ECE 212 Lab - Introduction to Microprocessors Department of Electrical and Computer Engineering University of Alberta

Lab 3: Introduction to Subroutines

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Contents

1	ntroduction	2
	1 Part A	2
	2 Part B	2
	3 Part C	2
2	Design	3
	1 Part A	3
	2 Part B	4
	2.2.1 Sample Calculations	6
	3 Part C	
3	esting	8
_	1 Part A	
	2 Part B	
	3 Part C	
4	uestions	12
	1 Question 1	12
	2 Question 2	
	3 Question 3	
5	conclusion	13
6	ppendix	14
	1 Part A Assembler Code	14
	2 Part A Flowchart Diagram	
	3 Part B Assembler Code	
	4 Part B Flowchart Diagram	
	5 Part C Assembler Code	
	6 Part C Flowchart Diagram	
7	Marking Sheet	30
8	farking Sheet	30

1 Introduction

This lab deals with stack operation (push and pop), segmenting a long program/function into several smaller and simpler subroutines/sub-functions. Each part uses these subroutines in order to produce the final output of the program which is statistical information about the input values, including the mean, max, min and number of entries divisible by another specified value. The subroutines must be run in order, however each is designed to be transparent and to flow seamlessly together. In addition this lab introduces the use of strings for user prompts, allowing the user to get feedback from the program as well as information about what they should enter at each step. These two combined are major aspects of programming in assembly. Subroutines allow a longer task to be subdivided into smaller pieces that can be called multiple times if required, and text prompts help not only with the execution of the program by the users, but also with the debugging by the developers.

1.1 Part A

In part A, FILLINTHEINTRO

1.2 Part B

For part B, FILLINTHEINTRO

1.3 Part C

FILLINTHEINTRO

The purpose of this lab is to learn and test with the Assembly language in a hands on environment in order to solidify the concepts learned in class and to improve our skill in the language. In addition, we will be learning how to handle the Netburner ColdFire boards directly, manipulating the contents of their memory and data structures. Finally, we are going to learn how to work inside the Eclipse IDE environment and how to properly use the powerful tools that come alongside it.

The code will be developed for the Netburner ColdFire Platform, which has some parameters that should be kept in mind throughout testing. There are multiple Data and Address registers, and the memory is indexed by hexadecimal codes. The data and the stored locations can each be modified directly, values can be compared and the code can branch into different sections depending on the values of the CCR (Condition Code Register) bits, which store information about the outcome from the last comparison or valid operation, such as if a value is negative or zero. These will be used to execute code conditionally.

The lab will be split into two sections, each with a different goal but with similar implementations. For one part we will be taking in an ASCII value and if the character it represents is a character included in the symbols for hexadecimal numbers then that hexadecimal value is output, otherwise it returns an error. For the second part an ASCII value is taken in, and if the character it represents is a letter in the English language (A-Z) then the ASCII code for the character in the opposite case is output. Thus, valid uppercase English letters are converted to their lowercase equivalents in ASCII and vice versa.

These experiments will introduce implementing high level programming practices of loops, if - then - else statements, using the Assembly language. More specifically, this will introduce the movement of memory and data to and from different parts of the Netburner chip, using techniques

such as referencing memory addresses and copying data to local data registers. The debugger tools of the IDE will be used to closely examine this movement and to analyze all changes to the data in order to solve issues in development as well as to test the code. This is all building off the concepts explored in Lab 0.

The computing science practice of Pair Programming was also introduced, where two people develop and test code in tandem. The partners are divided into the Driver and the Navigator. In this structure the Driver is the one responsible for the physical typing of the code into the computer, and the Navigator reviews this code and clarifies the meaning of each passage in order to find bugs faster and to improve efficiency in testing. The two partners should communicate constantly and switch in order to maximize the efficiency of this working model. This will not only decrease time needed for development but it will also improve the quality of code from each partner.

2 Design

2.1 Part A

For the first part of the lab, we begin by writing our own subroutine. Some code was already supplied such that:

- Stack Entry Condition =
 - 1. Space allocated for the number of entries on stack (long word)
 - 2. Space allocated for the divisor number on stack (long word)
- Stack Exit Condtion =
 - 1. Number of entries on stack (long word)
 - 2. Divisor number on stack (long word)

We started by editing Lab3a.s, and chose to back up address registers a2-a6 and the data registers d2-d7. Since the stack pointer was a7, this was accomplished by subtracting 44 from the stack pointer, then using the *movem* command to push the values onto the stack. Then at the end, just before the rts command in the subroutine, we move the values we backed up to the registers as they were before running the subroutine. Address register a2 was chosen to hold the memory location of where the valid numbers which were entered would go, which was 0x2300000. Next, at the very end of the file, we defined several messages as strings that would be used as prompts. Labels are used to reference these strings, which are used for prompting the user in the MTTTY serial monitor to enter relevant data for each part.

With the initial setup for this subroutine completed, we moved on to actually coding the prompt messages for the program. We begin by pushing the Welcome Prompt onto the stack, and using the provided *iprintf* and *cr* subroutines to print the Welcome Message and a carriage return for the user. Once the message was displayed, we immediately clean up the stack by adding 4 to the stack pointer, since the address of the message was pushed onto the stack, and addresses take up 4 bytes of memory. For every message printed onwards, we immediately clean up the stack after the message is printed. We then branch to a label where the user is prompted to enter the number of entries. The subroutine getstr, which was provided captures input from the serial monitor, and puts it into data register *d0*. For this part, the only valid entries are the numbers 3-15, with anything

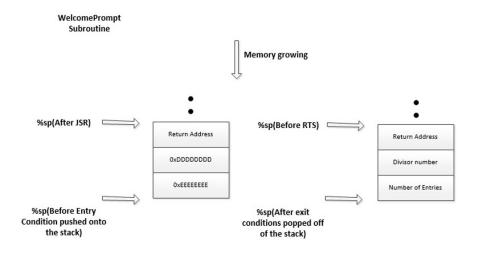


Figure 1: Visualization of the stack for Part A

else entered being rejected. This was accomplished by comparing the input to the number 15. If the number was greater than 15. we branch to a label to warn the user that an invalid entry was entered, and then the program returns to the label which prompts the user to enter the number of entries. Similarly, if the value entered was less than 3, the same thing would occur. If the entry was valid, we put the numentries value that the user entered in the spot reserved on the stack for us. This value was also copied to d7, which is later used as a counter for when the user enters the numbers. As a feature for the user, if the input was accepted, the value is printed to the serial monitor for the user to see. This is done in a similar method to how the messages are printed onto the stack.

In a very similar fashion, the user is then prompted to enter the divisor number. Here, the accepted values were from 2-5, with the data validation checks happening in almost exactly the same was as outlined above for entering the number of entries. The divisor also went into a spot on the stack, which was reserved for us.

Next, we enter a loop, where the user is prompted to enter numbers, until the user enters as many numbers as the user declared when prompted to enter the number of entries. The loop itself loops numentries-1 times, since there's a different prompt message for the last number. As mentioned earlier, we use d7 as our counter, since numentries was copied to this data register. The user then goes through the same process of entering numbers, and is prompted to re-enter the number if it is not positive. This is done similarly to above, in that the number entered is compared to the number 1, and if the number entered is less than 1, we branch to a label where the user is warned about the mistake, and then branches back to continue the loop. If the number is accepted, then it is printed to the serial monitor, and then the number entered is moved to the address location pointed to by a2, and a2 is post incremented. On each successful loop iteration, the number 1 is subtracted from the counter d7, until the counter is 1.

At this point, the user needs to enter one last number, for which there is an appropriate prompt, and the data validation is the exact same as above, with only positive values being accepted. If the

number is valid, then it is printed to the serial monitor, and moved to the memory location where a2 points.

Finally, we reach the part where the backed up registers on the stack are restored, and the subroutine hits the rts command.

2.2 Part B

In Part B, another subroutine is written, which computes some stats on the numbers entered. Similar to Part A, some code was provided such that the stack entry and exit conditions are as follows:

- Stack Entry Condition =
 - 1. Space allocated for how many numbers are divisible by the divisor (long word)
 - 2. Number of entries (long word) on stack
 - 3. Divisor number (long word) on stack
- Stack Exit Condtion =
 - 1. How many numbers are divisible by the divisor (long word) on stack
 - 2. Number of entries on stack (long word)
 - 3. Divisor number on stack (long word)

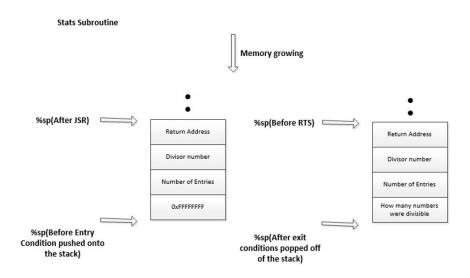


Figure 2: Visualization of the stack for Part B

We begin by backing up the registers a0-a6 and d0-d7 by subtracting 60 from the stack pointer, and then using the *movem* command to push those values onto the stack. Similarly, at the end of

the subroutine just before rts, we restore the values by using *movem* to restore the values, then add 60 to the stack pointer before returning to the previous subroutine.

We begin by reading the numerities and divisor off of the stack, which were put there by the subroutine in Part A. The addresses 0x2300000 and 0x2310000 were loaded into address registers a2 and a3 respectively. a2 kept track of the location were the numbers were entered, and a3 kept track of where the min, mean, max, and divisible numbers would be output.

To begin calculating the minimum number, we copy the numentries to another data register to use as a counter. The first number is read by indirectly addressing a2 with a post increment, and moved to a data register which will hold the final minimum value. Inside a loop, which loops numentries-1 times, (since the first number was already read), the next number is read by indirectly addressing a2 with a post increment, and moved into a temporary data register, where it is then compared to the data register which is supposed to hold the minimum value. If the temporary value is less than the minimum value, we move the temporary value to the data register which holds the minumum value. If the temporary value is greater than the current minimum, then the loop continues. At the end of every loop iteration, the loop counter is decremented by 1, and compared to the number 1, since it loops numentries-1 times. Once this loop is complete, the minumum value is stored in data register d5, so we move it to the memory location pointed at by a3 and post increment it.

The maximum number is calculated in a very similar manner compared to how the min was calculated. First, a2 is reset to 0x2300000, and the loop counter is reset to numerities, since we post incremented a2 when finding the min, and we decremented the counter when going through the numbers again to find the min. Then, we go through a loop that is almost the same as for calculating the minumum number. except that the current maximum number is held in one data register, and a temporary value is read and compared to the maximum number. If the temporary number is greater than the current maximum, then the temporary number becomes the maximum, and by the end of the loop, the maximum value is stored in data register d6, so we move it to the memory location pointed at by a3 and post increment it.

To find the mean, address register a2 was reset to 0x2300000, and the loop counter was reset to numentries again. A data register was used to store the cumulative sum of the numbers, which was accomplished by first clearing the data register, then looping over the list of numbers entered (pointed at by a2), and adding each number to the sum. This was done by indirectly addressing a2 with a post increment to copy the value to a data register, then adding that value to our cumulative sum. At the end of the loop, we have the cumulative sum, so to find the mean, we simply divide the cumulative sum by the numentries by using divs.l. When the division was complete, the mean was stored in data register d4, so so we move it to the memory location pointed at by a3 and post increment it.

Finally, to find the divisible numbers, we once more reset a2 to 0x2300000, and the loop counter to numentries. We go through one last loop, where a number is read by indirectly addressing a2, moving it to a data register, and then using divs.w to divide the number by the divisor, since divs.w provides the remainder stored in the 16 most significant bits. If the number is divisible, it should not have a remainder, so we check the remainder by bit shifting the division result 16 bits to the right, and check if the remainder is 0. If so, a counter which holds the number of divisible numbers is incremented by 1, and the divisible number is moved to the location pointed at by a3 and post incremented. The loop then continues until all the numbers have been processed. At the end of the loop, the number of divisible entries is stored in data register d7, so we move it to the spot reserved on the stack for us.

We then reach the end of the subroutine, where we restore the backed up register values on the stack, add 60 to the stack pointer, and return to the previous subroutine.

2.2.1 Sample Calculations

Numbers entered: 1,2,3,4,5

Numentries: 5 Divisor: 2

Min: 1 Max: 5

Mean: $\frac{1+2+3+4+5}{5} = 3$ Numdivisible: 2 Divisible numbers: 2,4

2.3 Part C

In Part C, the last subroutine is written, which prints the stats on the numbers entered. Similar to the previous parts, some code was provided such that the stack entry and exit conditions are as follows:

- Stack Entry Condition =
 - 1. How many numbers were divisible by the divisor
 - 2. Number of entries (long word)
 - 3. Divisor number (long word)
- Stack Exit Condtion =
 - 1. How many numbers are divisible by the divisor
 - 2. Number of entries (long word)
 - 3. Divisor number (long word)

Just like in part B, we begin by backing up the registers $a\theta$ -a6 and $d\theta$ -d7 by subtracting 60 from the stack pointer, and then using the *movem* command to push those values onto the stack. Similarly, at the end of the subroutine just before rts, we restore the values by using *movem* to restore the values, then add 60 to the stack pointer before returning to the previous subroutine. We begin by loading 0x2300000 into a2, 0x2310000 into a3, and the divisor, numentries, and numdivisible are read from the stack and moved into data registers.

At the very end of the program, we define strings just like in Part A, which are printed for the user to read. We begin by first printing the number of entries. A message is printed by pushing the address of a string defined earlier onto the stack, and using *iprintf*. Next, the actual number of entries is pushed onto the stack and printed to the serial monitor using the subroutine *value*. The stack is also cleaned up by adding to the stack pointer. The user is then prompted that the numbers will be printed, which is done in a very similar manner, using *iprintf*. Then, a loop iterates over the numbers, which are indirectly addressed with a2 with post increment, and printed with the subroutine *value*. The data register holding the numentries is decremented at the end of each

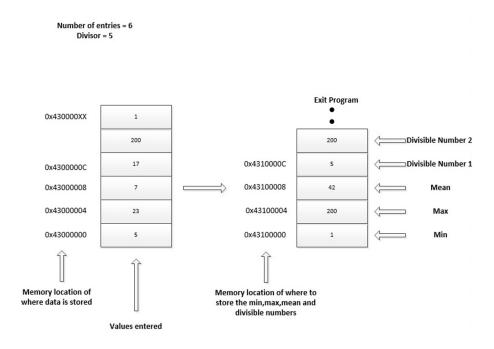


Figure 3: Visualization of the stack for Part C

iteration of the loop, until it reaches 0, which is when all the numbers have been printed. The min, max, mean, and number of divisible numbers are all printed in a similar manner to how the numentries was printed, with a message printed first using *iprintf*, and then the value being printed with the subroutine *value*. Whenever something is printed, the stack is immediately cleaned up afterwards by adding 4 to the stack pointer whenever a value is printed.

Finally, in a manner similar to how the numbers were printed, the divisible numbers are printed by indirectly addressing a3 with post increment, putting the number on the stack, and jumping to the subroutine value to print the numbers. This happens in a loop, with numdivisible being decremented at the end of each loop iteration, until all the divisible numbers are printed out.

At the end of the subroutine, the string "End of program" is printed to let the user know that the program has ended.

3 Testing

3.1 Part A

Initially, we visually tested our code by using the debugger in Eclipse IDE. While stepping through the code, we would check the values at relevant memory locations, and the data and address registers. Extra care had to be taken while monitoring the stack pointer, as it was very easy to miss an off-by-one error, and then have the program crash unexpectedly. On many occasions, the board had to be flashed with its default program when our would not upload. This was done by pressing the blue reset button, then typing Shift + A within 2 seconds. Next, fla was typed in and then

the F5 key was pressed where we could then send the default program to the microcontroller and re flash our program when trying to debug it. A lab TA showed us a helpful debugging practice, which was to insert a breakpoint at the very beginning of the subroutine, and at the very end. Using this method allowed us to see if our program at least produced the correct output before the trap error occurs. At least for this first part, a brute force approach also helped, where the value we wanted to replace on the stack was increased by 4 and re-uploaded to the microcontroller until the program worked. However, we quickly realized that this method was much less effective when we worked on Part B, which was more complicated and thus required a thorough understanding of how the stack works. When the bugs were ironed out our code was tested using the provided Lab3.s file. More specifically, this program was moved into the project folder along with main.cpp, and the program's functionality was verified. We intentionally entered invalid data such as negative numbers, or numbers outside of the accepted ranges at each step, and actually discovered a small bug, where our program accepted divisors ranging from 3-15 instead of 2-5. Luckily this was a trivial fix, since only two numbers had to be tweaked in the code. When we had a working solution for Part A, we moved on to Part B.

```
Testing Subroutines. Choose from the menu below
1 - Test First subroutine
 - Test Second subroutine
3 - Test Third subroutine
4 - Test All subroutine
Testing 1st subroutine.
Welcome to our amazing statistics program
Please enter the number (3min-15max) of entries followed by enter
Please enter the divisor (2min-5max) followed by enter
Please enter a number (positive only)
Please enter a number (positive only)
Invalid entry, please enter proper value.
Please enter a number(positive only)
Please enter a number (positive only)
Please enter a number (positive only)
Please enter the last number (positive only)
The Number of entries is 5
The Divisor number is 2
```

Figure 4: MTTTY output when testing our Part A solution

3.2 Part B

The procedure for testing our code for part B was similar to the process described above in Part A. We visually inspected our code in the Eclipse IDE, used the Eclipse debugger to step through our code, and monitor relevant memory addresses and registers. This part certainly proved to be more involved than Part A, in that a brute force approach to stop the program from crashing would take too long. Deeper inspection was required when monitoring the memory addresses and stack pointer, as well as a thorough understanding of what was happpening on the stack in order to catch errors. The default program also had to be re-uploaded many more times (compared to Part A) when the microcontroller froze. In the end, however, the program worked, and we verified the output by plugging numbers into our calculators to verify the mean and divisible numbers, and then visually inspecting the min, max, and numdivisible values. When we had a working solution

for Part A, we moved on to Part C.

```
Testing Subroutines. Choose from the menu below

1 - Test First subroutine

2 - Test Second subroutine

3 - Test Third subroutine

4 - Test All subroutine

2

Testing 2nd subroutine.

The entered values are:

5

4

3

2

1

The Min, Max, Mean are:

1

5

3

The Numbers divisible by 2 is/are 4 2
```

Figure 5: MTTTY output when testing our Part B solution

3.3 Part C

The procedure for testing our code for Part C felt more trivial than for Parts A and B, likely as a result from the practice we had debugging those parts. We took care to avoid the mistakes made when writing Parts A and B, such as forgetting to clean up the stack or not understanding where the values are on the stack. This resulted in producing a working solution for Part C in a efficient manner. When testing, we ran into a small bug where the program would enter an infinite loop displaying random numbers, when it was supposed to only display the numbers divisible by the divisor that the user enterered. Using our previous debugging experience from Parts A and B, we quickly discovered that we had mixed up the data registers which hold the numdivisible and the divisor values, and then verified our solution by checking the numbers, and manually calculating the mean ourselves to verify the program's output.

```
The Numbers divisible by 3 is/are 24 45 63 18
Testing Subroutines. Choose from the menu below
1 - Test First subroutine
2 - Test Second subroutine
3 - Test Third subroutine
4 - Test All subroutine
Testing 3nd subroutine.
The number of entries was: 10
The entered number(s) were:
24
17
14
63
10
19
18
Min number: 5
Max number: 63
Mean number: 24
There are 4 number(s) divisible by 3
45
63
18
```

Figure 6: MTTTY output when testing our Part C solution

Finally, our solution was verified by a lab TA. We initially entered values on our own, with invalid data as well to show how our program re-prompts the user to re-enter the data. Then, the TA tried several futher test cases to validate our solution.

```
58Z
Please enter a number (positive only)
Please enter a number (positive only)
Please enter the last number (positive only)
The number of entries was: 10
The entered number(s) were:
51
56
1000
150
16
582
54
80211
69
Min number: 16
Max number: 80211
Mean number: 8227
There are 6 number(s) divisible by 2
1000
150
16
582
The stack at beginning is set at SP = 0x200027C8
```

Figure 7: MTTTY output when testing all parts of our solution

4 Questions

4.1 Question 1

Is it always necessary to implement either callee or caller preservation of registers when calling a subroutine. Why?

A: Yes, because pretty much all subroutines need to modify registers to get anything meaningful done, whether the subroutine needs to store a temporary value, do some addition, or keep track of a memory location. Without the preservation of registers, the subroutine will modify registers that the caller may have been using, which will likely result in unexpected behavior in the caller or callee programs. However, not all registers need to be backed up. The only registers which need to be backed up and restored are the ones that the subroutine modifies.

but you don't want to overwrite the data in a register that the previous subroutine might need

4.2 Question 2

Is it always necessary to clean up the stack. Why?

A: Not necessarily but it is good practice, and you run the risk of memory leaks or trap errors from attempting to access a restricted memory location. Cleaning up the stack saves a lot of headaches, and it makes it easier to keep track of the values that your program uses. It also prevents your program from behaving unexpectedly.

4.3 Question 3

If a proper check for the getstring function was not provided and you have access to the buffer, how would you check to see if a valid # was entered? A detailed description is sufficient. You do not need to implement this in your code.

A: Since we have direct access to the buffer, we can read one byte of the input at a time, and compare it to the ASCII range for the characters '0' - '9', which happens to be the hex values 30-39. This would be done in a similar fashion to how we checked if the number entered was between 3-15 when validating the input for numentries. If any of the bytes is not within the ASCII range '0' - '9', then we know that the whole number entered is not valid. i.e. if the byte entered is greater than 39 or less than 30, the character is invalid, which makes the whole entry invalid. Example:

- ASCII(9) = 0x39 -valid
- ASCII(h) = 0x68 -invalid
- ASCII(0) = 0x30 -valid
- ASCII(4) = 0x34 -valid

Therefore, the entry 409 would be valid, and 4h9 would be invalid.

5 Conclusion

This lab demonstrated how to perform operations and modify data while moving it around using the Assembly language for the ColdFire architechture. In addition, the lab improved our understanding of the debugger software, a very powerful tool in the development of this kind of code. The main issue we found was related to the hardware itself, as there was some instances where the code did not execute properly and the board itself needed to be reset. The other issue we faced was mostly around getting used to the software and the workflow in the Eclipse IDE and the debugger. Once we understood the ways to use the tools we found our workflow sped up considerably, as we were able to check step by step and find bugs at the source. The last issue we had was with the syntax of the code, but that was solved quickly by reading over documentation and with the help of the TAs. Overall the lab went smoothly and has indeed succeeded at the goals of improving our familiarity and skill with the Netburner ColdFire system, Assembly code and Pair Programming practices.

6 Appendix

6.1 Part A Assembler Code

```
/* DO NOT MODIFY THIS -----*/
.text
. global WelcomePrompt
.extern iprintf
.extern cr
.extern value
.extern getstring
/* General Information ******************************/
/* Names of Students: Arun Woosaree and Navras Kamal
/* Date: March 5, 2018
                                              **/
/* General Description:
                                              **/
/* Takes in values with prompts for the statistics calculator
/* getstr stores value in d0*/
/* iprintf pops last thing on stack and prints it */
/* allocate 44 bytes because we are using up to 11 registers*/
/* 11*44 = 44 \text{ bytes*}/
suba.l #44, %sp
/* back up register contents onto the stack*/
movem. 1 \%a2-\%a6/\%d2-\%d7, (%sp)
lea 0x2300000, %a2
pea WelcomeMessage
/*print the welcome message*/
jsr iprintf
jsr cr
/*clean up stack*/
addq.1 #4 , %sp
bra numentries
```

```
failnumentries:
pea INVALIDENTRY
jsr iprintf
jsr cr
/*clean up stack*/
\operatorname{addq.l}\ \#4 , \%\mathrm{sp}
bra numentries
numentries:
/*prompt number of entries*/
pea Plztellusnumentries
jsr iprintf
jsr cr
/*clean up stack*/
addq.\,l\ \#4,\ \%sp
jsr getstring
/*sanitize input*/
cmp. 1 #15, %d0
bgt failnumentries
cmp.1 #3, %d0
blt failnumentries
/*success, replace value on stack*/
move. 1 \%d0, 52(\%sp)
/*counter used for loop*/
move. \ l \ \%d0 \ , \ \%d7
/*print what user entered*/
move. 1 \% d0, -(\% sp)
jsr value
jsr cr
addq.1 #4, %sp
bra divisor
faildivisor:
pea INVALIDENTRY
jsr iprintf
jsr cr
```

```
/*clean up stack*/
addq.l #4 , %sp
divisor:
/*divisor*/
pea Plztellusdivisor
jsr iprintf
jsr cr
/*cleanup stack*/
addq.l #4, %sp
/*get the string*/
jsr getstring
/*sanitize input*/
cmp.1 #5, %d0
bgt faildivisor
cmp.l #2, %d0
blt faildivisor
/*success, replace value on stack*/
move. 1 \%d0, 48(\% sp)
/*print what user entered*/
move. 1 \%d0, -(\%sp)
jsr value
jsr cr
addq.l #4, %sp
bra loopvals
failloopvals:
pea INVALIDENTRY
jsr iprintf
jsr cr
/*clean up stack*/
addq.l #4 , %sp
loopvals:
/*get first n-1 numbers*/
pea Plzenternumber
jsr iprintf
jsr cr
/*clean up stack*/
```

```
addq.1 #4, %sp
/*get the string and push onto stack*/
jsr getstring
/*sanitize input*/
/*positive only*/
cmp.1 #1, %d0
blt failloopvals
/*success, move to memory location*/
move. 1 \% d0, (\% a2)+
sub.1 #1, %d7
/*print what user entered*/
move. 1 \% d0, -(\% sp)
jsr value
jsr cr
addq.1 #4, %sp
cmp. l #1, %d7
beq lastnum
bra loopvals
faillastnum:
pea INVALIDENTRY
jsr iprintf
jsr cr
/*clean up stack*/
\operatorname{addq.l}#4 , %sp
lastnum:
/*get the last number*/
pea Plzenterlastnum
jsr iprintf
jsr cr
/*clean up stack*/
addq.l #4, %sp
/*get the string and push onto stack*/
jsr getstring
/*sanitize input*/
cmp. l #1, %d0
```

```
blt faillastnum
/*success, move to memory location*/
move.1 \%d0, (\%a2)+
/*print what user entered*/
move. 1 \% d0, -(\% sp)
jsr value
jsr cr
addq.1 #4, %sp
/* restore values */
movem. 1 (%sp), \%a2-\%a6/\%d2-\%d7
add.l #44, %sp
/* return to original program */
rts
. data
WelcomeMessage:
.string "Welcome to our amazing statistics program"
Plztellusnumentries:
.string "Please enter the number (3min-15max) of entries followed by enter"
Plztellusdivisor:
.string "Please enter the divisor (2min-5max) followed by enter"
Plzenternumber:
.string "Please enter a number(positive only)"
Plzenterlastnum:
.string "Please enter the last number (positive only)"
INVALIDENTRY:
.string "Invalid entry, please enter proper value."
```

6.2 Part A Flowchart Diagram

CVE-2017-12794:
Possible XSS in
traceback section of
technical 500 debug
page

Fixed an issue where we'd send out debug notifications to all users

Miscellaneous bug fixes and improvements

Further improvements to overall system stability and other minor adjustments have been made to enhance this user experience for your general satisfaction

Increased version number from 3.9.1 to 3.9.2

We improved the fuck outta yo user experience



6.3 Part B Assembler Code

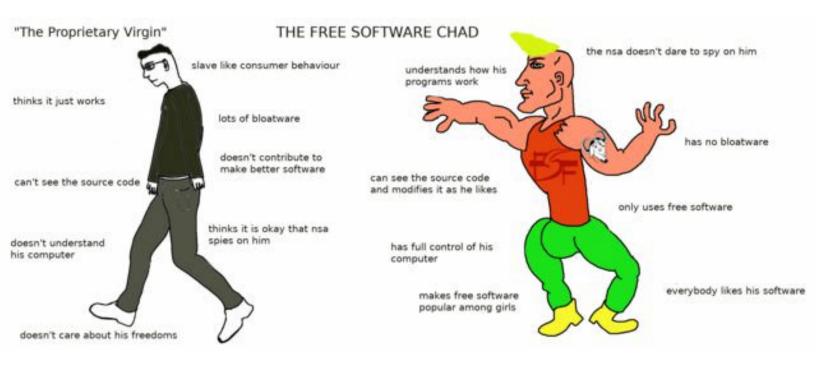
```
/* DO NOT MODIFY THIS ———
. text
. global Stats
.extern iprintf
.extern cr
.extern value
.extern getstring
/* Names of Students: Arun Woosaree and Navras Kamal **/
/* Date: March 5, 2018
                                               **/
/* General Description:
                                               **/
/* Computes stats on the numbers entered (min mean max)
                                               **/
/* backup values */
suba. 1 #60, %sp
movem. 1 \%a0-\%a6/\%d0-\%d7, (%sp)
/*divisor is at 4sp
jsr value
jsr cr
*/
/*numentries is at 8sp*/
move. 1 68(\% \text{sp}), -(\% \text{sp})
jsr value
jsr cr
addq.l #4, %sp
move.l 64(\% \text{sp}), \% d0 /* divisor */
move. 1 68(%sp) ,%d1 /* numentries */
lea 0x2300000, %a2 /* numbers are here */
lea 0x2310000, %a3 /*divisible numbers stored here*/
move.1 %d1, %d7 /*counter*/
move.l (\%a2)+, \%d5 /*first number*/
findmin:
```

```
loopmin:
        move. 1 (\%a2)+, \%d6
        cmp.1~\%d5\,,~\%d6
                 bge nochangemin
                 move.1 %d6, %d5 /* new min */
        nochangemin:
        subq.l #1, %d7
        cmp.l #1, %d7
        bne loopmin
/* min is stored in d5*/
move. 1 \% d5, (\% a3) +
/*
move. 1 \% d5, -(\% sp)
jsr value
jsr cr
addq.l #4, %sp
*/
lea 0x2300000, \%a2 /* numbers are here */
move.l %d1, %d7 /*counter*/
move.l (\%a2)+, \%d6 /*first number*/
findmax:
        loopmax:
        move. 1 (\%a2) + , \%d4
        cmp.1 %d6, %d4
                 ble nochangemax
                 move.\,l~\%d4\,,~\%d6~/*~new~max~*/
        nochange max:
        subq.l #1, %d7
        cmp.1 #1, %d7
        bne loopmax
/*max is stored in d6*/
move. 1 \% d6, (\% a3) +
/*
move. 1 %d6, -(\%sp)
jsr value
jsr cr
addq.l #4, %sp
*/
findmean:
         clr.l%d4
```

```
lea 0x2300000, %a2 /* numbers are here */
        move.1 %d1, %d7 /*counter*/
        loopaddvalues:
        move. 1 (\%a2) + , \%d3
        add.l %d3, %d4
        subq.l #1, %d7
        cmp.1 \#0, \%d7
        bne loopaddvalues
        /* sum in d4*/
        /*divide d4 by numertries*/
        divs.1 %d1, %d4
        /* mean should be in d4*/
        move. 1 \% d4, (\% a3) +
/*move.1 \%d4, -(\%sp)
jsr value
jsr cr
addq.l #4, %sp
finddivisible:
        lea 0x2300000, %a2 /* numbers are here */
        /* at this point, d6, d5, d4 store values and therefore should not be touched*
        /* this is the last time using the numentries, so will modify d1 directly */
        /* ok so d7, d3, and d2 are left free to use*/
        clr.1 %d7
        loopdivisible:
        move. 1 (\%a2) + , \%d3
        /* copy d3*/
        move. 1 %d3, %d4
        /* check if it's divisible, divisor is d0*/
        divs.w %d0, %d3
        /* get remainder*/
        lsr.l #8, %d3
        lsr.1 #8, %d3
        /*if 0, divisible */
        bne notdivisible
        /*it is divisible, move copy*/
        move. 1 \% d4 , (\% a3) +
        /*increment counter for divisible numbers*/
        addq.l #1, %d7
```

```
not divisible:
    subq.l #1, %d1
    cmp.l\ \#0,\ \%d1
    bne loopdivisible
    /* d7 holds number of divisible numbers*/
/*move.l \%d7, -(\%sp)
jsr value
jsr cr
addq.l #4, %sp
*/
move. 1 %d7, 72(\% \text{sp})
movem.l (%sp), \%a0-\%a6/\%d0-\%d7
adda.\,l\ \#60\ ,\ \%sp
rts
```

6.4 Part B Flowchart Diagram



6.5 Part C Assembler Code

```
/* DO NOT MODIFY THIS —
. text
. global Display
.extern iprintf
.extern cr
.extern value
.extern getstring
/* Names of Students: Arun Woosaree and Navras Kamal
/* Date: March 16 2018
                                                **/
/* General Description:
                                                **/
/* Prints the results of calculation from the stats subroutine
                                                **/
/* backup values */
suba.l #60, %sp
movem. 1 \%a0-\%a6/\%d0-\%d7, (%sp)
lea 0x2300000, %a2 /* numbers are here */
lea 0x2310000, %a3 /*divisible numbers stored here*/
move.\,l~72(\%\,\mathrm{sp}\,)~,~\%d7~/*~\mathrm{divisor}\,*/
move.1 76(\%sp), \%d6/* numentries*/
move.l 80(\% \text{sp}), \%d5/*numdivisible*/
jsr cr
/* tell user numentries*/
pea numentries
jsr iprintf
move. 1 %d6, -(\% sp)
jsr value
jsr cr
jsr cr
adda.l #8, %sp
/* print numbers */
pea numbers
jsr iprintf
```

```
adda.1 #4, %sp
jsr cr
loopnumbers:
move. 1 (\%a2)+, -(\%sp)
jsr value
jsr cr
adda.l #4, %sp
subq.l #1, %d6
bne loopnumbers
jsr cr
/*min*/
pea min
jsr iprintf
move. l (\%a3) + , -(\%sp)
jsr value
jsr cr
jsr cr
adda.l #8, %sp
/*max*/
pea max
jsr iprintf
move. l (\%a3) + , -(\%sp)
jsr value
jsr cr
jsr cr
adda.l #8, %sp
/*mean*/
pea mean
jsr iprintf
move. l (\%a3) + , -(\%sp)
jsr value
jsr cr
jsr cr
adda.l #8, %sp
/*print num divisible*/
pea numdivisible
jsr iprintf
move. 1 %d5, -(\% sp)
jsr value
pea numdivisible2
jsr iprintf
```

```
move. 1 \% d7, -(\% sp)
jsr value
jsr cr
adda.l #16, %sp
loopdivisible:
move. 1 (\%a3)+, -(\%sp)
jsr value
jsr cr
adda.l #4, %sp
subq.l #1, %d5
bne loopdivisible
/*restore values*/
movem. 1 (%sp), \%a0-\%a6/\%d0-\%d7
adda.1 #60 , %sp
rts
. data
numentries:
.string "The number of entries was: "
.string "The entered number(s) were: "
.string "Min number: "
max:
.string "Max number: "
mean:
.string "Mean number: "
numdivisible:
.string "There are "
numdivisible2:
.string "number(s) divisible by "
endprogram:
.string "End of program"
```

6.6 Part C Flowchart Diagram



8 Marking Sheet