

Unit 4: Relational Data Model

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OUTLINE

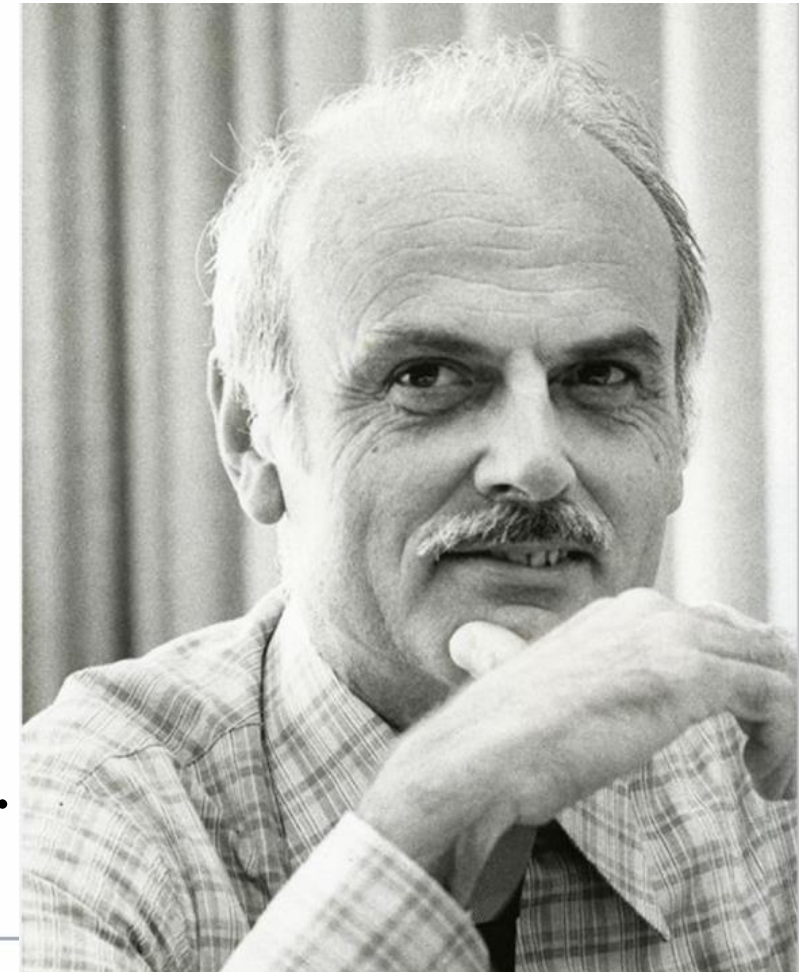
- Relational Data Model
- Constraints & Keys
- Relational Algebra Operations

Relational Data Model

The relational data model describes the world as
“a collection of inter-related relations (or tables).”

The relational data model was introduced by
Edgar F. Codd in **1970**, while working for IBM.

✓ Currently, it is the most widely used data model.



Relational Data Model

The relational data model has provided the basis for:

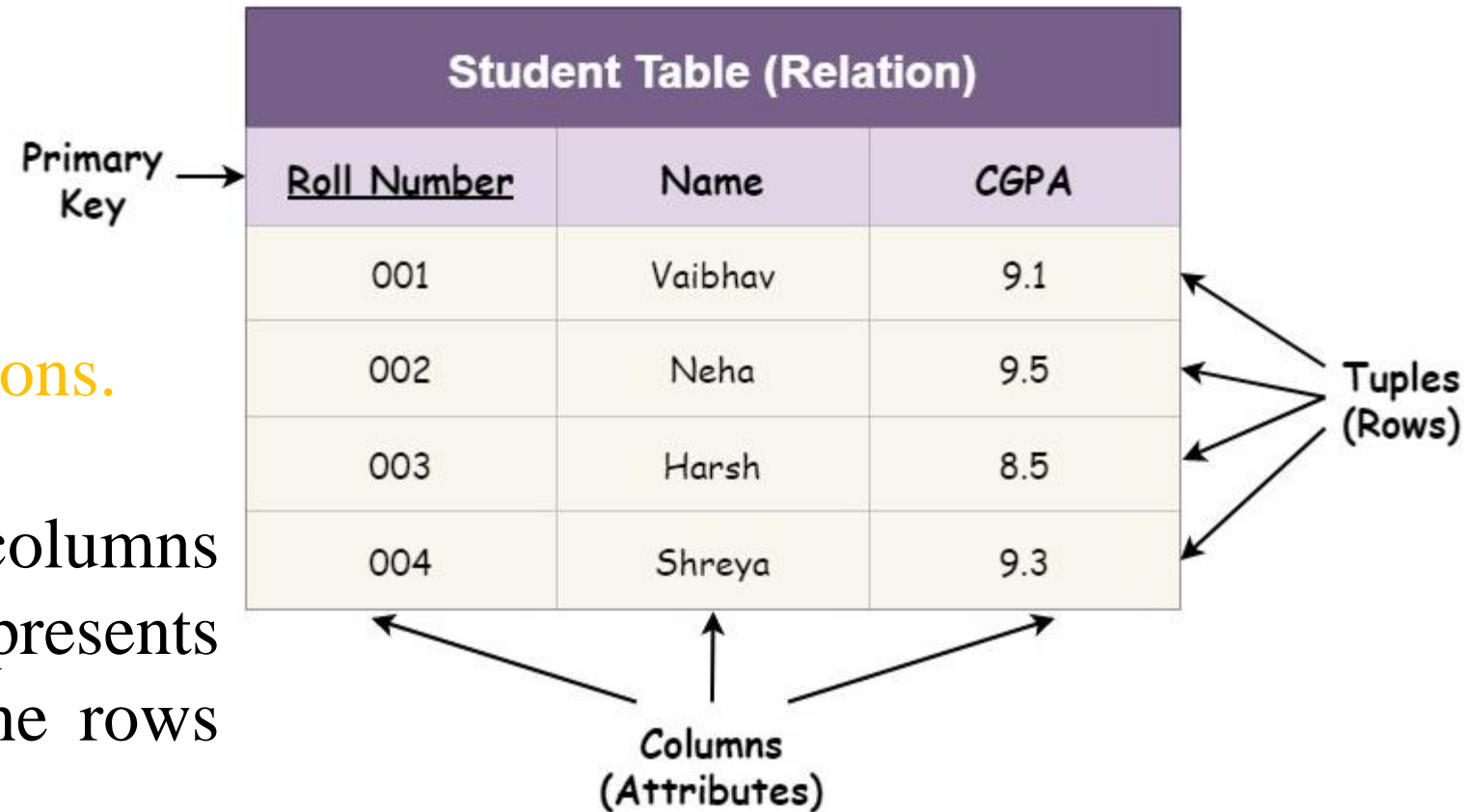
- Research on the theory of data/relationship/constraint
- Numerous database design methodologies
- The standard database access language called *structured query language (SQL)*
- Almost all modern commercial database management systems.

Relational Data Model

The relational model represents how data is stored in Relational Databases.

➤ Tables are also known as relations.

Each relation is a collection of columns and rows, where the column represents the attributes of an entity and the rows (or tuples) represent the records.



Relational Data Model

Attribute: Attributes are the properties that define an entity.

e.g.; ROLL_NO, NAME, ADDRESS

Tuple: Each row in the relation is known as a tuple.

ROLL_NO	NAME	ADDRESS	PHONE	AGE
1	RAM	DELHI	9455123451	18
2	RAMESH	GURGAON	9652431543	18
3	SUJIT	ROHTAK	9156253131	20
4	SURESH	DELHI		18

This relation contains 4 tuples.

Relational Data Model

Column: The column represents the set of values for a particular attribute.

Null Values: The value which is not known or unavailable. It is represented by blank space.

Degree: The number of attributes in the relation.

The STUDENT relation defined above has degree 5.

Cardinality: The number of tuples in a relation.

The STUDENT relation defined above has cardinality 4.

Relational Data Model

Advantages of the Relational Model

- Simple model:** It is simple and easy to use in comparison to other languages.
- Flexible:** It is more flexible than any other relational model present.
- Secure:** It is more secure than any other relational model.
- Data Accuracy:** Data is more accurate in the relational data model.
- Data Integrity:** The integrity of the data is maintained in the relational model.

Disadvantages of the Relational Model

- Relational Database Model is not very good for large databases.
- Sometimes, it becomes difficult to find the relation between tables.
- Because of the complex structure, the response time for queries is high.

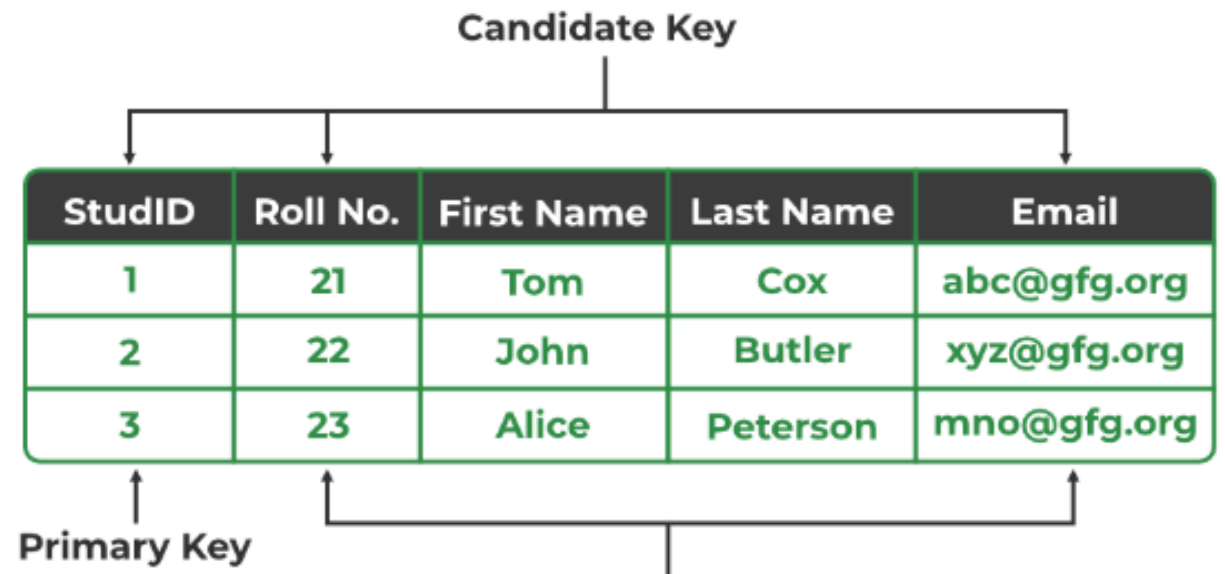
Relation Key

The keys that are used to identify the rows uniquely and also helps in identifying tables.

- Primary Key
- Foreign Key
- Super Key
- Candidate Key

Relation Key : Primary Key

A primary key is a column or a set of columns in a table whose values uniquely identify a row in the table.



Relation Key : Foreign Key

If an attribute can only take the values which are present as values of some other attribute, it will be a foreign key.

- It acts as a primary key in one table and it acts as secondary key in another table.
- It combines two or more relations (tables) at a time.
- They act as a cross-reference between the tables.

Primary Key

ID	Name	Course
2041	Tom	Java
2204	John	C++
2043	Alice	Python
2032	Bob	Oracle

Student Details

Foreign Key

ID	Marks
2041	65
2204	55
2043	73
2032	62

Student Marks

Relation Key : Super Key

The set of attributes that can uniquely identify a tuple is known as Super Key.

Example:

1. STUD_NO + STUD_NAME
2. STUD_NO + PHONE

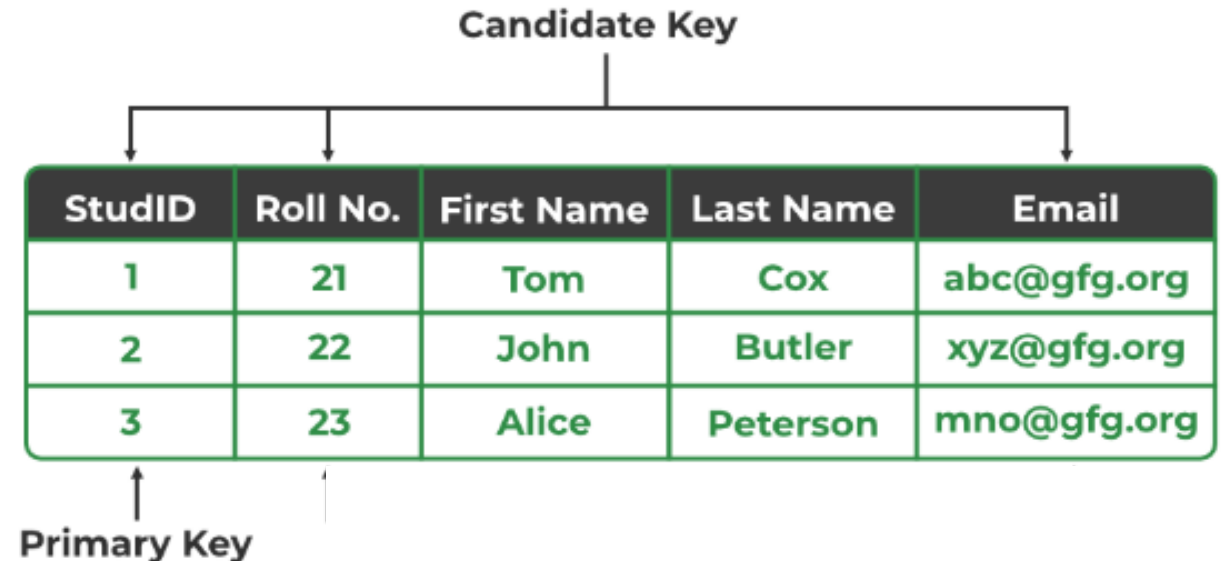
A super key is a group of single or multiple keys that identifies rows in a table.

STUD_NO	SNAME	ADDRESS	PHONE
1	Shyam	Delhi	123456789
2	Rakesh	Kolkata	223365796
3	Suraj	Delhi	175468965

Relation Key : Candidate Key

The minimal set of attributes that can uniquely identify a tuple is known as a candidate key.

Example: Stud ID, Roll No, Email



Not Null Constraint

The constraint is typically used for **columns** that must contain essential information.

Not Null Constraint: It ensures that a specific column in a table cannot contain null values.

Syntax: `CREATE TABLE table_Name(column1 data_type(size) NOT NULL);`

Example: `CREATE TABLE Doctor(Doctor_ID INT NOT NULL,
Doctor_Name VARCHAR(100) NOT NULL);`

- ✓ When you insert or update data in a column with “not null constraint”, you are required to provide a valid value, and leaving it empty is not allowed.

Check Constraint

It allows you to ensure a value meets a certain condition before it is inserted into a table.

Syntax: `CREATE TABLE table_name (column1 datatype CONSTRAINT constraint_name CHECK (expression));`

Example: `CREATE TABLE employees (name varchar(25) CONSTRAINT age_check CHECK (age > 18));`

Relational Algebra Operations

It is used to perform different operations in relational databases.

It consists of a set of operations enabling users to manipulate and handle data stored in relational databases.

These are similar to the operations of set theory.

for example- union intersection and difference

And some personalized operations for relational databases.

Relational Algebra Operations

Types of Relational Algebra Operations:

1. Basic

Selection (σ)

Projection (π)

Rename (ρ)

Cross Product(X)

Set operators (**Union,**
Intersection, Set Difference)

2. Derived/advanced

Join (\bowtie) (Natural & Outer Join)

Aggregate Functions

Relational Algebra Operations

Here are two tables/relations which will be used for demonstrating the relational algebra operations examples:

1. STUDENT (ROLL, NAME, AGE)

ROLL	NAME	AGE
1	Aman	20
2	Atul	18
3	Baljeet	19
4	Harsh	20
5	Prateek	21
6	Prateek	23

2. EMPLOYEE (EMPLOYEE_NO, NAME, AGE)

EMPLOYEE_NO	NAME	AGE
E-1	Anant	20
E-2	Ashish	23
E-3	Baljeet	25
E-4	Harsh	20
E-5	Pranav	22

Selection (σ)

Selection Operator which is represented by "sigma"(σ).

It is used to retrieve tuples(rows) from the table where the given condition is satisfied.

Example: Suppose we want the row(s) from STUDENT Relation where "AGE" is 20

σ AGE=20 (STUDENT)

ROLL	NAME	AGE
1	Aman	20
4	Harsh	20

Projection (Π)

It pulls out some specific columns (attributes) from a relation.

It is represented by "pi"(Π).

Example: Suppose we want the names of all students from STUDENT Relation.

Π NAME(STUDENT)

NAME
Aman
Atul
Baljeet
Harsh
Prateek

Rename (ρ)

It is denoted by "Rho"(ρ).

It is used to rename the output relation.

Example: If we want to rename the STUDENT relation as STUDENT_NAME

$\rho(\text{STUDENT_NAME}, \Pi \text{ NAME}(\text{STUDENT}))$

Cross Product(X)

It is denoted by symbol ‘X’.

It is used to perform operation on two relations.

Example: Consider two relations STUDENT(SNO, FNAME, LNAME) and
DETAIL(ROLLNO, AGE) below:

S. No.	FNAME	LNAME
1	Albert	Singh
2	Nora	Fatehi

ROLLNO	AGE
5	18
9	21

Cross Product(X)

On applying CROSS PRODUCT on STUDENT and DETAIL:

STUDENT × DETAILS

SNO	FNAME	LNAME	ROLLNO	AGE
1	Albert	Singh	5	18
1	Albert	Singh	9	21
2	Nora	Fatehi	5	18
2	Nora	Fatehi	9	21

Set operators: Union

Union Operator which is represented by "union"(\cup), same as the union operator from set theory.

It selects all tuples/rows from both relations/tables but with the exception that; both relations/tables must have the same set of Attributes.

Example: Suppose we want all the names from STUDENT and EMPLOYEE relation.

$\Pi \text{ NAME}(\text{STUDENT}) \cup \Pi \text{ NAME}(\text{EMPLOYEE})$

NAME
Aman
Anant
Ashish
Atul
Baljeet
Harsh
Pranav
Prateek

Set operators: Intersection & Set Difference

Intersection is represented by (\cap) .

It selects all the tuples which are present in both relations.

Set difference gives the difference between two relations.

It is denoted by a "Hyphen"(-) and it returns all the tuples(rows) which are in relation R but not in relation S.

Set operators: Intersection

Example: If we want the names which are present in STUDENT as well as in EMPLOYEE relation.

$\Pi \text{NAME}(\text{STUDENT}) \cap \Pi \text{NAME}(\text{EMPLOYEE})$

NAME
Baljeet
Harsh

Set operators: Set Difference

Example: To know the names of students who are in STUDENT Relation but not in EMPLOYEE Relation.

$\Pi \text{NAME}(\text{STUDENT}) - \Pi \text{NAME}(\text{EMPLOYEE})$

NAME
Aman
Atul
Prateek

Join (⋈)

A join is an operation that combines the rows/tuple of two or more tables based on common attribute/columns.

The main purpose of Join is to retrieve the data from multiple tables.

It is denoted by ⋈.

Frequently used JOIN types:

1. Inner Join

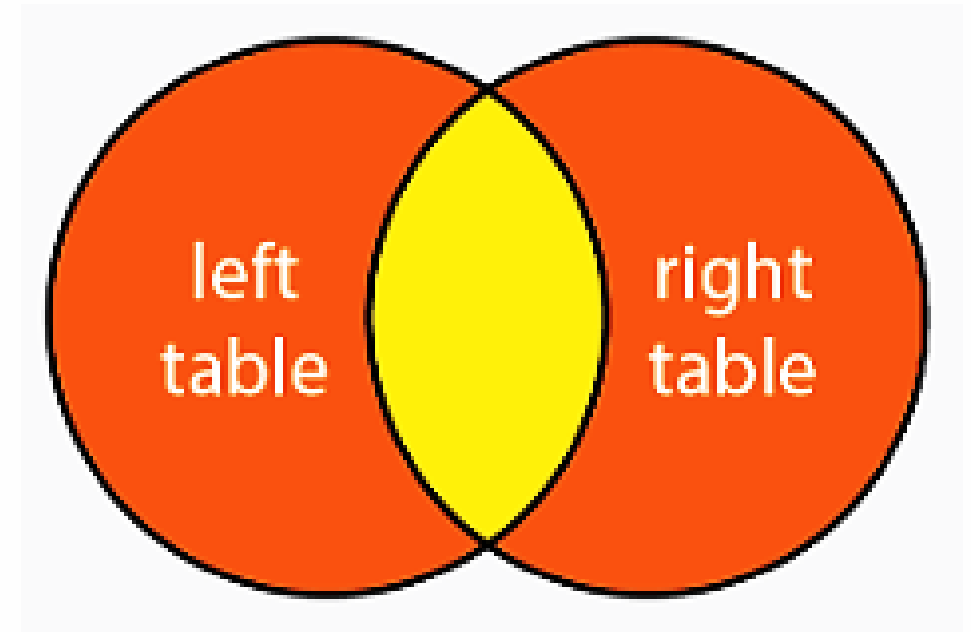
2. Outer Join

Inner Join

Inner Join is a join operation that combines two or more tables based on related columns.

It returns only rows that have matching values among tables.

- Equi Join
- Natural Join



➤ **Natural Join** joins two tables based on the same attribute name and datatypes.

Natural Join

Department

DEPT_NAME	MANAGER_NAME
IT	ROHAN
SALES	RAHUL
HR	TANMAY
FINANCE	ASHISH
MARKETING	SAMAY

Employee



EMP_ID	EMP_NAME	DEPT_NAME
1	SUMIT	HR
2	JOEL	IT
3	BISWA	MARKETING
4	VAIBHAV	PU
5	SAGAR	SALES

Output:

SELECT * FROM department
NATURAL JOIN employee;

DEPT_NAME	MANAGER_NAME	EMP_ID	EMP_NAME
HR	TANMAY	1	SUMIT
IT	ROHAN	2	JOEL
MARKETING	SAMAY	3	BISWA
SALES	RAHUL	5	SAGAR

Outer Join

Outer joins retrieves matching as well as non-matching records from related tables.

Types:

- | | | | |
|---------------------|---|------------|------------------|
| 1. Left outer join | / | Left Join | $\Join\lrcorner$ |
| 2. Right outer join | / | Right Join | $\lrcorner\Join$ |
| 3. Full outer join | / | Full Join | $\Join\ltimes$ |

Left Join

It retrieves all records from the left table and retrieves matching records from the right table.

Example: Table A

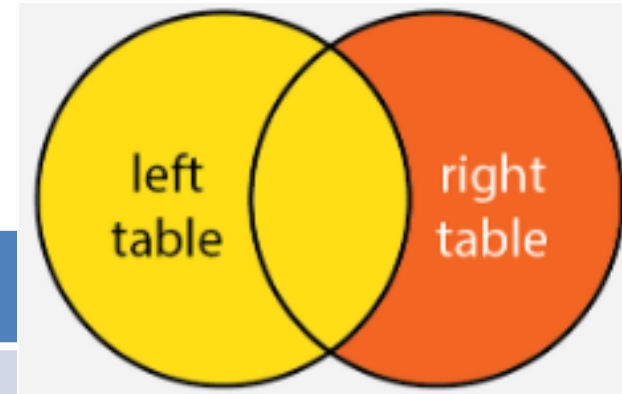
Number	Square
2	4
3	9
4	16

Table B

Number	Cube
2	8
3	27
5	125

$A \bowtie B$

Number	Square	Cube
2	4	8
3	9	27
4	16	-



Left Join

Department

DEPT_NAME	MANAGER_NAME
IT	ROHAN
SALES	RAHUL
HR	TANMAY
FINANCE	ASHISH
MARKETING	SAMAY

Employee



EMP_ID	EMP_NAME	DEPT_NAME
1	SUMIT	HR
2	JOEL	IT
3	BISWA	MARKETING
4	VAIBHAV	PU
5	SAGAR	SALES

Output:

```
SELECT * FROM department  
LEFT OUTER JOIN employee ON  
department.DEPT_NAME =  
employee.DEPT_NAME;
```

DEPT_NAME	MANAGER_NAME	EMP_ID	EMP_NAME	DEPT_NAME
HR	TANMAY	1	SUMIT	HR
IT	ROHAN	2	JOEL	IT
MARKETING	SAMAY	3	BISWA	MARKETING
SALES	RAHUL	5	SAGAR	SALES
FINANCE	ASHISH	-	-	-

Right Join

It retrieves all records from the right table and retrieves matching records from the left table.

Example: Table A

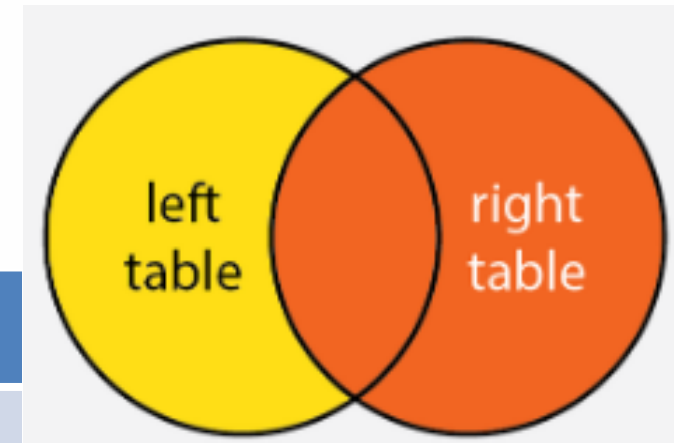
Number	Square
2	4
3	9
4	16

Table B

Number	Cube
2	8
3	27
5	125

$A \bowtie B$

Number	Square	Cube
2	4	8
3	9	27
5	-	125



Right Join

Department

DEPT_NAME	MANAGER_NAME
IT	ROHAN
SALES	RAHUL
HR	TANMAY
FINANCE	ASHISH
MARKETING	SAMAY

Employee

EMP_ID	EMP_NAME	DEPT_NAME
1	SUMIT	HR
2	JOEL	IT
3	BISWA	MARKETING
4	VAIBHAV	PU
5	SAGAR	SALES

Output:

```
SELECT * FROM department  
RIGHT OUTER JOIN employee ON  
department.DEPT_NAME =  
employee.DEPT_NAME;
```

DEPT_NAME	MANAGER_NAME	EMP_ID	EMP_NAME	DEPT_NAME
IT	ROHAN	2	JOEL	IT
SALES	RAHUL	5	SAGAR	SALES
HR	TANMAY	1	SUMIT	HR
MARKETING	SAMAY	3	BISWA	MARKETING
-	-	4	VAIBHAV	PU

Full Join

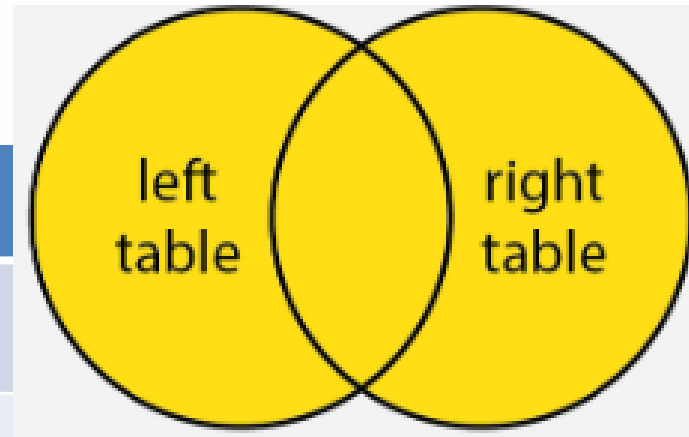
It combines the results of both LEFT JOIN and RIGHT JOIN. The result set will contain all the rows from both tables.

Example: Table A

Table B

$A \bowtie B$

Number	Square	Number	Cube	Number	Square	Cube
2	4	2	8	2	4	8
3	9	3	27	3	9	27
4	16	5	125	4	16	NULL
				5	NULL	125



Full Join

Department

DEPT_NAME	MANAGER_NAME
IT	ROHAN
SALES	RAHUL
HR	TANMAY
FINANCE	ASHISH
MARKETING	SAMAY

Employee

EMP_ID	EMP_NAME	DEPT_NAME
1	SUMIT	HR
2	JOEL	IT
3	BISWA	MARKETING
4	VAIBHAV	PU
5	SAGAR	SALES



Output:

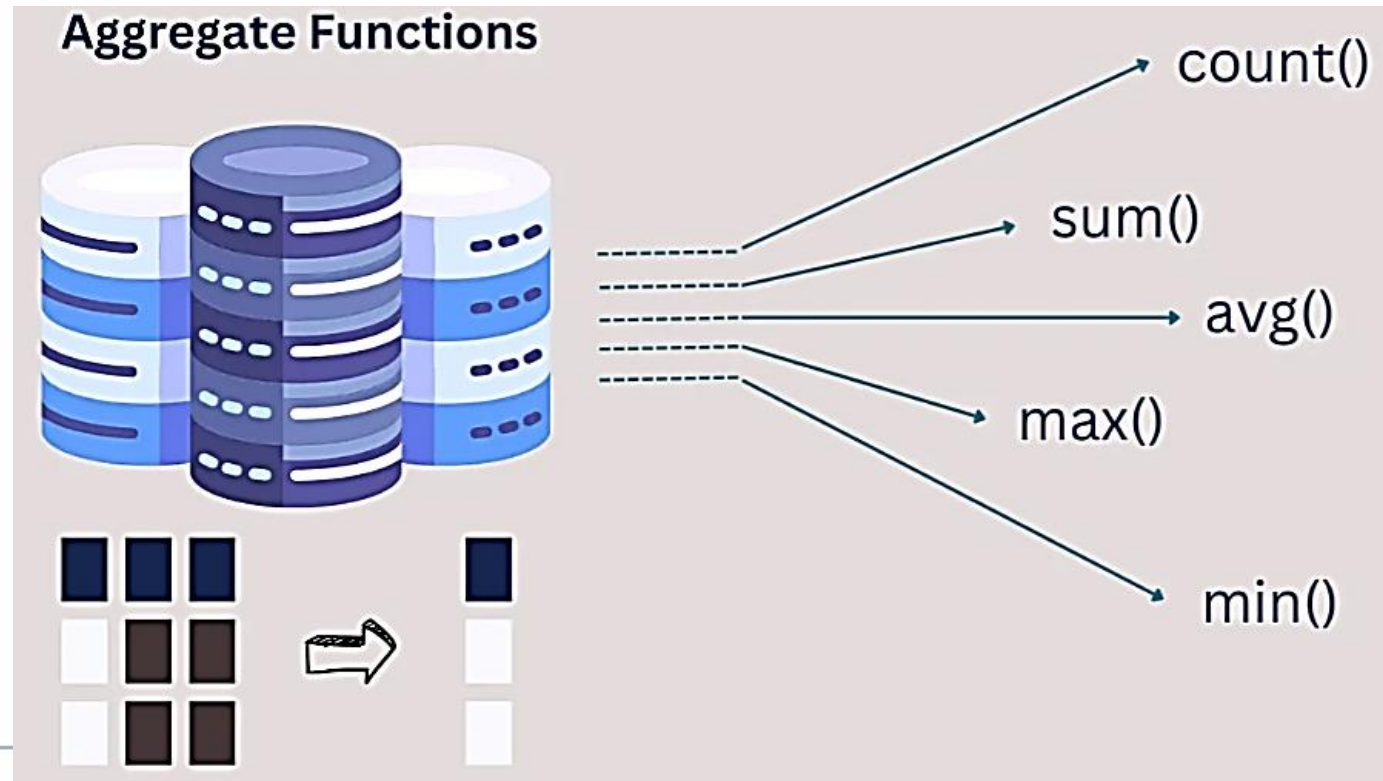
```
SELECT * FROM department
FULL OUTER JOIN employee ON
department.DEPT_NAME =
employee.DEPT_NAME;
```

DEPT_NAME	MANAGER_NAME	EMP_ID	EMP_NAME	DEPT_NAME
HR	TANMAY	1	SUMIT	HR
IT	ROHAN	2	JOEL	IT
MARKETING	SAMAY	3	BISWA	MARKETING
-	-	4	VAIBHAV	PU
SALES	RAHUL	5	SAGAR	SALES
FINANCE	ASHISH	-	-	-

Aggregate Function

Aggregation function is used to perform the calculations on multiple rows of a single column of a table.

It returns a single value.



Aggregate Function

```
CREATE TABLE Employee (Id INT, Name CHAR(1),  
Salary DECIMAL(10));  
INSERT INTO Employee (Id, Name, Salary)  
VALUES      (1, 'A', 802),  
            (2, 'B', 403),  
            (3, 'C', 604),  
            (4, 'D', 705),  
            (5, 'E', 606),  
            (6, 'F', NULL);
```

Id	Name	Salary
1	A	802
2	B	403
3	C	604
4	D	705
5	E	606
6	F	NULL

Aggregate Function

Aggregate Function Example:

--Count the number of employees

```
SELECT COUNT(*) AS Total Employees FROM Employee;
```

-- Calculate the total salary

```
SELECT SUM(Salary) AS Total Salary FROM Employee;
```

-- Find the average salary

```
SELECT AVG(Salary) AS Average Salary FROM Employee;
```

-- Get the highest salary

```
SELECT MAX(Salary) AS Highest Salary FROM Employee;
```

-- Determine the lowest salary

```
SELECT MIN(Salary) AS Lowest Salary FROM Employee;
```

Total Employees: 6

Total Salary: 3120

Average Salary: 624

Highest Salary: 802

Lowest Salary: 403

Thank You!!!

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