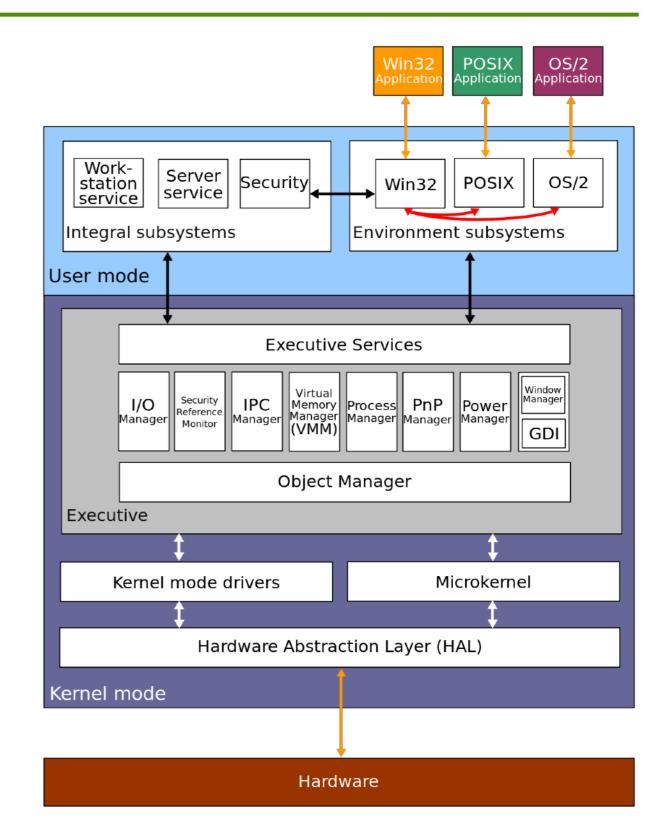


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- Microkernel structure
 - a small kernel providing
 - interprocess communication (message passing usually through ports)
 - basic functionality (e.g., scheduling and virtual memory management)
 - other OS functionality (device drivers, file system, networking, user interface, etc.) implemented as user-space processes



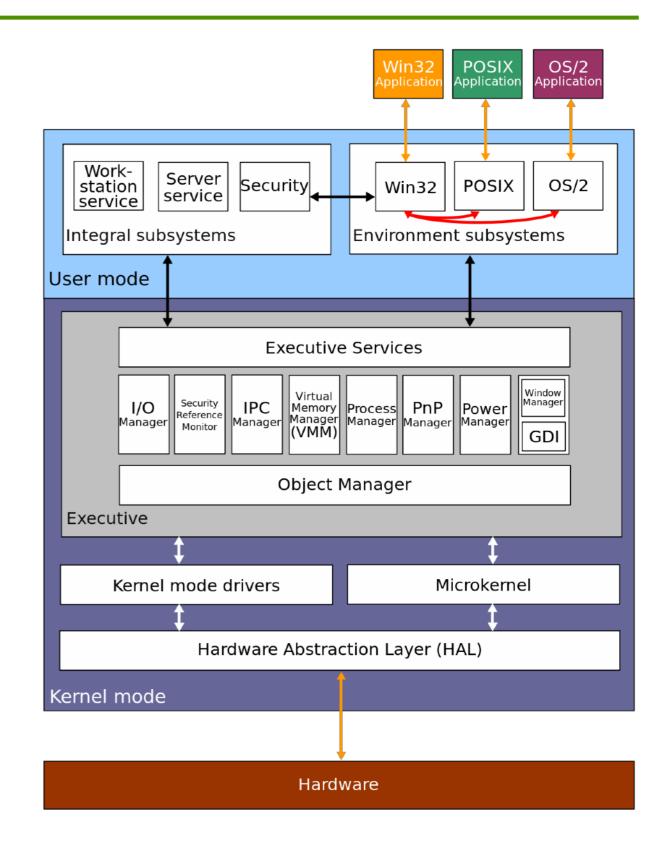
Windows NT

- Layered architecture with several modules
- the Windows executive provides core OS services (Ntoskrnl.exe)
- the kernel itself resides between HAL and the executive layer
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 - POSIX, OS/2 and Win32 APIs are documented and can be used by applications
- the hardware abstraction layer (HAL.dll) provides portability across a variety of hardware platforms
 - device drivers call functions in the HAL to interface with the hardware

