

DSA A2 Q4 Explanation
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First we scan "n" which denotes the number of nodes and then declare an array `lvOrder` and scan the elements in them.

Then we create the Root Node (T) and enqueue it in a Queue Q and then pass it the function `superInsert` which has arguments T, `lvOrder` array, n (number of elements), Q.

Then we start a while loop till we reached the end of the `lvOrder` array (ie till all elements are inserted). Then we start inserting each element in the following manner :-

1. If the element can be inserted left of the node in the front of the Q we insert it to the left of it and then enqueue the node in the queue Q. Then we increment the index for the `lvOrder` array.
2. If the element can be inserted right of the node in the front of the Q we insert it to the right of it and then enqueue the node in the queue Q and dequeue the Q because the node in the front can no longer have any nodes inserted left/right of it because top-down order has been provided as input. Then we increment the index for the `lvOrder` array.
3. If the element cannot be inserted to left/right of the node in the front of the queue we dequeue the node from Q.

The time complexity of the function `superInsert` is $O(n)$ because each node will be enqueued once in the Q and also dequeued after either it has had a no children or it has had a right child or if it has had both right and left children.

To figure that if a node can go under a node I have first modified the struct node to include a `greaterThan` and `lessThan` value which denotes the range of keys which can go under it. When a node has been added to the left of the parent node the `greaterThan` value is retained from the parent and the `lessThan` value becomes key of the parent. Similarly when the node has been inserted to the right.

The time complexity of the function `canItGoUnderIt` is $O(1)$.

Then the function `superInorder` is called which uses `inorder` to travel the BST in ascending order and instead of printing the node updates its value to `value+carrySum` and then updates the `carrySum` to the new value of the node. `carrySum` is a long long int declared in main and passed by reference to the `superInorder` function.

The time complexity of the function `superInorder` is $O(n)$ as every node is visited once.

The function `LvlOrderWithQueue` is called which performs level order traversals using queues as discussed in lectures and tutorials in addition to it, it also adds the value of the node to the `totalSum` variable which is passed by reference to the function to calculate sum of all nodes in the BST.

The time complexity of the function `LvlOrderWithQueue` is $O(n)$ as every node is visited once.