CMPUT 379 Lab

ETLC E1003: Tuesday, 5:00 – 7:50 PM.

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CAB 311: Thursday, 2:00 – 4:50 PM.

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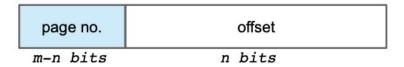
Today's Lab

- More examples on memory management
- Assignment 3

Memory management examples

Recap - Paging

- The memory space of a process is divided into fixed-size pages.
- The logical address is divided into page number and offset



Recap - Paging

- The memory space of a process is divided into fixed-size pages.
- The logical address is divided into page number and offset
- The physical memory is divided into fixed-size frames that holds pages

Physical memory

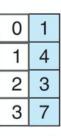
page 1 of process 5	Frame 0
(empty)	Frame 1
page 3 of process 2	Frame 2
(empty)	Frame 3

. . .

(empty)	Frame n-2
page 4 of process 2	Frame n-1
(empty)	Frame n

Recap - Paging

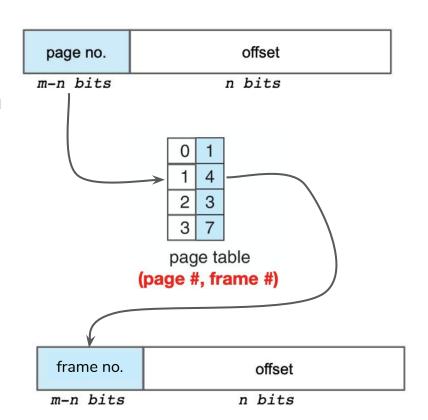
- The memory space of a process is divided into fixed-size pages.
- The logical address is divided into page number and offset
- The physical memory is divided into fixed-size frames that holds pages
- The page table of each process stores in which frame each page locates



page table (page #, frame #)

Recap - Paging translation

- 1. Read the **page no.** and **offset** from a VA
- Look up the frame no. of that page no. in the process' page table
- 3. Replace the **page no.** in the VA with **frame no.** to get the **PA**



Page table size

page no. offset

m-n bits n bits

- Assuming we can address 1 byte increments
- Page / frame size: 2ⁿ bytes
- Number of pages: 2^(m-n)
- Large page will cause internal fragmentation
- Small page will incur large page table ...

- Example:
 - o 32-bit address space
 - 4KB (2^12 byte) page size
 - 4-byte page-table entry
- Number of page table entries: 2^32 / 2^12 = 2^20
- Page table size: 2^20 * 4 bytes = 4MB (per process)
- My laptop currently has 499 running processes
- RAM for storing page tables: ~2GB
- How to reduce page table size?

Multi-level page tables

page r	number	page offset
p_1	p_2	d
10	10	12

- Look up the outer page table to get the frame number for the inner page table
- Not all inner page tables need to be allocated at first save space!
- Aliases
 - Outer page table page directory
 - Inner page table page table

known as forward-mapped page table because address translation works from the outer page table inward

able

age of the

couter page table

page of page table

 p_1 is an index into the outer page table p_2 is the displacement within the page of the inner page table d is the page offset

Multi-level page tables - have a try

- Download the example from eClass
- Go to **HW-Paging-MultiLevelTranslate** folder
- Run paging-multilevel-translate.py to get example page directory, page table, memory dump and some VAs
- Try to locate the page table for each VA (by looking up the page directory), translate VA to PA (by looking up the page table), and check your answers

Page replacement policy

- RAM are fast but small
 - Physical memory < memory used by processes
 - Virtual memory: move some pages to the disk
- Which one to replace? (cache replacement policies)
 - First in first out (FIFIO)
 - Least recently used (LRU)
 - Least frequently used (LFU)
 - Most recently used (MRU)
 - o etc.

Page replacement policy - have a try

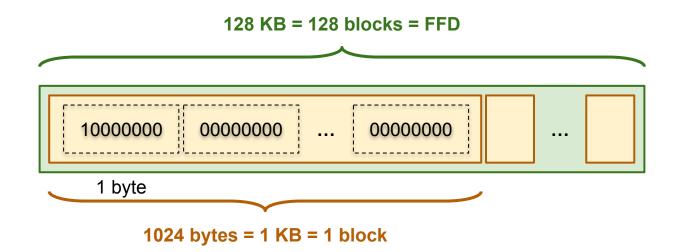
- Download the example from eClass
- Try ./paging-policy.py --addresses=0,1,2,0,1,3,0,3,1,2,1
 --policy=LRU --cachesize=3 -c with different policies, and compare the number of misses

UNIX-like File System

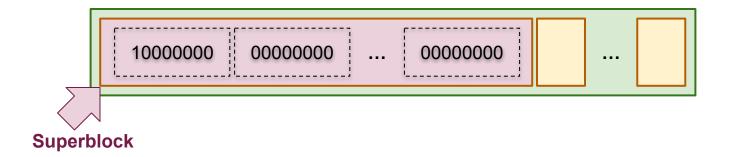
Implementing a Unix-like FS

Assignment 3

- Our toy FFD (fake file disk) has a capacity of 128 KB, which contains only one partition and one simulated file system.
- Blocks on our simulated file system are 1024 bytes (or 1 KB) in length.



The first block is the superblock, that stores free-space list and inodes.



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10000000
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00000000
                                                                 00000000
                                                               Superblock
```

8 bytes 16 bytes 24 bytes 25 -32 bytes 33 -40 bytes 48 bytes 56 bytes 64 bytes 72 bytes 80 bytes 73 -88 bytes 96 bytes 104 bytes 112 bytes 120 bytes - 128 bytes 128 -1024 bytes

Free-block list

几

ſ											
	10000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	11111111			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
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	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000			
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	Superblock										
- 1								•			

8 bytes 16 bytes 24 bytes 32 bytes 25 -40 bytes 41 -48 bytes 49 -56 bytes 64 bytes 65 - 72 bytes 73 - 80 bytes 81 - 88 bytes 89 - 96 bytes 97 - 104 bytes 105 - 112 bytes 113 - 120 bytes 121 - 128 bytes 128 -1024 bytes

Free-block list

53rd block is available

								l
10000000	00000000	00000000	00000000	00000000	00000000	000000	00000000	
00000000	00000000	00000000	00000000	00000000	00000000	00000000	11111111	
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	1
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00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	
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00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	6
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	7
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	8
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	8
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	2
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	16
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	11
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	12
						Su	perblock	12
							-	1

8 bytes 16 bytes 24 bytes 32 bytes 40 bytes 48 bytes 56 bytes 64 bytes 65 - 72 bytes 73 - 80 bytes 81 - 88 bytes 96 bytes 97 - 104 bytes .05 - 112 bytes 13 - 120 bytes 21 - 128 bytes 28 -1024 bytes

Free-block list

128th block is occupied

10000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	1 -	8	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	11111111	9 -	16	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	17 -	24	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	25 -	32	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	33 -	40	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	41 -	48	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	49 -	56	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	57 -	64	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	65 -	72	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	73 -	80	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	81 -	88	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	89 -	96	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	97 -	104	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	105 -	112	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	113 -	120	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	121 -	128	bytes
						Su	perblock	128 -1	L024	bytes

Free-block list

128th block is occupied

10000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	1 .	- 8	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	11111111	9 -	- 16	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	17 -	- 24	bytes
00000000	00000000	00000000	00000000	20000000	00000000	00000000	00000000	25 -	- 32	bytes
00000000	00000000	00000000		-	00000000	00000000	00000000	33 -	- 40	bytes
00000000	00000000	00000000	1st i	node	00000000	00000000	00000000	41 -	- 48	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	49 -	- 56	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	57 -	- 64	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	65 -	- 72	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	73 -	- 80	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	81 -	- 88	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	89 -	- 96	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	97 -	- 104	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	105 -	- 112	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	113 -	- 120	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	121 -	- 128	bytes
						Su	perblock	128 -	-1024	bytes

Free-block list

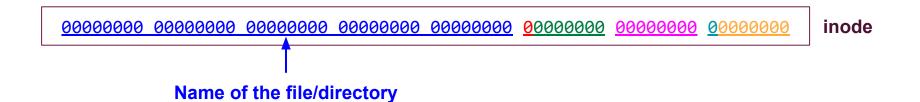
128th block is occupied

								1 /			
10000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	1	-	8	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	11111111 <mark>1</mark>	9	-	16	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	17	-	24	bytes
00000000	00000000	00000000	00000000	20000000	00000000	00000000	00000000	25	-	32	bytes
00000000	00000000	00000000			00000000	00000000	00000000	33	-	40	bytes
00000000	00000000	00000000	1st i	node	00000000	00000000	00000000	41	-	48	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	49	-	56	bytes
00000000	00000000	00000000	9th i	node	00000000	00000000	00000000	57	-	64	bytes
00000000	00000000	00000000	5111	Touc	00000000	00000000	00000000	65	-	72	bytes
00000000	00000000	00000000	0000000	0000000	00000000	00000000	00000000	73	-	80	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	81	-	88	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	89	-	96	bytes
00000000	00000000	00000000	14th	inode	00000000	00000000	00000000	97	-	104	bytes
00000000	00000000	00000000		1	00000000	00000000	00000000	105	-	112	bytes
00000000	00000000	00000000	0000000	0000000	00000000	00000000	00000000	113	-	120	bytes
00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000	121	-	128	bytes
						Su	perblock	128	-1	1024	bytes

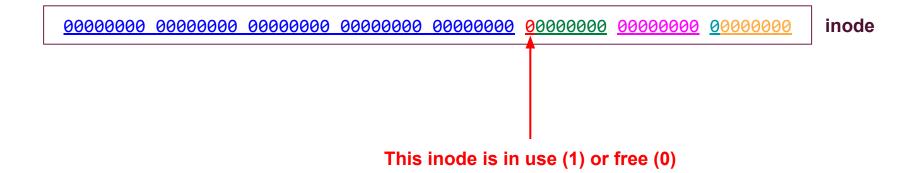
 The inode (short for index node) contains the metadata for each file or directory in the file system, such as name, size, etc.

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• The **inode** (short for index node) contains the metadata for each file or directory in the file system, such as name, size, etc.

Size of the file $(1\sim127)$ or directory (0)

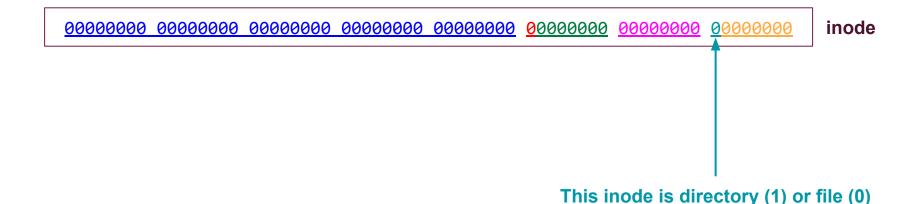
inode

 The inode (short for index node) contains the metadata for each file or directory in the file system, such as name, size, etc.

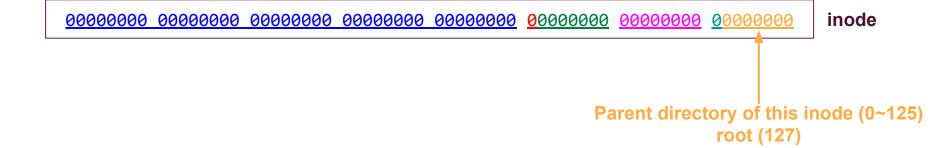


File start block index (1~127) or directory (0)

 The inode (short for index node) contains the metadata for each file or directory in the file system, such as name, size, etc.



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 The inode (short for index node) contains the metadata for each file or directory in the file system, such as name, size, etc.

This inode is directory (1) or file (0)

- The first block is the superblock, that stores free-space list and inodes.
- Free-space list occupy the first 128 bits (16 bytes) in the superblock, each bit represents the corresponding block is available or occupied.
- Each inode (short for index node) occupy 8 bytes in the superblock. We have 126 inodes in the superblock, which occupy 126 × 8 = 1008 bytes in the superblock in total.

In superblock: 16 bytes (free-block list) + 1008 bytes (inodes) = 1024 bytes = 1 KB

Your program also have a buffer, which is 1 KB. Zero-out it at the beginning.

```
void fs_mount(char *new_disk_name);
void fs_create(char name[5], int size);
void fs_delete(char name[5]);
void fs_read(char name[5], int block_num);
void fs_write(char name[5], int block_num);
void fs_buff(uint8_t buff[1024]);
void fs_ls(void);
void fs_resize(char name[5], int new_size);
void fs_defrag(void);
void fs_cd(char name[5]);
```

```
void fs_mount(char *new_disk_name);
```

- Mount FFD to your program. Remember we are mounting FFD with our simulated FS, so we are not mounting a real disk with real FS, which means you should not use system call "mount".
- Do not forget to check the consistency of the file system. In this assignment, we only care about 6
 potential errors in the FFD caused by unknown reasons. You need to check them one by one and
 follow the given order.
- If the FFD passes the test, then mount it into the program by reading the superblock of this FFD into the memory, and unmount the previous FFD. Otherwise, keep using the last FFD.
- A success fs_mount will change the current working directory to root.

```
void fs_create(char name[5], int size);
```

- Create a file on current mounted FFD with given name and a fixed size.
- When assign blocks to the file, the file must be allocated a number of contiguous blocks.
- A size of 0 means the user wants to create a directory.
- Remember that you cannot have two files with the same name under the same path.
- Name . and .. are reserved and cannot be used.
- In this assignment, user can have symbols in the name as well (such as ...). We will only test the characters that is on your keyboard (i.e. will not test '\n', '\r', etc.).

```
void fs_delete(char name[5]);
```

- Delete the file on mounted FFD by its name.
- If a directory is selected, then remove the directory and all files/directories in it.

```
void fs_read(char name[5], int block_num);
```

Put 1 KB data in the specified block to buffer.

```
void fs_write(char name[5], int block_num);
```

Put 1 KB data in buffer to the specified block.

```
void fs_buff(uint8_t buff[1024]);
```

Put user input to buffer.

```
void fs_ls(void);
```

- Same thing as when you type Is in your terminal.
- List files and directories based on the order they stored in inodes.
- Format in description.

```
void fs_resize(char name[5], int new_size);
```

- Change the size of the file into new_size.
- If you have to shift the file, remember to clean-up data in the old position. We do not want any dummy data in your FFD.

```
void fs_defrag(void);
```

- Organize the contents of all files in FFD.
- You should shift all files to the lowest start block index, in the meantime maintain the order of those files (i.e., if file1 locate physically after file4, then after defragmentation, file1 should still physically after file4).

```
void fs_cd(char name[5]);
```

- Change the current working directory into the new directory.
- Same as cd in your terminal.

General comments:

- 1. Concepts are relatively easy. This is a coding assignment, and the implementation is straight forward.
- 2. Test your code with errors. We listed all errors you need to handle, and we assigned priorities to those errors.
- 3. You are encouraged to make your test cases and submit them.
- 4. Start early.