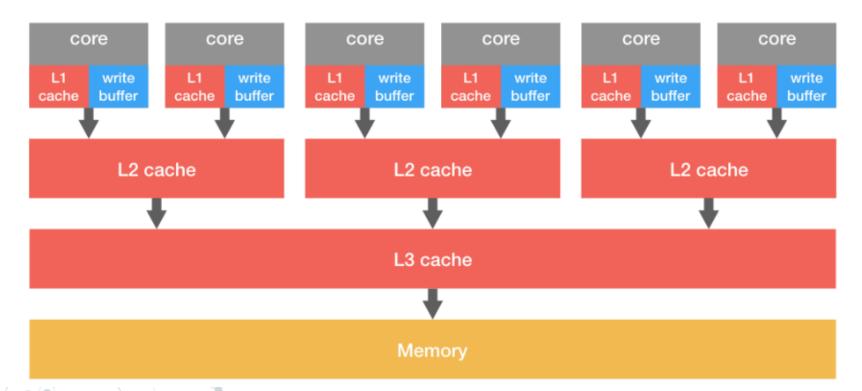


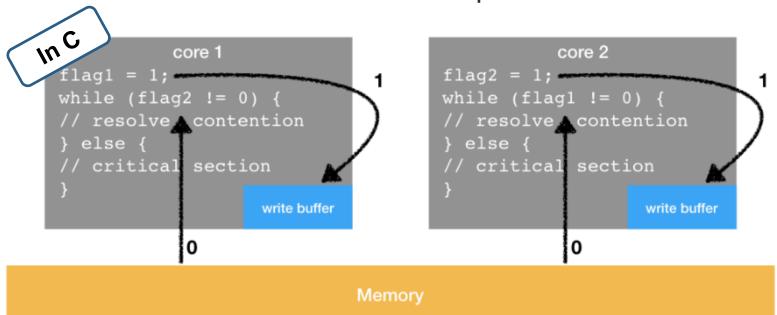
## Concurrency



- The hardware, the OS, and the compilers are too complex.
- Since 1970 processor cores do not work with the memory directly and instead use a complex hierarchy of caches and write buffers.



- Let us remove all the cache and consider the write buffers only.
- Write buffers are absolutely essential for the performance of the processor since writes to the memory are expensive and we want to batch them as much as possible.



 Another way of thinking about the problem is to think that the write buffer has reordered the order of operations by executing

flag2 != 0 before flag1 = 1.

- Similarly, we can think that caches also reorder the operations.
- Operations are reordered by the compiler that does optimizations.
- As a result, the order of operations in the source code will be different from the order executed by a specific core.
- The same code can have a different order of operations when executed in parallel on two separate cores.

- The order of operations would not be a problem if we were not using threads running on different cores to collaborate with each other.
- Collaborating threads require us to argue that operation X
  happened on thread A before operation Y on thread B
- Multithreading requires us to be able to talk about causality between operations across threads.

Without special tools that wouldn't be possible!

# **Threads**



```
use std::thread::spawn;
fn main() {
        spawn(||{
            for i in 0..10{
                println!("{}",i);
            }
        });
        for i in 10..20{
            println!("{}",i);
        }
        println!("Done");
}
```

- Spawn creates the new thread.
- It is a function which takes a function as a parameter.
- We can write normal functions or use closures.

```
//output:
10
11
12
13
14
15
16
17
18
19
Done
//possible to see some output from the other thread, but
completely random
```

• The program has only run the latter part of that. We've not seen the print from the threads because when the program reaches its end, all other threads are killed.

• If we don't want that to happen then we can call **h.join**, and that will wait for the first thread to finish before completing.

```
use std::thread::spawn;
fn main() {
    let h=spawn(||{
        for i in 0..10{
            println!("{}",i);
     });
     for i in 10..20{
        println!("{}",i);
     h.join().expect("Error");
     println!("Done");
```

```
//output:
10.. 19
0..9
Done
```



• Because **h.join** returns a result, we can pull information out of that result. What if our function for exam was to return 5?

```
use std::thread::spawn;
fn main() {
    let h=spawn(||{
        for i in 0..10{
            println!("{}",i);
         return 5;
     });
     for i in 10..20{
        println!("{}",i);
     let r = h.join().unwrap();
     println!("Done R= {}",r);
```

```
//output:
10..19
0..9
Done R= 5
```



- Now, every time we've run this so far, we've always finished the one thread before the other really got to get started.
- So, let's perhaps make the one run a little bit slower:

```
use std::thread::{spawn,sleep};
use std::time::Duration;
fn main(){
    let h=spawn(||{
        for i in 0..10{
            sleep(Duration::from millis(10));
            println!("{}",i);
         }return 5;
     });
     for i in 10..20{
        sleep(Duration::from millis(10));
        println!("{}",i);
     let r = h.join().unwrap();
     println!("Done R= {}",r);
```

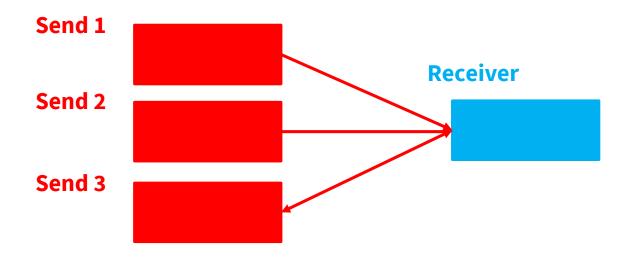
 When we run the program, hopefully we'll find that the two sets of numbers interlace as they both wait and take it in turns.

```
//output:
0
10
1
11
12
2
3
13
4
14
```

```
15
5
16
6
17
7
18
8
19
9
Done R= 5
```



• <u>multiple-producer</u>, <u>single-consumer</u> channel (mpsc).





Channels allow us to pass data between threads safely.

```
use std::sync::mpsc::{Sender, Receiver};
use std::thread;
static NTHREADS: i32 = 3;
fn main() {
    let (tx, rx): (Sender< >, Receiver< >) = mpsc::channel();
    let mut children = Vec::new();
    for id in 0..NTHREADS {
        let thread tx = tx.clone();
        let child = thread::spawn(move || {
               thread tx.send(id).unwrap();
               println!("thread {} finished", id);
        });
        children.push(child);
    let mut ids = Vec::with capacity(NTHREADS as usize);
    for in 0..NTHREADS {ids.push(rx.recv());}
    for child in children {child.join().expect("oops! the
child thread panicked");}
    println!("{:?}", ids);
```

```
Compiling
Finished
Running
thread 0 finished
thread 2 finished
thread 1 finished
[Ok(0), Ok(2), Ok(1)]
```



```
fn main() {
    // Channels have two endpoints: the `Sender<T>` and the
`Receiver<T>`,
    // where `T` is the type of the message to be transferred
    // (type annotation is superfluous)

let (tx, rx): (Sender<_>, Receiver<_>) = mpsc::channel();

let mut children = Vec::new();
```



```
for id in 0..NTHREADS {
        // The sender endpoint can be copied
        let thread tx = tx.clone();
        // Each thread will send its id via the channel
        let child = thread::spawn(move || {
            // The thread takes ownership over `thread tx`
            // Each thread queues a message in the channel
            thread tx.send(id).unwrap();
            // Sending is a non-blocking operation, the thread
will continue
            // immediately after sending its message
            println!("thread {} finished", id);
        });
        children.push(child);
```

```
// Back in the main thread
// Here, all the messages are collected
    let mut ids = Vec::with capacity(NTHREADS as usize);
    for in 0..NTHREADS {
        // The `recv` method picks a message from the channel
        // `recv` will block the current thread if there are
no messages available
        ids.push(rx.recv());
```

```
// Wait for the send threads to complete any remaining work

for child in children {
   child.join().expect("oops! the child thread panicked");
 }

// Show the order in which the messages were sent,
remember, receiving is in one thread.

println!("{:?}", ids);
}
```

